



单元6.1 图的基本概念

第二编 图论 第七章 图

7.1 图的基本概念



北京大学



内容提要

- 图、无向图、有向图、简单图
- 相邻、关联
- 度、度序列、握手定理
- 图同构
- 图族
- 图运算





无序积

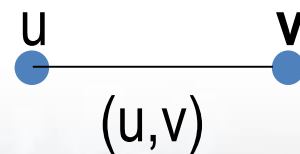
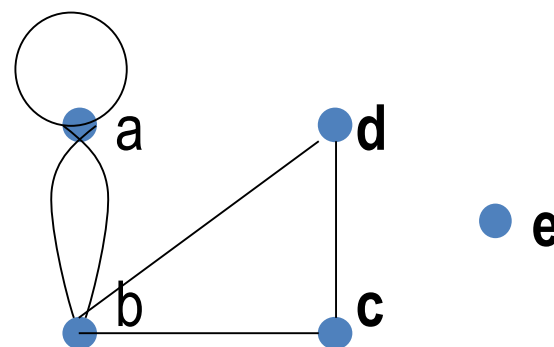
$$A \& B = \{ \{a, b\} \mid x \in A \wedge y \in B \}$$

- 记 $\{a, b\} = (a, b)$
- 允许 $a = b$
- $(a, b) = (b, a)$



无向图

- 图 $G=\langle V,E \rangle$
 - $V \neq \emptyset$ 顶点集 $V(G)$
 - $E \subseteq V \times V$ 边集(多重集) $E(G)$
- 例: $G=\langle V,E \rangle$
 - $V=\{a,b,c,d,e\}$
 - $E=\{(a,a),(a,b),(a,b),(b,c),(c,d),(b,d)\}$
- 顶点、边



有向图

- 有向图 $D=\langle V,E \rangle$

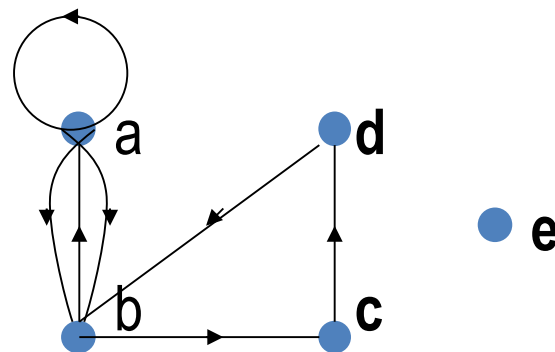
- $V \neq \emptyset$, 顶点集 $V(D)$

- $E \subseteq V \times V$, 边集(多重集) $E(D)$

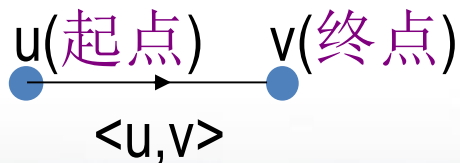
- 例: $D=\langle V,E \rangle$

- $V=\{a,b,c,d,e\}$

- $E=\{ \langle a,a \rangle, \langle a,b \rangle, \langle a,b \rangle, \langle b,a \rangle, \langle b,c \rangle, \langle c,d \rangle, \langle d,b \rangle \}$



- 有向边

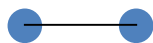


n阶图、有限图、零图、平凡图、空图

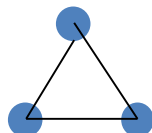
- n阶图: $|V(G)|=n$



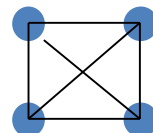
1阶图



2阶图



3阶图



4阶图



- 有限图: $|V(G)|<\infty$
- 零图 N_n : $E=\emptyset$
- 平凡图: 1阶零图 N_1
- 空图: $V=E=\emptyset$



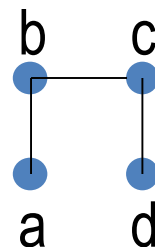
N_1



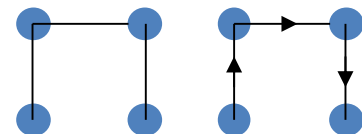
N_4

标定图、非标定图、底图

- 标定图：顶点或边带标记的图



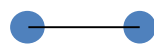
- 非标定图：顶点和边不带标记的图



- 底图(基图)：有向图去掉边的方向后得到的无向图

相邻、邻接、关联

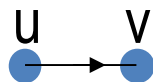
- 有边相连的两个顶点是相邻的



- 有公共顶点的两条边是相邻的



- u 邻接到 v , v 邻接于 u



- 一条边的端点与这条边是关联的

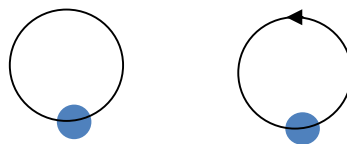


- 关联次数

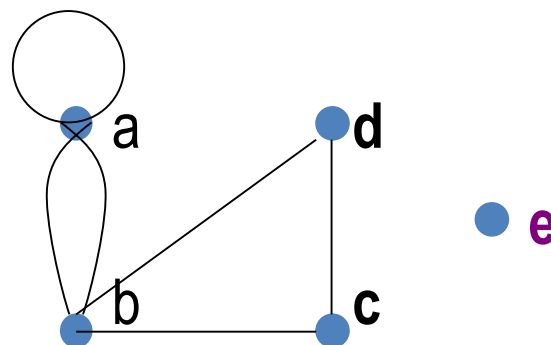


环、孤立点、平行边

- 环：只与一个顶点关联的边

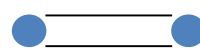


- 孤立点：不与任何边关联的顶点

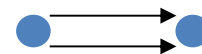


- 平行边

– 端点相同的两条无向边是平行边



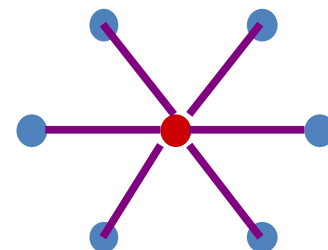
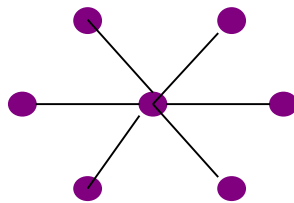
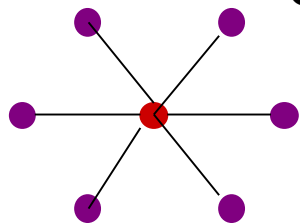
– 起点与终点相同的两条有向边是平行边



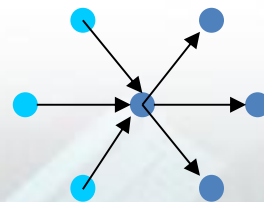
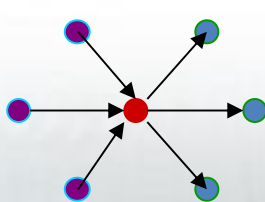
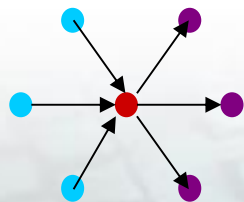
非平行边

邻域、闭邻域、关联集


- 邻域: $N_G(v) = \{u \in V(G) \mid (u, v) \in E(G) \wedge u \neq v\}$
- 闭邻域: $N_G(v) = N_G(v) \cup \{v\}$
- 关联集: $I_G(v) = \{e \mid e \text{ 与 } v \text{ 关联}\}$

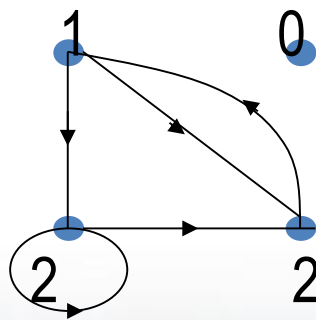
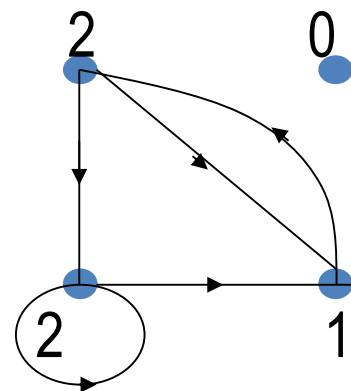
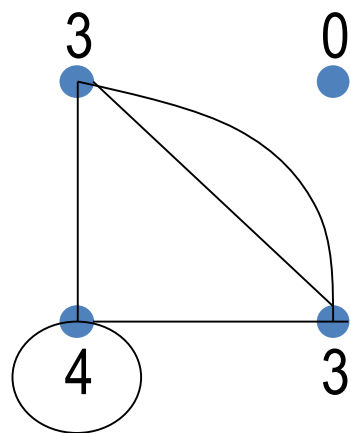
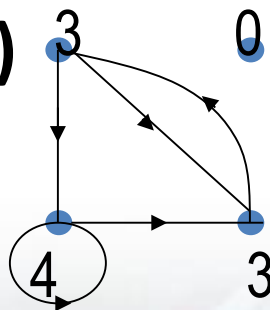


- 后继: $\Gamma_D^+(v) = \{u \in V(D) \mid \langle v, u \rangle \in E(D) \wedge u \neq v\}$
- 前驱: $\Gamma_D^-(v) = \{u \in V(D) \mid \langle u, v \rangle \in E(D) \wedge u \neq v\}$
- (闭)邻域: $N_D(v) = \Gamma_D^+(v) \cup \Gamma_D^-(v)$ $N_D(v) = N_D(v) \cup \{v\}$





- 度: $d_G(v)$ = 与 v 关联的边的次数之和
- 出度: $d_D^+(v)$ = 与 v 关联的出边的次数之和
- 入度: $d_D^-(v)$ = 与 v 关联的入边的次数之和
- 度: $d_D(v) = d_D^+(v) + d_D^-(v)$ 



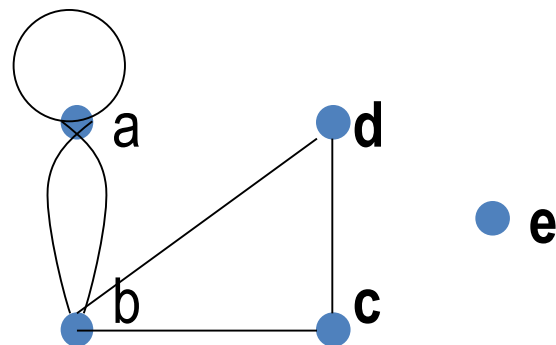


最大度、最小度

- 最大度 $\Delta(G) = \max\{ d_G(v) \mid v \in V(G) \}$
- 最小度 $\delta(G) = \min\{ d_G(v) \mid v \in V(G) \}$
- 最大出度 $\Delta^+(D) = \max\{ d_D^+(v) \mid v \in V(D) \}$
- 最小出度 $\delta^+(D) = \min\{ d_D^+(v) \mid v \in V(D) \}$
- 最大入度 $\Delta^-(D) = \max\{ d_D^-(v) \mid v \in V(D) \}$
- 最小入度 $\delta^-(D) = \min\{ d_D^-(v) \mid v \in V(D) \}$

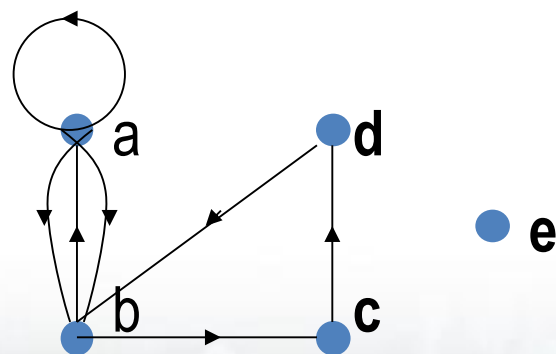
举例

- $\Delta=4, \delta=0$



- $\Delta^+=3, \delta^+=0$

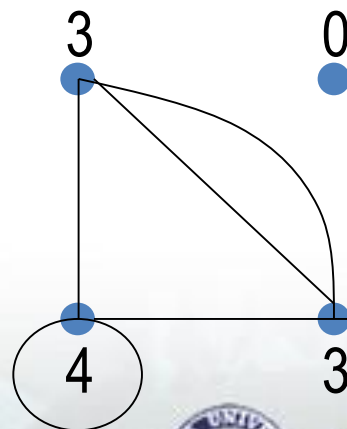
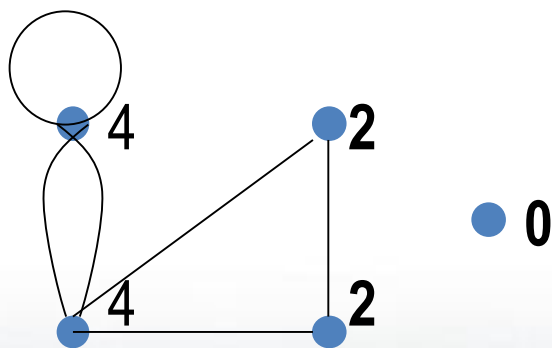
$$\Delta^-=3, \delta^-=0$$



图论基本定理(握手定理)

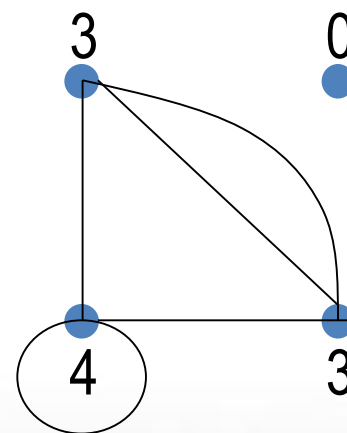
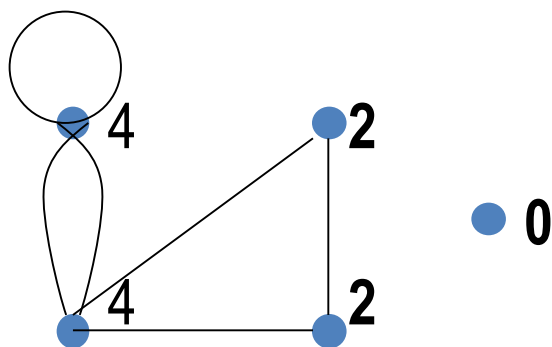
- 设 $G=\langle V, E \rangle$ 是无向图,
 $V=\{v_1, v_2, \dots, v_n\}$, $|E|=m$, 则

$$d(v_1)+d(v_2)+\dots+d(v_n)=2m. \quad \#$$



握手定理推论

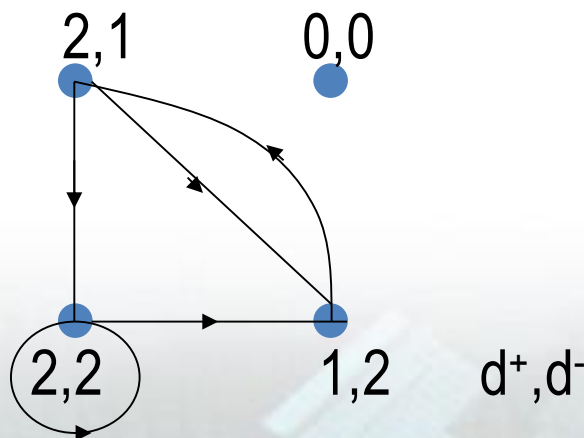
任何图中奇度顶点的个数是偶数. #



图论基本定理(握手定理)

- 设 $D=\langle V,E \rangle$ 是有向图, $V=\{v_1, v_2, \dots, v_n\}$, $|E|=m$, 则

$$\begin{aligned} & d^+(v_1) + d^+(v_2) + \dots + d^+(v_n) \\ &= d^-(v_1) + d^-(v_2) + \dots + d^-(v_n) = m. \quad \# \end{aligned}$$



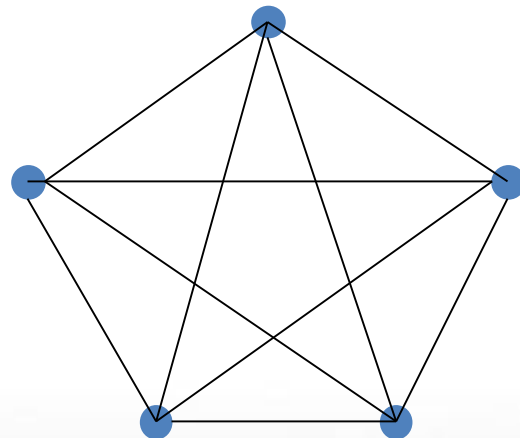
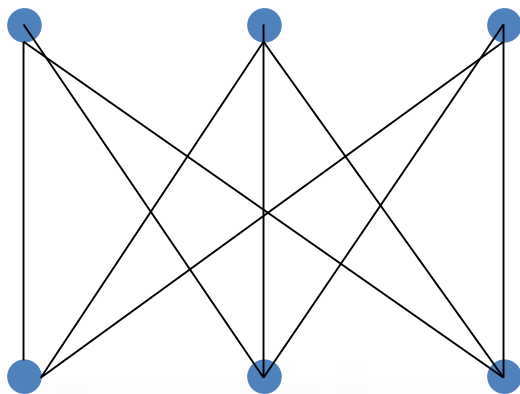
d^+, d^-

简单图、k-正则图

- 简单图：既无环也无平行边的图

$$\Rightarrow 0 \leq \Delta(G) \leq n-1$$

- k-正则图：所有顶点的度都是k



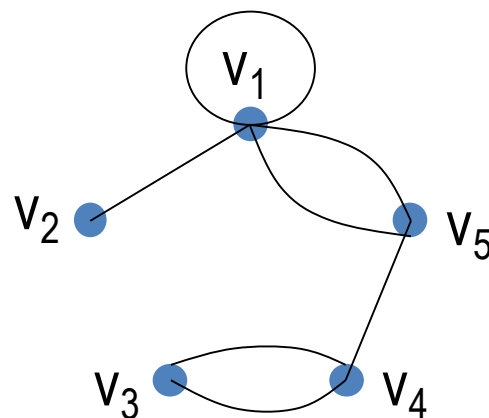
度数列

- 设 $G = \langle V, E \rangle$, $V = \{v_1, v_2, \dots, v_n\}$, 称

$$d = (d(v_1), d(v_2), \dots, d(v_n))$$

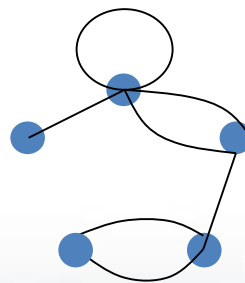
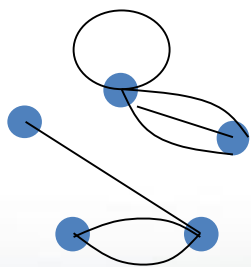
为 G 的度数列

- 例: $d = (5, 1, 2, 3, 3)$



可图化

- 设非负整数列 $d=(d_1, d_2, \dots, d_n)$, 若存在图 G , 使得 G 的度数列是 d , 则称 d 为可图化的
- 可图化举例: 度数列 $d=(5, 3, 3, 2, 1)$





可图化充要条件

非负整数列 $d=(d_1, d_2, \dots, d_n)$ 可图化 \Leftrightarrow
 $d_1 + d_2 + \dots + d_n \equiv 0 \pmod{2}.$

- 证: (\Rightarrow) 握手定理
 (\Leftarrow) 奇数度点两两之间连一边,
剩余度用环来实现. #



可简单图化

- 设非负整数列 $d=(d_1, d_2, \dots, d_n)$, 若存在简单图 G , 使得 G 的度数列是 d , 则称 d 为可简单图化的
- 例: $d=(5, 3, 3, 2, 1)$ 不可简单图化
 - $\Delta=n$ ($\Delta \leq n-1$)
- 例: $d=(4, 4, 3, 2, 1)$ 不可简单图化
 - $(n-1, n-1, \dots, 1)$



可简单图化充要条件(Havel定理)

- 设非负整数列 $d=(d_1, d_2, \dots, d_n)$ 满足:

$$d_1 + d_2 + \dots + d_n \equiv 0 \pmod{2},$$

$$n-1 \geq d_1 \geq d_2 \geq \dots \geq d_n \geq 0,$$

则 d 可简单图化 \Leftrightarrow

$$d'=(d_2-1, d_3-1, \dots, d_{d_1+1}-1, d_{d_1+2}, \dots, d_n)$$

可简单图化

- 例: $d=(4, 4, 3, 3, 2, 2)$, $d'=(3, 2, 2, 1, 2)$



Havel定理举例

- 判断下列非负整数列是否可简单图化. (1)
 $(5, 5, 4, 4, 2, 2)$ (2) $(4, 4, 3, 3, 2, 2)$
- 解: (1) $(5, 5, 4, 4, 2, 2)$, $(4, 3, 3, 1, 1)$,
 $(2, 2, 0, 0)$, $(1, -1, 0)$, 不可简单图化.
- (2) $(4, 4, 3, 3, 2, 2)$, $(3, 2, 2, 1, 2)$, $(3, 2, 2, 2, 1)$,
 $(1, 1, 1, 1)$, $(0, 1, 1)$, $(1, 1)$, 可简单图化. #

可简单图化充要条件

- 定理7.4(P.Erdős,T.Gallai,1960):

设非负整数列 $d=(d_1,d_2,\dots,d_n)$ 满足:

$$n-1 \geq d_1 \geq d_2 \geq \dots \geq d_n \geq 0,$$

则 d 可简单图化 \Leftrightarrow

$$d_1+d_2+\dots+d_n \equiv 0 \pmod{2}$$

并且对 $r=1,2,\dots,n-1$ 有

$$d_1+d_2+\dots+d_r \leq r(r-1) + \min\{r,d_{r+1}\} + \min\{r,d_{r+2}\} + \dots + \min\{r,d_n\}.$$

#

定理7.4等价形式

- **定理7.4'** (P.Erdős,T.Gallai,1960):非负整数列 $d=(d_1,d_2,\dots,d_n)$ 可简单图化 \Leftrightarrow
 $d_1+d_2+\dots+d_n=0(\bmod 2)$
并且对 $r=1,2,\dots,n$ 有
$$d_1+d_2+\dots+d_r \leq r(r-1)+\min\{r,d_{r+1}\}+\min\{r,d_{r+2}\}+\dots+\min\{r,d_n\}. \quad \#$$
- **说明:** $n-1 \geq d_1 \geq d_2 \geq \dots \geq d_n \geq 0 \Rightarrow$
$$d_1+d_2+\dots+d_n \leq n(n-1).$$



定理7.4举例

• 下列非负整数列是否可简单图化?

(1) (5,4,3,2,2,1)

(2) (5,4,4,3,2)

(3) (3,3,3,1)

(4) (6,6,5,4,3,3,1)

(5) (5,5,3,3,2,2,2)

(6) (d_1, d_2, \dots, d_n) , $d_1 > d_2 > \dots > d_n \geq 1$



定理7.4举例

- (1) $5+4+3+2+2+1=17 \neq 0 \pmod{2}$.

不可(简单)图化.

- (2) $5+4+4+3+2=18=0 \pmod{2}$.

但是 $d_1=5 > n-1=4$, 不满足 $n-1 \geq d_1$, 不可简单图化.

(或者: 但是 $r=1$ 时, $d_1=5 > 1(1-1) + \min\{1,4\} + \min\{1,4\} + \min\{1,3\} + \min\{1,2\} = 4$, 不可简单图化.)

- (3) $3+3+3+1=10=0 \pmod{2}$. $d_1=3=n-1$, 满足 $n-1 \geq d_1$,

但是 $r=2$ 时, $d_1+d_2=6 > 2(2-1) + \min\{2,3\} + \min\{2,1\} = 5$,
不可简单图化.

定理4举例(4)

- (4) $6+6+5+4+3+3+1=28=0(\text{mod } 2)$.

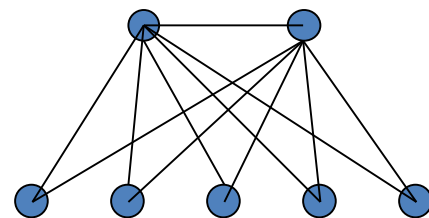
$d_1=6=n-1$, 满足 $n-1 \geq d_1$. $r=1, 2$ 时,

$d_1=6 \leq 1(1-1) + \min\{1, 6\} + \min\{1, 5\} + \dots = 6$,

$d_1+d_2=12 > 2(2-1) + \min\{2, 5\} + \dots = 11$,

不可简单图化.

- 或: $6, 6, *, *, *, *, 1$ 不可简单图化





定理4举例(5)

• (5) $5+5+3+3+2+2+2=22=0(\text{mod } 2)$.

$d_1=5 < n-1$, 满足 $n-1 \geq d_1$. $r=1, 2, \dots, 7$ 时,

$d_1=5 < 1(1-1)+\min\{1,5\}+\min\{1,5\}+ \dots=6$,

$d_1+d_2=10 < 2(2-1)+\min\{2,3\}+ \dots=12$,

$d_1+d_2+d_3=13 < 3(3-1)+\min\{3,3\}+ \dots=15$,

$d_1+d_2+d_3+d_4=16 < 4(4-1)+\min\{4,2\}+ \dots=18$,



定理4举例(5)

$$d_1+d_2+d_3+d_4=16<4(4-1)+\min\{4,2\}+\dots=18,$$

$$d_1+d_2+\dots+d_5=18<5(5-1)+\min\{5,2\}+\dots=24,$$

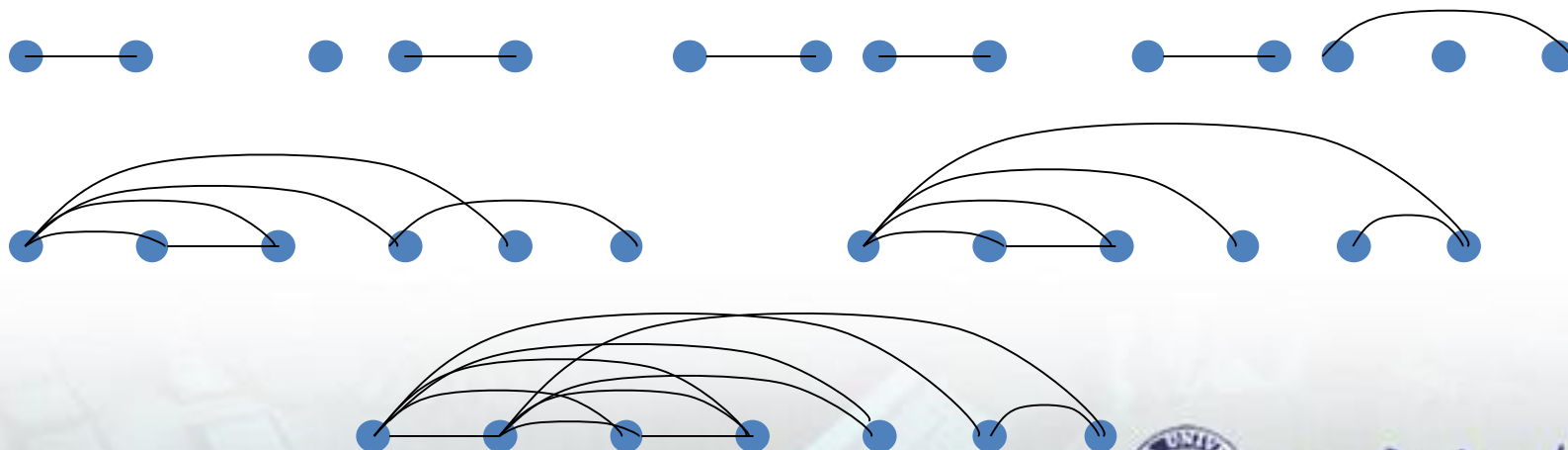
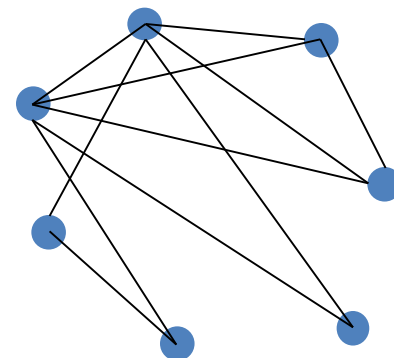
$$d_1+d_2+\dots+d_6=20<6(6-1)+\min\{6,2\}=32,$$

$$d_1+d_2+\dots+d_7=22<7(7-1)=42,$$

可简单图化.

定理4举例(5)

$(5,5,3,3,2,2,2), (4,2,2,1,1,2), (4,2,2,2,1,1),$
 $(1,1,1,0,1), (1,1,1,1), (0,1,1), (1,1)$





定理4举例(6)

- (6) $d_1 > d_2 > \dots > d_n \geq 1$,
 $d_{n-1} \geq 2, d_{n-2} \geq 3, \dots, d_1 \geq n$,
不满足 $n-1 \geq d_1$,
不可简单图化. #

Paul Erdős(1913-1996)

- 保罗•爱尔特希, 匈牙利人
- 廿世纪数学界的传奇人物
- “Another roof, another proof.”





Paul Erdős



“My mother said, ‘Even you, Paul, can be in only one place at one time.’

Maybe soon I will be relieved of this disadvantage.

Maybe, once I've left, I'll be able to be in many places at the same time.

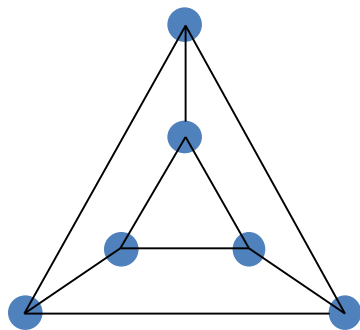
Maybe then I'll be able to collaborate with Archimedes and Euclid.”



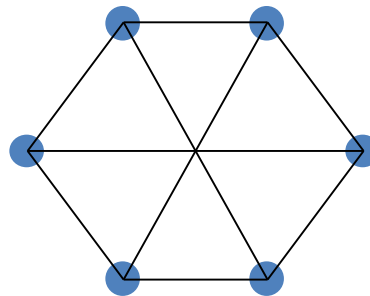
图同构

- 无向图 $G_1 = \langle V_1, E_1 \rangle$, $G_2 = \langle V_2, E_2 \rangle$, 若存在双射 $f: V_1 \rightarrow V_2$ 满足 $\forall u, v \in V_1, (u, v) \in E_1 \leftrightarrow (f(u), f(v)) \in E_2$, 则称 G_1 与 G_2 同构, 记作 $G_1 \cong G_2$
- 同构的图, 其图论性质完全一样
- NAUTY 算法
- 有向图 $D_1 = \langle V_1, E_1 \rangle$, $D_2 = \langle V_2, E_2 \rangle$, 若存在双射 $f: V_1 \rightarrow V_2$, 满足 $\forall u, v \in V_1, \langle u, v \rangle \in E_1 \leftrightarrow \langle f(u), f(v) \rangle \in E_2$, 则称 D_1 与 D_2 同构, 记作 $D_1 \cong D_2$

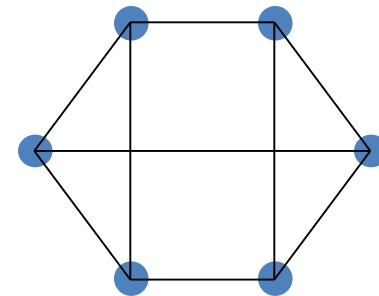
图同构(举例)



G_1



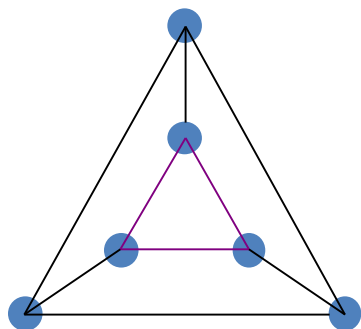
G_2



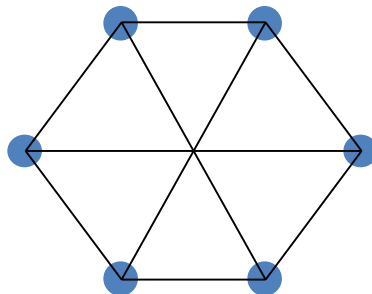
G_3

$$G_1 \cong G_3, \quad G_1 \not\cong G_2$$

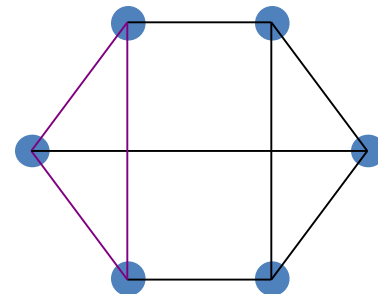
图同构(举例)



G_1

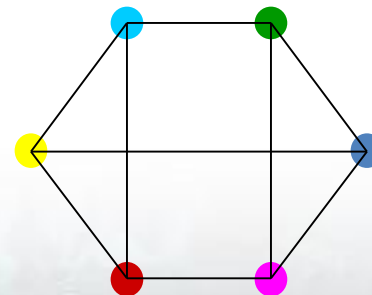
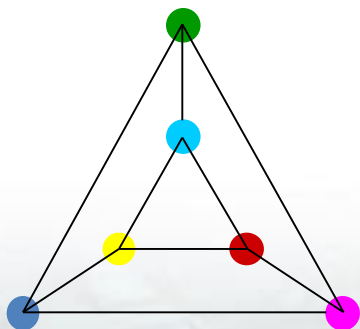


G_2

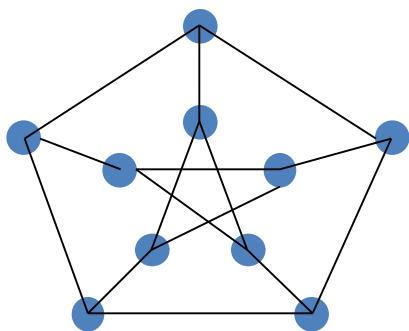


G_3

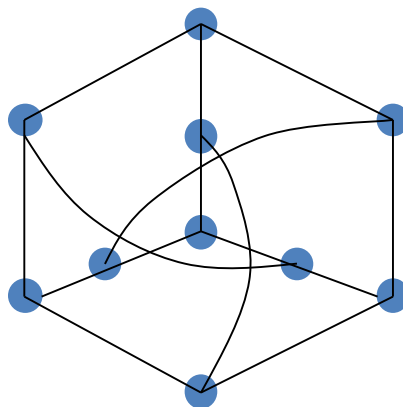
$$G_1 \cong G_3, \quad G_1 \not\cong G_2$$



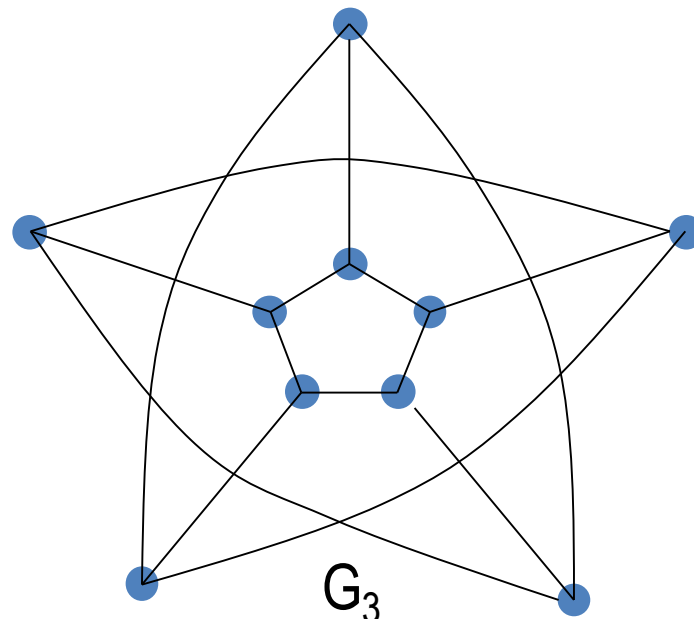
图同构(举例)



G_1



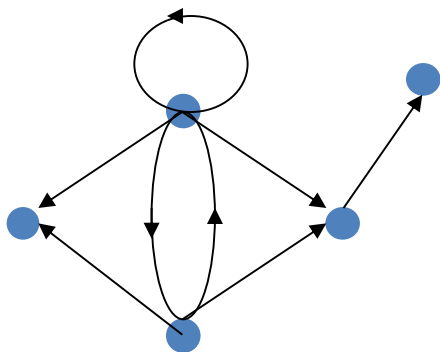
G_2



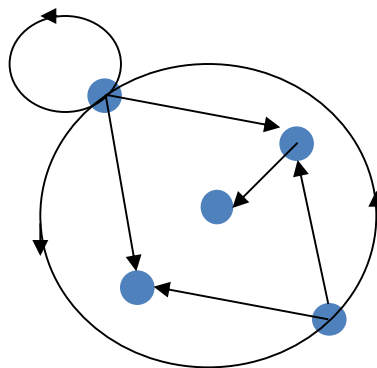
G_3

$$G_1 \cong G_2 \cong G_3$$

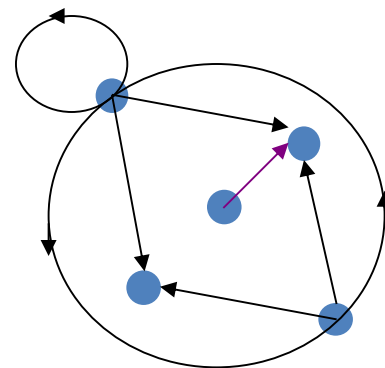
图同构(举例)



D_1



D_2



D_3

$$D_1 \cong D_2, \quad D_2 \not\cong D_3$$



图族(graph class)

- 完全图,有向完全图,竞赛图
- 柏拉图图,彼德森图,库拉图斯基图
- r 部图,二部图(偶图),完全 r 部图
- 路径,圈,轮,超立方体



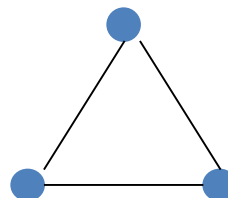
完全图



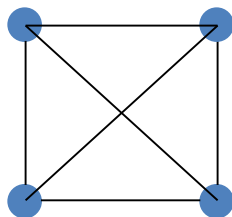
K_1



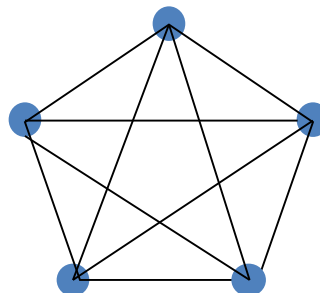
K_2



K_3

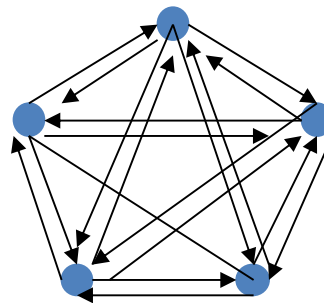
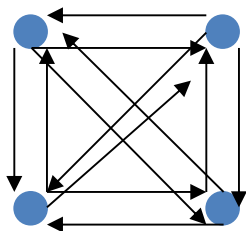
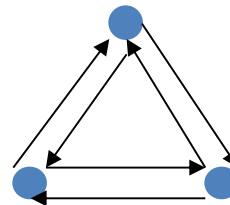


K_4



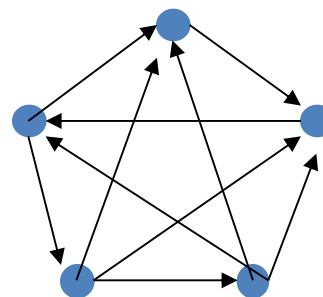
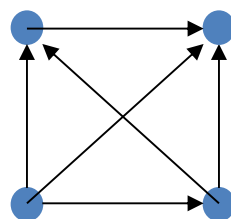
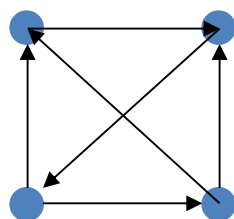
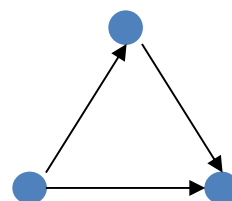
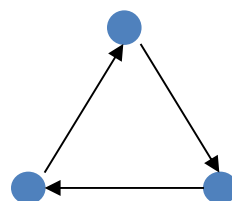
K_5

有向完全图

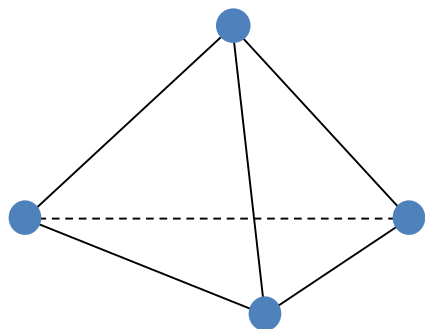




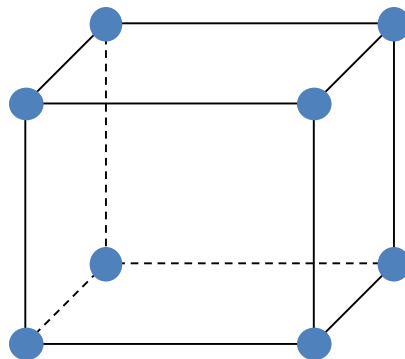
竞赛图



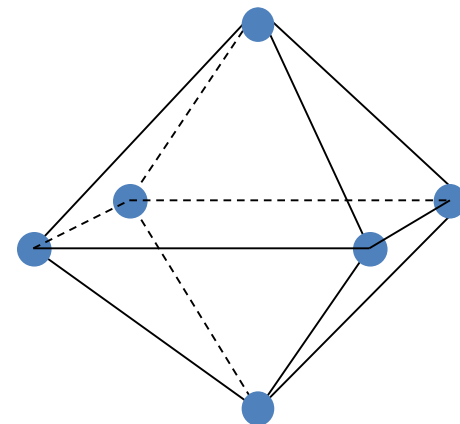
柏拉图图



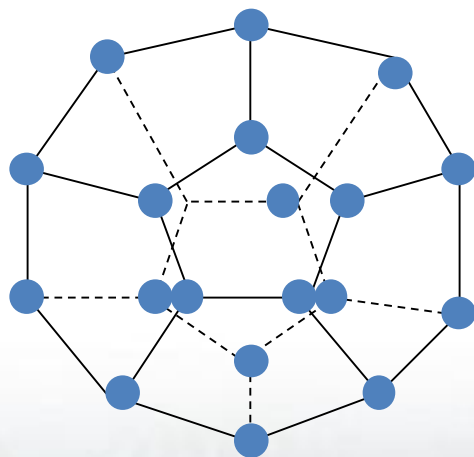
正四面体图



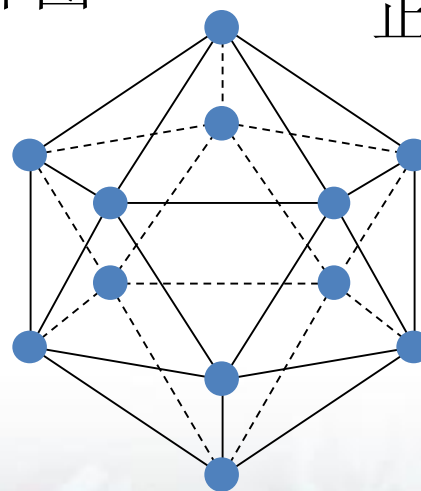
正六面体图



正八面体图

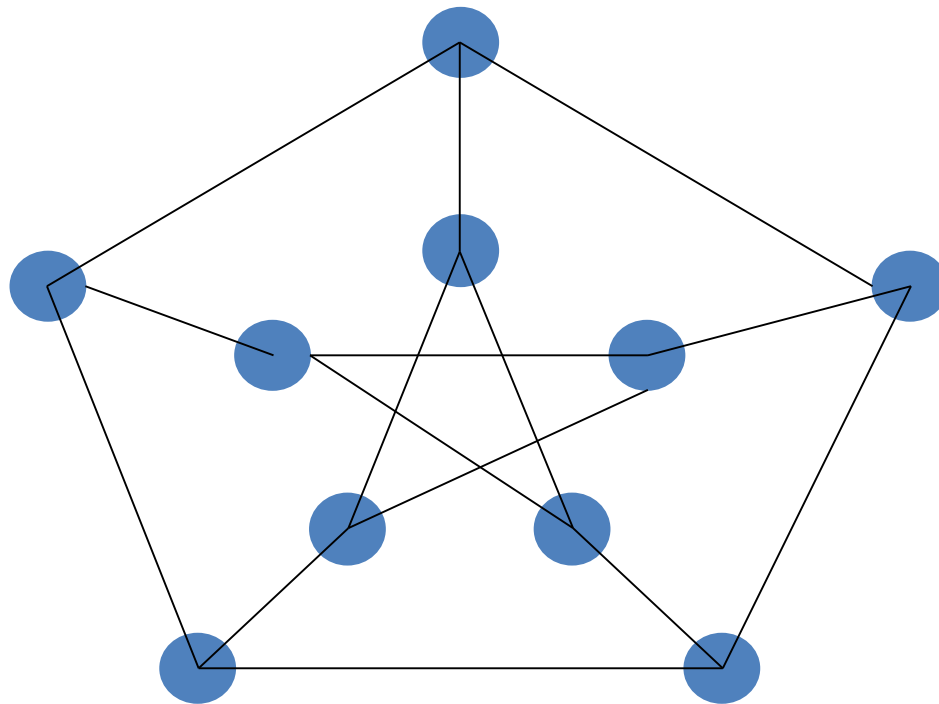


正十二面体图



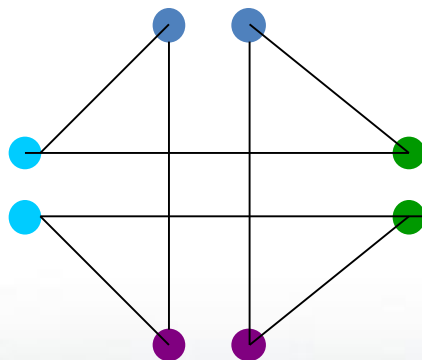
正二十面体图

彼德森图



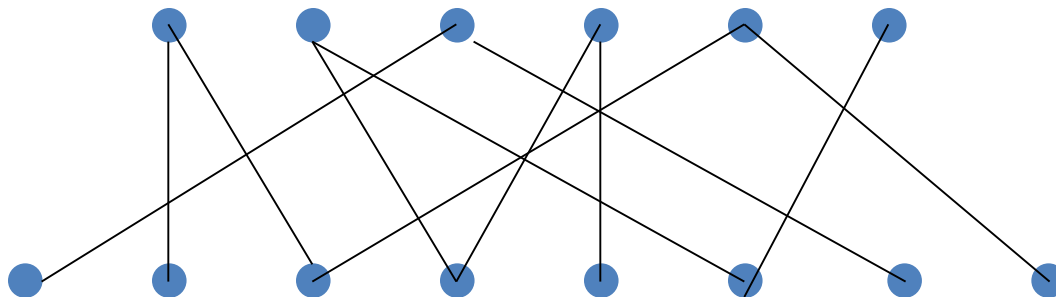
r部图

- **r部图**: $G=\langle V, E \rangle, V=V_1 \cup V_2 \cup \dots \cup V_r,$
 $V_i \cap V_j = \emptyset \ (i \neq j), E \subseteq \bigcup_{i \neq j} (V_i \times V_j),$
- 也记作 $G=\langle V_1, V_2, \dots, V_r; E \rangle.$

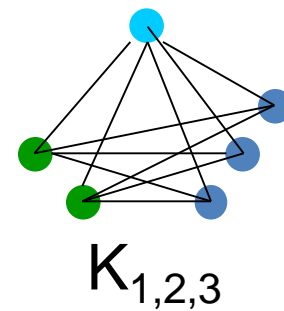
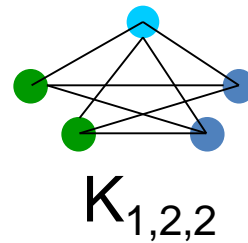
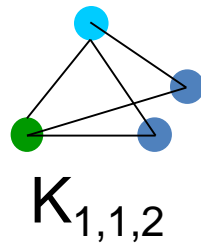
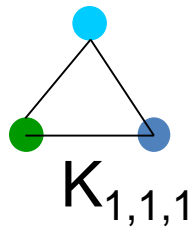
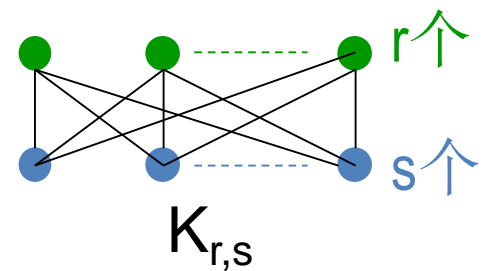
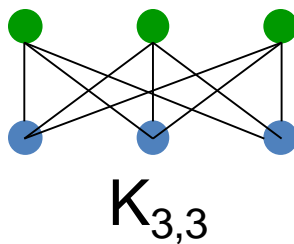
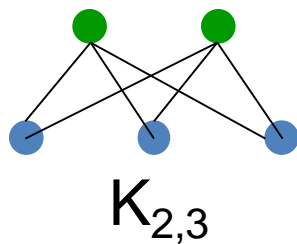


二部图(偶图)

- 二部图: $G=\langle V_1, V_2; E \rangle$, 也称为偶图



完全r部图





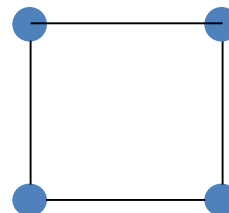
超立方体



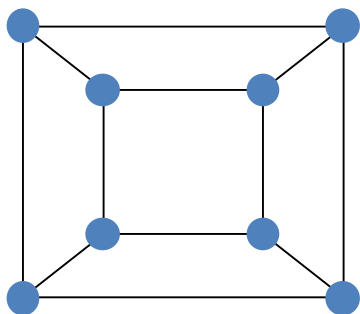
Q_0



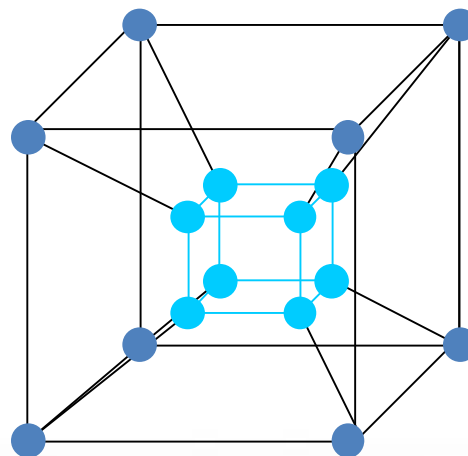
Q_1



Q_2



Q_3



Q_4



子图,生成子图

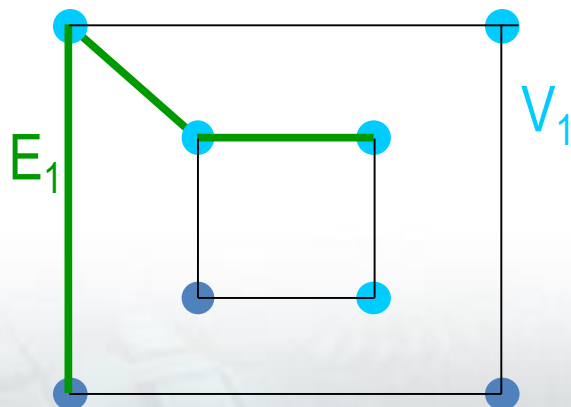
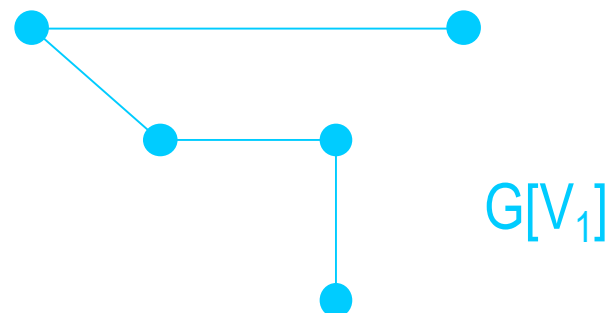
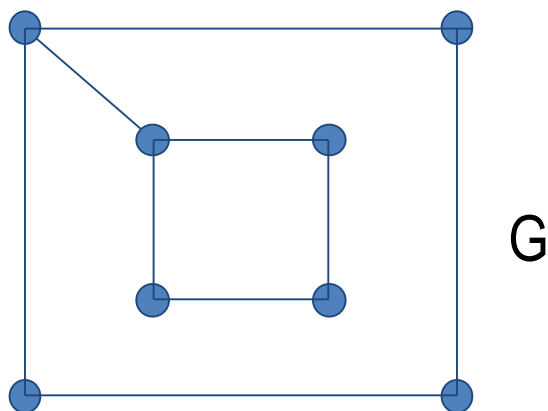
- 子图(subgraph): $G=\langle V,E\rangle$, $G'=\langle V',E'\rangle$,
 $G'\subseteq G \Leftrightarrow V'\subseteq V \wedge E'\subseteq E$
- 真子图(proper subgraph):
 $G'\subset G \Leftrightarrow G'\subseteq G \wedge (V'\subset V \vee E'\subset E)$
- 生成子图(spanning subgraph):
 G' 是 G 的生成子图 $\Leftrightarrow G'\subseteq G \wedge V'=V$



导出子图

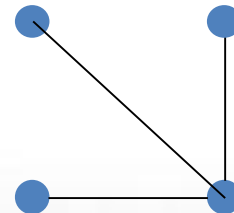
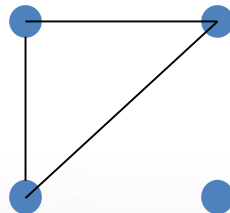
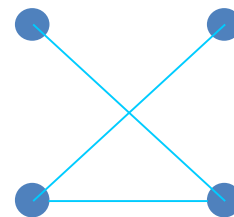
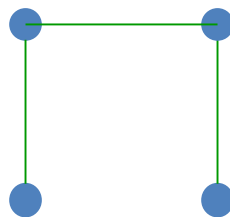
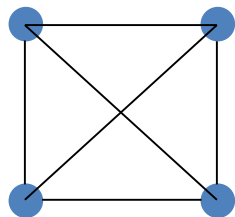
- 导出子图: $G = \langle V, E \rangle$,
- 若 $V_1 \subset V, E_1 = E \cap V_1 \& V_1$, 则称
$$G[V_1] = \langle V_1, E_1 \rangle$$
为由 V_1 导出的子图
- 若 $\emptyset \neq E_1 \subset E, V_1 = \{v \mid v \text{ 与 } E_1 \text{ 中的边关联}\}$, 则称
$$G[E_1] = \langle V_1, E_1 \rangle$$
为由 E_1 导出的子图

导出子图(举例)



补图

- 补图: $G=\langle V,E\rangle$, $\overline{G}=\langle V,E(K_n)-E\rangle$
- 自补图(self-complement graph): $G\cong\overline{G}$



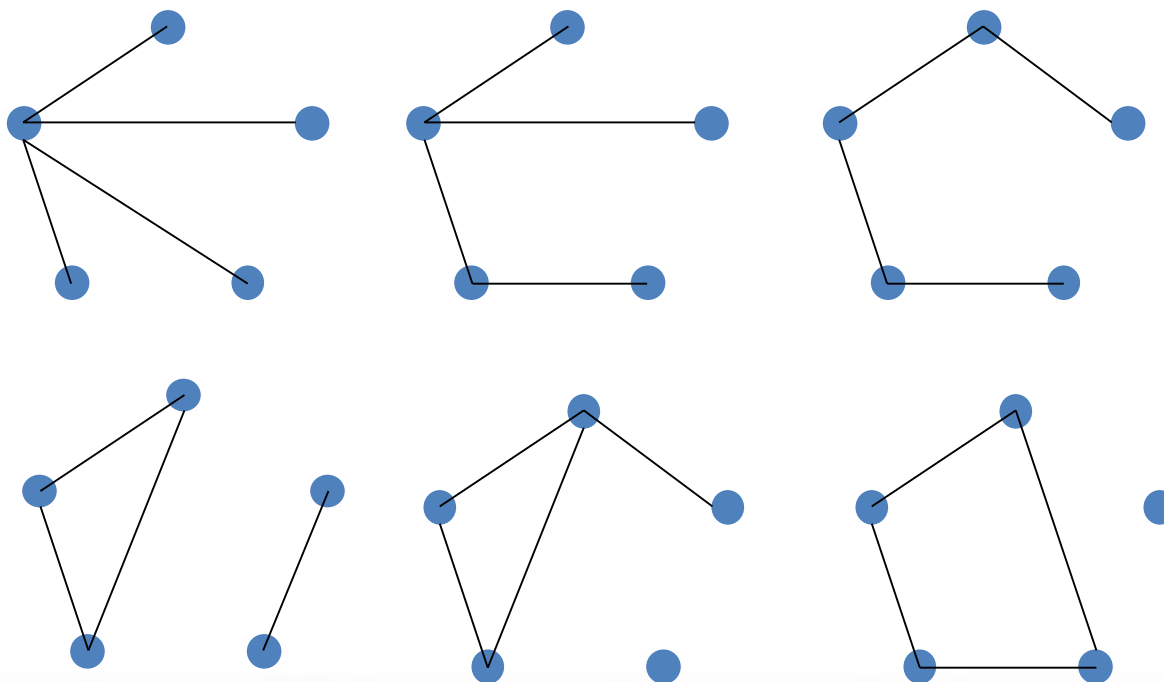


例7.5

- (1) 画出5阶4条边的所有非同构的无向简单图; (2) 画出4阶2条边的所有非同构的有向简单图.
- (1) $\sum d(v)=2m=8, \Delta \leq n-1=4,$
 $(4,1,1,1,1), (3,2,1,1,1), (2,2,2,1,1),$
 $(3,2,2,1,0), (2,2,2,2,0)$
- (2) $\sum d^+(v)=\sum d^-(v)=m=2, \sum d(v)=2m=4,$
 $(2,1,1,0), (1,1,1,1), (2,2,0,0)$

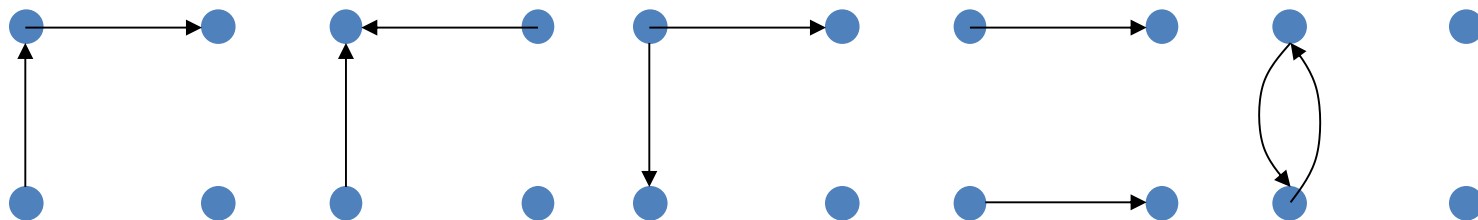
例7.5(1)

- 5阶4条边的所有非同构无向简单图



例7.5(2)

- 4阶2条边的所有非同构有向简单图.





图的运算

- 删除,收缩,加新边,不交
- 并图,交图,差图,环和
- 联图,积图

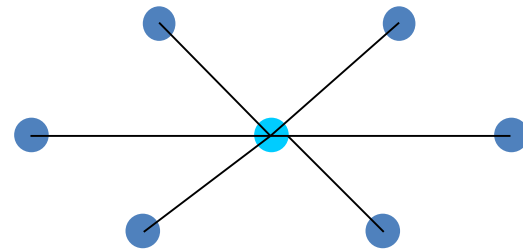
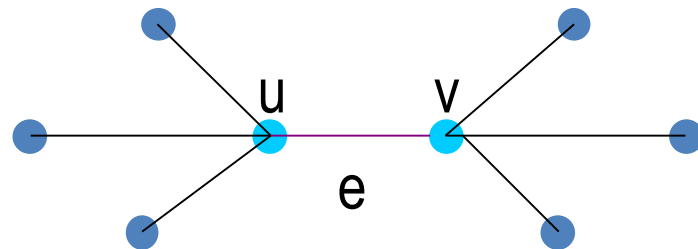


删除(delete)

- $G-e = \langle V, E-\{e\} \rangle$, 删除边 e
- $G-E' = \langle V, E-E' \rangle$, 删除边集 E'
- $G-v = \langle V-\{v\}, E-I_G(v) \rangle$, 删除顶点 v 以及 v 所关联的所有边
- $G-V' = \langle V-V', E-I_G(V') \rangle$, 删除顶点集 V' 以及 V' 所关联的所有边

收缩、加新边

- $G \setminus e$: $e=(u,v)$, 删除 e , 合并 u 与 v
 - $G \cup (u,v) = \langle V, E \cup \{(u,v)\} \rangle$
 - 在 u 与 v 之间加新边
- 也写成 $G+(u,v)$





不交

- $G_1 = \langle V_1, E_1 \rangle, G_2 = \langle V_2, E_2 \rangle,$
- G_1 与 G_2 不交 $\Leftrightarrow V_1 \cap V_2 = \emptyset$
- G_1 与 G_2 边不交(边不重) $\Leftrightarrow E_1 \cap E_2 = \emptyset$



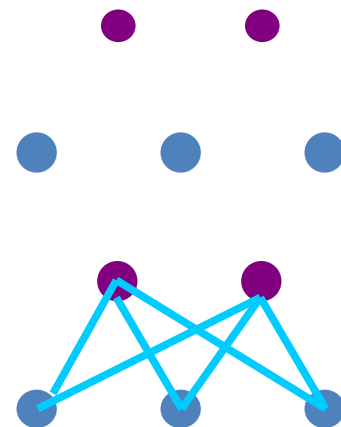
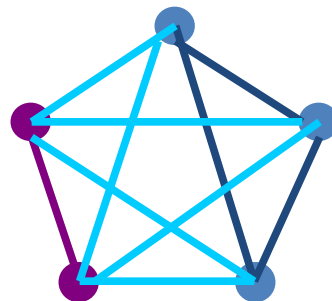
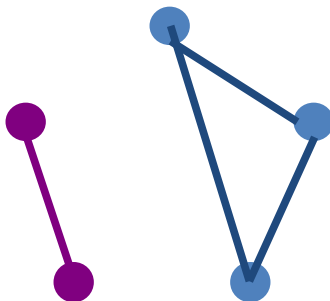
并图,交图,差图,环和

- $G_1 = \langle V_1, E_1 \rangle$, $G_2 = \langle V_2, E_2 \rangle$, 都无孤立点
- 并图: $G_1 \cup G_2 = \langle V(E_1 \cup E_2), E_1 \cup E_2 \rangle$
- 交图: $G_1 \cap G_2 = \langle V(E_1 \cap E_2), E_1 \cap E_2 \rangle$
- 差图: $G_1 - G_2 = \langle V(E_1 - E_2), E_1 - E_2 \rangle$
- 环和: $G_1 \oplus G_2 = \langle V(E_1 \oplus E_2), E_1 \oplus E_2 \rangle$

联图

- $G_1 = \langle V_1, E_1 \rangle$, $G_2 = \langle V_2, E_2 \rangle$, 不交无向图
- $G_1 + G_2 = \langle V_1 \cup V_2, E_1 \cup E_2 \cup (V_1 \& V_2) \rangle$

- $K_r + K_s = K_{r+s}$
- $N_r + N_s = K_{r,s}$
- $n = n_1 + n_2$, $m = m_1 + m_2 + n_1 n_2$

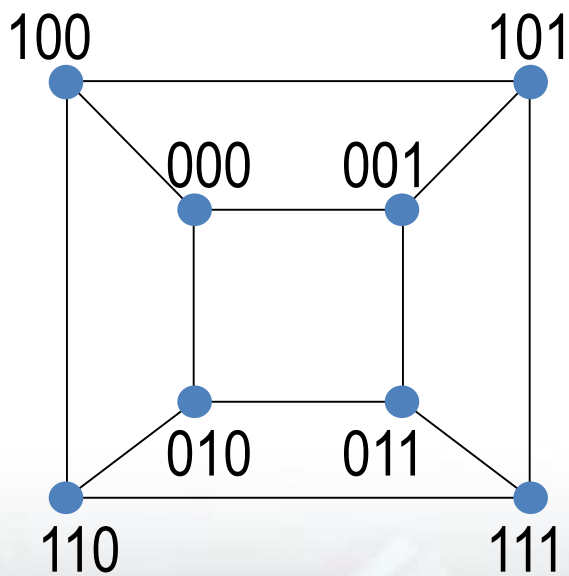
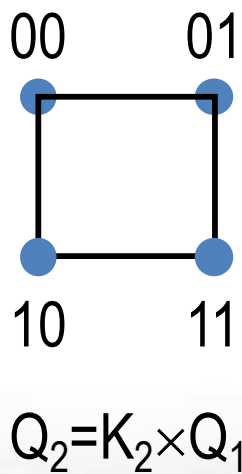
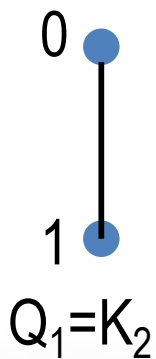
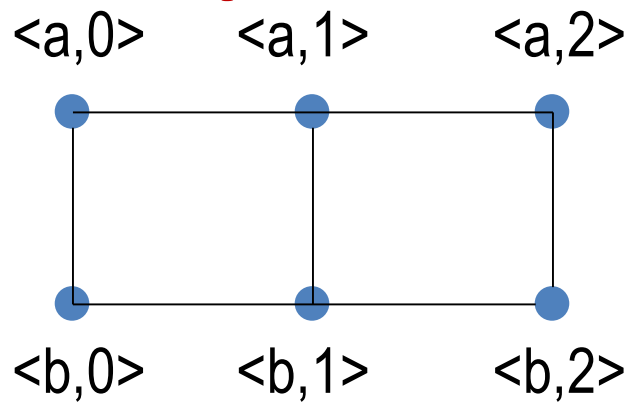
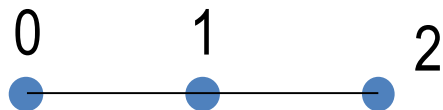
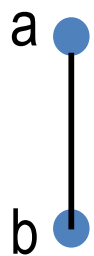




积图

- $G_1 = \langle V_1, E_1 \rangle$, $G_2 = \langle V_2, E_2 \rangle$, 无向简单图
- $G_1 \times G_2 = \langle V_1 \times V_2, E \rangle$, 其中
$$E = \{ (\langle u_i, u_j \rangle, \langle u_k, u_s \rangle) \mid$$
$$(\langle u_i, u_j \rangle, \langle u_k, u_s \rangle \in V_1 \times V_2) \wedge$$
$$((u_i = u_k \wedge u_j \text{ 与 } u_s \text{ 相邻}) \vee (u_j = u_s \wedge u_i \text{ 与 } u_k \text{ 相邻})) \}$$
- $n = n_1 n_2$, $m = n_1 m_2 + n_2 m_1$

积图(举例)



$$Q_k = K_2 \times Q_{k-1}$$



小结

- 无向图,有向图,简单图,相邻,关联
- 度,握手定理
- 图同构
- 图族(完全图,零图, ...)
- 图运算(并图,联图,积图, ...)

