Question 1 [50 points]: "If it's 'dynamic', it must be better."

The following rules define an environment-based semantics for lexically-scoped functions, Lam, and dynamically-scoped functions, Ds-lam.

 $env \vdash e \Downarrow v$ Under environment env, expression e evaluates to value v

$$\frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Num \ n) \Downarrow (Num \ n)} \xrightarrow{Env-num} \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} \xrightarrow{Env-add} \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} = \frac{env \vdash e1 \Downarrow (Num \ n1+n2)}{env \vdash (Add \ e1 \ e2) \Downarrow (Num \ n1+n2)} =$$

$$\frac{\textit{env} \vdash \textit{e1} \Downarrow \textit{v1} \qquad \textit{x=v1}, \textit{env} \vdash \textit{e2} \Downarrow \textit{v2}}{\textit{env} \vdash (\mathsf{Let} \; \textit{x} \; \textit{e1} \; \textit{e2}) \Downarrow \textit{v2}} \; \mathsf{Env\text{-let}}$$

$$\frac{lookup(env,x)=e}{env \vdash (\operatorname{Id} x) \Downarrow e} \text{ Env-id } \frac{lookup(env,x) \text{ undefined}}{env \vdash (\operatorname{Id} x) \text{ unknown-id-error}} \text{ Env-unknown-id-error}$$

$$\frac{env \vdash (\mathsf{Lam} \ x \ e1) \Downarrow \big(\mathsf{Clo} \ env \ (\mathsf{Lam} \ x \ e1)\big)}{env \vdash (\mathsf{Clo} \ env_{old} \ e) \Downarrow v} \ \mathsf{Env\text{-}clo}$$

$$\frac{\textit{env} \vdash \textit{e1} \Downarrow (\textit{Clo} \; \textit{env}_{old} \; (\textit{Lam} \; \textit{x} \; \textit{eB})) \qquad \textit{env} \vdash \textit{e2} \Downarrow \textit{v2} \qquad \textit{x=v2}, \textit{env}_{old} \vdash \textit{eB} \Downarrow \textit{v}}{\textit{env} \vdash (\textit{App} \; \textit{e1} \; \textit{e2}) \Downarrow \textit{v}} \; \texttt{Env-app}$$

$$\frac{env \vdash (\mathsf{Ds\text{-}lam} \ x \ e1) \Downarrow (\mathsf{Ds\text{-}lam} \ x \ e1)}{env \vdash (\mathsf{Ds\text{-}lam} \ x \ e1)} \underbrace{\mathsf{Env\text{-}ds\text{-}lam}}_{\mathsf{Env\text{-}ds\text{-}app}} \underbrace{\frac{env \vdash e1 \Downarrow (\mathsf{Ds\text{-}lam} \ x \ eB)}{env \vdash (\mathsf{App} \ e1 \ e2) \Downarrow v}}_{\mathsf{Env\text{-}ds\text{-}app}} \underbrace{\frac{env \vdash e1 \Downarrow (\mathsf{Ds\text{-}lam} \ x \ eB)}{env \vdash (\mathsf{App} \ e1 \ e2) \Downarrow v}}_{\mathsf{Env\text{-}ds\text{-}app}} \underbrace{\frac{env \vdash eB \Downarrow v}{env \vdash eB \Downarrow v}}_{\mathsf{Env\text{-}ds\text{-}app}}$$

Assume that lookup(env, x) returns the **leftmost** binding of x. For example:

$$lookup((x=(Num 2), x=(Num 1), \emptyset), x) = (Num 2)$$

Consider the following expression, shown in concrete syntax (left) and in abstract syntax (right).

$$\left(\text{Let y (Num 100)} \right)$$

$$\left(\text{Let f (Let y (Num 10)} \right)$$

$$\left(\text{Lam x (Add (Id x) (Id y)))} \right)$$

$$\left(\text{Let y (Num 2)} \right)$$

$$\left(\text{App (Id f) (Id y))} \right)$$

Q1a [10 points] Complete the abstract syntax tree for the above expression.

(Writing Add I) is also okay) extra lines