

Database Concepts: Basic concepts, Database Modeling, The Relational Data Model, SQL, Integrity and Security, Relational Database Design, Object-Oriented Databases, Distributed and Cloud, Databases, Big Data.

Unit 1: Introduction to Database Concepts

- What is a database? DBMS vs RDBMS
 - Components of a DBMS
 - Advantages and disadvantages of DBMS
 - Data independence (logical and physical)
 - Data abstraction levels: external, conceptual, internal
 - Database users and database system architecture
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Unit 2: Database Modeling

- Entity-Relationship (ER) model: entities, attributes, relationships
 - Enhanced ER model: generalization, specialization, aggregation
 - ER to relational mapping
 - UML diagrams (basic overview)
 - Conceptual, logical, and physical data models
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Unit 3: The Relational Data Model

- Structure of relational databases
 - Keys: super key, candidate key, primary key, foreign key
 - Relational algebra (select, project, union, difference, join, etc.)
 - Relational calculus (tuple and domain)
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Unit 4: Structured Query Language (SQL)

- DDL (CREATE, ALTER, DROP)
- DML (SELECT, INSERT, UPDATE, DELETE)
- DCL (GRANT, REVOKE)
- Aggregate functions, GROUP BY, HAVING, subqueries
- Joins (inner, outer, self join, natural join)
- Views, indexes, transactions in SQL

Unit 5: Integrity and Security

- Integrity constraints: domain, entity, referential, user-defined
- Triggers and assertions
- Transaction management: ACID properties
- Concurrency control (locks, timestamp, serializability)
- Database recovery techniques (log-based, checkpointing)
- Database security mechanisms (authentication, authorization, encryption)

Unit 6: Relational Database Design

- Functional Dependencies (FDs)
- Normal Forms: 1NF, 2NF, 3NF, BCNF, 4NF, 5NF
- Anomalies in database design
- Decomposition: lossless and dependency-preserving
- Multivalued dependencies and join dependencies

Unit 7: Object-Oriented Databases

- Object-oriented concepts in DBs: classes, objects, inheritance
- Object-relational databases (vs. pure object-oriented DBs)
- Persistent objects, encapsulation, polymorphism
- Querying object-oriented databases (OQL basics)

Unit 8: Advanced Topics

a. Distributed Databases

- Characteristics and architecture
- Fragmentation, replication, and allocation
- Distributed query processing
- Distributed transaction management

b. Cloud Databases

- Cloud database architecture
- Database-as-a-Service (DBaaS)

- Scalability and elasticity
- Popular cloud database platforms (e.g., Google BigQuery, Amazon RDS)

c. Big Data

- Characteristics: Volume, Variety, Velocity, Veracity
- NoSQL databases: key-value, document, column, graph
- MapReduce and Hadoop basics
- Comparison of traditional vs. big data storage systems

✓ 1. What is a Database? DBMS vs RDBMS

- **Database:** A structured collection of data that can be easily accessed, managed, and updated. Think of it as an organized digital filing cabinet.
 - **DBMS (Database Management System):** Software that allows users to define, create, maintain, and control access to the database. Example: Microsoft Access.
 - **RDBMS (Relational DBMS):** A type of DBMS that stores data in **tables** (relations) and allows relationships between data using **keys**. Example: MySQL, PostgreSQL.
 - ✓ *All RDBMS are DBMS, but not all DBMS are relational.*
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✓ 2. Components of a DBMS

A DBMS has several essential components:

- **Database Engine:** Core service that manages data storage, retrieval, and processing.
 - **Query Processor:** Translates user queries (e.g., SQL) into low-level instructions.
 - **Metadata Catalog:** Stores data about the data (e.g., table definitions).
 - **Transaction Manager:** Ensures that all operations are completed successfully or not at all.
 - **Concurrency Control Manager:** Handles multiple users accessing the same data.
 - **Recovery Manager:** Restores the database to a consistent state after a crash.
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✓ 3. Advantages and Disadvantages of DBMS

Advantages:

- Reduces **data redundancy** (no duplicate data)
- Ensures **data integrity and accuracy**
- Supports **data sharing** across multiple users
- Enforces **security and access control**
- Enables **backup and recovery**

Disadvantages:

- Complex to set up and manage
 - High **hardware/software cost**
 - Requires **trained personnel**
 - May lead to **performance issues** if not optimized
-

✓ 4. Data Independence (Logical and Physical)

Data independence means you can change the **data structure** without affecting the **application programs**.

- **Physical Data Independence:** Changing how data is stored (e.g., from SSD to HDD) doesn't affect higher-level applications.
- **Logical Data Independence:** Changing the schema (e.g., adding a column) without affecting existing queries or programs.

□ *This makes databases flexible and easy to maintain.*

✓ 5. Data Abstraction Levels

Data abstraction hides complex details and shows only relevant data. There are **3 levels**:

1. **Internal Level (Physical):** How data is actually stored (files, indexes, pointers).
2. **Conceptual Level (Logical):** The overall logical structure (tables, fields, relationships).
3. **External Level (View):** How individual users see the data (customized views).

✂ *Each user sees only what they need, not the entire database.*

✓ 6. Database Users and System Architecture

Types of Database Users:

- **End Users:** Use applications to interact with the database (e.g., customers on a shopping site)
- **Application Programmers:** Write programs to access data
- **DBA (Database Administrator):** Maintains and secures the database
- **System Analysts:** Design the database structure and rules

Database Architecture:

- **1-tier:** All in one system (e.g., personal DB in MS Access)
- **2-tier:** Client (UI) and Server (DBMS) separated
- **3-tier:** Client → Application Server → Database Server

☞ *This is most common in real-world web applications*

Unit 2: Database Modeling

✓ 1. Entity-Relationship (ER) Model

This is a visual way to design databases using **entities**, **attributes**, and **relationships**.

- **Entity:** A real-world object or concept.
 - ◆ Examples: Student, Teacher, Book
- **Attributes:** Properties that describe an entity.
 - ◆ Student may have: Student_ID, Name, DOB
- **Primary Key:** An attribute that uniquely identifies each entity.
 - ◆ Example: Student_ID
- **Relationship:** How entities are related to each other.
 - ◆ A Student *borrow*s a Book.
- **Types of Relationships:**
 - One-to-One (1:1): A person has one passport.
 - One-to-Many (1:N): A teacher teaches many students.
 - Many-to-Many (M:N): Students enroll in many courses; courses have many students.

✎ We draw ER diagrams using rectangles (entities), ovals (attributes), and diamonds (relationships).

✓ 2. Enhanced ER (EER) Model

The Enhanced ER model adds **more details** to the ER model to support real-world complexities:

- **Generalization:** Top-down approach.
 - ◆ *Combine similar entities into a single general entity.*
Example: Car and Bike → generalized to Vehicle.
- **Specialization:** Bottom-up approach.
 - ◆ *Split an entity into more specific entities.*
Example: Employee → specialized into Teacher and Admin.
- **Aggregation:** Treat a relationship as an entity.
 - ◆ Example: If Student *borrow*s Book, we may treat Borrow as an entity if it has attributes like Borrow_Date.

✎ These help model complex systems more accurately.

✓ 3. ER to Relational Mapping

This means converting the ER diagram into **tables**, which can be created in SQL.

- **Step 1:** Convert entities → tables.
- **Step 2:** Convert attributes → columns.
- **Step 3:** Add primary keys to uniquely identify rows.
- **Step 4:** Convert relationships → foreign keys or new tables (for M:N).

□ Example:

- Student (Student_ID, Name)
 - Book (Book_ID, Title)
 - Borrows (Student_ID, Book_ID, Borrow_Date) ← for M:N relationship
-

✓ 4. UML Diagrams (Basic Overview)

- **UML = Unified Modeling Language**, often used in software engineering.
- In database context, it's like ER diagrams but more standardized.
- UML class diagrams are similar to ER diagrams:
 - Class = Entity
 - Attributes = Fields
 - Associations = Relationships
 - Multiplicity = 1:1, 1:N, M:N

✎ UML is helpful when designing a database as part of a software project.

✓ 5. Conceptual, Logical, and Physical Data Models

These are **levels of database design**:

- **Conceptual Model** (What to store?)
 - High-level view, usually an ER diagram.
 - Focus on business rules, entities, and relationships.
- **Logical Model** (How to store?)
 - More technical.
 - Adds data types, constraints, and normalization.
 - Ready to be translated into SQL.
- **Physical Model** (Where to store?)
 - Focus on how data is stored on disk.
 - Includes indexes, file structures, and performance tuning.

□ Analogy:

- Conceptual: *Sketch on paper*
 - Logical: *Technical drawing*
 - Physical: *Actual construction with materials*
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Unit 3: The Relational Data Model — this unit forms the *core theory* behind how modern relational databases like MySQL, PostgreSQL, and Oracle work.

✓ 1. Structure of Relational Databases

A **relational database** organizes data into **tables** (also called **relations**).

- Each **table** represents an **entity** (e.g., Students, Courses).
- Each **row** in a table is called a **tuple** — it represents a record.
- Each **column** is an **attribute** — it stores one kind of information.

◆ Example: STUDENT Table

Student_ID Name Age Department

101 Aditi 20 CS

102 Ravi 21 IT

Key concepts:

- Tables must have a **unique identifier**: a **Primary Key**.
 - Data is organized in **relations** with **no duplicate rows**.
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✓ 2. Keys: Super Key, Candidate Key, Primary Key, Foreign Key

These are special columns used to **identify rows** and **connect tables**:

- **Super Key**: Any combination of attributes that uniquely identifies a row.
◆ Example: {Student_ID}, {Student_ID, Name}
- **Candidate Key**: A **minimal** super key (no extra fields).
◆ Example: {Student_ID} (if it's enough on its own)
- **Primary Key**: One candidate key is chosen as the main identifier.
◆ Must be **unique** and **not null**.
- **Foreign Key**: A column that creates a **relationship** by pointing to the **Primary Key** in another table.
◆ Ensures **referential integrity**.

□ Example:

If `Enrollment(Student_ID, Course_ID)` references `Student(Student_ID)` and `Course(Course_ID)` — then `Student_ID` is a **foreign key** in `Enrollment`.

✓ 3. Relational Algebra

This is a **theoretical language** for querying relational databases — it forms the **foundation** of SQL.

Here are some common operations:

- **σ (SELECT):** Filters rows based on condition.
 - ◆ Example: `σ Age > 20 (Student)`
- **π (PROJECT):** Selects specific columns.
 - ◆ Example: `π Name, Department (Student)`
- **\cup (UNION):** Combines rows from two tables with the same schema.
 - ◆ Example: `Student_2023 \cup Student_2024`
- **$-$ (DIFFERENCE):** Finds rows in one table not in another.
 - ◆ Example: `All_Students - Graduated_Students`
- **\times (CARTESIAN PRODUCT):** Combines all rows from two tables.
 - ◆ Rarely used directly.
- **\bowtie (JOIN):** Combines tables based on a condition (usually matching keys).
 - ◆ Example: `Student \bowtie Enrollment`
- **Rename (ρ):** Renames a table or column.

✂ These operations are used to build complex queries step by step.

✓ 4. Relational Calculus (Tuple and Domain)

This is another theoretical way to query data, but more **declarative** — you describe *what* you want, not *how* to get it.

- **Tuple Relational Calculus (TRC):**
 - Uses variables that represent **rows (tuples)**.
 - ◆ Example:
`{ t | t \in Student AND t.Age > 20 }`
(*t is a tuple from Student where Age > 20*)
- **Domain Relational Calculus (DRC):**
 - Uses variables that represent **column values (domains)**.

◆ Example:

$\{ \langle n, d \rangle \mid \exists id (\langle id, n, age, d \rangle \in \text{Student AND } age > 20) \}$

□ Think of TRC like saying “find the students as rows,” and DRC like “find the names and departments as values.”

Unit 4: Structured Query Language (SQL) is the most practical and applied part of your syllabus — it's the language we use to talk to relational databases.

Let's go point-by-point with clear and simple explanations:

✓ 1. DDL – Data Definition Language

Used to **define or modify** the structure of database objects like tables.

- **CREATE:** Creates a new table or database

```
sql

CREATE TABLE Student (
    ID INT PRIMARY KEY,
    Name VARCHAR(50),
    Age INT
);
```

- **ALTER:** Modifies an existing table (add/remove/change column)

```
sql

ALTER TABLE Student ADD Email VARCHAR(100);
```

- **DROP:** Deletes a table or database permanently

```
sql

DROP TABLE Student;
```

✓ 2. DML – Data Manipulation Language

Used to **work with data inside tables** (not the structure).

- **SELECT:** Retrieves data

```
sql

SELECT Name, Age FROM Student;
```

- **INSERT:** Adds a new record

```
sql

INSERT INTO Student (ID, Name, Age) VALUES (1, 'Amit', 20);
```

- **UPDATE:** Modifies existing records

sql

```
UPDATE Student SET Age = 21 WHERE ID = 1;
```

- **DELETE:** Removes records

sql

```
DELETE FROM Student WHERE ID = 1;
```

✓ 3. DCL – Data Control Language

Used to **control access** to the database.

- **GRANT:** Gives permission

sql

```
GRANT SELECT, INSERT ON Student TO user1;
```

- **REVOKE:** Removes permission

sql

```
REVOKE INSERT ON Student FROM user1;
```

✂ *This is how we manage user security and access rights.*

✓ 4. Aggregate Functions, GROUP BY, HAVING, Subqueries

◆ Aggregate Functions:

Used to **summarize data**

- COUNT(), SUM(), AVG(), MIN(), MAX()

sql

```
SELECT AVG(Age) FROM Student;
```

◆ GROUP BY:

Groups rows based on one or more columns

sql

```
SELECT Department, COUNT(*) FROM Student GROUP BY Department;
```

◆ **HAVING:**

Filters groups (like WHERE does for rows)

sql

```
SELECT Department, COUNT(*)  
FROM Student  
GROUP BY Department  
HAVING COUNT(*) > 10;
```

◆ **Subqueries:**

A query **inside another query**

sql

```
SELECT Name FROM Student  
WHERE Age > (SELECT AVG(Age) FROM Student);
```

✓ **5. Joins**

Used to **combine data from multiple tables** based on a related column.

◆ **Inner Join:**

Returns only matching rows.

sql

```
SELECT Student.Name, Course.CourseName  
FROM Student  
INNER JOIN Course ON Student.CourseID = Course.ID;
```

◆ **Left Outer Join:**

Returns all from the left table + matching from the right.

sql

```
SELECT Student.Name, Course.CourseName  
FROM Student  
LEFT JOIN Course ON Student.CourseID = Course.ID;
```

◆ Right Outer Join:

Returns all from the right table + matching from the left.

◆ Full Outer Join:

Returns all records when there is a match in one of the tables.

◆ Self Join:

A table joined with itself.

sql

```
SELECT A.Name, B.Name
FROM Employee A, Employee B
WHERE A.ManagerID = B.ID;
```

◆ Natural Join:

Auto-matches columns with the same name and type.

✓ 6. Views, Indexes, Transactions in SQL

◆ View:

A virtual table based on the result of a SELECT query.
Used for **security** or **simplifying queries**.

sql

```
CREATE VIEW SeniorStudents AS
SELECT * FROM Student WHERE Age > 21;
```

◆ Index:

Improves **query performance** by speeding up data retrieval.

sql

```
CREATE INDEX idx_name ON Student (Name);
```

◆ Transactions:

A group of SQL operations that execute **together** — either all succeed or none.

- ACID properties: **Atomicity, Consistency, Isolation, Durability**

sql

```
BEGIN;  
UPDATE Account SET Balance = Balance - 500 WHERE ID = 1;  
UPDATE Account SET Balance = Balance + 500 WHERE ID = 2;  
COMMIT;  
-- Or use ROLLBACK to undo
```

Unit 5: Integrity and Security in simple and clear terms. This unit focuses on making sure your data is **correct, consistent, secure, and protected during crashes or misuse**.

✓ 1. Integrity Constraints

These are rules that ensure **data stays valid and consistent**.

◆ Domain Constraint

Ensures data is within a valid range or type.

→ Example: Age should be a number between 0 and 120.

◆ Entity Integrity

Every table must have a **primary key**, and it **can't be NULL**.

→ No two rows should have the same ID.

◆ Referential Integrity

Ensures that **foreign keys** must match an existing value in the related table.

→ A student can't be enrolled in a course that doesn't exist in the Course table.

◆ User-Defined Constraints

Custom business rules.

→ Salary must be greater than minimum wage, or enrollment date can't be in the future.

✓ 2. Triggers and Assertions

◆ Trigger

An **automatic action** executed when an event occurs in the database (like insert, update, delete).

→ Example: When a new student is added, automatically insert their email into the contact table.

sql

```
CREATE TRIGGER before_insert_check
BEFORE INSERT ON Student
FOR EACH ROW
BEGIN
```

```
IF NEW.Age < 0 THEN
    SIGNAL SQLSTATE '45000' SET MESSAGE_TEXT = 'Age cannot be negative';
END IF;
END;
```

◆ Assertion

A **condition** that must always be true in the database.

☞ Example: Total enrollment in a class must not exceed 60 students.

Note: Not all databases support assertions directly, but the logic can often be implemented using triggers or constraints.

✓ 3. Transaction Management: ACID Properties

A **transaction** is a group of SQL operations that should be treated as a **single unit**.

ACID stands for:

- **Atomicity:** All operations in a transaction happen, or none do.
- **Consistency:** The database goes from one valid state to another.
- **Isolation:** Transactions don't interfere with each other.
- **Durability:** Once a transaction is committed, it stays even after a crash.

☞ Example: Transferring ₹500 from Account A to Account B — both debit and credit must succeed or both must fail.

✓ 4. Concurrency Control

When **multiple users** access the database at the same time, we need rules to **avoid conflicts**.

◆ Locking

- **Shared lock:** Allows read but not write.
- **Exclusive lock:** Allows read and write, but no one else can access the data.

◆ Timestamp Ordering

- Each transaction is given a timestamp.
- Older transactions get priority.
- Used to ensure no overlapping writes/reads.

◆ Serializability

- Ensures that the **final result** of concurrent transactions is the same as if they were run **one at a time**.

□ Goal: Prevent problems like **lost updates**, **dirty reads**, or **inconsistent data**.

✓ 5. Database Recovery Techniques

Used when there's a **crash**, power failure, or error.

◆ Log-Based Recovery

- Every action is written to a **log file** before it's actually done.
- If a crash happens, logs are used to redo or undo operations.

◆ Checkpointing

- At intervals, the database takes a **snapshot** of its state.
- During recovery, it can start from the last checkpoint instead of the beginning.

□ Combined, these techniques make sure your database can bounce back to a safe state.

✓ 6. Database Security Mechanisms

Used to **protect the database from unauthorized access** or misuse.

◆ Authentication

- Verifies the user's identity (e.g., username & password).

◆ Authorization

- Defines what **actions** a user can perform (e.g., read-only access to some tables, full access to others).

◆ Encryption

- Converts sensitive data into a **secure format** so even if someone steals it, they can't understand it.

- Used for storing passwords, credit card numbers, etc.

□ *Just like securing your house with a lock, security ensures only the right people can open the right doors — and safely.*

Unit 6: Relational Database Design is all about designing databases that are **efficient**, **non-redundant**, and **easy to update without errors**. Let's break this down step by step.

✓ 1. Functional Dependencies (FDs)

A **Functional Dependency (FD)** is a rule that shows how one piece of data **determines** another.

◆ Notation:

$A \rightarrow B$ means **if you know A, you can determine B**.

◆ Example:

If $\text{Student_ID} \rightarrow \text{Name}$, then:

- Knowing the `Student_ID`, you can find the `Name`.
- But knowing the `Name` doesn't guarantee a unique `Student_ID`.

□ *FDs help us decide how to split or join tables during normalization.*

✓ 2. Normal Forms (1NF to 5NF)

Normalization is the process of organizing tables to **remove redundancy and prevent anomalies**.

◆ 1NF (First Normal Form)

No repeating groups. Each cell must have a single value.

✗ Bad Table:

ID Name Phone Numbers

1 Asha 1234, 5678

✓ 1NF:

ID Name Phone Number

1 Asha 1234

1 Asha 5678

◆ 2NF (Second Normal Form)

Remove partial dependencies (applies to composite keys).

☞ *Every non-key attribute must depend on the full primary key.*

Example:

StudentID CourseID StudentName

Problem: StudentName depends **only** on StudentID, **not** CourseID.

✓ Split into two tables:

- Student(StudentID, StudentName)
 - Enrollment(StudentID, CourseID)
-

◆ 3NF (Third Normal Form)

Remove transitive dependencies.

☞ *Non-key columns must depend only on the key — not on other non-key columns.*

Example:

| ID | Name | Department | DeptHead |

Here, DeptHead depends on Department, **not** on ID.

✓ Split:

- Student(ID, Name, Department)
 - Department(Department, DeptHead)
-

◆ BCNF (Boyce-Codd Normal Form)

Stricter version of 3NF.

⇒ *For every functional dependency $A \rightarrow B$, A must be a super key.*

◆ 4NF (Fourth Normal Form)

Deals with **multivalued dependencies**.

⇒ If two attributes are **independent** of each other but both depend on the same key, split them.

Example:

A person can have **multiple skills** and **multiple hobbies** — but they're not related.

✓ Break into:

- `Person_Skills`
 - `Person_Hobbies`
-

◆ 5NF (Fifth Normal Form)

Deals with **join dependencies** — break down data so it can be reconstructed by joining without redundancy.

Used in **very advanced modeling**, like handling complex business rules.

✓ 3. Anomalies in Database Design

Poorly designed databases can lead to three types of anomalies:

- **Insertion Anomaly:** Can't add data unless related data is also available.
Example: Can't add a course unless a student is enrolled.
 - **Update Anomaly:** Updating in one place but forgetting in another creates inconsistency.
 - **Deletion Anomaly:** Deleting one piece of data causes **loss of other valuable data**.
Example: Deleting a student who's the only one enrolled in a course may delete the course info too.
-

✓ 4. Decomposition: Lossless and Dependency-Preserving

◆ Decomposition

Splitting a table into two or more smaller tables.

- **Lossless:** No data is lost when tables are joined back.
 - Must have a common key.
- **Dependency-Preserving:** All original FDs should still hold in the new tables.
 - Helps maintain data rules without extra joins.

□ *Good decomposition = no redundancy + no data loss + rules preserved.*

✓ 5. Multivalued and Join Dependencies

◆ Multivalued Dependency (MVD)

When one attribute depends on a key, and another also depends on the same key, but they're **independent of each other**.

→ Example:

- A student has multiple **phone numbers** and multiple **hobbies**.
- But phone numbers and hobbies are **not related**.

◆ Join Dependency

A table can be split into smaller tables and later **reconstructed (joined) without data loss**.

→ 5NF ensures that **all join dependencies are preserved**.

Unit 7: Object-Oriented Databases — this is where database design meets object-oriented programming (OOP). It's perfect for advanced applications like CAD systems, multimedia, and AI where complex data types and relationships are common.

✓ 1. Object-Oriented Concepts in Databases: Classes, Objects, Inheritance

◆ Class

A blueprint or template for an object — just like in OOP.

→ In a database, a class defines a structure (like a table) and behavior (methods).

Example:

text

```
Class: Student
Attributes: ID, Name, Age
Methods: enroll(), updateProfile()
```

◆ Object

An instance of a class — a real data record.

→ Object = row in a table, but with behavior included.

◆ Inheritance

Just like in OOP, a subclass can inherit properties and methods from a parent class.

→ Example:

Employee class → Professor and Admin inherit ID, Name, Salary, etc.

□ This allows **code reuse** and **data consistency**.

✓ 2. Object-Relational Databases vs Pure Object-Oriented Databases

◆ Object-Oriented Database (OODB)

- Stores data as **objects**, not just rows and columns.
- Supports full OOP: classes, objects, inheritance, polymorphism.
- Data + behavior (methods) are stored together.

◆ Examples: db4o, ObjectDB

◆ Object-Relational Database (ORDBMS)

- A **hybrid**: Combines relational databases with object-oriented features.
- Extends SQL to handle complex data types (arrays, multimedia, user-defined types).
- Still table-based, but supports classes and inheritance.

◆ Examples: PostgreSQL, Oracle (with object support)

□ *ORDBMS is more practical in real-world systems; OODB is more powerful for complex apps but less common.*

✓ 3. Persistent Objects, Encapsulation, Polymorphism

◆ Persistent Objects

Objects that remain stored in the database even after the program ends.

⇒ Think of them as "saved" objects — they don't disappear when the app closes.

◆ Encapsulation

Data + behavior (methods) are bundled together inside an object.

⇒ You can't directly change the data from outside — must use methods like `updateSalary()`.

◆ Polymorphism

Same method name can behave differently in different classes.

⇒ Example: `calculateSalary()` behaves differently for `Teacher` and `Admin`.

□ *These features bring the power of OOP into data management.*

✓ 4. Querying Object-Oriented Databases (OQL Basics)

OQL = Object Query Language

It's like SQL but designed to query objects instead of just rows.

Example:

```
sql
```

```
SELECT s.name  
FROM Student s
```

```
WHERE s.age > 20;
```

Advanced:

```
sql
```

```
SELECT c.name  
FROM Course c  
WHERE c.professor.department.name = 'Computer Science';
```

□ OQL supports **navigating object relationships** using dot notation (just like accessing attributes in programming).

Summary Table:

Concept	Meaning
Class	Template/structure of an object
Object	Real instance of a class
Inheritance	One class gets attributes from another
Persistent Object	Stored permanently in the DB
Encapsulation	Hide data, expose only methods
Polymorphism	One function, multiple behaviors
OQL	SQL-like language for querying objects

Would you like a diagram comparing RDBMS vs ORDBMS, a real-world example using OQL, or a quiz on object-oriented DB concepts?

Unit 8a: Distributed Databases, one of the most important real-world topics for large-scale, global systems.

✓ What is a Distributed Database?

A **distributed database** is a database that is **spread across multiple computers** (called nodes), which may be in the same location or across the world — but all appear to users as a **single database**.

→ Example: Google stores your emails in data centers across continents, but you access it as one seamless inbox.

✓ 1. Characteristics of Distributed Databases

- **Data is stored on multiple sites (servers)**, often geographically far apart.
 - Users feel they are working with a **single database**.
 - Data can be **replicated** or **fragmented** across sites.
 - **Local autonomy**: Each site can manage its own data.
 - **Fault tolerance**: If one server fails, others can take over.
-

✓ 2. Architecture of Distributed Databases

There are 3 major types of architectures:

◆ Client-Server

Clients send queries to a central or distributed set of servers.

→ Simple and common in enterprise apps.

◆ Peer-to-Peer

All nodes (servers) are equal — each can act as a client or a server.

→ Used in modern scalable systems like blockchain or file-sharing apps.

◆ Multi-Database Systems

Each site has its **own local database**, but they're linked through a coordination system.

→ Best for big corporations with departments in different cities.

✓ 3. Fragmentation, Replication, and Allocation

◆ Fragmentation

Splitting a large table into smaller pieces (fragments) across sites.

- **Horizontal Fragmentation:** Rows are divided
 ☞ Example: Customers from India in one fragment, from USA in another
- **Vertical Fragmentation:** Columns are divided
 ☞ Example: Names and Emails in one fragment, Orders and Payments in another

◆ Replication

Copying the same data to multiple sites for availability and faster access.

☞ If one site goes down, others still have the data.

- **Full replication:** Every site has a copy of the entire DB.
- **Partial replication:** Only selected data is copied to certain sites.

◆ Allocation

Deciding **where to store which data** to optimize performance.

☞ Based on:

- Access frequency
- Cost of storage
- Network speed

☐ *Goal: Keep data close to where it's used the most.*

✓ 4. Distributed Query Processing

How to **execute a query** that may need data from multiple locations?

Steps:

1. **Query is parsed** and optimized.
2. **Subqueries** are sent to relevant sites.
3. **Results** from each site are combined.

4. Final result is returned to the user.

Challenges:

- Data transfer takes time.
 - Must minimize the number of site accesses and network traffic.
 - Ensuring results are **correct and fast**.
-

✓ 5. Distributed Transaction Management

A **transaction** is a set of operations that must be completed fully or not at all — just like in centralized databases, but harder in distributed ones.

◆ Problem:

A transaction may touch data from **multiple locations** — we must ensure:

- No site is left in an inconsistent state.
- All sites **agree** to commit.

◆ Two-Phase Commit (2PC) Protocol:

1. **Prepare Phase:** Coordinator asks all sites — “Can you commit?”
2. **Commit Phase:** If all say “yes,” it tells them to **commit**. If even one says “no,” it **rolls back** everywhere.

□ It ensures **Atomicity** and **Consistency**, but can be slow.

Summary Table:

Concept	Meaning
Fragmentation	Split data by rows/columns across sites
Replication	Copy of data kept on multiple sites
Allocation	Decide where each piece of data should go
Query Processing	Execute queries efficiently across sites
Transaction Mgmt	Ensure consistency & commit across sites (2PC, etc.)

Great! Let's explore **Unit 8b: Cloud Databases** — this is all about how databases run over the internet using cloud platforms. It's a must-know topic in today's data-driven world.

✓ 1. Cloud Database Architecture

A **cloud database** is hosted on **cloud infrastructure** (like AWS, Azure, or GCP) and accessed **remotely over the internet**, not installed on your own machine.

◆ Key Features:

- **Accessible via web interfaces or APIs**
- **Managed by the provider** (no manual installation, backup, scaling)
- **Highly available and fault-tolerant**
- Can be **relational** (like MySQL) or **NoSQL** (like MongoDB)

◆ Architecture Layers:

1. **Frontend** – Web app or dashboard to access data
2. **API Layer** – Communication between app and database
3. **Compute Layer** – Executes queries and processing
4. **Storage Layer** – Actual disk/data storage (automatically scalable)
5. **Security Layer** – Access control, encryption, authentication

□ Think of it as "Database on Rent" — no hardware, only usage fees.

✓ 2. Database-as-a-Service (DBaaS)

DBaaS is a cloud model where you get a **fully managed database service** — you don't need to worry about setup, patching, scaling, or backups.

◆ Features:

- Pay-as-you-go pricing
- Auto backup & recovery
- Monitoring and performance tuning
- Easy integration with apps

◆ Examples:

- **Amazon RDS (Relational)**
- **MongoDB Atlas (NoSQL)**

- **Google Cloud SQL / Firebase**
- **Azure SQL Database**

□ *DBaaS = You focus on data & queries, the provider handles everything else.*

✓ 3. Scalability and Elasticity

These are **core advantages** of cloud databases:

◆ Scalability:

Ability to **grow** with demand.

- **Vertical scaling:** Add more power (CPU/RAM) to a single server.
- **Horizontal scaling:** Add more servers to handle more data/users (sharding, clustering).

◆ Elasticity:

System automatically **scales up or down** based on real-time need.

→ Saves cost when demand is low, boosts power when demand is high.

□ Imagine your database **expanding during festive sales** and **shrinking after** — without you doing anything.

✓ 4. Popular Cloud Database Platforms

Here are some widely used cloud DB services:

Platform	Type	Key Features
Amazon RDS	Relational	Supports MySQL, PostgreSQL, Oracle, SQL Server
Amazon DynamoDB	NoSQL	Fast, scalable key-value and document DB
Google BigQuery	Relational (Analytical)	Petabyte-scale SQL analytics engine
Google Cloud SQL	Relational	Fully managed PostgreSQL/MySQL/SQL Server
Microsoft Azure SQL	Relational	Scalable, secure cloud SQL database
MongoDB Atlas	NoSQL	Fully managed MongoDB in the cloud

Platform	Type	Key Features
Firestore	NoSQL	Real-time syncing for mobile/web apps

Summary Table:

Concept	Meaning
Cloud DB	A database hosted online using cloud infrastructure
DBaaS	Managed database service provided by cloud providers
Scalability	Ability to handle increasing data or user load
Elasticity	Auto-adjust resources based on demand
BigQuery / RDS Examples of popular cloud-based database platforms	

Let's break down **Unit 8c: Big Data** in clear and simple terms. This unit introduces the world of **huge, fast, and complex datasets** that traditional databases can't easily handle.

✓ 1. Characteristics of Big Data: 4 Vs

Big Data is often described using **four main characteristics**:

◆ Volume

The **amount of data** is massive — from terabytes to petabytes.

→ Example: Facebook generates over 4 petabytes of data per day.

◆ Variety

Data comes in **different formats**: structured (tables), semi-structured (XML/JSON), and unstructured (videos, images, social media posts).

◆ Velocity

Data is created and processed at **high speed**.

→ Example: Stock prices updating every second.

◆ Veracity

Refers to the **trustworthiness or quality** of the data.

→ Is the data accurate? Does it have errors or noise?

□ These 4 Vs make Big Data **complex** but also **valuable** when processed correctly.

✓ 2. NoSQL Databases

Traditional relational databases (RDBMS) are not great for big, varied, fast-changing data. NoSQL databases are built to solve this.

Here are **4 main types of NoSQL databases**:

◆ Key-Value Store

Simple format: a key and a value

→ Example: UserID: 101 → {"name": "Amit", "age": 22}

- ✓ Very fast for lookups
- ◆ Examples: Redis, Amazon DynamoDB

◆ Document Store

Stores data as **documents**, often in JSON or XML.

- ☞ Flexible schema
- ◆ Examples: MongoDB, CouchDB

◆ Column Store

Stores data **by columns** instead of rows — great for big analytical queries.

- ☞ Used in data warehousing.
- ◆ Examples: Apache Cassandra, HBase

◆ Graph Database

Stores data as **nodes and relationships** — great for social networks or recommendation systems.

- ◆ Examples: Neo4j, Amazon Neptune

☐ *NoSQL* = “Not Only SQL” → *suitable for Big Data problems*.

✓ 3. MapReduce and Hadoop Basics

◆ MapReduce

A programming model used to process big data in parallel across many computers.

- **Map Step:** Breaks data into chunks and processes them separately.
 - ☞ Example: Count words in multiple documents.
- **Reduce Step:** Combines results into a final answer.
 - ☞ Total up word counts from each document.

◆ Hadoop

An open-source framework that supports MapReduce and provides a **distributed file system (HDFS)**.

- Stores big data **across multiple machines**
- Works even if some machines fail
- Handles large-scale data processing

□ Hadoop = HDFS (storage) + MapReduce (processing)

✓ 4. Traditional vs. Big Data Storage Systems

Feature	Traditional DB (RDBMS)	Big Data System (e.g., Hadoop, NoSQL)
Data Size	GBs to low TBs	TBs to PBs
Schema	Fixed schema (strict)	Flexible or schema-less
Scalability	Vertical (add CPU/RAM)	Horizontal (add more machines)
Data Type	Mostly structured	Structured + Unstructured + Semi-structured
Query Language	SQL	NoSQL / Custom APIs
Processing Model	Single node	Distributed (MapReduce, Spark, etc.)

□ *Big Data systems are built to scale, handle messy data, and process it fast.*