

# Approximate Algorithms

## Minimum VERTEX COVER

Given an undirected graph  $G = (V, E)$  find the smallest set of vertices  $S \in V$  such that all edges in  $G$  have at least one of their endpoints in  $S$ .

This is an NP-hard problem.

### Claim

MVC is a 2-approximation for VERTEX COVER.

We don't know how big  $OPTVC(G)$  is, but we can lower bound its size.

$$|MVC(G)| \leq 2|OPTVC(G)|$$

Because edges selected by MVC form a matching, no vertex covers more than one edge in a matching.

## Matching Vertex Cover

Pick edge  $(u,v)$ , remove the edge and all edges adjacent to  $u$  or  $v$ . Add  $u, v$  to vertex cover. Repeat

## List Scheduling Approximation 1966 Graham

Given  $n$  jobs, job  $i$  must execute uninterruptedly for  $P_i$  time units.

$m$  machines (identical) each machine can work on one job at a time

Find a schedule of jobs that minimizes completion time (time when last machine finishes).