

## CPSC 421 - Homework 5

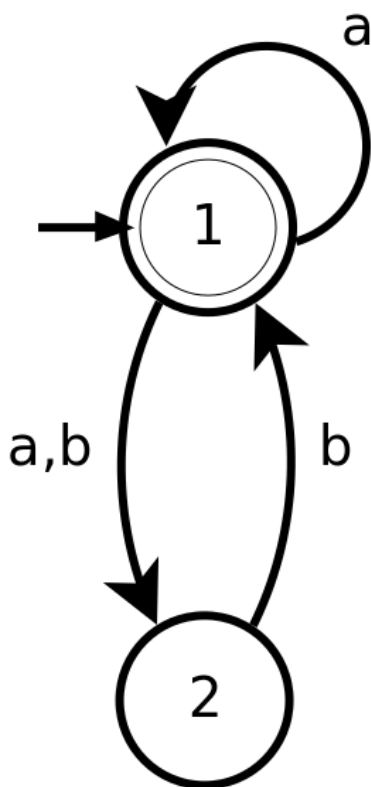
Tristan Rice, q7w9a, 25886145

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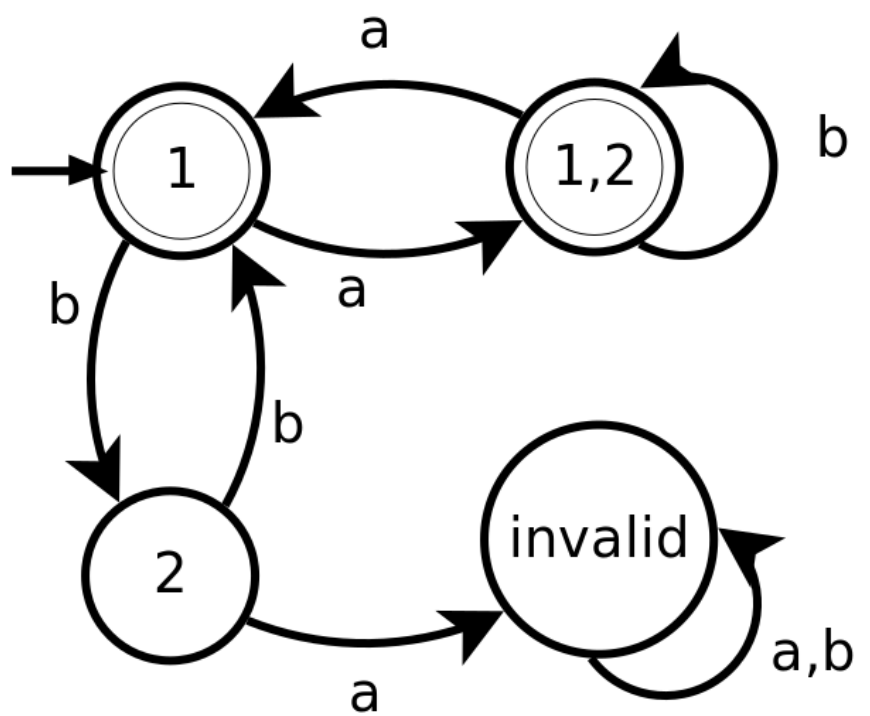
Possible states:

- empty  $\{1\}$
- a  $\{1, 2\}$
- b  $\{2\}$
- aa  $\{1\}$
- ab  $\{1, 2\}$
- ba  $\{\}$
- bb  $\{1\}$

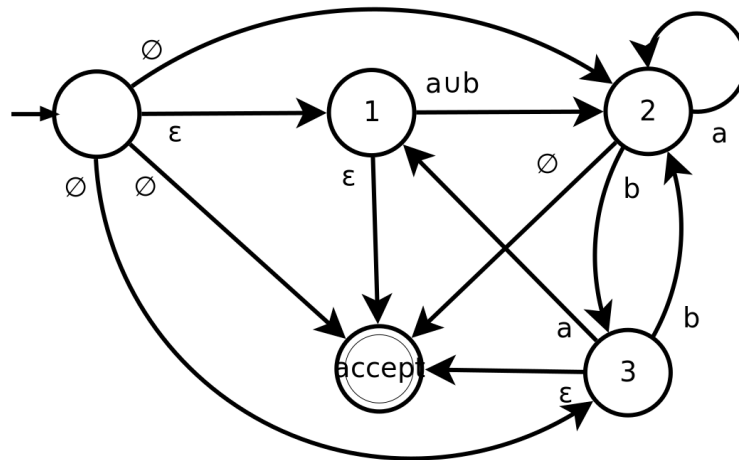
Original



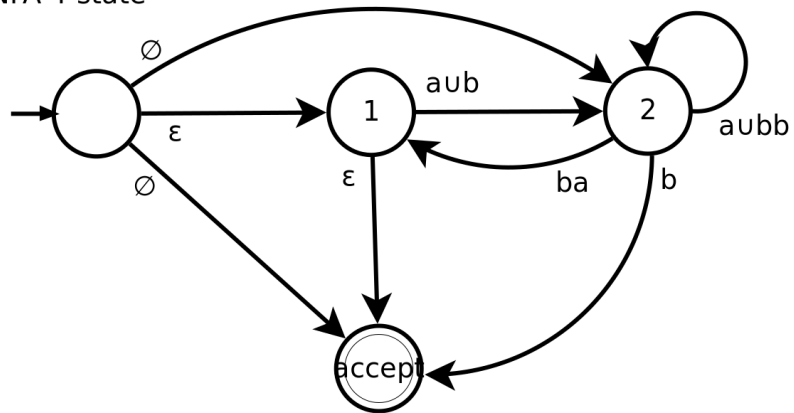
DFA



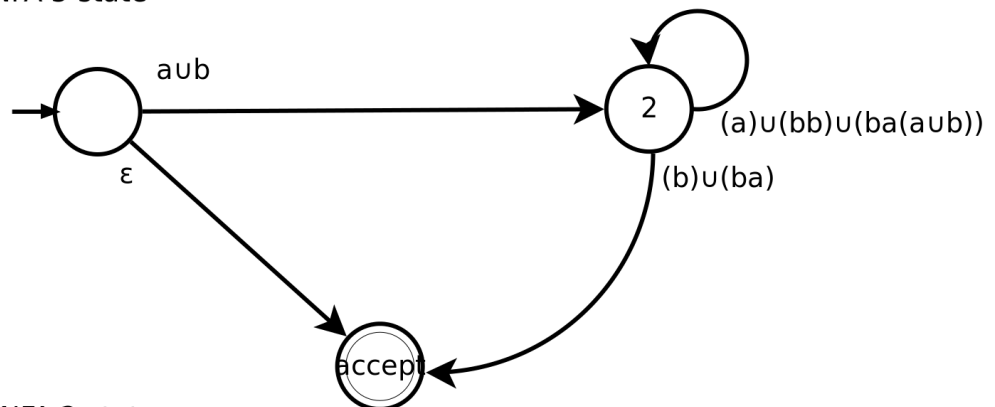
GNFA 5-state



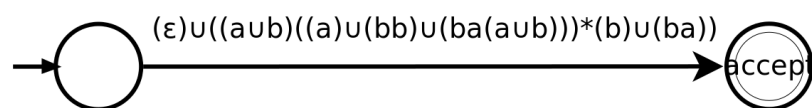
GNFA 4-state



GNFA 3-state



GNFA 2-state



### 3

#### 3.a

Claim:  $\binom{2n}{n} \sim \frac{2^{2n}\gamma}{\sqrt{n}}$ .

Proof:

$$n! \sim \sqrt{2\pi n}(n/e)^n$$

$$\begin{aligned} \binom{2n}{n} &= \frac{(2n)!}{n!(2n-n)!} \\ &= \frac{(2n)!}{n!n!} \\ &\sim \frac{\sqrt{2\pi 2n}(2n/e)^{2n}}{\sqrt{2\pi n}(n/e)^n \sqrt{2\pi n}(n/e)^n} \\ &= \frac{(2n/e)^{2n}}{(n/e)^{2n} \sqrt{\pi n}} \\ &= \frac{2^{2n}}{\sqrt{\pi n}} \end{aligned}$$

This is clearly equal to  $\frac{2^{2n}\gamma}{\sqrt{n}}$  where  $\gamma = \frac{1}{\sqrt{\pi}}$ .

Thus,  $\binom{2n}{n} \sim \frac{2^{2n}\gamma}{\sqrt{n}}$ .

#### 3.b

Claim:  $\binom{2n+1}{n} \sim \frac{2^{2n+1}\gamma}{\sqrt{n}}$ .

Proof:

We know that  $\binom{2n}{n} \sim \frac{2^{2n}\gamma}{\sqrt{n}}$ . Thus, if we can show that the asymptotic ratio is the same for  $\frac{\binom{2n+1}{n}}{\binom{2n}{n}}$  and  $\frac{2^{2n+1}\gamma/\sqrt{n}}{2^{2n}\gamma/\sqrt{n}}$ , we know will that  $\binom{2n+1}{n} \sim \frac{2^{2n+1}\gamma}{\sqrt{n}}$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{\binom{2n+1}{n}}{\binom{2n}{n}} &= \lim_{n \rightarrow \infty} \frac{(2n+1)!n!}{(2n)!n!(2n+1-n)!} \\ &= \lim_{n \rightarrow \infty} \frac{(2n+1)!n!}{(2n)!(n+1)!} \\ &= \lim_{n \rightarrow \infty} \frac{2n+1}{n+1} \\ &= 2 \end{aligned}$$

$$\begin{aligned} \lim_{n \rightarrow \infty} \frac{2^{2n+1}\gamma/\sqrt{n}}{2^{2n}\gamma/\sqrt{n}} &= \lim_{n \rightarrow \infty} \frac{2^{2n+1}}{2^{2n}} \\ &= \lim_{n \rightarrow \infty} 2 \\ &= 2 \end{aligned}$$

Thus, since the asymptotic ratios are the same, we know that  $\binom{2n+1}{n} \sim \frac{2^{2n+1}\gamma}{\sqrt{n}}$  must be proportional.

#### 3.c

$$L = \{s \in \{0, 1\}^* \mid s \text{ has } n \text{ zeros and } m \text{ ones, and } m = n \text{ or } m = n + 1\}$$

If a DFA were to exist, it would have a walk counting function  $f(l)$  equal to the number of words of length  $l$ .

When length  $l$  is even,  $m = n$  and there are  $\binom{2n}{n}$  words. When length  $l$  is odd,  $m = n + 1$  and there are  $\binom{2n+1}{n}$  words.

For the even case, we know  $\binom{2n}{n}$  is proportional to  $\sqrt{2\pi n}\gamma n^{-\frac{1}{2}}$ . This means that  $f(l) = \Theta(2^l n^r)$ ,  $r = -\frac{1}{2}$ . Since  $r$  is not an integer,  $f(l)$  must not be a walk count due to Theorem 7.1 from handout 1.

For the odd case, we know  $\binom{2n+1}{n}$  is proportional to  $\sqrt{2^{2n+1}}\gamma n^{-\frac{1}{2}}$ . This means that  $f(l) = \Theta(2^l n^r)$ ,  $r = -\frac{1}{2}$ . Since  $r$  is not an integer,  $f(l)$  must not be a walk count due to Theorem 7.1 from handout 1.

Thus, since the function  $f(l)$  is not a valid walk count and no DFA can be constructed for the language. Since all regular languages have a corresponding DFA,  $L$  must not be a regular language.

## 4

### 4.a

Claim:  $L = \{1^m \mid 7 \text{ divides } m\}$  is not recognized by a DFA with fewer than 7 states.

Proof: We can apply the Myhill-Nerode Theorem.

Say that for some  $w \in \sum^*$ ,  $xw \in L$ ,  $yw \notin L$ , then  $x, y$  are in different states.

$$L \in \sum^*, x \in \sum^*$$

$$\text{AcceptingFuture}(L, x) = \{w \in \sum^*\}$$

$$L = \{1^m \mid 7 \text{ divides } m\}$$

$$\text{AcceptingFuture}(L, \epsilon) = (1^7)^*$$

$$\text{AcceptingFuture}(L, 1) = 1^6(1^7)^*$$

$$\text{AcceptingFuture}(L, 1^2) = 1^5(1^7)^*$$

$$\text{AcceptingFuture}(L, 1^3) = 1^4(1^7)^*$$

$$\text{AcceptingFuture}(L, 1^4) = 1^3(1^7)^*$$

$$\text{AcceptingFuture}(L, 1^5) = 1^2(1^7)^*$$

$$\text{AcceptingFuture}(L, 1^6) = 1(1^7)^*$$

Thus, since there are at least 7 distinct accepting futures, by the Myhill-Nerode Theorem, the DFA representing this language must have at least 7 states.

### 4.b

Pumping Lemma: Say that  $L$  is regular and accepted by a DFA of  $p$  states or fewer. Then if  $s \in L$  and  $|S| \geq p$  we can write  $S = xyz$  such that

1.  $xz, xyz, xy^2z, xy^3z, \dots \in L$
2.  $y \neq \epsilon$
3.  $|xy| \leq p$

Claim:  $L = \{1^m \mid m \text{ is a perfect square}\}$  is not regular.

Proof: Say  $L$  is regular and accepted by a DFA of  $p$  states.

Now consider  $S = 1^p \in L$ . Using the pumping lemma, let  $s = xyz$  such that 1-3 above hold. Since  $S = 1^{2p} = xyz$  and  $|xy| \leq p$ , we have  $x = 1^a, y = 1^b, z = 1^{2p-a-b}$

Thus,  $\forall i \in \mathbb{N}$ , such that  $1^{a+bi}1^{2p-a-b} \in L$ .

$1^{2p+b(i-1)}$  cannot be in  $L$  since there is no constant  $b$  such that  $\forall i \in \mathbb{N}$ ,  $2p+b(i-1)$  is a perfect square. That would imply perfect squares are separated by a constant factor. Proof by contradiction.  $L$  must not be regular.