# Algebraic Topology

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# **Contents**

| 1 Category Theory |     |                          | 3  |  |
|-------------------|-----|--------------------------|----|--|
|                   | 1.1 | Lecture 1 8/26           | 3  |  |
|                   | 1.2 | Lecture 2 8/28           | 5  |  |
| 2                 | Hon | Homologies, Cohomologies |    |  |
|                   | 2.1 | Lecture 3 9/4            | 8  |  |
|                   | 2.2 | Lecture 4 9/9            | 1  |  |
|                   | 2.3 | Lecture 5 9/11           | 13 |  |
|                   | 2.4 | Lecture 6 9/16           | 16 |  |

# Chapter 1

# **Category Theory**

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### 1.1 Lecture 1 8/26

**Definition 1.1** (Category). A category C consists of the following data:

- 1. A collection of objects denoted as Ob(C)
- 2. Given two objects  $X, Y \in Ob(\mathcal{C})$ , a collection of morphisms between  $X, Y, f : X \to Y$ , denoted as  $mor_{\mathcal{C}}(X, Y)$ .
- 3. (Composition) We have  $mor_{\mathcal{C}}(X,Y) \times mor_{\mathcal{C}}(Y,Z) \to mor_{\mathcal{C}}(X,Z)$  that satisfies associativity

$$f \circ (g \circ h) = (f \circ g) \circ h$$

4. (Identity) There is a distinguished morphism for each X,  $Id_{\mathcal{C}}(X,X)$  such that given any  $f \in mor(X,Y)$ , we have  $f \circ id_X = id_Y \circ f = f$ .

In this course, we will make the assumption that in all the categories that we work with, Ob(C) need not be a set, but given any  $X, Y \in Ob(C)$ , mor(X, Y) will always be a set. Now we talk bout some examples of categories.

**Example 1.1** (Sets). Let Ob(Sets) be all the sets in the universe. Given X, Y sets, mor(X, Y) be all the set maps from X to Y, and  $id_X$  is the identity map.

**Example 1.2** (Top). Let Ob(Top) be all the topological spaces, and mor(X, Y) be all the continuous maps from X to Y.

**Example 1.3** (Vect<sub> $\mathbb{F}$ </sub>). Let  $\mathbb{F}$  be a field, and let Ob be all the  $\mathbb{F}$ -vector spaces. Then mor(V, W) is all the  $\mathbb{F}$ -linear homomorphisms from V to W, where  $id_V$  is the identity homomorphism.

**Example 1.4** (Posets). Fix a poset P, let Ob(P) be the collection of elements in P, and given p,q we define

$$mor(p,q) = \begin{cases} *, \text{ if } q \leq p \\ \varnothing, \text{ otherwise} \end{cases}$$

### Problem 1.1. HW(Q1): check this is a category

**Example 1.5** (Opposite category). Given a category C, there is another category called the opposite category, denoted as  $C^{op}$ , where

- 1. The objects are the same as C
- 2. Given  $X, Y \in Ob(C^{op})$ , we have  $mor_{op}(X, Y) := mor_{\mathcal{C}}(Y, X)$ .
- 3. Moreover, given  $f \in mor_{op}(X,Y), g \in mor_{op}(Y,Z)$ , then  $g \circ f$  in  $C^{op}$  is  $f \circ g : Z \to X$ .

Naturally, we define isomorphisms now.

**Definition 1.2** (isomorphism). Given a category C, and a morphism  $f \in mor_C(X,Y)$ , we say f is an isomorphism if there exists  $g \in mor_C(Y,X)$  such that

$$f \circ g = Id_Y, g \circ f = Id_X$$

Now we introduce maps between categories.

**Definition 1.3** (functor). Given categories C, D, a functor  $F: C \to D$  is the following;

- 1. Given an object X in C, F(X) is an object in D.
- 2. Given a morphism  $f: X \to Y$ , F(f) is a functor  $F(f): F(X) \to F(Y)$ . Moreover, it satisfies the following:
  - (a)  $F(id_X) = id_{F(X)}$
  - (b)  $F(f \circ g) = F(f) \circ F(g)$ . Alternatively, we can rewrite this condition as the following:

$$\begin{array}{ccc} mor(X,Y) \times mor(Y,Z) & \longrightarrow & mor(X,Z) \\ & & \downarrow^{mor(F) \times mor(F)} & & \downarrow^{mor(F)} \\ mor(F(X),F(Y)) \times mor(F(Y),F(Z)) & \longrightarrow & mor(F(X),F(Z)) \end{array}$$

such that this diagram commutes.

#### Problem 1.2. HW(Q2): functors take isomorphisms to isomorphisms.

Now we talk about some examples of functors.

**Example 1.6.**  $F: Top \rightarrow Set$ , where  $X \mapsto X$ , where the latter is a set, and  $f \mapsto f$  as set maps.

**Example 1.7.** Let  $\mathbb{F}$  be a field, and  $F: Sets \to \text{Vect}_{\mathbb{F}}$ , where  $X \mapsto \mathbb{F}\langle X \rangle$ , where  $\mathbb{F}\langle X \rangle$  is the free vector space over  $\mathbb{F}$  on the set X.

Problem 1.3. HW(Q3): extend this to a functor by defining mor(f) and show this is a functor.

**Example 1.8.** Let  $\mathbb{F}$  be a field, then the following is a functor,  $F: Sets^{op} \to Vect_{\mathbb{F}}$ , where

$$hF: X \mapsto Maps(X, \mathbb{F})$$

1.2. LECTURE 2 8/28 5

**Problem 1.4. HW(Q4)**: show this extends to a functor by defining F(f), and show it is a functor.

### 1.2 Lecture 2 8/28

**Definition 1.4** (contravariant functor). Let  $F: \mathcal{C} \to \mathcal{D}$  is a contravariant functor from  $\mathcal{C}^{op} \to \mathcal{D}$ , (equivalently,  $\mathcal{C} \to \mathcal{D}^{op}$ ).

**Problem 1.5. HW(Q5):** Show that the following functor F from  $Vect_{\mathbb{F}}$  to  $Vect_{\mathbb{F}}$  extends to a contravariant functor, where

$$Ob_F: V \mapsto V^* = Hom(V, \mathbb{F})$$

i.e., define the morphism function and show it is a contravariant functor.

We remark that we can define a category of categories: let Cat be the category of categories, with morphisms as functors, and note that objects or morphisms in this case are both not sets!

**Definition 1.5** (natural transformation). Given functors  $F, G : \mathcal{C} \to \mathcal{D}$ , a natural transformation T from F to G is the following:  $T : F \Rightarrow G$ :

- 1. given object  $X \in Ob(\mathcal{C})$ ,  $T(X) \in mor(F(X), G(X))$
- 2. Given  $f \in mor(X, Y)$ , the following diagram commutes:

$$F(X) \xrightarrow{F(f)} F(Y)$$

$$T(X) \downarrow \qquad \qquad \downarrow T(Y)$$

$$G(X) \xrightarrow{G(f)} G(Y)$$

where  $mor_F$ ,  $mor_G$  is the identification function on morphisms by functors F, G

If for all X, T(X) is an isomorphism, then this natural transformation is called a natural isomorphism.

In other words, this natural transformation is how one takes a functor F and turn it to another functor G. We will (in a homework) show there exists natural transformation between the following two functors.

**Example 1.9.** Consider  $F, G : Vect_{\mathbb{F}} \to Vect_{\mathbb{F}}$ , define

$$F(V) = V \otimes_{\mathbb{F}} V/_{\langle a \otimes b - b \otimes a \rangle} = V \otimes_{\mathbb{F}} V/\Sigma_2, G(V) = (V \otimes_F V)^{\Sigma_2} = \{\alpha \in V \otimes_{\mathbb{F}} V : \sigma(\alpha) = \alpha\}$$

Both are vector spaces are fixed under "swaps." Then a natural transformation can be defined as follows T(V):

$$T(V): a \otimes b \mapsto a \otimes b + b \otimes a$$

**Problem 1.6. HW(Q6):** For the above F, G

- 1. Show that T defines a natural transformation from F to G.
- 2. Find conditions on  $\mathbb{F}$  for T being a natural isomorphism.

Next we define limits and colimits. Let C, D be categories, d be an object in D, then we can define a functor  $F_d : C \to D$  such that for any object c in C,

$$F_d(c) = d, F_d(f) = Id_d$$

In other words, this is the "constant functor" on  $\mathcal{D}$ , i.e., every object is sent to d, and every morphism is sent to  $id_d$ .

**Definition 1.6** (colimit). Given any functor  $F: \mathcal{C} \to \mathcal{D}$ , the colimit of F, denoted as  $\operatorname{colim}(F)$  is an object in  $\mathcal{D}$  endowed with a natural transformation:

$$\varphi_F: F \Rightarrow F_{\operatorname{colim}(F)}$$

such that given any other object d in D and a natural transformation

$$\varphi: F \Rightarrow F_d$$

there exists a unique morphism in  $\mathcal{D}$ ,  $f:\operatorname{colim}(F)\to d$  making the following diagram commute: for any X,Y,g:



Next we prove some facts about colimits and give an example, where colim(F) exists.

**Proposition 1.1.** If colim F exists, then colim F is unique up to isomorphisms.

*Proof.* Let  $\operatorname{colim}(F)$ ,  $\operatorname{colim}(F)'$  be two  $\operatorname{colimits}$  that satisfy the criteria. They are both objects in  $\mathcal{D}$ , then we get a morphism  $f:\operatorname{colim}(F)\to\operatorname{colim}(F)'$ , and likewise  $g:\operatorname{colim}(F)\to\operatorname{colim}(G)'$ , then

$$f \circ g : \operatorname{colim}(F)' \to \operatorname{colim}(F)'$$

is the only morphism, and is the identity morphism. Similarly for  $g \circ f$ .

Next we demonstrate a fact via an example.

**Theorem 1.1.** Let  $\mathcal{C}$  be a category where  $Ob(\mathcal{C}), mor(X,Y)$  are all sets. Let  $F: \mathcal{C} \to \mathsf{Top}$  be any functor, then  $\mathsf{colim}(F)$  exists.

*Proof.* Define  $\operatorname{colim}(F) := \bigsqcup_{c} F(c) / \sim$ , where  $\sim$  is induced by the equivalence relation given by

$$y \sim F(f)y$$

where  $y \in F(C_1)$ ,  $f: C_1 \to C_2$ ,  $F(f)x \in F(C_2)$ . The natural transformation we endow on F as  $\varphi_F: F \Rightarrow F_{\text{colim}(F)}$ :

$$\varphi_F: F(C) \mapsto \bigsqcup_{C \in Ob(C)} F(C) / \sim$$

1.2. LECTURE 2 8/28 7

## **Problem 1.7. HW(Q7):** Show that colim(F), $\varphi_F$ is indeed a colimit.

We note that colimits also exist (the same argument goes through) if we replace Top with groups, sets, but with slightly different constructions, replacing disjoint unions with products, etc.

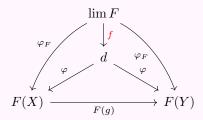
**Definition 1.7** (limit). Given a functor  $F: \mathcal{C} \to \mathcal{D}$ , the limit of F, denoted as  $\lim(F)$  is an object of  $\mathcal{D}$ , endowed with a natural transformation:

$$\varphi_F: F_{\lim(F)} \Rightarrow F$$

such that given any other object  $d \in Ob(\mathcal{D})$  and a natural transformation

$$\varphi: F_d \to F$$

there exists a unique  $f: \lim F \to d$  such that the following diagram commutes:



Just like colimits, limits are unique up to isomorphisms.

**Problem 1.8. HW(Q8):** Given  $F: \mathcal{C} \to \mathcal{D}$ , consider  $F^{op}: \mathcal{C}^{op} \to \mathcal{D}^{op}$ , then

$$\lim F = \operatorname{colim} F^{op}$$

The above problem is interpretation of diagrams and essentially we just reverse all the maps.

## **Chapter 2**

# Homologies, Cohomologies

### 2.1 Lecture 3 9/4

Today we define (co)chain complexes: let R be a commutative ring, let  $Mod_R$  denote the category of R-modules and R-module maps.

**Definition 2.1** (chain complex). A chain complex of *R*-modules is a collection of *R*-modules and *R*-modules maps

$$\cdots \to M_{i+1} \xrightarrow{\partial_{i+1}} M_i \xrightarrow{\partial_i} M_{i-1} \xrightarrow{\partial_{i-1}} \cdots$$

such that  $\partial_i \circ \partial_{i+1} = 0$  for all i. In other words, the image of previous map is contained in the kernal of the subsequent map. In short, we have

$$\partial^2 = 0$$

We will denote a chain complex by  $\{M.; \partial.^M\}$ .

Next we introduce morphisms between chain complexes.

**Definition 2.2** (morphism between complexes). Let  $\{M.; \partial.^M\}, \{N.; \partial.^N\}$ , a morphism  $\{f.\}$  between chain complexes is a "ladder" such that the following commutes:

$$\dots \longrightarrow M_{i+1} \xrightarrow{\partial_{i+1}^{M}} M_{i} \xrightarrow{\partial_{i}^{M}} M_{i-1} \xrightarrow{\partial_{i-1}^{M}} \dots$$

$$f_{i+1} \downarrow \qquad f_{i} \downarrow \qquad f_{i-1} \downarrow \qquad \vdots$$

$$\dots \longrightarrow N_{i+1} \xrightarrow{\partial_{i+1}^{N}} N_{i} \xrightarrow{\partial_{i}^{N}} N_{i-1} \xrightarrow{\partial_{i-1}^{N}} \dots$$

Moreover, we define composition of morphisms:

$$\{f.\} \circ \{g.\} := \{(f \circ g).\}$$

where  $\{g.\}:\{M.;\partial.^M\}\to\{N.;\partial.^N\}$ , and  $\{f.\}:\{N.;\partial.^N\}\to\{L.;\partial.^L\}$ , which is simply vertical stacking.

**Problem 2.1. HW(Q9):** Prove that chain complexes of R-modules form a category  $ch_R$ .

There are interesting functors  $F: \operatorname{ch}_R \to Mod_R$ , and we begin with the following one:

**Definition 2.3** ( $H_n$ , nth-homology). Given  $n \in \mathbb{Z}$ , there is a functor

$$H_n: \operatorname{ch}_R \to Mod_R$$

defined as follows:

$$H_n(\lbrace M.; \partial.^M \rbrace) := \ker \partial_n^M / Im \partial_{n+1}^M$$

and for  $f: \{M.; \partial.^M\} \to \{N.; \partial.^N\}$ , we define:  $H_n(f): H_n(\{M.; \partial.^M\}) \to H_n(\{N.; \partial.^N\})$ ,

$$H_n(f)[x] := [f_n(x)]$$

where  $[x] \in H_n(\{M.; \partial.^M\})$ .

*Proof.* We need to show  $H_n$  is well-defined on objects and morphisms. We need to check that  $Im\partial_{n+1} \subset \ker \partial_n$ . This is a consequence of  $\partial^2 = 0$ .

On morphisms: for  $x \in \ker \partial_n^M$ , we have  $f_n(x) \in \ker \partial_n^N$ . This is we have

$$\partial_n^N (f_n(x) = f_{n+1})(\partial_n^M(x)) = 0$$

Moreover, we need to check that this desn't depend on the choice of representatives, i.e., we can check that

$$Im\partial_{n+1}^M \mapsto 0$$

Let  $x = \partial_{n+1}^M(y)$ , we have

$$f_n(x) = f_n(\partial_{n+1}^M(y)) = \partial_{n+1}^N(f_{n+1}(y)) = 0$$

$$M_{n+1} \xrightarrow{\partial_{n+1}^M} M_n$$

$$f_{n+1} \downarrow \qquad \qquad \downarrow f_n$$

$$N_{n+1} \xrightarrow{\partial_{n+1}^N} N_n$$

Next we talk about homotopy between morphisms between chain complexes.

**Definition 2.4** (homotopy). Given two morphisms,  $f.,g.:M.\to N.$ , a chain homotopy h. between them is a collection of R-modules maps, for all  $n\in\mathbb{Z}$ ,

$$h_n:M_n\to N_{n+1}$$

such that

$$\partial_{n+1} \circ h_n + h_{n-1} \circ \partial_n = f_n - g_n$$

denoted as  $\partial h + h\partial = f - g$ .

$$M_{n+1} \longrightarrow M_n \xrightarrow{\partial_n^M} M_{n-1}$$

$$f_{n+1} \downarrow \qquad \qquad \downarrow f_n/g_n \downarrow \qquad \downarrow f_{n-1}$$

$$N_{n+1} \xrightarrow{\partial_{n+1}^N} N_n \longrightarrow N_{n-1}$$

**Problem 2.2. HW(Q10):** Show that homotopy is an equivalence relation between morphisms. Hint: replace  $h_n$  with  $-h_n: M_n \to N_{n+1}$ .

*Proof.* Reflexive is shown by defining  $h_n$  to be the zero map. For symmetry, we choose  $-h_n$ . Transitive is a ladder.

**Proposition 2.1.** Let h. be a chain homotopy between f. and g., then we have an equality

$$H_n(f.) = H_n(g.)$$

where  $H_n(f.), H_n(g.): H_n(M.) \to H_n(N.)$ .

*Proof.* Given  $[x] \in H_n(M.)$ , we have

$$H_n(f)[x] = [f_n(x)]$$

$$= [g_n(x) + \partial h.(x) + h.\partial(x)]$$

$$= [g_n(x) + \partial h.(x)]$$

$$= [g_n(x)]$$

$$= H_n(g)[x]$$

Next we define a new category.

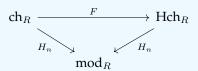
**Definition 2.5** ( $Hch_R$ ). Define the category  $Hch_R$  as follows:

- 1.  $Ob(Hch_R) = Ob(ch_R)$
- 2.  $mor_{\mathsf{Hch}_R}(M., N.) = mor_{\mathsf{ch}_R}(M., N.) / \sim$ , where  $\sim$  is the homotopy equivalence.

**Problem 2.3. HW(Q11):** Show that  $Hch_R$  is a category, admitting a functor

$$F: ch_R \to Hch_R$$

such that the following diagram commutes:



Next we introduce long and short exact sequences.

Definition 2.6 (exactness). Firstly, given a pair of *R*-module maps,

$$X_1 \xrightarrow{f} X_2 \xrightarrow{g} X_3$$

we say that the above is exact at  $X_2$  if  $\ker(g) = \operatorname{im}(f)$ . Hence given a sequence of R-module maps,

$$\cdots \to X_{i+1} \to X_i \to X_{i-1} \to \ldots$$

this is called a long exact sequence if it is exact at all  $X_i$ . Finally, given a pair of R-module maps,

$$0 \to X_i \xrightarrow{f} X_2 \xrightarrow{g} X_3 \to 0$$

This is a short exact sequence, and f is injective, g is surjective.

2.2. LECTURE 4 9/9

#### Problem 2.4. HW(Q12): Prove the following:

1. Given LES,

$$\cdots \to X_{i+1} \xrightarrow{f_{i+1}} X_i \xrightarrow{f_i} X_{i-1}$$

show the following is a short exact sequence:

$$0 \to \ker(f_i) \xrightarrow{i} X_i \xrightarrow{f_i} \ker(f_{i-1}) \to 0$$

2. Prove the 5-lemma. Given the below sequence, exact at positions  $X_i, Y_i$ , where i = 2, 3, 4, and assume the diagram commutes and if  $t_1, t_2, t_4, t_5$  are isomorphisms, show that  $t_3$  is also an isomorphism.

Next we state the most important theorem in chain complexes:

**Theorem 2.1** (The snake lemma). Let  $A : \xrightarrow{f} B : \xrightarrow{g} C$  be a SES of chain complexes, i.e.,

$$A_n \xrightarrow{f_n} B_n \xrightarrow{g_n} C_n$$

is a short exact sequence of all n. Then there exists a LES of homology groups.

$$H_{n}(A) \xrightarrow{\delta_{n-1}} H_{n}(C)$$

$$H_{n}(A) \xrightarrow{\delta_{n}} H_{n}(B) \xrightarrow{\delta_{n}} H_{n}(C)$$

$$H_{n-1}(A) \xrightarrow{\delta_{n-1}} H_{n-1}(B) \xrightarrow{\delta_{n-1}} H_{n-1}(C)$$

$$H_{n-2}(A)$$

## 2.2 Lecture 4 9/9

Today we prove the snake lemma. We will refer to this following diagram throughout the proof.

$$\begin{array}{c|cccc} A_{n+1} & \xrightarrow{f_{n+1}} & B_{n+1} & \xrightarrow{g_{n+1}} & C_{n+1} \\ \delta^A \Big\downarrow & \delta^B \Big\downarrow & \delta^C \Big\downarrow \\ A_n & \xrightarrow{f_n} & B_n & \xrightarrow{g_n} & C_n \\ \delta^A \Big\downarrow & \delta^B \Big\downarrow & \delta^C \Big\downarrow \\ A_{n-1} & \xrightarrow{f_{n-1}} & B_{n-1} & \xrightarrow{g_{n-1}} & C_{n-1} \\ \delta^A \Big\downarrow & \delta^B \Big\downarrow & \delta^C \Big\downarrow \\ A_{n-2} & \xrightarrow{f_{n-2}} & B_{n-2} & \xrightarrow{g_{n-2}} & C_{n-2} \end{array}$$

*Proof.* First we define the map  $\delta_n: H_n(C) \to H_{n-1}(A)$ . Let  $[x] \in H_n(C)$ , then  $x \in \delta^C$ , where  $\delta^C: C_n \to C_{n-1}$ . We define

$$\delta[x] = [y], y \in A_{n-1}$$

as follows: for  $x \in C_n$ ,  $g_n : B_n \to C_n$  is surjective, hence there exists  $b \in B_n$  such that  $g_n(b) = x$ . Then consider  $d = \delta^B(b)$ , since the diagram commutes, we have

$$d \in \ker g_{n-1} \Rightarrow d \in \operatorname{im} f_{n-1}$$

Let  $y \in A_{n-1}$  be this unique y such that  $f_{n-1}(y) = d$ , where uniqueness is by  $f_{n-1}$  is injective. This is indicated in the below diagram:

We first need to check that [y] does not depend on the choice of b. Let  $g_n(b_1) = g_n(b_2) = x$ , then

$$g(b_1 - b_2) = 0 \Rightarrow b_1 - b_2 = f_n(a), a \in A_n$$

let  $y_1, y_2$  be those determined by  $b_1, b_2$ , then

$$f_{n-1}(y_1 - y_2) = \delta^B(b_1 - b_2) = \delta^B(f_n(a)), a \in A_n$$

Because the following diagram commutes,

$$\begin{array}{ccc}
\mathbf{a} \in A_n & \xrightarrow{f_n} B_n \\
\delta^A \downarrow & & \downarrow \delta^B \\
A_{n-1} & \xrightarrow{f_{n-1}} B_{n-1}
\end{array}$$

we then have

$$y_1 - y_2 = \delta^A(a)$$

i.e.,  $[y_1] = [y_2]$ , as they only differ by im  $\delta$ .

**Problem 2.5. HW(Q13):** Check that if  $x \in \text{im } \delta^C$ , then  $\delta_n[x] = 0$ .

the proof is not finished, too lazy to tex it up

Next we review the tensor products of *R*-modules. We first review *R*-bilinear maps

**Definition 2.7** (bilinear maps). Let M, N, P be R-modules, an R-bilinear map  $f: M \times N \to P$  is a map such that

- 1. f is linear in both coordinates, we have  $f(m_1 + m_2, n) = f(m_1, n) + f(m_2, n)$ , and similarly,  $f(m, n_1 + n_2) = f(m, n_1) + f(m, n_2)$ .
- 2. For all  $r \in R$ , we have f(rm, n) = f(m, rn) = rf(m, n).

Next we define tensor products.

2.3. LECTURE 5 9/11 13

**Definition 2.8** (tensor product). A tensor product of  $M \times N$  is an R-module denoted by  $M \otimes_R N$  such that

1.  $M \otimes_R N$  comes endowed with an R-bilinear map

$$M \times N \xrightarrow{\varphi} M \otimes_R N$$

2. given any other R-bilinear map  $f: M \times N \to P$ , there exists a unique R-module map  $\psi$  such that the following diagram commutes:

$$M \times N \xrightarrow{\varphi} M \otimes_R N$$

$$\downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \downarrow \qquad \qquad \qquad \qquad \qquad \qquad$$

It is not clear that  $M \otimes_R N$  exists or not. In fact, they exist!

**Theorem 2.2** ( $M \otimes_R N$  exists). Define  $M \otimes_R N = R\langle M \times N \rangle / K$ , where  $R\langle M \times N \rangle$  is the free R-module on the set  $M \times N$ . We define K as the submodule generated by the following four relations:

- 1.  $\langle (m_1+m_2,n)\rangle \langle (m_1,n)\rangle \langle (m_2,n)\rangle$
- 2.  $\langle (m, n_1 + n_2) \rangle \langle (m, n_1) \rangle \langle (m, n_2) \rangle$
- 3.  $r\langle (m,n)\rangle \langle (rm,n)\rangle$
- 4.  $r\langle (m,n)\rangle \langle (m,rn)\rangle$

Moreover, the map  $\varphi: M \times N \to M \otimes_R N$  given by

$$(m,n) \mapsto \langle (m,n) \rangle := m \otimes_R n$$

**Problem 2.6. HW(Q14):** show that  $M \otimes_R N$  is a tensor product.

### 2.3 Lecture 5 9/11

We continue with the tensors of R-modules. Let  $f:A\to B$  an an R-module map, let N be some fixed R-module, then f induces maps:  $f\otimes id:A\otimes_R N\to B\otimes_R N$ ,

$$f \otimes id : a \otimes n \mapsto f(a) \otimes n$$

and  $id \otimes f : N \otimes f : N \otimes_R A \to N \otimes_R B$ :

$$id \otimes f : n \otimes a \mapsto n \otimes f(a)$$

Problem 2.7. HW(Q15(a)): Show that the following maps induce functors:

1.  $-\otimes_R N: Mod_R \to Mod_R$ , where

$$A \mapsto A \otimes_R N, f \mapsto f \otimes id$$

2.  $N \otimes_R -: Mod_R \to Mod_R$ , where

$$A \mapsto N \otimes_R A, f \mapsto id \otimes f$$

Problem 2.8. HW(Q15(b)): Show that one has the following natural isomorphisms:

- 1.  $0 \otimes_R M \cong 0$ , and  $0 \otimes_R \cong F_0$  (recall the definition of  $F_0$  as a functor).
- 2.  $R \otimes_R M \cong M$ , and  $R \otimes_R \cong id$ .
- 3.  $M \otimes_R N \cong N \otimes_R M$ , and  $M \otimes_R \cong \otimes_R M$ .
- 4.  $M \otimes_R (N \otimes_R K) \cong (M \otimes_R N) \otimes_R K$ .
- 5.  $(M \oplus N) \otimes_R K \cong (M \otimes_R K) \oplus (N \otimes_R K)$ .

For convenience, we introduce the following definition:

**Definition 2.9** (positively graded chain complex). A positively graded chain complex  $\{M.; \partial.^M\}$  is a chain complex so that  $M_i = 0$  for all i < 0. The category of positively graded chain complexed is denoted as  $ch_R^+$ .

We have our first important theorem for  $ch_R^+$ .

**Theorem 2.3.** There exists a functor  $\otimes_R$  and a natural transformation X such that the following diagram of functors commutes up to some X:

$$ch_{R}^{+} \times ch + R^{+} \xrightarrow{\otimes_{R}} ch_{R}^{+}$$

$$H_{i} \times H_{j} \downarrow \qquad \downarrow H_{i+j}$$

$$Mod_{R} \times Mod_{R} \xrightarrow{\otimes_{R}} Mod_{R}$$

where  $X: \otimes_R \circ (H_i \times H_j) \Rightarrow H_{i+j} \circ \otimes_R$  is a natural transformation.

We note that the existence of X means this diagram doesn't commute "on the nose," but these two composition functors are the same up to some natural transformation. Before we given the proof, we recall that  $Ob(C \times D) = Ob(C) \times Ob(D), mor((X,Y),(X',Y')) = mor(X,Y) \times mor(X',Y').$ 

*Proof.* We define  $\otimes_R$  of positively graded chain complexes as follows: let  $\{M.; \partial.^M\}, \{N.; \partial.^N\}$  be two PGCC. Define  $\{M \otimes_R N.; \partial.^M\}$ :

$$(M \otimes_R N) = \bigoplus_{i+j=n} (M_i \otimes_R N_j)$$

note that the RHS is always a finite sum. Moreover,  $\partial^{\otimes}$  is defined as follows:

 $\partial^{\otimes}: (M \otimes_R N)_n \to (M \otimes_R N)_{n-1}$  is defined on the component  $M_i \otimes_R N_j$  (from the RHS)

and

$$\partial^{\otimes}(m_i \otimes n_j) := \partial^M(m_i) \otimes n_j + (-1)^i m_i \otimes \partial^N(n_j)$$

2.3. LECTURE 5 9/11 15

It is easy to check that  $\partial^{\otimes} \circ \partial^{\otimes} = 0$ .

Now we've show  $ch_R^+ \otimes_R ch_R^+$  is well-defined, it remains to define X, the natural transformation. We define

$$X: H_i(M.) \otimes_R H_j(N.) \to H_{i+j}(M. \otimes_R N.)$$

again, it suffices to define X on elementary tensors.

$$X : [\alpha] \otimes [\beta] \mapsto [\alpha \otimes \beta]$$

we need to check that

- 1.  $\partial^{\otimes}(\alpha \otimes \beta) = 0$  if  $\partial^{M}(\alpha) = 0$  and  $\partial^{N}(\beta) = 0$ .
- 2. If  $\alpha = \partial(r)i$ , then notice that  $\partial^{\otimes}(r \otimes \beta) = \alpha \otimes \beta$ , similarly for  $\beta$ . This would show that X is well-defined.

It is straightforward to check that X commutes with morphisms in  $ch_R^+ \times ch_R^+$ .

Next we define cochain complexes and cohomologies.

**Definition 2.10** (cochain). A cochain of *R*-modules  $(M^{\bullet}, \partial_{M}^{\bullet})$  is a sequence of *R*-module maps:

$$\ldots \longrightarrow M^i \xrightarrow{\partial^i} M^{i+1} \xrightarrow{\partial^{i+1}} M^{i+2} \longrightarrow \ldots$$

such that  $\partial \circ \partial = 0$ .

Cochain complexes form a category, with morphisms  $\{f^{\bullet}\}\$  form a ladder:

The *n*-th cohomology of a cochain complex  $\{M^{\bullet}; \partial_{M}^{\bullet}\}$  is defined as:

$$H^n(M^{\bullet};\partial_M^{\bullet}) := \frac{\ker \partial^i : M^i \to M^{i+1}}{\operatorname{im} \partial^{i-1} : M^{i-1} \to M^i}$$

We remark that there is nothing unexpected here from what we learned about chain complexes. Namely, if we reindex  $\{M^{\bullet}; \partial_{M}^{\bullet}\}$ , this defines a chain complex with  $M'_{i} = M^{-i}$ . This implies that the snake lemme holds! (with  $\partial^{i}: H^{i}(C) \to H^{i+1}(A)$ ).

**Theorem 2.4.** There is a functor D and a natural transformation  $\beta$  such that the following diagram of functors commute up to the natural transformation  $\beta$ :

$$\begin{array}{ccc} ch_R^{op} & \xrightarrow{D} coch_R \\ H_n^{op} & \xrightarrow{\beta} & \downarrow H^n \\ Mod_R^{op} & \xrightarrow{\overline{D}} Mod_R \end{array}$$

where  $\overline{D}(M) = Hom_R(M, R)$ , and

$$D(\{M_{\bullet}; \partial_{\bullet}^{M}\})$$
 is defined as  $\{DM^{\bullet}; \partial^{\bullet}\}$ 

where

$$DM^n:=Hom_R(M_n,R), \partial^n:DM^n\to DM^{n+1}$$
 is the map induced by  $\partial_{n+1}:M_{n+1}\to M_n$ 

We observe that  $\partial^{n+1}\partial^n = 0$  since  $\partial_{n+2}\partial_{n+1} = 0$ .

## **Problem 2.9. HW(Q16):** Show that D is a functor.

Next we define the natural transformation  $\beta$ . We have  $\beta: H^n(DM) \to Hom_R(H_n(M_{\bullet}), R)$ , such that

$$\beta: [\varphi] \mapsto \beta[\varphi]$$

let  $[x] \in H_n(M_{\bullet})$ , where  $\beta[\varphi]([x]) = \varphi(x)$  (where  $\varphi \in Hom_R(M_n, \mathbb{R}), x \in M_n$ ).

*Proof.* We first need to show that  $\beta$  is well-defined. If  $x = \partial_{n+1}(y)$ , then consider

$$\beta[\varphi][x] = \varphi(x) = \varphi(\partial_{n+1}(y)) = \partial^n(\varphi)(y) = 0, x \in \ker \varphi$$

Conversely, if  $\varphi = \delta^{n-1}(\psi)$ , we have

$$\beta[\varphi][x] = \varphi(x) = \delta^{n-1}\psi(x) = \psi(\partial_n(x)) = 0$$

It remains to check that  $\beta$  commutes with morphisms in  $ch_R^{op}$  (which we will do next time).

### 2.4 Lecture 6 9/16

Today we continue our discussion of homological algebra. Let M be an R-module.

**Definition 2.11** (resolution). A resolution of M is a positively graded chain complex  $\{P_{\bullet}, \partial_{\bullet}\}$  such that

- 1.  $H_n(P_{\bullet}) = 0$  for all n > 0
- 2.  $H_0(P_{\bullet}) = \frac{P_0}{\operatorname{im} \partial_1} \cong M$ , where  $\partial_1 : P_1 \to P_0$ .

We say  $\{P_{\bullet}, \partial_{\bullet}\}$  is a free resolution if  $P_i$  is a free R-module for each i.

For resolutions, we prove the following two things: first, free resolutions always exist; second, every *R*-module map can be extended to a map between their resolutions (with extra assumptions) and these extensions are unique up to homotopies.

**Proposition 2.2.** For any M, a free resolution for M exists.

*Proof.* We proceed this inductively. Defien  $P_0$  to be  $R\langle M \rangle$ , where it is the free R-module defined on the set M. Note that

$$R\langle M\rangle \to M$$
 is surjective :  $\langle m\rangle \mapsto m$ 

Let *K* be the kernel of this map, we have an isomorphism:

$$\varepsilon: P_0/K \cong M$$

Define  $P_1$  as  $R\langle K \rangle$ , note that  $P_1 \rightarrow K$ , then we define

$$\partial_1: P_1 \to P_0$$

to be the composite

$$P_1 \twoheadrightarrow K \subset P_0$$

Now we consider  $P_2$ : let  $K_1 \subset P_1$  be the kernel of  $\delta_1$ , define  $P_2 = R\langle K_1 \rangle$ , then define  $\partial_2; P_2 \to P_1$  to be the composite"

$$P_2 \twoheadrightarrow K_1 \subset P_1$$

note that  $\ker \delta_1 / \operatorname{im} \delta_2 = K_1 / K_1 = 0$ . Then we define  $K_2 = \ker \delta_2$ , define  $P_3 = R \langle K_2 \rangle, \dots$ 

Just like the above proposition, the next theorem uses induction.

2.4. LECTURE 6 9/16 17

Theorem 2.5 (extension theorem). Let  $\{P_{\bullet}^M, \delta_{\bullet}^M, \varepsilon_M\}$  be a free resolution on M, and let  $\{P_{\bullet}^N, \delta_{\bullet}^N, \varepsilon^N\}$  be an arbitrary resolution of N. Then given a map of R-modules  $f: M \to N$ , we may extend it to a map of chain complexes:

$$f.: \{P^M_{ullet}, \delta^M_{ullet}\} \to \{P^N_{ullet}, \delta^N_{ullet}\}$$

such that the following diagram commutes:

$$\begin{array}{ccc} H_0(P_{\bullet}^M) & \xrightarrow{H_0(f_{\bullet})} & H_0(P_{\bullet}^N) \\ \downarrow_{\varepsilon_M} & & \downarrow_{\varepsilon_N} \\ M & \xrightarrow{f} & N \end{array}$$

Moreover, given any two extension  $f^1_{\bullet}, f^2_{\bullet}$  of f, we have a chain homotopy  $h_{\bullet}$  between  $f^1_{\bullet}, f^2_{\bullet}$ .

Remark: if  $f_{\bullet}$  makes the diagram commute, and  $g_{\bullet}$  is homotopic to  $f_{\bullet}$ , then  $g_{\bullet}$  also makes the diagram commutes: homotopy classes work the same on homologies (they are the same on the nose).