

# Algorithm Design and Analysis (Fall 2022)

## Midterm Exam

### Instructions:

1. The exam contains 5 pages. Please write your name and student ID on *every* page.
2. Do not start to read the questions until the examination starts.
3. This is an open book exam.
4. Exam duration: 8:00AM - 9:30AM.
5. For each question, you can continue your solution on the back of the page.

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Name: \_\_\_\_\_

Student ID: \_\_\_\_\_

1. (35 Points) Consider the special case of the Knapsack problem where each item  $i$ 's value  $v_i$  equals to its weight  $w_i$ . Suppose  $v_i = w_i$  is a (positive) integer power of 2 for each  $i = 1, \dots, n$ , and the capacity constraint  $K$  is an integer. Consider the following greedy algorithm:

- sort the items by descending order of weights (or values);
- for each item  $i$ , if putting  $i$  into the knapsack does not make the total weight more than  $K$ , put  $i$  in.

Does this algorithm always output an optimal solution? If so, prove it. If not, give a counterexample.

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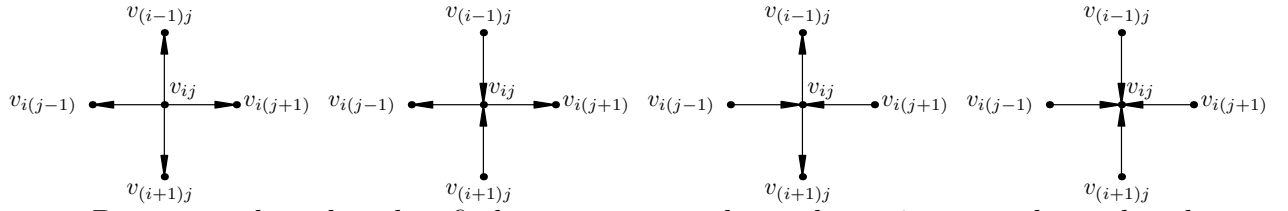
2. (a) (15 Points) Let  $G$  be a *directed* graph with  $n$  vertices labeled  $v_1, \dots, v_n$ . The graph contains  $n - 1$  edges  $e_1, \dots, e_{n-1}$ , and it is known that  $e_i$  can be either  $(v_i, v_{i+1})$  or  $(v_{i+1}, v_i)$ . That is, if we ignore the directions of the edges,  $G$  is a line. Design an algorithm that finds *one* vertex with out-degree 0. Your algorithm must run in  $O(\log n)$  time. Prove the correctness of your algorithm, and analyze its time complexity.

The adjacency list is used to store the graph, and the indices of the  $n$  vertices are stored in an array.

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- (b) (15 Points) Let  $G$  be a *directed* graph with  $n^2$  vertices labeled  $\{v_{ij}\}_{i=0,1,\dots,n-1;j=0,1,\dots,n-1}$ . Each vertex  $v_{ij}$  has four neighbors  $v_{(i-1)j}$ ,  $v_{(i+1)j}$ ,  $v_{i(j-1)}$ ,  $v_{i(j+1)}$ , where we adopt the convention that  $0-1 = n-1$  and  $(n-1)+1 = 0$  (for example,  $v_{(n-1)j}$  is a neighbor of  $v_{0j}$ ). It is known that the directions of the edges between  $v_{ij}$  and the “left” and the “right” neighbors of  $v_{ij}$  are the same, i.e., either both edges point towards  $v_{ij}$  or both edges point outwards  $v_{ij}$ , and the directions of the edges between  $v_{ij}$  and the “upper” and the “lower” neighbors of  $v_{ij}$  are the same. Specifically, the four edges incident to  $v_{ij}$  must be in one of the following four configurations.



Design an algorithm that finds *one vertex with out-degree 0 or one directed cycle*. Your algorithm must run in  $O(\log n)$  time. Prove the correctness of your algorithm, and analyze its time complexity.

The adjacency list is used to store the graph, and the indices of the  $n^2$  vertices are stored in an array.

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3. (35 Points) Given an *undirected* (unweighted) graph  $G = (V, E)$  and a vertex  $s$ , design an efficient algorithm to find a cycle with minimum number of edges that contains  $s$ . Notice that a cycle cannot contain an edge more than once. In particular, for a neighbor  $u$  of  $s$ ,  $s-u-s$  is not a valid cycle. Prove the correctness of your algorithm and analyze its time complexity. Try to design an algorithm with time complexity as lower as you can.