

# Graphs

Graphs are built from

- ▶ Vertices (aka nodes)
- ▶ Edges (perhaps with extra data)

Two kinds of graphs:

- ▶ directed: edge is from source  $u$  to target  $v$  ( $u = v$  possible)
- ▶ undirected: between  $u$  and  $v$  ( $u = v$  not possible)

We do not model

- ▶ hypergraphs (edges with  $> 2$  end points)
- ▶ multigraphs

We often study paths in graphs; a cycle is a path that returns to itself without following the same edge twice.

# Size Measures

The size of a graph depends on

- ▶  $n$ , the number of nodes
- ▶  $a$ , the number of edges

$n$  and  $a$  may not be related, except

$$a \in O(n^2)$$

Many graph algorithms run in time  $\Theta(n + a)$  which

- ▶ is in  $\Theta(n)$  if  $a \in O(n)$
- ▶ is in  $\Theta(n^2)$  if  $a \in \Theta(n^2)$

# Connectivity

For an **undirected** graph, two nodes are **connected** if there is a path between them.

For a **directed** graph, two nodes  $u, v$  are

- ▶ **strongly connected** if path from  $u$  to  $v$  **and** path from  $v$  to  $u$
- ▶ **unilaterally connected** if path from  $u$  to  $v$  **or** path from  $v$  to  $u$
- ▶ **weakly connected** if path **between**  $u$  and  $v$  in corresponding **undirected** graph

# Trees

A **tree** is an **undirected** graph that is

- ▶ **connected**, and
- ▶ **acyclic**

In a tree, **any** node may be picked as the **root**. Then one can define various notions:

- ▶ **parent** (exactly one for each non-root node)
- ▶ **child** (each node may have arbitrarily many)
- ▶ **leaf** (node with no children)
- ▶ **depth** (distance from root)
- ▶ **height** (maximum depth).

Special case: **binary** trees.

# Graph Operations

- ▶ **Get** to check if a given edge exists (and if so perhaps get extra data)
- ▶ **Put** to add an edge to the graph (we do not require it checks for duplicates)
- ▶ **Delete**
- ▶ **AllFrom** returns a list of **all** edges from given node

# Representations

An **Adjacency Matrix** has an entry in row  $i$ , column  $j$  iff the graph has an edge from  $i$  to  $j$

- ▶ **get**, **put**, **delete** all run in **constant** time
- ▶ **allFrom** runs in **linear** time (even if resulting list very short)

With **Adjacency Lists**, each node has a pointer to a **list** of the nodes to which there is an **edge from** that node

- ▶ **put** runs in **constant** time (as no check for duplicates)
- ▶ **allFrom** runs in **constant** time (though the returned list may be long)
- ▶ **get** (and **delete**) do **not** run in **constant** time.

# Assessment

## An Adjacency Matrix

- + allows for quick `get` (and `delete`) even in `dense` graph
- but for a `sparse` graph,
  - ▶ `allFrom` is unnecessarily slow
  - ▶ wastes space