

1. (20) Multiple choice, 1pt each. For each question circle the **best** answer.

Q1. Numerical algorithms should be

- a. efficient
- b. accurate
- c. reliable

Xd. all of the above

Q2. In numerical analysis we have to analyze the following types of errors:

- a. design errors
- Xb. roundoff errors
- c. hardware errors
- d. timing errors

Q3. The error in  $x$  has to be given in absolute form if:

- a.  $x$  is unknown
- Xb.  $x$  is zero
- c.  $x$  can be negative
- d.  $x$  can overflow

Q4. If  $x = 10$  with relative error 0.01, what is the absolute error in  $x^2/2$ ?

- Xa. 1
- b. 0.5
- c. 0.0005
- d. 0.1

Q5. Which of the following iterative methods covers fastest under normal circumstances:

- Xa. Newton's method
- b. Secant method
- c. Bisection method
- d. Golden section search method

Q6. For some root finding method the error  $e_n$  satisfies  $e_n = 0.1 \times e_{n-1}^3$ . Which statement is true:

- a. The order of the method is  $0.1^3$
- b. The convergence rate is  $-\log_{10}(0.1)$
- c. The method is strictly convergent

Xd. The method is superlinearly convergent

- Q7. The bisection method
- a. converges faster than the secant method
  - b. is in between Newton's and the secant method in speed
  - c. is used for optimization
  - Xd. can converge when Newton's method does not
- Q8. In large scale realistic applications the dominant computational cost in a root finding method usually is
- a. the selection of the initial bracketing interval  $[a \ b]$
  - Xb. evaluating the function and possibly its derivative
  - c. computing the termination criterium
  - d. switching to the bisection method when an iteration fails
- Q9. A matrix  $A$  is singular if
- a. it is not square
  - b. Gaussian elimination with partial pivoting fails
  - c. it is its own inverse
  - Xd. it has no inverse
- Q10. The p-norm of vector  $(x_1, \dots, x_n)$  is
- a.  $|x_1 + x_2 + \dots + x_p|$
  - b.  $\text{length}(x)^p$
  - c.  $|x_1|^p + |x_2|^p + \dots + |x_n|^p$
  - Xd.  $(|x_1|^p + |x_2|^p + \dots + |x_n|^p)^{1/p}$
- Q11. The induced matrix norm ( $\|\cdot\|$ ) satisfies for matrix  $A$  and vector  $x$
- a.  $Ax \leq \|Ax\|$
  - b.  $\|Ax\| = \|A\|\|x\|$
  - c.  $\|A\| + \|x\| > 0$
  - Xd.  $\|Ax\| \leq \|A\|\|x\|$
- Q12. For a general nonsingular matrix  $A$  of moderate size the equation  $Ax = b$  can best be solved using:
- Xa. Gaussian elimination with scaled partial pivoting
  - b. Gaussian elimination with complete pivoting
  - c. Gaussian elimination with partial pivoting
  - d. Direct inversion

- Q13. If Gaussian elimination for a matrix equation  $Ax = b$  without pivoting encounters a division by zero
- a.  $A$  is singular
  - b. the system has multiple solutions
  - c. the equation is ill-conditioned
  - Xd. nothing can be concluded about  $A$
- Q14. The operation count of solving an upper or lower triangular system of linear equations size  $n$  (i.e.,  $n \times n$  matrix) is about
- Xa.  $n^2$
  - b.  $2n^2$
  - c.  $n^3$
  - d.  $n$
- Q15. The complexity of Gaussian elimination with partial pivoting is
- a.  $O(n)$
  - b.  $O(n^2)$
  - Xc.  $O(n^3)$
  - d.  $O(n^4)$
- Q16. The complexity of computing the  $LU$  factorization of a matrix  $A$  is
- a.  $O(n)$
  - b.  $O(n^2)$
  - Xc.  $O(n^3)$
  - d.  $O(n^4)$
- Q17. If we have to solve  $Ax = b$  many times for different vectors  $b$  but the same  $A$  we
- a. compute the inverse of  $A$
  - Xb. use  $LU$  decomposition once, then forward backward substitution for the different  $b$ 's
  - c. use iterative methods
  - d. use Gaussian elimination
- Q18. If  $A$  has no  $LU$  factorization then
- a.  $A$  is singular
  - b. every matrix has a  $LU$  factorization
  - c.  $A$  requires complete pivoting
  - Xd. we can't conclude anything about  $A$

- Q19. The condition number of a matrix  $A$  is used to
- a. speed up pivoting
  - b. determine how many steps pivoting takes
  - c. determine the relation between error in  $A$  and error in  $b$
  - Xd. estimate the accuracy of Gaussian elimination
- Q20. Would you like a free point?
- Xa. Yes
  - b. No
  - c. Don't care
  - d. None of the above

2. (10) Consider the following finite difference approximation to the derivative of a function  $f$  at  $x$

$$f'(x) \approx \frac{f(x+h) - f(x-h)}{2h}. \quad (1)$$

- a. (3) Assuming that the (roundoff) absolute error in computing  $f$  is  $\epsilon$ , what is the resulting AE in equation 1?

Each term in the numerator has error  $\epsilon$  so the sum has error  $2\epsilon$  so we have roundoff error

$$\epsilon/h$$

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- b. (4) Derive a formula for the absolute discretization error of the approximation, valid for small  $h$ .

Taylor expansion of  $f(x+h)$  and  $f(x-h)$  gives

$$f(x+h) \approx f(x) + f'(x)h + \frac{1}{2}f''(x)h^2 + \frac{1}{6}f'''(x)h^3 + O(h^4)$$

and

$$f(x-h) \approx f(x) - f'(x)h + \frac{1}{2}f''(x)h^2 - \frac{1}{6}f'''(x)h^3 + O(h^4).$$

Substitute in eq. refeq and we get

$$\frac{f(x+h) - f(x-h)}{2h} = f'(x) + \frac{1}{6}f'''(x)h^2 + O(h^3).$$

So discretization error is about

$$\frac{1}{6}f'''(x)h^2.$$

- c. (3) Describe (in words) what happens to the overall accuracy of the approximation if we make  $h$  smaller and smaller.

$$\text{Total error} = \epsilon/h + \frac{1}{6}f'''(x)h^2.$$

If I decrease  $h$  the discretization error gets smaller but the roundoff error gets bigger. If  $\epsilon$  is small the decrease in discretization error will make the total error go down at first, until  $h$  gets so small that the roundoff error becomes important, after which the error will increase.