

# CIS 575. Introduction to Algorithm Analysis

## Material for April 17, 2024

### Algorithms Based on Depth-First Search

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The topic of this note is part of what is covered in *Cormen's* Section 20.3.

## 1 Generic Depth-First Search

We shall *augment* our algorithm for depth-first search so that it can take various kinds of action at various points: when a node has just been discovered; when a tree edge is about to be explored; when a tree edge has been explored; when a non-tree edge is encountered; when a node is finished. We get the algorithm

```
DFS( $u$ )  //  $u$  is white
     $color[u] \leftarrow \text{gray}; d[u] \leftarrow \text{current\_time}$ 
    PreNode action
    foreach  $w$  where  $E$  contains an edge from  $u$  to  $w$ 
        if  $color[w] = \text{white}$ 
            PreEdge action
             $T \leftarrow T \cup \{u \rightarrow w\}$ 
            DFS( $w$ )
            PostEdge action
        else
            OtherEdge action
     $color[u] \leftarrow \text{black}; f[u] \leftarrow \text{current\_time}$ 
    PostNode action
```

By suitable instantiation of the 5 kinds of action, numerous useful graph algorithms may be constructed. In the subsequent notes, we shall present two such instantiations: **Topological sort** for **directed** graphs (the topic of *Cormen's* Section 20.4), and **Articulation points** for **undirected** graphs (as mentioned in *Cormen's* Problem 20-2).

**Running Time.** Assuming that the various actions each run in *constant* time, we see that the total time spent within the **foreach** loop, not counting the recursive calls, is in  $\Theta(|E|)$ , and that the total time spent elsewhere is in  $\Theta(|V|)$ . Thus the generalized DFS algorithm runs in time  $\Theta(|V| + |E|)$  which for a **connected** undirected graph, or a **strongly connected** directed graph, amounts to  $\Theta(|E|)$ .