CS536 Abstract Syntax Tree (AST) and Syntax Directed Translation (SDT)

A Sahu
CSE, IIT Guwahati

http://jatinga.iitg.ac.in/~asahu/cs536/

Outline

- Basic of AST
- Basic of Syntax Directed Translation
- Intermediate Representation

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Doubt of Last Class

```
main(){
    char *A="H";
    printf("%s",A);
    }

Text Data Bss

1578 600 8
```

LLVM Installation

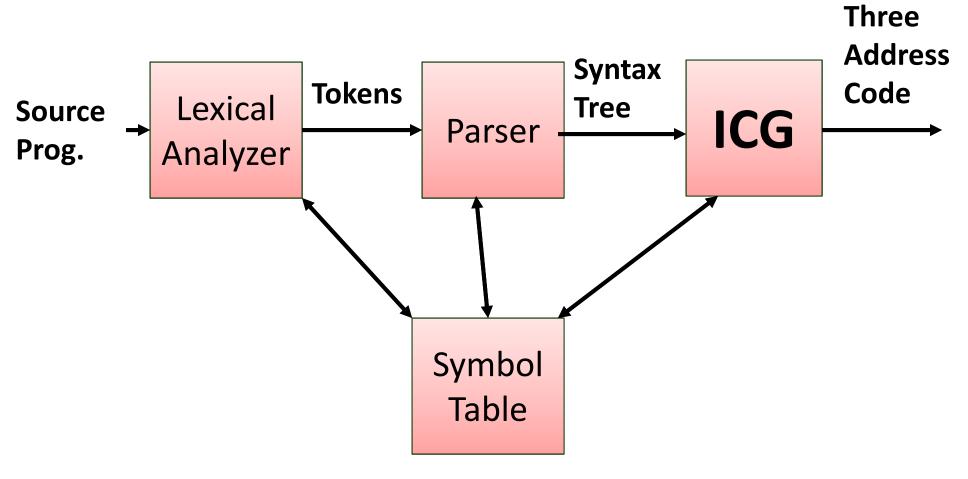
- \$sudo apt-get install cmake clang swapsapce
- Download Ilvm-project-Ilvmorg-12.0.0.tar.gz (or recent version) from LLVM official website "Source code (tar.gz)"
- \$mkdir ~/LLVM; cd ~/LLVM; mkdir install; mv ~/Download/ll*tgz.; tar -xvzf ll*tgz
- \$cd llvm-project-llvmorg-12.0.0/
- \$mkdir build; cd build
- \$cmake --clean-first -G "Unix Makefiles"
 DLLVM_PARALLEL_LINK_JOBS=1 DCMAKE_INSTALL_PREFIX=\$(HOME)/LLVM/install/ ../Ilvm
- \$make -j1 //This may take some hours
- \$make install

Test Installation of LLVM

- How test a in built pass/transform
- cd \$(HOME)/LLVM/llvm-project-llvmorg-12.0.0/llvm/test/Transforms/HelloNew
- \$\$(HOME)/LLVM/install/bin/opt -disableoutput -passes=helloworld helloworld.ll
- \$\$(HOME)/LLVM/llvm-project-llvmorg-12.0.0/build/bin/opt -disable-output passes=helloworld helloworld.ll

Basic of Syntax Directed Translation

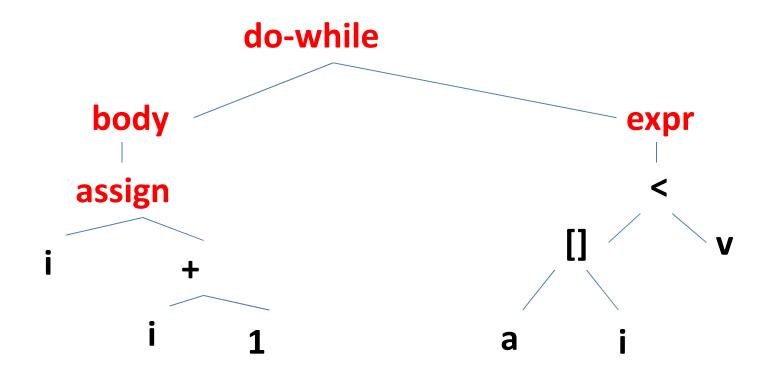
A model of Compiler Front end



ICG: Intermediate Code Generator

Abstract Syntax Tree: AST

- Syntax tree: hierarchical syntactic structure of the source program
- AST for: do i=i+1; while (a[i] < v);



Syntax Definition

 A grammar naturally describe the hierarchical structure of the most program

if (expression) statement else statement

Production rule: Can have the form

stmt → if (expr) stmt else stmt

- In a production: if, else, (,) are terminals
 - The term expr, stmt are non-terminal
 - Can have the form (again)

Definition of Context Free Grammar

- CFG has four components
- A set of terminal symbols (referred as token)
 - Elementary Symbols of the Grammar
- S set of non-terminals (NT/syntactic variables)
- A set of production rules
 - Each production of NT called head/left side
 - Arrow and a sequence of T and/or NT called body/right side production
- A designation of one NT as Start symbol

Production Example

Expression: list of digits separated by plus and minus signs

```
list → list + digit
list → list – digit
list → digit
digit → 0|1|2|3|4|5|6|7|8|9
```

```
list \rightarrow list + digit | list – digit | digit digit \rightarrow 0|1|2|3|4|5|6|7|8|9
```

Derivation

- A grammar derives strings by beginning with start symbols
- And repeated replacing a non terminals by body of the production for that non terminals
- The terminal strings can be derived from the start symbol

```
9-5+2 is list, Can be derived as follows

list \rightarrow list +2  // list \rightarrow list + digit

\rightarrow list -5 + 2  // list \rightarrow list - digit

\rightarrow 9 - 5 - 2  // list \rightarrow digit
```

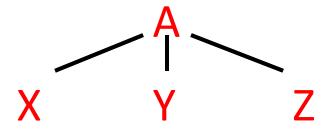
Another Production Example

```
call \rightarrow id (optparams)
optparams \rightarrow params | \epsilon
params \rightarrow params, param | param
```

- The term ϵ specifies the empty string
- This analogous/similar to earlier production

Parsing Trees

 A parse tree pictorially shows: How the start symbol of a grammar derives strings in language



- 1. The root is labeled by start symbol
- 2. Each leaf is labeled by a terminal or by ϵ
- Each interior node is labeled by a non-terminals
- 4. If NT A $X_{1,} X_{2,} ... , X_n$ are labeled children of A from left to right then there must be production A \rightarrow $X_{1,} X_{2,} ... , X_{n,}$ where X1, X2, ...Xn are eithe NT or T,
- 5. If A $\rightarrow \epsilon$, then A may have single child ϵ

Parsing Trees: Properties

- A tree consists of one or more nodes
- Exactly one root node in a Tree
 - Root have no-parent, it is top node
 - Other node have exactly one parent
- Leaf: node with no children
- N is parent of M, M is child of N, Children of one node is Siblings, Ordered from left to right
- Descendent (self, child*), Ancestor (self, parent*)

Thanks