





Overview of advanced storage technologies and storage virtualization

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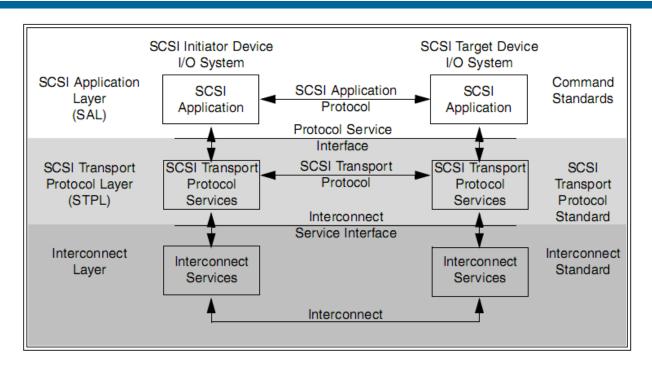
SCSI Interface

(SCSI-3 specification)

SCSI-3 History

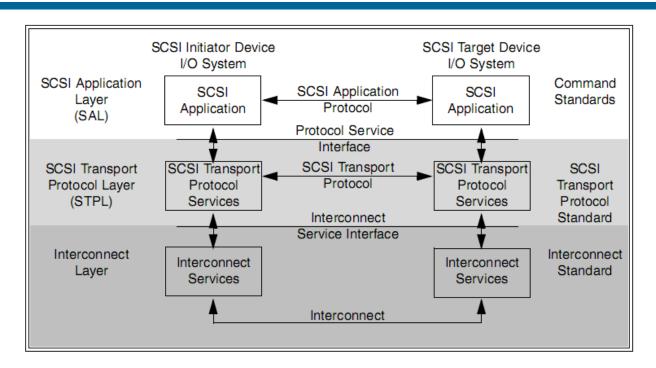
- SCSI interface created in 1978 by Shugart Associates, first standarized 1981
 - **SCSI** = Small Computers System Interface
- Developed as logical interface to drive high-performance disks over 8-bit SCSI parallel bus
- Very successful and widely accepted
 - Extended to control all kinds of storage peripherals
- Around 1995 begins development of SCSI-3 standard, published in 2003
- SCSI-3 is no longer a single standard, but a "standards family" (set of interrelated standards)
 - Suffix "-3" has been dropped, so nowadays "SCSI" means "SCSI-3"

SCSI layered protocol stack



- SCSI-3 models SCSI interface as objects, protocol layers and service interfaces
 - Concept very similar to OSI computer networks reference model
- SCSI objects (protocol entities) exchange frames between peers
 - Frames get encapsulated when travelling downwards through protocol stack and de-encapsulated when travelling upwards

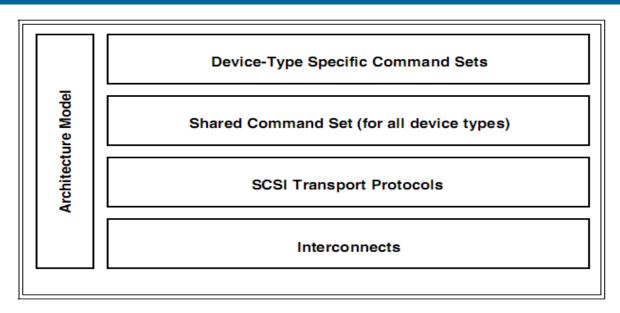
SCSI layered protocol stack



Model has three layers:

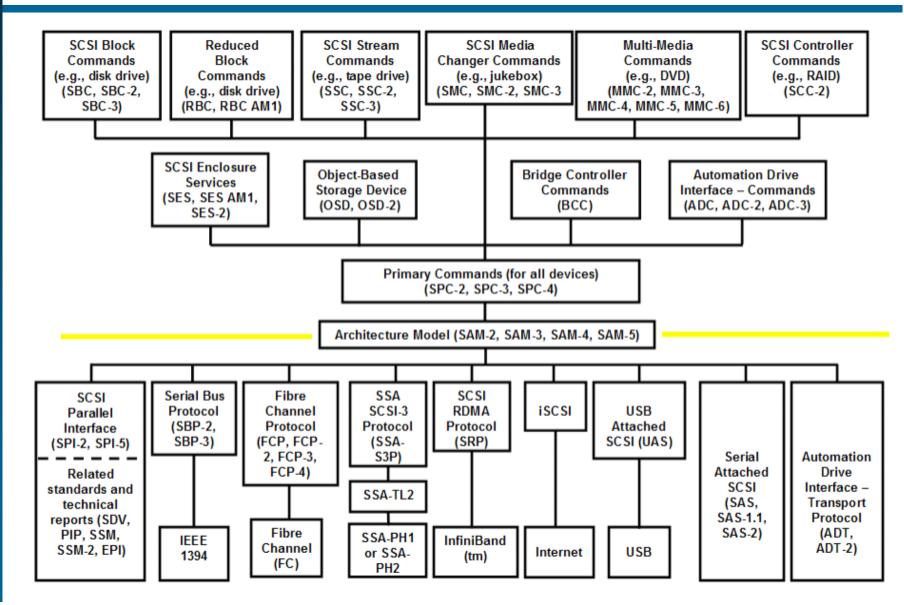
- Application layer: Commands sent from initiator to target, answers from initiator with data and status info
- <u>Transport layer:</u> Protocols to wrap SCSI payload within link-level frames
- <u>Interconnect layer:</u> Essentially, equivalent to link and physical layers of OSI model

SCSI standards family

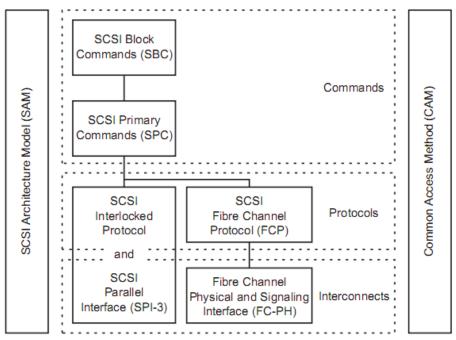


- SCSI-3 standards family defines:
 - A common architecture model, standardizing concepts and behaviors in all layers
 - Set of commands for the application layer
- Commands organized in two different subsets:
 - Primary commands: MUST be implemented by any SCSI device
 - Basic device identification and management
 - Device-specific sets of commands

SCSI standards family

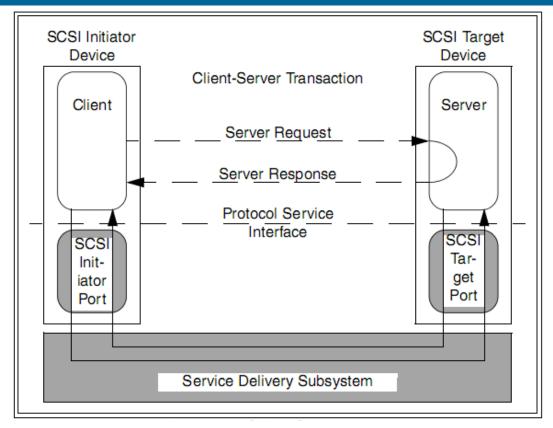


SCSI standards family



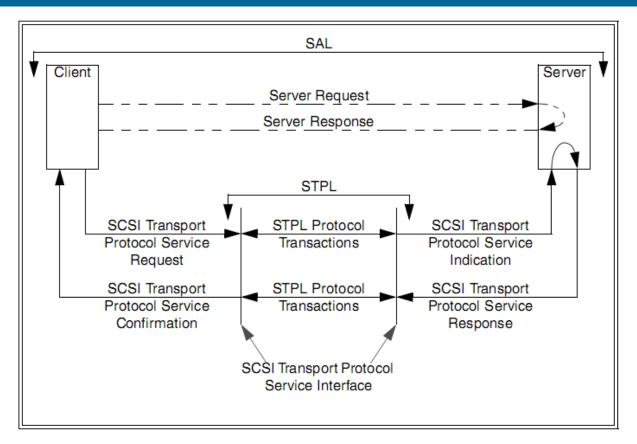
- SCSI-3 protocol stack model allows sending exactly the same commands over different buses or storage networking interconnects
 - Just transport-level wrapping of commands/answers, and link-level behavior, need to be changed
 - Allows also bidirectional payload transfer through two or more different transports between initiator and target

SCSI client-server model



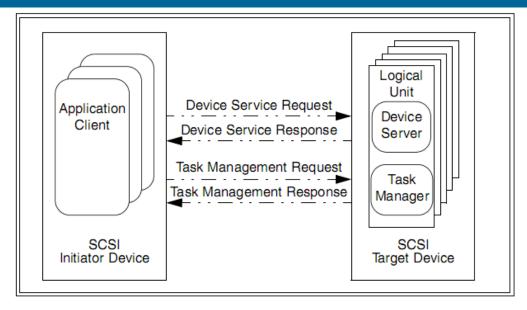
- SCSI I/O transaction between initiator and target is modelled as a client-server transaction
 - Server (= target = storage device) offers service to client (= initiator = HBA)
 - Service = access to write or retrieve data, get status or manage device

SCSI client-server model



- SCSI command is request from service sent from client (initiator) to server (target)
- Transport and interconnect layer perform physical transport of request and response

LUNs and tasks



- SCSI target (server) exposes services (functionality) through Logical Units
 - Logical Unit = object that processes commands (ex: read from media)
- Application client (driver in initiator) issues request (command) to Logical Unit
 - Command defines unit of work to be performed by Logical Unit
 - This unit of work is named task in SCSI parlance
 - Thus, initiator issues tasks (= commands) to Logical Unit, which processes them

Command Descriptor Block (CDB)

Typical CDB for 6-byte commands

Bit Byte	7	6	5	4	3	2	1	0	
0		OPERATION CODE							
1	Miscella	Miscellaneous CDB information (MSB)							
2									
3		LOGICAL BLOCK ADDRESS (if required) (LSB)							
4	TRANSFER LENGTH (if required) PARAMETER LIST LENGTH (if required) ALLOCATION LENGTH (if required)								
5		CONTROL							

- Initiator sends command to Logical Unit issuing a CDB (Command Descriptor Block)
 - Structured set of bytes, typically of 6, 10, 12 or 16 bytes
 - Interpreted by target as request to execute a given command, and parameters to it

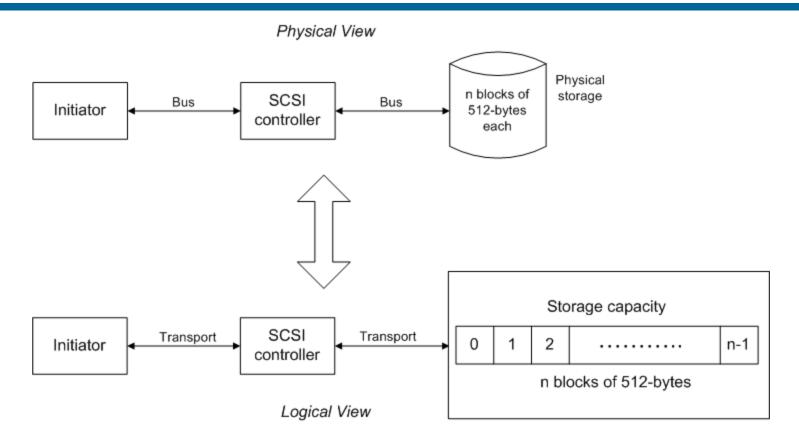
Command Descriptor Block (CDB)

Typical CDB for 10-byte commands

Bit Byte	7	6	5	4	3	2	1	0	
0	OPERATION CODE								
1	Miscellaneous CDB information SERVICE ACTION (if required)								
2	(MSB)	(MSB)							
3		LOCICAL BLOCK ADDDECC (if an entire d)							
4	LOGICAL BLOCK ADDRESS (if required)								
5		(LS						(LSB)	
6			M	iscellaneous (CDB information	on			
7	(MSB)								
8		-	PARAMETER LIST LENGTH (if required) ALLOCATION LENGTH (if required) (LSE						
9		CONTROL							

- Different size CDBs for same command differ in length of parameters
 - Larger address allows addressing larger storage volumes

Logical Block Addressing



- SCSI controller hides physical details of storage organization
 - Storage capacity is shown just as a linear array of individuallyaddressable blocks

Logical Block Addressing

Typical CDB for 6-byte commands

Bit Byte	7	6	5	4	3	2	1	0	
0		OPERATION CODE							
1	Miscellar	Miscellaneous CDB information (MSB)							
2	LOCION PLOCK APPERENCE (II								
3		LOGICAL BLOCK ADDRESS (if required) (LSB)							
4	TRANSFER LENGTH (if required) PARAMETER LIST LENGTH (if required) ALLOCATION LENGTH (if required)								
5		CONTROL							

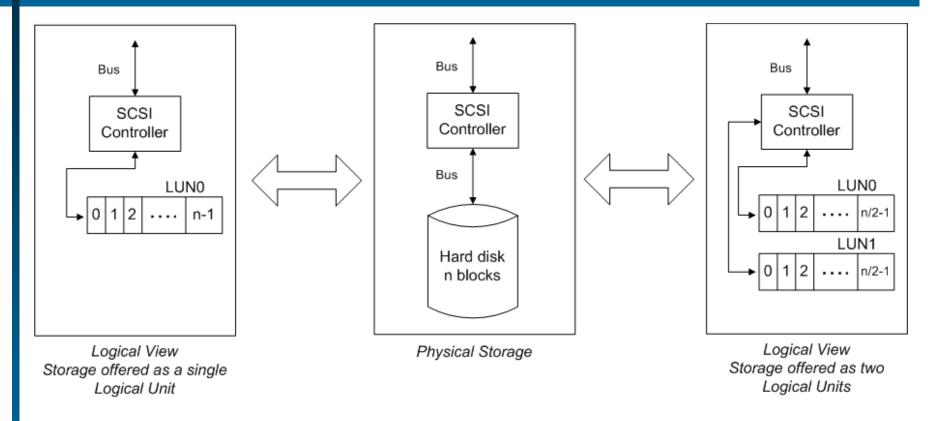
- Blocks to be accessed are addressed through 2 parameters in CDB:
 - Address of first logical block of transfer
 - Length, in blocks, of transfer
 - > Typical block size = 512 bytes

Logical Block Addressing

Bit Byte 7 6 5 4 3 2 1 0 0 OPERATION CODE 1 Miscellaneous CDB information 2 (MSB) 3 4 5 6 7 8 9 (LSE 10 (MSB) TRANSFER LENGTH (If required) PARAMETER LIST LENGTH (If required)		Typical CDB for long LBA 16-byte commands										
1		7	6	5	4	3	2	1	0			
2 (MSB) 3 4 5 6 7 8 9 (LSE 10 (MSB) TRANSFER LENGTH (If required)	0		OPERATION CODE									
3 4 5 6 COGICAL BLOCK ADDRESS 10 (LSE 10 (MSB) TRANSFER LENGTH (If required)	1		Miscellaneous CDB information									
4	2	(MSB)										
5 LOGICAL BLOCK ADDRESS 7 8 9 (LSE 10 (MSB) 11 TRANSFER LENGTH (If required)	3		_									
6 7 8 9 (LSE 10 (MSB) TRANSFER LENGTH (If required)	4]										
6	5											
8 9 (LSE 10 (MSB) TRANSFER LENGTH (If required)	6		LOGICAL BLOCK ADDRESS									
9 (LSE 10 (MSB) 11 TRANSFER LENGTH (If required)	7											
10 (MSB) 11 TRANSFER LENGTH (If required)	8											
11 TRANSFER LENGTH (If required)	9								(LSB)			
Trouter Ett EEtte Tri (in required)	10	(MSB)										
PARAMETER LIST LENGTH (if required)	11		_									
12 ALLOCATION LENGTH (if required)	12											
13 (LSE	13								(LSB)			
14 Miscellaneous CDB information	14			M	liscellaneous (CDB information	on					
15 Control	15				Cor	ntrol						

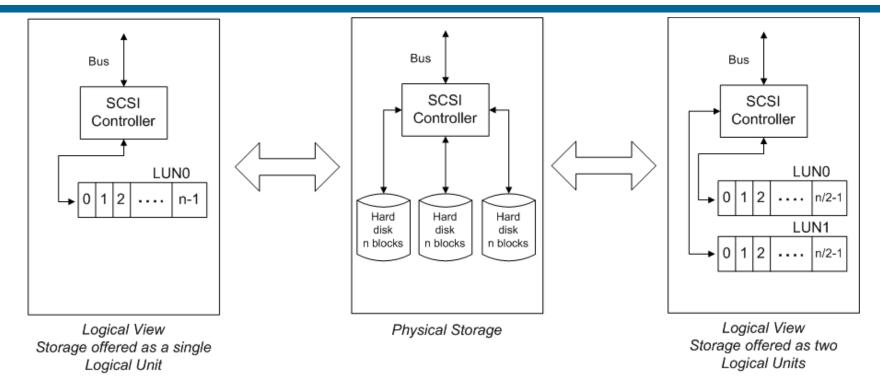
- Allowing larger address fields, SCSI interface copes with increasingly larger storage volumes:
 - CDB-6: 2^{21} blocks = 1 GByte
 - CDB-10: 2^{32} blocks = 2 TBytes
 - > Typical LUN size limit for many SCSI controllers and virtualization software
 - CDB-16: 2⁶⁴ blocks = ~ 9 ZetaBytes

LUNs and LBA give flexibility



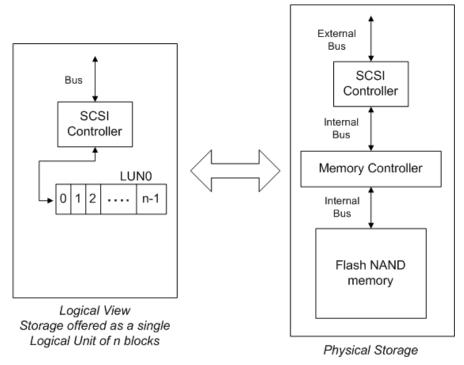
- Logical Units allow greater flexibility to organize storage
 - Controller just needs to map physical blocks to logical addresses
 - Required unique numerical identifier for each Logical Unit
 - > **LUN** = Logic Unit Number

LUNs and LBA hide complexity



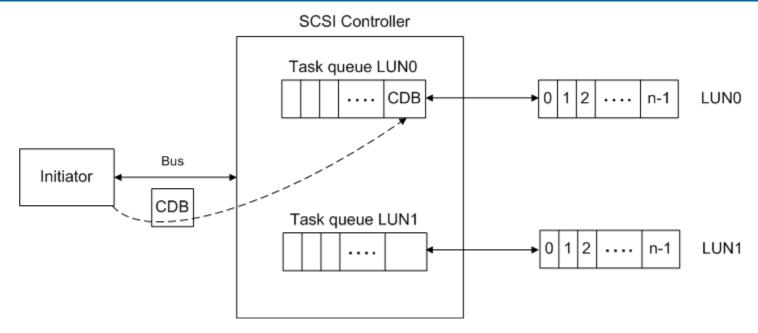
- Logical Units hide complexity of hardware setup
 - Left: controller shows several physical disks as single, larger disk
 - LBA hides physical placement of data
 - Controller free to transparently optimize placement
 - Right: controller sums total storage as split in several LUNs
 - Allow transparency of complex placement strategies (RAID, replication)

LUNs and LBA hide complexity



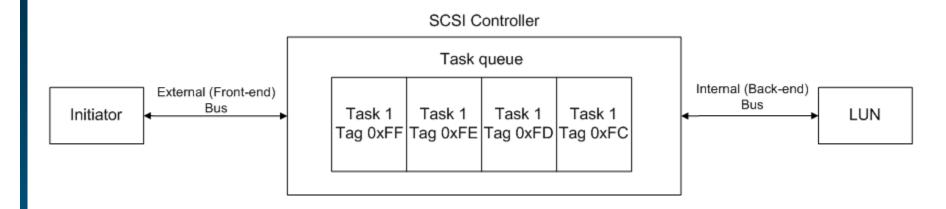
- Logical Units also hide nature of storage
 - Example: Flash storage (SSD, USB drives) has a complex internal architecture
 - Blocks are dynamically migrated for wear-leveling
 - Simple external view as a single LUN of n blocks
 - Memory controller manages mapping between logical LBA address and physical memory page

LUNs and task queues



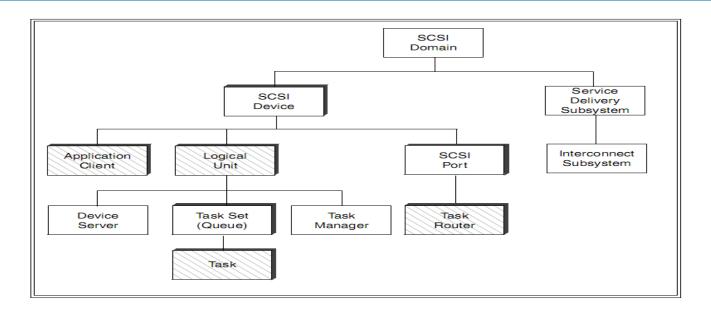
- SCSI controller can offer for each LUN none, one or several task queues to manage pending commands (*tasks*)
 - CDB inserted into queue, and can be executed asynchronously
 - If several commands in queue, can be reordered before execution, if hardware supports it
- If no queue, commands must execute using synchronous I/O (only a single command pending at any time)

LUNs and task queues



- Use of task queue requires identifying each individual task
 - Each task is assigned numerical identifier, unique within queue
 - Tag provides context reference for task when executing
- Tags are not required in synchronous I/O ("untagged commands")

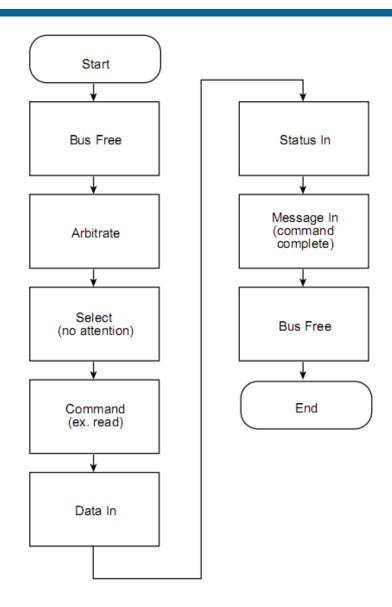
SCSI task addressing



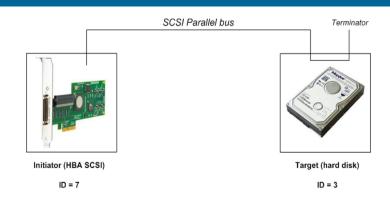
- Issuing or answering SCSI task requires use of several numerical identifiers:
 - Initiator ID and Target ID = Identifiers of physical devices involved
 - Logical Unit Number (LUN) = Identifier or logical unit, of target device,
 which must process this command
 - Task queue = Identifier of queue, within LU, in which command is inserted
 - Task tag = unique identifier for this task within queue

Simplest way to understand operation of SCSI interface

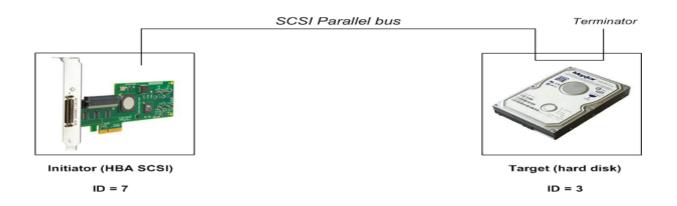
- SCSI interface executes command as an ordered sequence of phases
 - Task will end with sending of status info on outcome of command



- Figure contains decoding
 of real traffic capture (with
 logic analyzer) for SCSI
 WRITE(10) command
 between initiator and
 target shown above
 - Write: data goes from initiator to disk



Time		Time		STAT	DATA-8	
4554214772	ns	151.61	us	ARBITRATION	80	
4554215356	ns	0.584	us	SELECT	88	
4554218924	ns	3.568	us	MSG OUT	CO	
4554220922	ns	1.998	us	MSG OUT	20	
4554223122	ns	2.2	us	MSG OUT	FE	
4554228472	ns	5.35	us	COMMAND	2 A	
4554228896	ns	0.424	us	COMMAND	00	
4554229172	ns	0.276	us	COMMAND	00	
4554229448	ns	0.276	us	COMMAND	03	
4554229722	ns	0.274	us	COMMAND	22	
4554229998	ns	0.276	us	COMMAND	FF	
4554230272	ns	0.274	us	COMMAND	00	
4554230548	ns	0.276	us	COMMAND	00	
4554230822	ns	0.274	us	COMMAND	80	
4554231096	ns	0.274	us	COMMAND	00	
4554236148	ns	5.052	us	DATA OUT	00	
4554236248	ns	0.1	us	DATA OUT	00	
4554236350	ns	0.102	us	DATA OUT	20	
4560918946	ns	6682.596	us	STATUS	00	
4560920940	ns	1.994	us	MSG IN	00	



4554214772 ns 151.61 us ARBITRATION 80 = 1000 0000 4554215356 ns 0.584 us SELECT 88 = 1000 1000

- SCSI transaction starts addressing physical devices involved
 - For parallel SCSI, addressing done in ARBITRATION stage
 - Parallel SCSI is multi-drop, so bus-master and bus-slave must be chosen
 - HBA (ID = 7) wins arbitration and becomes bus master
 - HBA selects then hard disk (ID = 3) as bus slave (peer for this transfer)

IDENTIFY message format

Bit Byte	7	6	5	4	3	2	1	0
0	IDENTIFY	DISCPRIV			LU	JN		

Task attribute message codes

		Suj			
Code	IU Transfers Disabled		IU Transfers Enabled		Message name
	Iniator	Target	Initiator Target		
24h	0	0	N/A	N/A	ACA
21h	Q	Q	N/A	N/A	Head of Queue
22h	Q	Q	N/A N/A		Ordered
20h	Q	Q	N/A	N/A	Simple

4554218924 ns 3.568 us MSG OUT CO LUN = 0
4554220922 ns 1.998 us MSG OUT 20 Simple queue
4554223122 ns 2.2 us MSG OUT FE TAG

- SCSI transaction needs now address LU, queue and command tag
 - Initiator addresses LUN0 of target device (hard disk)
 - Command is inserted into "Simple" queue
 - Command is uniquely identified within simple queue by tag 0xFE

	WRITE (10) command						
Bit Byte	7	6	5	4	3	2	1	О
0		•		OPERATION	CODE (2Ah)		•	•
1		WRPROTECT	WRPROTECT DPO FUA Reserved FUA_NV					Obsolete
2	(MSB)							
5			LOGICAL BLOCK ADDRESS (LSB)					
6		Reserved	Reserved GROUP NUMBER					
7	(MSB)	_	TRANSFER LENGTH —					
8			TRANSFER LENGTH					
9		CONTROL						
45542284	554228472 ns 5.35 us COMMAND 2A 🚈 WRITE					(10) codop		
45542288			24 us		MAND	00	*********************************	-(10) codop
45542291			76 us		MAND	00 <==	─ LBA (N	/ISB)
45542294	448 ns	0.2	76 us	COM	MAND	03		
45542297	722 ns	0.2	74 us	COM	MAND	22		
45542299	998 ns	0.2	76 us	COM	MAND	FF <	─ LBA (l	_SB)
45542302	272 ns	0.2	74 us	COM	MAND	00 `		
45542305	548 ns	0.2	76 us	COM	MAND	00 <==	Transfe	er Length (MS
45542308	322 ns	0.2	74 us	COM	MAND	80 ⊱	Transfe	er Length (LS
45542310	096 ns	0.2	74 us	COM	MAND	00 `		

- Initiator sends now CDB to target
 - Target interprets CDB and gets ready to write 0x80 blocks (64 KB) data in media, starting from specified LBA

Status byte code bit values

Status byte	Status represented	Task Ended
00h	Good	Yes

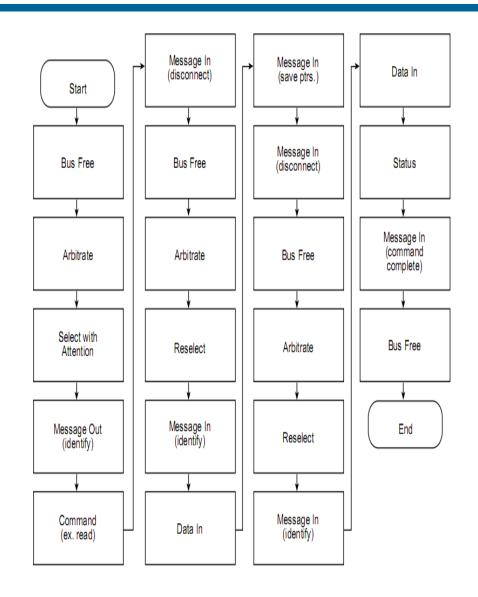
Message format

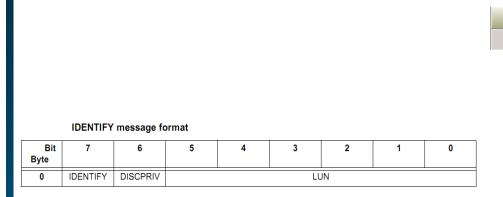
Message code	Message format
00h	One-byte message (TASK COMPLETE)

ns 5.052	นธ	DATA OUT	00 \ Start of data transfer
ns 0.1	นธ	DATA OUT	00
ns 0.102	นธ	DATA OUT	20 End of data transfer
ns 6682.596	นธ	STATUS	00 Command ended OK
ns 1.994	us	MSG IN	00 Cash Task completed
	ns 0.1 ns 0.102 ns 6682.596	ns 0.1 us ns 0.102 us ns 6682.596 us	ns 0.1 us DATA OUT ns 0.102 us DATA OUT ns 6682.596 us STATUS

- Initiator sends data to target (disk)
- Target ends transaction by:
 - Reporting command executed OK (Status Good)
 - Declaring command as finished (Task Complete)

- SCSI interface allows split transactions
 - Target can disconnect after command, and reconnect to start data transfer
 - Data transfer itself can be interrupted and resumed, if needed
- Coupled with tagged queuing, provides high flexibility for optimizing command execution
 - Allows command reordering
 - Avoids wasting time idling the bus





Time		Time		STAT	DATA-8
4547006302	ns	0.584	us	SELECT	88
4547009868	ns	3.566	us	MSG OUT	CO
4547011968	ns	2.1	us	MSG OUT	20
4547014268	ns	2.3	us	MSG OUT	FE
4547019518	ns	5.25	us	COMMAND	2 A
4547019918	ns	0.4	us	COMMAND	00
4547020194	ns	0.276	us	COMMAND	00
4547020468	ns	0.274	us	COMMAND	03
4547020794	ns	0.326	us	COMMAND	22
4547021118	ns	0.324	us	COMMAND	7F
4547021418	ns	0.3	us	COMMAND	00
4547021718	ns	0.3	us	COMMAND	00
4547021994	ns	0.276	us	COMMAND	80
4547022268	ns	0.274	us	COMMAND	00
4547024850	ns	2.582	us	MSG IN	04
4547339692	ns	314.842	us	ARBITRATION	08
4547340602	ns	0.91	us	RESELECT	88
4547342134	ns	1.532	us	MSG IN	80
4547356896	ns	14.762	us	MSG IN	20
4547358492	ns	1.596	us	MSG IN	FE
4547378766	ns	20.274	us	DATA OUT	00
4547378868	ns	0.102	us	DATA OUT	00
4547378968	ns	0.1	us	DATA OUT	00
4554061068	ns	6682.1	us	STATUS	00
4554063162	ns	2.094	us	MSG IN	00

- Figure shows single SCSI task (command), split in two I/O transactions
- Note that LUN addressing from initiator (MSG OUT C0) allows also DISCONNECT privilege
 - Essentially, tells target that initiator can store in RAM CDB, pointers and status info to resume this command, if needed

	ı	1						_
		Support						
		IU tran disabl		IU transfers enabled				
	Code	Init	Targ	Init	Targ	Message name		_
	04h	0	0	0	0	DISCONNECT		-
454702	4850 ns	2	.582	นธ	-	MSG IN	04	
454733	9692 ns	314	1.842	นธ	ARBIT	RATION	08 = 00	00 1000
454734	10602 ns		0.91	us	RE	SELECT	88 = 10	00 1000
454734	12134 ns	1	1.532	us		MSG IN	80 💳	LUN = 0
454735	6896 ns	14	1.762	us		MSG IN	20 <	Simple queue
454735	8492 ns	1	L.596	us		MSG IN	FE ⊨	TAG
454737	78766 ns	20	.274	us	DA	TA OUT	00 `	

- After last byte of CDB, target disconnects from initiator
 - Task still active and pending
- When ready for transfer, target reconnects
 - Wins arbitration and becomes bus-master (!)
 - Re-selects initiator, which becomes bus-slave (!).
 - > SPI functional equivalent to raising interrupt
 - Note that only <u>bus role</u> is reversed, NOT interface role

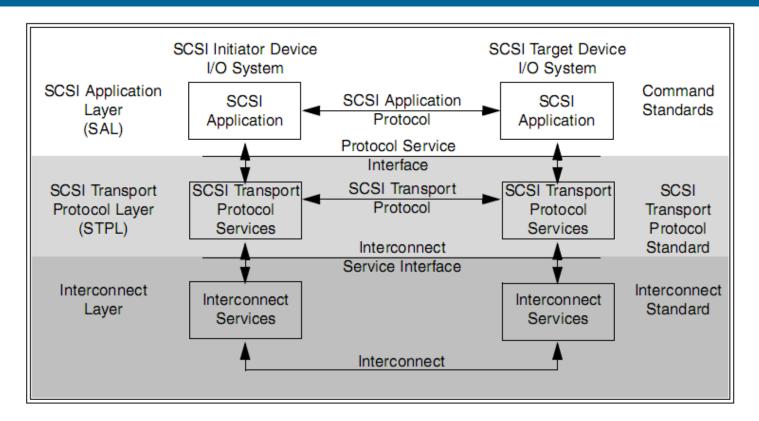
Link Control message codes

	Suppo	ort			
	IU tran		IU transfers enabled		
Code	Init	Targ	Init	Targ	Message name
04h	0	0	0	0	DISCONNECT

	04	MSG IN	นธ	2.582	ns	4547024850
0000 1000	08 = 0000	ARBITRATION	นธ	314.842	ns	4547339692
1000 1000	88 = 1000	RESELECT	นธ	0.91	ns	4547340602
LUN = 0	80 <	MSG IN	us	1.532	ns	4547342134
Simple queue	20 ⊨	MSG IN	us	14.762	ns	4547356896
TAG	FE <	MSG IN	us	1.596	ns	4547358492
	00 `	DATA OUT	นธ	20.274	ns	4547378766

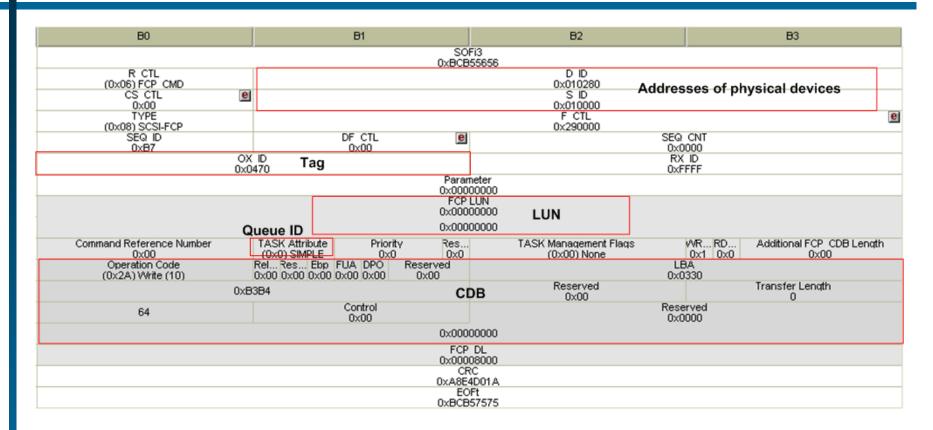
- Target sends now messages to identify task being resumed
 - Addresses LUN, queue and tag
- Disk accepts now data from initiator
- Command ends with Status Good (see end of complete transaction)

Example: WRITE over other transports



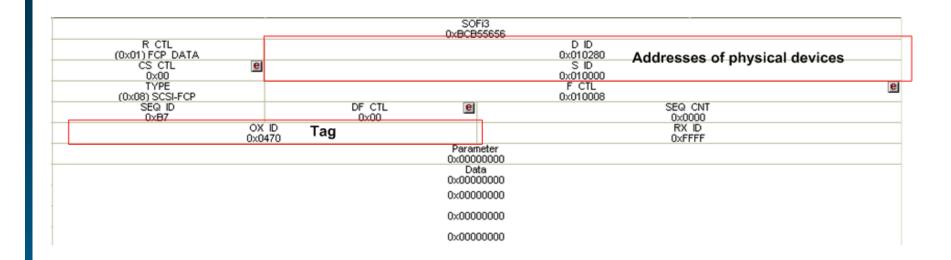
- Layered stack model means that SCSI protocol works the same way over any transport
 - Once understood over parallel SCSI (SPI), it is easy to understand its operation over any other bus

Example: WRITE over Fibre Channel



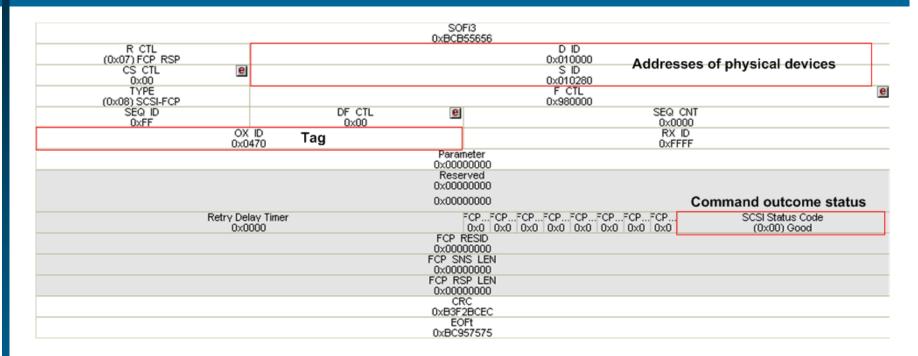
- Figure shows SCSI WRITE(10) command encapsulated within a Fibre Channel frame
 - SCSI protocol payload is straightforward to identify and interpret

Example: WRITE over Fibre Channel



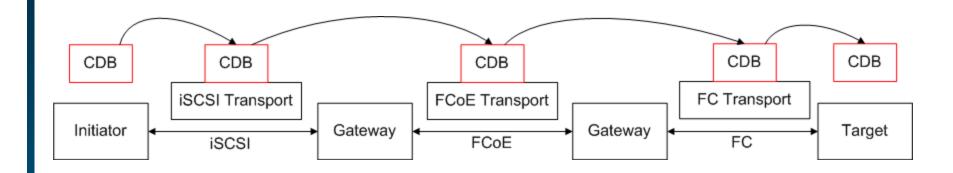
- Figure shows data transfer belonging to that same split transaction
 - Physical identifiers plus tag are enough to recover context of current task

Example: WRITE over Fibre Channel



- Figure shows frame that completes SCSI protocol for WRITE command
 - Again, physical identifiers plus tag allow to recover context of current task
 - Status code is same as in parallel SCSI

Summary: SCSI gives great flexibility



- SCSI protocol payload can be hopped from transport to transport with no changes
 - Encapsulation/De-encapsulation is straightforward and reasonably fast
- SCSI allows transaction between initiator and target while traversing several different buses/storage networking infrastructures
 - Ideal for accessing real storage from virtual machines, or for virtualizing storage