

Quick intro to Relational Algebra

(More after chapter 3)

2.4

Relational Algebra defines the set of operations on relations.

- **Operands**: relations
- **Operators**: Π , σ , \times , \bowtie , etc.
- **Results**: relations.

Selection σ

$t \in \sigma_p R$ iff p is true for t .

p is a predicate on attrs of R

Ex. R :

a	b
3	x
2	y

$\sigma_{a>2} R$

a	b
3	x

What is $\sigma_{a>2} R$ schema?

Projection Π

Projects a subset of the attr of R

$\Pi_{\langle \text{att list} \rangle} R$ is a relation that only includes the attributes in $\langle \text{att list} \rangle$

We will extend this definition later.

Ex:

$R(a,b)$

a	b
3	x
3	y

$\Pi_a R$

a
3

What is $\Pi_a R$ schema?

Cross product: \times

Given relations R and S .

$(r, t) \in R \times S$ iff $r \in R$ and $s \in S$

$R(a, b)$

a	b
1	x
2	y

$S(c, d)$

c	d
5	8
2	12

$T = R \times S$

a	b	c	d
1	x	5	8
1	x	2	12
2	y	5	8
2	y	2	12

What is schema of T ?

Natural Join \bowtie

Given relations R and S

C is set of attributes of both S and R
with the same name

• if C is empty.

$$R \bowtie S = R \times S$$

• otherwise

$$\pi_{\underbrace{\text{Attr}(R), \text{Attr}(S) - C}}_{\uparrow\uparrow}$$

Do not project
both common
attributes (only
the first).

$$\sigma_{\bigwedge_{a_i \in C} R_{a_i} = S_{a_i}} (R \times S)$$

$\uparrow\uparrow$

match tuples
with same value in
common attributes.
conjunction over
all common attributes

Ex:

$R(a, b)$

a	b
1	x
2	y

$S(a, c)$

a	c
5	8
2	12

Common attributes = {a}

$$T = R \bowtie S = \Pi_{a,b,c} \sigma_{R.a=S.a}(R \times S)$$

$R \times S$

R.a	R.b	S.a	S.c
1	x	5	8
1	x	2	12
2	y	5	8
2	y	2	12

} $R.a = S.a$

$R \bowtie S$

a	b	c
2	y	12