CS536 Prog. Language, Memory, Mutifile and AST

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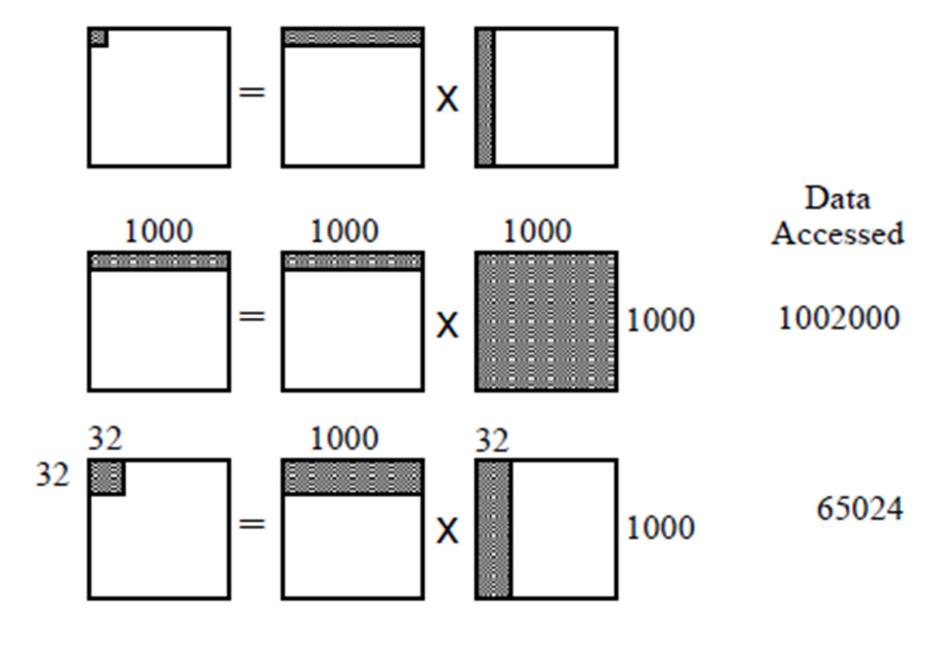
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Outline

- Last Class: Blocking in Matrix Multiplication
- Programming Language Basic
- Memory Layout of a C Program
- Compiling Multiple C Files
- Basic of Syntax Directed Translation
- Intermediate Representation

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Blocking for Matrix Multiplication



Blocking for Matmul: Original Code

```
for (i= 0; i< n; i++) {
  for (j = 0; j < n; j++) {
    for (k = 0; k < n; k++) {
      Z[i,j] = Z[i,j] + X[i,k]*Y[k,j];
```

Row major access

Column major access

Blocking for Matmul: Original Code

```
for (i= 0; i< n; i++) {
  for (k = 0; k < n; k++) {
    for (j = 0; j < n; j++) {
      Z[i,j] = Z[i,j] + X[i,k]*Y[k,j];
```

Row major access

Column major access

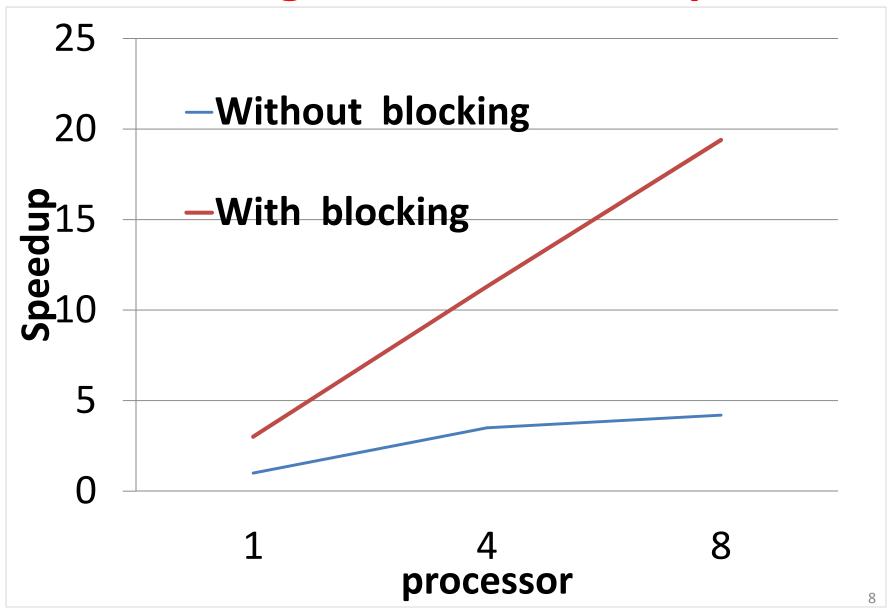
Blocking for Matmul: Stripmine 2 outerloop

```
for (ii = 0; ii < n; ii = ii+B) {
  for (i= ii; i< min(n,ii+B); i++) {
    for (jj= 0; jj< n; jj= jj+B) {
       for (j = jj; j < min(n, jj+B); j++) {
         for (k = 0; k < n; k++) {
           Z[i,j] = Z[i,j] + X[i,k]*Y[k,j];
```

Blocking for Matmul: permute

```
for (ii = 0; ii < n; ii = ii+B) {
  for (jj= 0; jj< n; jj= jj+B) {
    for (k = 0; k < n; k++) {
       for (i= ii; i< min(n,ii+B); i++) {
         for (j = jj; j < min(n,jj+B); j++) {
             Z[i,j] = Z[i,j] + X[i,k]*Y[k,j];
```

Blocking for Matmul: Impact



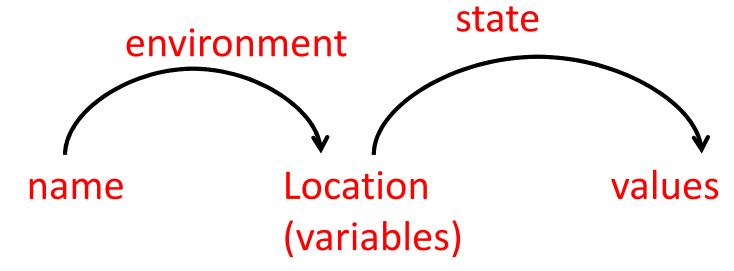
Programming Language Basic

Static Vs Dynamic Variable

- Static vs Dynamic Distinction
- Static Scope: Lexical/Compilation scope for a variable x
 - -Global x,
 - static int x
- Dynamic Scope: as the program run same used of x could refer to any several value of x

Environment and Scope

- Environment: mapping name to location in the memory (I-value of variable)
- States: mapping location to their values (r-value)



• int x=6; x have location &x, value =6

Scope and Block Structure

- Scope: Public, private, protected
- Static scope based on block: { ..}, begin..end
- C program consist of
 - top level declaration of variables
 - and functions
 - functions: may have variable declarations within them, scope is restricted with in that functions

Scope and Block Structure

```
main() {
                                             B1
int a=1;
int b=1;
                                  B2
      int b=2;
            int a=3;
            cout << a << b;
                                B3
           int b=4;
            cout << a << b;
                                B4
      cout << a << b;
cout << a << b;
```

Function Parameter Passing

- Actual parameter: used in the call of fun
- Formal parameter: used in fun definition
- Call by value: F(int A)
- Call by Reference: F(int *A), F(int &A)
- Call by name: F(int A)
- Aliasing: can refer to same location

```
F(int *A, int *B)//overlapped A & B
F(int * __restrict A, int *__restrict B)
```

Dynamic memory allocation

- Reduce wastage of memory
- Useful when data size is unknown before hand
- Array Declaration

```
int A[100];
```

- Easy, Not to use pointer, small size, known before
- Array Creation:
 - Not easy, use of pointer, typecast, lager size, necessary size

Memory management C: APIs

- Application program interfaces (APIs)
- Available function/APIs to manage memory
- Create/allocate/reserve space
 - malloc : memory allocation
 - calloc: memory allocation + initialization to 0
- Move a reserved space to another location
 - realloc: move the space to another location
- Destroy/de-allocate/free space
 - free:

Memory Allocation

- Memory can be allocated
- Declaring a variable

```
int A[100];
```

Explicitly requesting space

```
int *A;
A=(int*)malloc(sizeof(int)*100);
```

Example: Dynamic Array Allocation

- Given N persons (with their IQ level) in order
 - N may be dynamic, variable
- A person decide He/She is intelligent or dumb
- Decides locally:
 - If his/her IQ level is greater than equal to average of IQ level of both neighbors
 - Left neighbor and right neighbor

Example: Dynamic Array Allocation

```
main(){
   int *IQScore, *Intelligent, i, N;
  printf("Input N:"); scanf("%d", &N);
   IQScore=(int*)malloc(N*sizeof(int));
   Intelligent = (int*) calloc(N*sizeof(int));
   for (i=0; i < N; i++) scanf("%d", & IQScore[i]);</pre>
   for (i=1; i<N-1; i++) {</pre>
      if(IOScore[i] > = (IOScore[i-1] + IOScore[i+1])/2)
        Intelligent[i]=1; else Intelligent[i]=0;
       printf(" I am %d person %s\n", i,
           Intelligent[i]?"YES":"NO");
   free(IQScore); free(Intelligent);
```

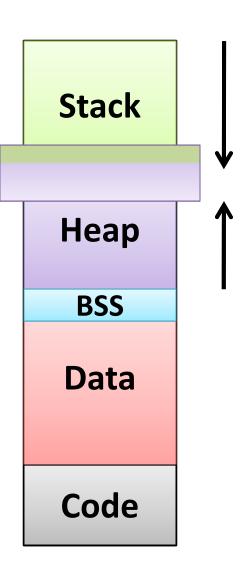
- Program: Input, Output, Processing
- Code (Instruction), Data (Stack, Heap)
- To store: Require memory
 - Input data, output data, intermediate data
- Memory can be allocated
 - Declaring a variable

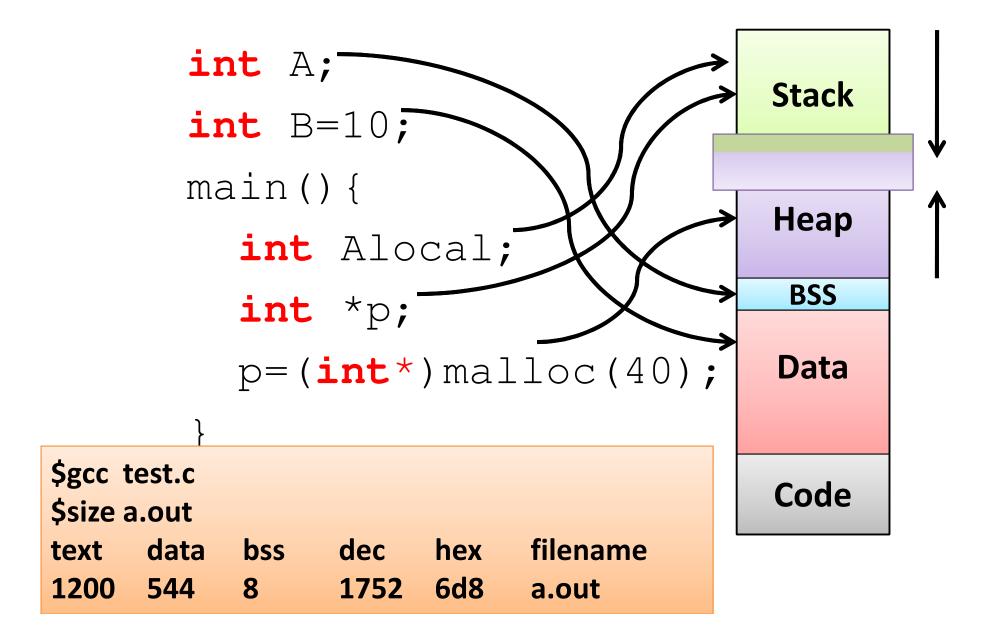
```
int A[100];
```

Explicitly requesting space

```
int *A;
A=(int*)malloc(sizeof(int)*100);
```

- Stack
 - automatic (default), local
 - Initialized/uninitialized
- Data
 - Global, static, extern
 - BSS: Block Started by Symbol
 - BBS: Uninitialized Data Seg.
- Code : program instructions
- Heap
 - malloc, calloc





Compiling Multiple C Files

Compiling multiple Files

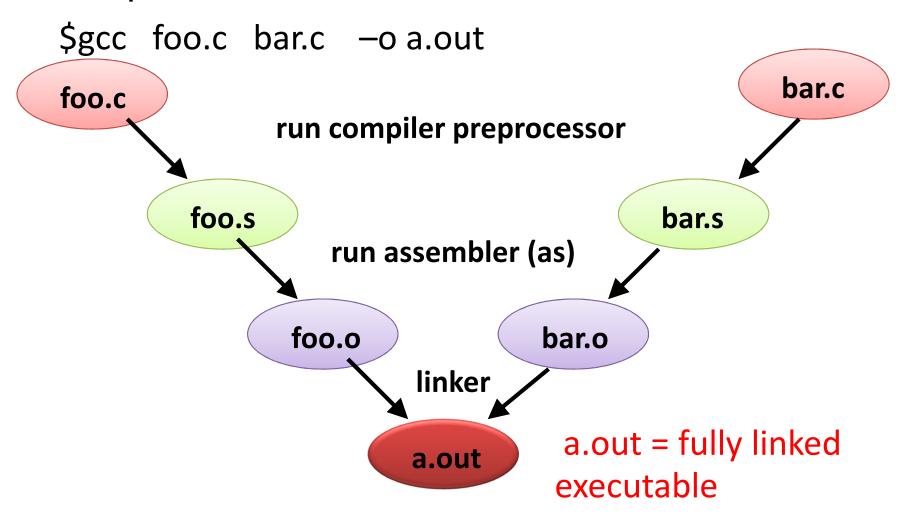
```
//foo.c
int foo3x(int x){
    return 3*x;
}
```

```
//bar.c
int main(){
    int x;
    x=foo3x(10);
    printf("%d",x);
    return 0;
}
```

- \$ gcc –c foo.c
- \$ gcc –c bar.c
- \$ gcc foo.o bar.o
- \$./a.out

Linker and Loader

Compiler in Action...



What is Linker?

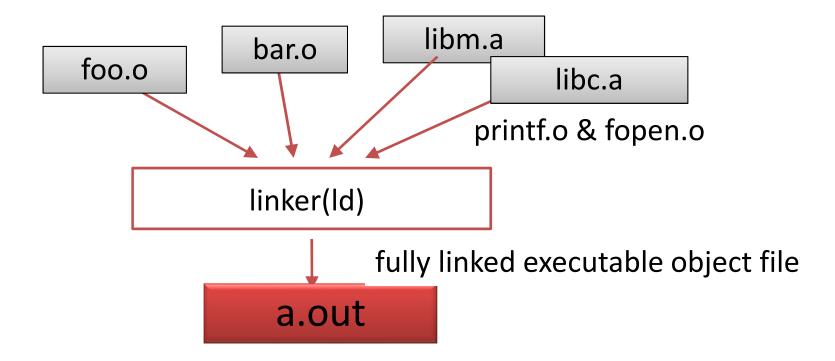
- Combines multiple relocatable object files
- Produces fully linked executable directly loadable in memory
- How?
 - Symbol resolution associating one symbol definition with each symbol reference
 - Relocation relocating different sections of input relocatable files

Object files

- Types
 - Relocatable : Requires linking to create executable
 - Executable : Loaded directly into memory for execution
 - Shared Objects : Linked dynamically, at run time or load time

Linking with Static Libraries

- Collection of concatenated object files stored on disk in a particular format – archive
- An input to Linker
 - Referenced object files copied to executable



Creating Static Library

```
//foo.c
int foo3x(int x){
    return 3*x;
}
```

- \$ gcc -c foo.c
- \$ ar rcs libfoo.a foo.o
- \$ gcc bar.c -L. -Ifoo
- \$./a.out

```
int main(){//bar.c
    int x;
    x=foo3x(10);
    printf("%d",x);
    return 0;
}
```

//it create libfoo.a

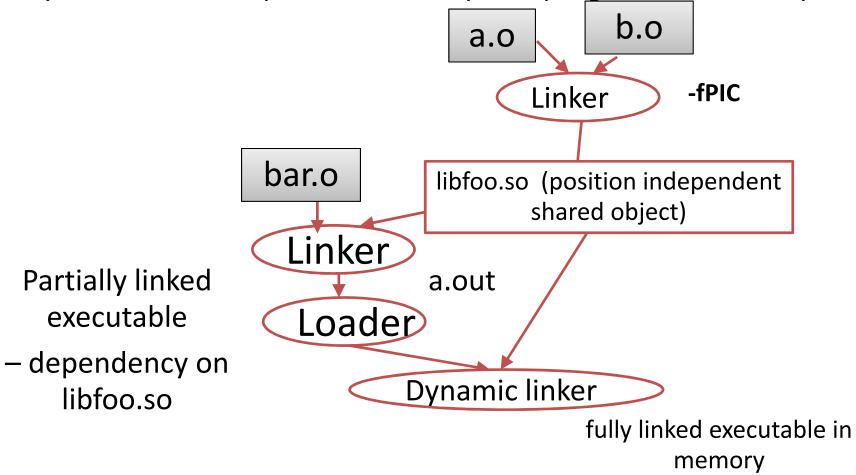
<u>Dynamic Linking – Shared Libraries</u>

- Addresses disadvantages of static libraries
 - Ensures one copy of text & data in memory
 - Change in shared library does not require executable to be built again
 - Loaded at run-time by dynamic linker, at arbitrary memory address, linked with programs in memory
 - On loading, dynamic linker relocates text & data of shared object

Shared Libraries ..(Cntd)

- Linker creates libfoo.so (PIC) from a.o b.o
- a.out partially executable dependency on libfoo.so

Dynamic linker maps shared library into program's address space



Creating Dynamic Library

```
//foo.c
int foo3x(int x){
    return 3*x;
}
```

```
int main(){//bar.c
    int x;
    x=foo3x(10);
    printf("%d",x);
    return 0;
}
```

- \$gcc -c -fPIC foo.c
- \$gcc -shared -Wl,-soname,libfoo.so.1 -o libfoo.so.1 foo.o
- \$ gcc bar.c -L. —Ifoo
- \$ export LD_LIBRARY_PATH=.
- \$./a.out

Thanks