

# **Real Lecture 22: More Clusters more Graphs**

# Parenthesis on Clusters

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Code to cluster all the data can be found here:

- <http://bit.ly/10nQ54Q>
- The code requires the code for clustering (<http://bit.ly/15cApcE>) and for reading the data (<http://bit.ly/ZCwSQg>) to both be in the same folder

Resulting clusters can be found here:

- <http://bit.ly/10F0V5A>
- Feel free to play with the code and experiment with the effect of different distance metrics

# What is a graph

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A set of Vertices together with a set of Edges connecting those vertices  $(V, E)$

- The set of vertices is an arbitrary set
  - for this class we assume it is finite

Edges are different depending on the type of graph

- Undirected: Edge is a pair  $(v_i, v_j)$ .

- The pair  $(v_i, v_j)$  is equivalent to  $(v_j, v_i)$



- Directed: Edge is a pair  $(v_i, v_j)$ .

- The pair  $(v_i, v_j)$  is NOT equivalent to  $(v_j, v_i)$

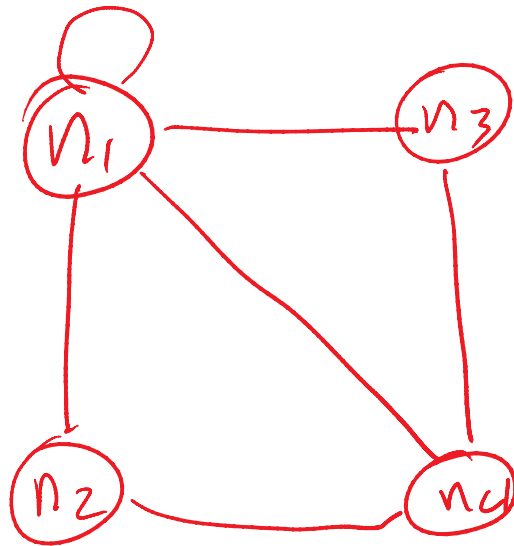
- Weighted: Edge is a triple  $(v_i, v_j, w)$

- Can be directed or undirected

# Example 1: Undirected

## • Example 1: Undirected

- Vertices  $\{n_1, n_2, n_3, n_4\}$
- Edges  $\{(n_1, n_2), (n_3, n_1), (n_1, n_4), (n_2, n_4), (n_1, n_1), (n_3, n_4)\}$

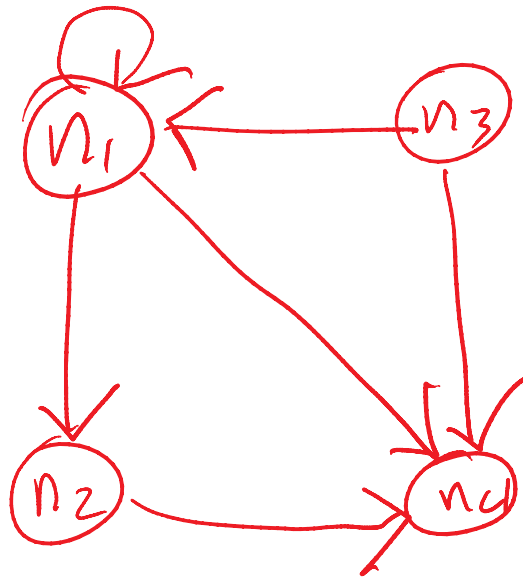


All nodes are reachable from each other

# Example 2: Directed

## • Example 2: Directed

- Vertices  $\{n_1, n_2, n_3, n_4\}$
- Edges  $\{(n_1, n_2), (n_3, n_1), (n_1, n_4), (n_2, n_4), (n_1, n_1), (n_3, n_4)\}$



There is no path from  $n_4$  to  $n_1$  for example

Shortest path from  $n_3$  to  $n_4$  takes one edge

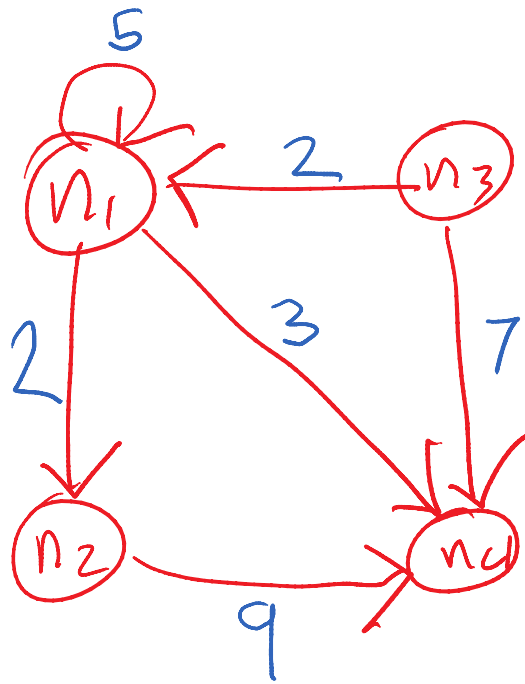
# Example 3: Weighted Directed

## • Example 3: Weighted Directed

- Vertices  $\{n_1, n_2, n_3, n_4\}$

- Edges

$\{(n_1, n_2, 2), (n_3, n_1, 2), (n_1, n_4, 3), (n_2, n_4, 9), (n_1, n_1, 5), (n_3, n_4, 7)\}$



Shortest path from  $n_3$  to  $n_4$  is not the one-edge path, but the path that goes through  $n_1$

Shortest path between any two nodes will never go through a loop

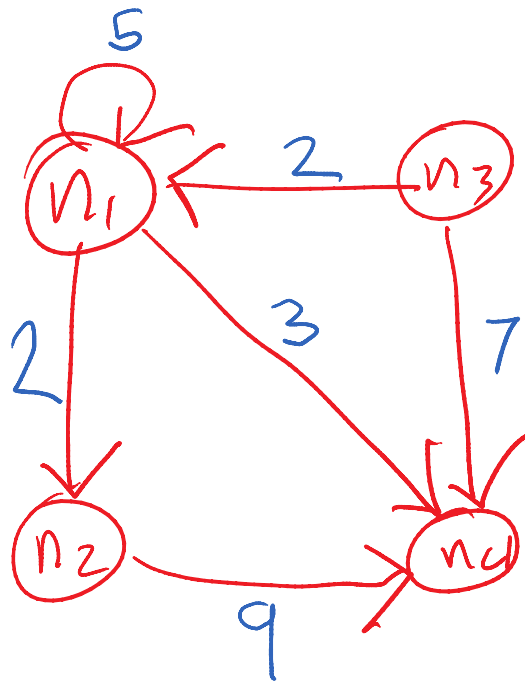
# Example 2: Weighted Directed

## • Example 2: Weighted Directed

- Vertices  $\{n_1, n_2, n_3, n_4\}$

- Edges

$\{(n_1, n_2, 2), (n_3, n_1, 2), (n_1, n_4, 3), (n_2, n_4, 9), (n_1, n_1, -1), (n_3, n_4, 7)\}$



Shortest path is now  
undefined for many nodes  
(it is  $-\infty$ )

For any path that goes through  
 $n_1$ , you can get a shorter path  
by going around  $n_1$  one more  
time.

# Representing Graphs

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## Representation 1: Set based representation

- Many important operations are expensive

## Representation 2: Object oriented

- Node object contains list of neighboring nodes
- List of neighbors contains references pointing to nodes

## Representation 3: Matrix

- Very inefficient for sparse matrices

## Representation 4: Implicit

- Graph represented as a function
- Sometimes the only mechanism that works when the graph is too big for any other representation



# dot: Graph Visualization Language

Very simple syntax:

```
digraph name{  
    v1 -> v2[label="something"];  
    v2->v3;  
}
```

You can find many visualizers on the web

- GraphViz is probably the most popular one

- <http://www.graphviz.org/>

- There is a simple web viewer here:

- <http://sandbox.kidstrythisathome.com/erdos/>

but it only supports a small number of nodes and edges if you want to play with this, it's better to download the stand alone version

# **Key algorithmic questions**

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Is there a path between two nodes?

What is the shortest path between two nodes?

Can we order the nodes in a graph so all nodes come before their successors in the graph?