MASARYK UNIVERSITY
FACULTY OF INFORMATICS

Parameter Synthesis from Hypotheses Formulable in CTL Logic

BACHELOR THESIS

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Declaration

Hereby I declare, that this paper is my original authorial work, which I have worked out by my own. All sources, references and literature used or excerpted during elaboration of this work are properly cited and listed in complete reference to the due source.

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Advisor: John Foo, Ph.D.

Acknowledgement

I would like to thank my supervisor . . .

Abstract

The aim of the bachelor work is to provide . . .

Keywords

keyword1, keyword2 . . .

Contents

1 Introduction

Lorem Ipsum [1]

2 Algorithm

In this chapter, we describe the distributed algorithm that computes the assumption function $\mathcal{A}$.

2.1 Distributed Environment

In this section, we briefly describe the distributed environment assumed by our algorithm, in order to prevent any possible confusion.

We assume a distributed environment with fixed number of reliable processes connected by reliable, order-preserving channels (The order preservation can be relaxed to some extent). We also assume that each process has a fixed identifier and the set of all process identifiers is equal to the result set of the partition function. Each process can communicate directly (using the function SEND) with any other process (assuming it knows other process's identifier) and all messages that can't be processed directly are stored in a queue until they can be processed.

Several parts of the algorithm do not have explicit termination (they terminate by reaching deadlock - no messages are exchanged between processes). In such cases, suitable termination detection algorithm is employed to detect this deadlock and terminate computation properly. Our implementation uses Safra's algorithm [CITATION] for this purpose, but the related code has been skipped for easier readability.

We also divide algorithm description into three main parts: Process variables, Initialization and Message handler. First section describes data structures stored in process memory available during whole computation. Initialization section is executed exactly once and no messages can be received until it's finished. Message handler defines what should happen when message is received.

2.2 Algorithm outline

The main idea of the algorithm is described in REFERENCE and resembles other CTL model checking algorithms.

```
1: procedure CHECKCTL(φ, K = (id, f, P, S, I, -p->, L))
2:     A ← {(p, s, α, tt) | p ∈ P ∧ α ∈ L(s)}
3:     for all i < |φ| do
4:         for all ψ in cl(φ) where |ψ| = i do
5:             A ← CHECKFORMULA(ψ, K, A)
6:         end for
7:     end for
8: end procedure
```

The algorithm starts by initializing the assumption function using the labeling function defined in kripke fragment. After that, it traverses the structure of formula, starting from smallest formulas and uses computed results to process more complex formulas. Function CHECKFORMULA computes all states and colors where formula ψ holds and returns assumption function updated accordingly. This is done using the local information contained in given kripke fragment, assumptions previously computed for smaller formulas and also by communicating with other processes. Note that only assumptions relevant for particular process are computed and returned (each process has information only about it's own state space).

2.3 Common operations

In this section, we define functions used to simplify the algorithm description. Let us fix a formula ϕ and a parametrised kripke fragment $\mathcal{K} = (id, f, \mathcal{P}, S, I, \xrightarrow{p}, L)$ as an input of the algorithm.

Intuitively, function $validStates : cl(\phi) \times AS_{\mathcal{K}}^{\phi} \rightarrow S \times 2^{\mathcal{P}}$ computes a set of states and parameters where truth of given formula is assumed. It is also responsible for handling of boolean operators, since these can be computed without any inter-process communication.

$$validStates(\phi_1 \wedge \phi_2, \mathcal{A}) = \{(s, p) \mid \mathcal{A}(s, p, \phi_2) = \mathtt{tt} \wedge \mathcal{A}(s, p, \phi_2) = \mathtt{tt}\}$$

$$validStates(\phi_1 \vee \phi_2, \mathcal{A}) = \{(s, p) \mid \mathcal{A}(s, p, \phi_2) = \mathtt{tt} \vee \mathcal{A}(s, p, \phi_2) = \mathtt{tt}\}$$

$$validStates(\neg\phi, \mathcal{A}) = \{(s,p) \mid \mathcal{A}(s,p,\phi) = \texttt{ff}\}$$

$$validStates(\texttt{tt}, \mathcal{A}) = \{(s,p) \mid s \in S \wedge p \in \mathcal{P}\}$$

$$validStates(\texttt{ff}, \mathcal{A}) = \emptyset$$

$$validStates(\phi, \mathcal{A}) = \{(s,p) \mid \mathcal{A}(s,p,\phi) = \texttt{tt}\}$$

Another useful function is $validColours : cl(\phi) \times S \times AS_{\mathcal{K}}^{\phi} \to 2^{\mathcal{P}}$ which returns a set of parameters for which given formula is assumed to be true in given state.

$$validColours(\phi, state, \mathcal{A}) = \{p | (state, p) \in validStates\}$$

Function $predecessors : S \to S \times 2^{\mathcal{P}}$ computes set of direct predecessors of given node including the color sets labeling the appropriate transitions.

$$predecessors(to) = \{(from, P) \mid P = \{p \mid from \xrightarrow{p} to\}\}$$

Symmetrically, function $successors : S \to S \times 2^{\mathcal{P}}$ computes set of direct successors of given node including the color sets labeling the appropriate transitions.

$$predecessors(from) = \{(to, P) \mid P = \{p \mid from \xrightarrow{p} to\}\}$$

We also define function $update : AS_{\mathcal{K}}^{\phi} \times S \times 2^{\mathcal{P}} \times tcl(\phi) \to AS_{\mathcal{K}}^{\phi}$ which takes current assumptions, a state, set of parameters and a formula and returns assumptions updated so that for all parameters of the given set, formula holds in given state.

$update(\mathcal{A}, state, colours, \phi)$:
for all $p \in colours$ **do**
  Set $\mathcal{A}(state, p, \phi) = \texttt{tt}$
end for

2.4 Temporal Operators

In this section, we describe how the CHECKFORMULA is implemented for each of the temporal operators. Note that all of the following algorithms has implicit termination and therefore needs a proper termination detection algorithm to correctly terminate.

2.4.1 Exist Next Operator

1: **Process variables:**
2: $\mathcal{K} = (id, f, \mathcal{P}, S, I, \xrightarrow{p}, L)$ ▷ Kripke fragment
3: $\phi = \mathrm{EX}\phi_1$ ▷ CTL formula
4: $\mathcal{A}$ ▷ Initial assumption function
5: **procedure** INIT
6: **for all** $(state, colSet)$ **in** $validStates(\phi_1, \mathcal{A})$ **do**
7: **for all** $(pred, tranCol)$ **in** $predecessors(state)$ **do**
8: SEND$(f(state), (pred, colSet \cap tranCol))$
9: **end for**
10: **end for**
11: **end procedure**
12: **procedure** RECEIVE$(colSet, to)$
13: $\mathcal{A} \leftarrow update(\mathcal{A}, to, colSet, \phi)$
14: **end procedure**

The simplest of temporal operators is the EX operator. During initialization, all states and colors where ϕ holds are computed. For each of such states, all predecessors are considered and appropriate message that will cause assumption update is sent.

2.4.2 Exist Until Operator

1: **Process variables:**
2: $\mathcal{K} = (id, f, \mathcal{P}, S, I, \xrightarrow{p}, L)$ ▷ Kripke fragment
3: $\phi = \mathrm{E}\phi_1\mathrm{U}\phi_2$ ▷ CTL formula
4: $\mathcal{A}$ ▷ Initial assumption function
5: **procedure** INIT
6: **for all** $(state, colSet)$ **in** $validStates(\phi_2, \mathcal{A})$ **do**
7: $\mathcal{A} \leftarrow update(\mathcal{A}, state, colSet, \phi)$
8: **for all** $(pred, tranCol)$ **in** $predecessors(state)$ **do**
9: SEND$(f(state), (pred, colSet \cap tranCol))$
10: **end for**
11: **end for**
12: **end procedure**
13: **procedure** RECEIVE$(colSet, state)$
14: $colSet \leftarrow colSet \cap valid(\phi_1, to, \mathcal{A})$
15: **if** $colSet \neq \emptyset$ **and** $colSet \setminus valid(\phi, to, \mathcal{A}) \neq \emptyset$ **then**

16: $\mathcal{A} \leftarrow update(\mathcal{A}, to, colSet, \phi)$
17: **for all** $(pred, tranCol)$ **in** $predecessors(to)$ **do**
18: SEND$(f(pred), (pred, colSat \cap tranCol))$
19: **end for**
20: **end if**
21: **end procedure**

The EU operator is a little more complex, but again fairly simple. The algorithm starts by computing all states and colours where ϕ_2 is true. Starting from these states, a backpropagation of parameter sets along the reversed transitions is performed. During the computation, the propagated parameter set is updated to reflect the validity of ϕ_1 and the validity of transitions used along the path. Note that backpropagation is stoped as soon as there is no new information computed (*colSat* is either empty or equal to already computed assumptions).

2.4.3 All Until Operator

1: **Process variables:**
2: $\mathcal{K} = (id, f, \mathcal{P}, S, I, \xrightarrow{p}, L)$ ▷ Kripke fragment
3: $\phi = \mathrm{A}\phi_1\mathrm{U}\phi_2$ ▷ CTL formula
4: $T = \xrightarrow{p}$ ▷ Uncovered edges
5: $\mathcal{A}$ ▷ Initial assumption function
6: **procedure** INIT
7: **for all** $(state, colSet)$ **in** $validStates(\phi_2)$ **do**
8: **for all** $(pred, tranCol)$ **in** $predecessors(state)$ **do**
9: SEND$(f(state), (state, pred, \mathcal{P} \cap tranCol))$
10: **end for**
11: **end for**
12: **end procedure**
13: **procedure** RECEIVE$(colSet, from, to)$
14: $T \leftarrow T \setminus \{(to, p, from) | p \in colSet\}$
15: $colSet \leftarrow \{p | p \in colSet \wedge \forall s_2 \in S.(to, p, s_2) \notin T\}$
16: $colSet \leftarrow colSet \cap valid(\phi_1, to, \mathcal{A})$
17: **if** $colSet \neq \emptyset$ **and** $colSet \setminus valid(\phi, to, \mathcal{A}) \neq \emptyset$ **then**
18: $\mathcal{A} \leftarrow update(\mathcal{A}, to, colSet, \phi)$
19: **for all** $(pred, tranCol)$ **in** $predecessors(to)$ **do**

```
20:             SEND(f(pred), (to, pred, colSet ∩ tranCol))
21:         end for
22:     end if
23: end procedure
```

The AU operator is the most complex one to handle. As opposed to EX, which requires at least one valid successor to be true, AU requires that all successors of the specific node are valid. In order to compute such information, we create a copy of transition relation and call it T.

During the computiation, T is modified in such way, so that we can guarantee that if edge is not present in T, this edge leads to a state where either ϕ_2 or Aϕ_2Uϕ2 holds.

A First appendix

Hello, here is some text without a meaning. This text should show what a printed text will look like at this place. If you read this text, you will get no information. Really? Is there no information? Is there a difference between this text and some nonsense like “Huardest gefburn”? Kjift – not at all! A blind text like this gives you information about the selected font, how the letters are written and an impression of the look. This text should contain all letters of the alphabet and it should be written in of the original language. There is no need for special content, but the length of words should match the language.

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B Another appendix

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[1] S. author, "Assumption-based distribution of CTL model checking," *STTT*, vol. 7, no. 1, pp. 61–73, 2005.