

### Question 4

• Definition: Given  $x = x_1 \dots x_T \in \Sigma^*$  and  $A \in N$

$$e(A, i, j) \stackrel{\text{def}}{=} P_0(A \Rightarrow x_i \dots x_j)$$

• Initialization:  $\forall A \in N; \forall i: 0 \dots T-1;$

$$e(A, i, i+1) = p(A \rightarrow b) \cdot \delta(b, x_{i+1})$$

• Recursion:  $\forall A \in N; \forall l: 2 \dots T; \forall i: 0 \dots T-l;$

$$e(A, i, i+l) = \left\{ \sum_{B, C \in N} p(A \rightarrow BC) \cdot e(B, i, i+l-1) \cdot e(C, i+l-1, i+l) \right\}$$

• Final Result: Probability of a sentence:  $P_0(x) = e(S, 0, T)$

• Temporal Cost:  $O(N^3)$  Being  $N$  the length of the sentence

### Question 1

• Outside algorithm's initial expression:

$$\forall A \in N; f(A, 0, T) = \delta(A, S)$$

Kronecker delta is a function which returns 1 if both variables are equal and 0 otherwise.

$$\delta(A, S) = \begin{cases} 1 & A = S \\ 0 & A \neq S \end{cases}$$

This is done so the substate  $A$  is equal to the initial state of the tree. Therefore, its probability is maximum and set to 1.

• Outside algorithm's final expression:

$$P_0(x) = \sum_{A \in N} f(A, i, i+1) \cdot p(A \rightarrow b) \cdot \delta(b, x_{i+1}) \quad 0 \leq i \leq T-1$$

The final result is the sum of all right and left productions multiplied by their own  $A$  production probability. Recursively, the algorithm computes all outside production's probabilities to determine the probability of sentence  $x$ .

## Question 2

• Definition: Optimal score and sequence for  $x_1, \dots, x_t$  ending at  $y_t = s \in Y$

$$\varphi_t(s) \stackrel{\text{def}}{=} \max_{y_1^t; y_t=s} \prod_{i=1}^t \Psi_i(y_{i-1}, y_i, x_i)$$

$$b_t(s) \stackrel{\text{def}}{=} \arg \max_{y_1^t; y_t=s} \prod_{i=1}^t \Psi_i(y_{i-1}, y_i, x_i)$$

• Initialization:  $\forall s \in Y; \varphi_1(s) = \Psi(y_0 = \text{null}, y_1 = s, x_1)$

$$b_1(s) = 0$$

• Recursion:  $\forall t = 2 \dots T; \forall s \in Y;$

$$\varphi_t(s) = \max_{s' \in Y} \left\{ \varphi_{t-1}(s') \cdot \Psi_t(y_{t-1} = s', y_t = s, x_t) \right\}$$

$$b_t(s) = \arg \max_{s' \in Y} \left\{ \varphi_{t-1}(s') \cdot \Psi_t(y_{t-1} = s', y_t = s, x_t) \right\}$$

Final Result: Optimal score for  $x$ :

$$\max_s \varphi_T(s)$$

Optimal sequence,  $\hat{y}$ , for  $x$ :

$$y_T = \arg \max_s \varphi_T(s)$$

$$\forall t = T \dots 2; y_{t-1} = b_t(y_t)$$

$$\hat{y} = y_1 \dots y_T$$