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Christopher Zúñiga, C28730

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Problem

We want to transfer a 10 Mb file from "Host 1" to "Host 2," through 3 intermediate routers (so there are 4 links in total). We compare the total time (from the very start of sending until the last bit arrives at Host 2) under:

- Circuit switching
- Message switching
- Datagram switching

The following parameters are given:

- File size: 10 Mb (megabits).
- Number of routers: 3 (hence 4 links total).
- Distance between Host 1 and Host 2: 2000 km (each of the 4 links is 500 km).
- Propagation speed: 200,000 km/s.
- Transmission bandwidth: 1 Mbps (i.e. 1 Mb/s).
- Node (router) processing delay: 100 ms per router for each arriving unit (message or packet).
- Datagram size: 75 kb (assume 1 kb = 1000 bits, so 75 kb = 75000 bits).

1 Key Quantities

Per-Link Propagation Delay

Each link is $\frac{2000}{4} = 500 \,\mathrm{km}$, and the propagation speed is $200,000 \,\mathrm{km/s}$. Hence the one-way propagation delay per link is:

$$t_{\text{prop}} = \frac{500 \,\text{km}}{200,000 \,\text{km/s}} = 0.0025 \,\text{s} = 2.5 \,\text{ms}.$$

Transmission Times

For the entire 10 Mb file at 1 Mbps:

$$t_{\text{tx(file)}} = \frac{10 \,\text{Mb}}{1 \,\text{Mb/s}} = 10 \,\text{seconds}.$$

For a single 75 kb datagram at 1 Mbps:

$$t_{\text{tx(datagram)}} = \frac{75,000 \,\text{bits}}{1,000,000 \,\text{bits/s}} = 0.075 \,\text{s} = 75 \,\text{ms}.$$

Router Processing Delay

Each router introduces a 100 ms processing time for a *message* or *datagram* after it fully arrives but before forwarding.

2 Circuit Switching

Basic assumptions:

• We first perform an end-to-end circuit setup, incurring:

$$(4 \text{ links of propagation}) + (3 \text{ routers' processing}).$$

If we assume the setup message is short and bandwidth is high enough, the main times are $4 \times t_{\text{prop}} + 3 \times 100 \,\text{ms}$.

- Once the circuit is established, data flows as a continuous stream (no store-and-forward at intermediate routers).
- Finally, the last bit traveling across the 4 links experiences (roughly) $4 \times t_{\text{prop}}$ of propagation once it is sent.

Setup Time

$$t_{\rm setup} = (4 \times t_{\rm prop}) + (3 \times 0.1 \,\mathrm{s}) = 4 \times 0.0025 + 3 \times 0.1 = 0.01 + 0.3 = 0.31 \,\mathrm{s}.$$

Data Transmission + Final Propagation

Once the circuit is up, we send 10 Mb at 1 Mbps over 4 links (pipelined). The *time to pump* out the file is 10 s. The last bit, after being transmitted, still must propagate to the receiver over $4 \times t_{\text{prop}} = 0.01$ s. Hence,

$$t_{\text{data}} + t_{\text{last_bit_prop}} = 10 \,\text{s} + 0.01 \,\text{s} = 10.01 \,\text{s}.$$

Total (Circuit Switching)

$$t_{\text{total(circuit)}} = t_{\text{setup}} + t_{\text{data}} + t_{\text{last_bit_prop}} = 0.31 + 10.01 = 10.32 \,\text{s}.$$

3 Message Switching

In message switching, the entire file (treated as one large message) must be fully transmitted to each router, stored (and processed for $100 \,\mathrm{ms}$), then forwarded to the next router, and so on. We have 3 routers = 4 hops.

- 1. At hop 1: Send 10 Mb at 1 Mbps \Rightarrow 10 s, plus 2.5 ms propagation. Once it arrives, router processes for 100 ms.
- 2. At hop 2: Same 10s send, 2.5 ms prop, 100 ms process.
- 3. **At hop 3**: Same ...
- 4. At hop 4: Same 10 s send, 2.5 ms prop, but no extra router processing since the destination is a host.

Putting this together:

Time per hop =
$$10 s + 0.0025 s \approx 10.0025 s$$
.

We do that 4 times (4 links), but we have 3 router-processing delays in between:

$$t_{\text{total(message)}} = (4 \times 10.0025) \,\text{s} + (3 \times 0.1) \,\text{s}.$$

Numerically,

$$4 \times 10.0025 = 40.01 \,\mathrm{s}, \quad 3 \times 0.1 = 0.3 \,\mathrm{s}.$$

 $t_{\text{total(message)}} = 40.01 + 0.3 = 40.31 \,\mathrm{s}.$

4 Datagram (Packet) Switching

Here, the file is split into many 75 kb (kilobit) datagrams (store-and-forward per datagram). We have:

File size = 10,000,000 bits, Datagram size =
$$75,000$$
 bits.
Number of datagrams = $\frac{10,000,000}{75,000}$ = $133.\overline{3} \approx 134$.

(We typically round up to 134 datagrams.)

One Datagram's End-to-End Delay

For a *single* datagram:

- Transmit time per link: 0.075 s.
- Propagation per link: 0.0025 s.
- Processing at each router: 0.1 s.

Each datagram goes over 4 links (thus 3 routers in between). The store-and-forward model says each router must fully receive the datagram before forwarding it, incurring a 100 ms process time at *each* intermediate router. For each of the **first 3 links**, we have (tx + prop + process). On the **4th link**, we have tx + prop but *no extra router process* at the end. Hence,

$$t_{1 \text{ datagram}} = 3 \times (0.075 + 0.0025 + 0.1) \text{ s} + (0.075 + 0.0025) \text{ s}.$$

= $3 \times 0.1775 + 0.0775 = 0.5325 + 0.0775 = 0.61 \text{ s}.$

Total Delay for 134134 Datagrams (Pipelined)

In datagram switching, the sending host transmits the 134 datagrams back to back. Each datagram has a $0.075 \,\mathrm{s}$ transmit time on the first link. So after the first datagram, the second starts $0.075 \,\mathrm{s}$ later, etc. Because of pipelining, while the nth datagram is traveling through the next stages, the (n+1)th datagram can be on the earlier link, etc.

A common formula for store-and-forward pipelines is:

$$t_{\text{total(datagram)}} = t_{1 \text{ datagram}} + (\text{NumberOfDatagrams} - 1) \times t_{\text{tx (onfirstlink)}}.$$

$$= 0.61 \text{ s} + (134 - 1) \times 0.075 \text{ s}.$$

$$= 0.61 + 133 \times 0.075 = 0.61 + 9.975 = 10.585 \text{ s}.$$

So once the first datagram has fully traversed all 4 links $(0.61 \,\mathrm{s})$, each subsequent datagram finishes about $0.075 \,\mathrm{s}$ after the previous one in the steady pipeline. By the time we have sent all 134, the total is about $10.59 \,\mathrm{s}$.

Final Numerical Results

• Circuit Switching:

$$t_{\rm total} \approx 10.32 \, {\rm s}$$

• Message Switching:

$$t_{\rm total} \approx 40.31 \, {\rm s}$$

• Datagram Switching:

$$t_{\rm total} \approx 10.59 \, \rm s$$