

Approximate Principal Components Analysis of Large Data Sets

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Approximation-Regularization for Analysis of Large Data Sets

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LESSON OF THE TALK

Many statistical methods use (perhaps implicitly) a singular value decomposition (SVD).

The SVD is computationally expensive.

We want to understand the statistical properties of some approximations which speed up computation and save storage.

Spoiler: sometimes approximations actually **improve** the statistical properties

CORE TECHNIQUES

Suppose we have a matrix $\mathbf{X} \in \mathbb{R}^{n \times p}$ and vector $Y \in \mathbb{R}^n$

LEAST SQUARES:

Finding

$$\hat{\beta} = \underset{\beta}{\operatorname{argmin}} \|\mathbf{X}\beta - Y\|_2^2$$

PRINCIPAL COMPONENTS ANALYSIS (PCA):

(Or graph Laplacian or diffusion map or..)

Finding \mathbf{U} , \mathbf{V} orthogonal and Λ diagonal such that

$$\mathbf{X} = \tilde{\mathbf{X}} - \bar{\mathbf{X}} = \mathbf{U}\Lambda\mathbf{V}^\top$$

where

$$\bar{\mathbf{X}} = \mathbf{1}\mathbf{1}^\top \tilde{\mathbf{X}}$$

CORE TECHNIQUES

If \mathbf{X} fits into RAM, there exist excellent algorithms in LAPACK that are

- Double precision
- Very stable
- cubic complexity, $O(\min\{n, p\}^3)$, with small constants
- require extensive random access to matrix

There is a lot of interest in finding and analyzing techniques that extend these approaches to large(r) problems

OUT-OF-CORE TECHNIQUES

Many techniques focus on randomized compression

(This is sometimes known as **sketching**)

LEAST SQUARES:

- 1 Rokhlin, Tygert, “A fast randomized algorithm for overdetermined linear least-squares regression” (2008).
- 2 Drineas, Mahoney, et al., “Faster least squares approximation” (2011).
- 3 Woodruff “Sketching as a Tool for Numerical Linear Algebra” (2013).
- 4 Homrighausen, McDonald, “Preconditioned least squares” (2016).

SPECTRAL DECOMPOSITION:

- 1 Halko, et al., “Finding Structure with Randomness: Probabilistic Algorithms for Constructing Approximate Matrix Decompositions” (2011).
- 2 Gittens, Mahoney, “Revisiting the Nystrom Method for Improved Large-Scale Machine Learning” (2013).
- 3 Pourkamali, “Memory and Computation Efficient PCA via Very Sparse Random Projections” (2014).
- 4 Homrighausen, McDonald, “On the Nyström and Column-Sampling Methods for the Approximate PCA of Large Data Sets” (2015).

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COMPRESSION

BASIC IDEA:

- Choose some matrix $Q \in \mathbb{R}^{q \times n}$ or $\mathbb{R}^{p \times q}$.
- Use $Q\mathbf{X}$ or $\mathbf{X}Q$ and QY instead

Finding $Q\mathbf{X}$ for arbitrary Q and \mathbf{X} takes $O(qnp)$ computations

This can be expensive,

To get this approach to work, we need some structure on Q

THE Q MATRIX

- Gaussian

(Well behaved distribution and easier theory. Dense matrix)

- Fast Johnson-Lindenstrauss Methods

- Randomized Hadamard (or Fourier) transformation

(Allows for $O(np \log(p))$ computations.)

- $Q = \pi\tau$ for π a permutation of I and $\tau = [I_q \ 0]$.

($\mathbf{X}Q$ means “only read q columns”)

- Sparse Bernoulli

$$Q_{ij} \stackrel{i.i.d.}{\sim} \begin{cases} 1 & \text{with probability } 1/(2s) \\ 0 & \text{with probability } 1 - 1/s \\ -1 & \text{with probability } 1/(2s) \end{cases}$$

This means $Q\mathbf{X}$ takes $O\left(\frac{np}{s}\right)$ “computations” on average

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TYPICAL RESULTS

THE GENERAL PHILOSOPHY: Find an approximation that is as close as possible to the solution of the original problem

PCA:

A typical result would be to find an approximate \tilde{V} such that

$$\text{angle}(V, \tilde{V}) \leq \sqrt{\frac{p}{n}} \left(\frac{1}{\text{spectral gap}} \right)$$

(This is the same order of convergence as PCA [Homrighausen, McDonald (2015)])

TYPICAL RESULTS

THE GENERAL PHILOSOPHY: Find an approximation that is **as close as possible** to the solution of the original problem

Least Squares:

A typical result would be to find an $\tilde{\beta}$ such that

$$\|\mathbf{X}\tilde{\beta} - Y\|_2^2 \leq (1 + \epsilon) \left(\min_{\beta} \|\mathbf{X}\beta - Y\|_2^2 \right)$$

Here, $\tilde{\beta}$ should be ‘easier’ to compute than $\hat{\beta}$

COLLABORATORS & GRANT SUPPORT

Collaborator:

Darren Homrighausen, Colorado State

Assistant Professor, Department of Statistics

Grant Support:

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- INET Grant INO—14-00020



Typical result: Principal components analysis

NOTATION

- Write (again) the SVD of \mathbf{X} as

$$\mathbf{X} = \mathbf{U}\mathbf{\Lambda}\mathbf{V}^\top,$$

- For general rank r matrices \mathbf{A} we write

$$\mathbf{A} = U(\mathbf{A})\mathbf{\Lambda}(\mathbf{A})V(\mathbf{A})^\top$$

where

$$\mathbf{\Lambda}(\mathbf{A}) = \text{diag}(\lambda_1(\mathbf{A}), \dots, \lambda_r(\mathbf{A}))$$

- For some matrix \mathbf{A} , we use \mathbf{A}_d to be the first d columns of \mathbf{A} .

APPROXIMATIONS

Can't use \mathbf{X}

What about approximating it?

Focus on two methods of “approximate SVD”

- 1 Nyström extension
- 2 Column sampling

INTUITION (NOT THE ALGORITHM)

- Essentially, let

$$\mathbf{S} = \mathbf{X}^\top \mathbf{X} \in \mathbb{R}^{n \times n} \quad \text{and} \quad \mathbf{R} = \mathbf{X} \mathbf{X}^\top \in \mathbb{R}^{p \times p}$$

- Both of these are symmetric, positive semi-definite
- Randomly choose q entries in $\{1, \dots, n\}$ or $\{1, \dots, p\}$
- Then partition the matrix so the selected portion is \mathbf{S}_{11} or \mathbf{R}_{11}

$$\mathbf{S} = \mathbf{V} \Lambda^2 \mathbf{V}^\top = \begin{bmatrix} \mathbf{S}_{11} & \mathbf{S}_{12} \\ \mathbf{S}_{21} & \mathbf{S}_{22} \end{bmatrix} \quad \mathbf{R} = \mathbf{U} \Lambda^2 \mathbf{U}^\top = \begin{bmatrix} \mathbf{R}_{11} & \mathbf{R}_{12} \\ \mathbf{R}_{21} & \mathbf{R}_{22} \end{bmatrix}$$

APPROXIMATING THINGS WE DON'T CARE ABOUT

If we want to approximate \mathbf{S} (or \mathbf{R}), we have for example

Nystrom

$$\mathbf{S} \approx \begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \mathbf{S}_{11}^\dagger \begin{bmatrix} \mathbf{S}_{11} & \mathbf{S}_{12} \end{bmatrix}$$

Column sampling

$$\mathbf{S} \approx U \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right) \Lambda \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right) U \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right)^\top$$

Previous theoretical results have focused on the accuracy of these approximations (and variants) for a fixed computational budget.

APPROXIMATING THINGS WE DON'T CARE ABOUT

If we want to approximate \mathbf{S} (or \mathbf{R}), we have for example

Nyström

$$\mathbf{S} \approx \begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \mathbf{S}_{11}^\dagger [\mathbf{S}_{11} \quad \mathbf{S}_{12}]$$

Column sampling

$$\mathbf{S} \approx U \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right) \Lambda \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right) U \left(\begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix} \right)^\top$$

We don't care about these approximations. We don't want these matrices.

WHAT WE ACTUALLY WANT...

We really want \mathbf{U} , \mathbf{V} , and Λ .

Or even better, the population analogues

Then we can get the principal components, principal coordinates, and the amount of variance explained.

It turns out that there are quite a few ways to use these two methods to get the things we want.

Let

$$L(\mathbf{S}) = \begin{bmatrix} \mathbf{S}_{11} \\ \mathbf{S}_{21} \end{bmatrix}$$

$$L(\mathbf{R}) = \begin{bmatrix} \mathbf{R}_{11} \\ \mathbf{R}_{21} \end{bmatrix}$$

LOTS OF APPROXIMATIONS

After some reasonable algebra...

Quantity of interest	Label	Approximations
V	\mathbf{V}_{nys} \mathbf{V}_{cs}	$L(\mathbf{S})\mathbf{V}(\mathbf{S}_{11})\Lambda(\mathbf{S}_{11})^\dagger$ $\mathbf{U}(L(\mathbf{S}))$
U	\mathbf{U}_{nys} \mathbf{U}_{cs} $\widehat{\mathbf{U}}_{nys}$ $\widehat{\mathbf{U}}_{cs}$ $\widehat{\mathbf{U}}$	$L(\mathbf{R})\mathbf{V}(\mathbf{R}_{11})\Lambda(\mathbf{R}_{11})^\dagger$ $\mathbf{U}(L(\mathbf{R}))$ $\mathbf{X}\mathbf{V}_{nys}\Lambda_{nys}^{\dagger/2}$ $\mathbf{X}\mathbf{V}_{cs}\Lambda_{cs}^{\dagger/2}$ $\mathbf{U}(\mathbf{x}_1)$

NOTES OF INTEREST

Method	Complexity:	
	Computational	Storage
Standard	$O(n^2p + pn^2)$	$O(np)$
Nyström	$O(nq^2 + q^3)$	$O(q^2)$ [$O(nq)$]
Column sampling	$O(qnp + qp^2)$	$O(pq)$

- Column sampling results in orthogonal \mathbf{U} and \mathbf{V} , Nyström doesn't
- $\hat{\mathbf{U}} = U(\mathbf{x}_1)$ where $\mathbf{X} = [\mathbf{x}_1 \quad \mathbf{x}_2]$ seems reasonable if we think that only the first l selected covariates matter (gets used frequently by new statistical methods)

SO WHAT DO WE USE?

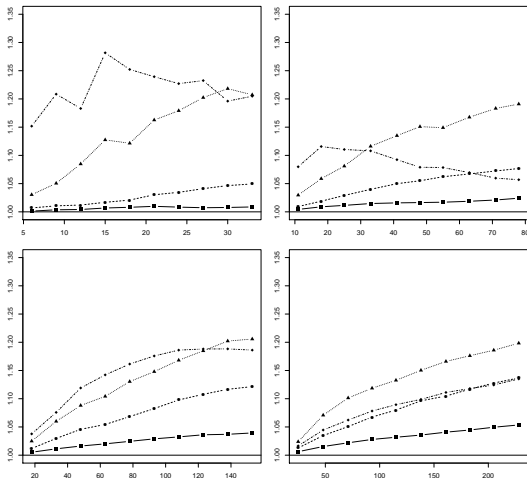
Well...

It depends. We did some theory which gives a way to calculate how far off your approximation might be. You could use these bounds if you like to use your data and make a choice.

Did some simulations too.

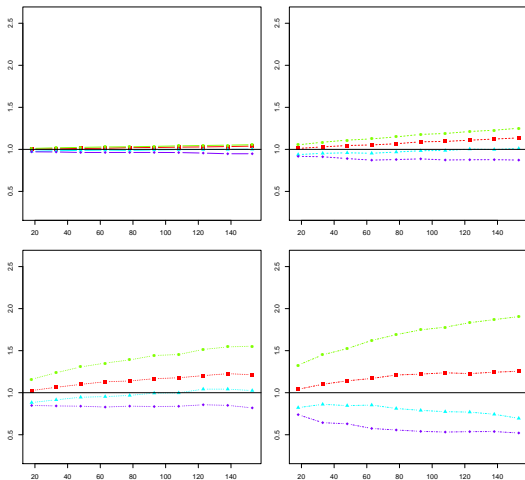
- Draw $\tilde{X}_i \stackrel{iid}{\sim} N_p(0, \Sigma)$ for $n = 5000$, $p = 3000$.
- 4 different conditions.
- For each method report: $\frac{\text{Projection error of estimator}}{\text{Projection error of baseline estimator}}.$

APPROXIMATIONS TO V



- y-axis, relative performance Nystrom to column sampling.
- $< 1 \rightarrow$ better; $> 1 \rightarrow$ worse.
- $d \in \{2, 5, 10, 15\}$.
- x-axis, approximation parameter $q \in [3d/2, 15d]$
- Random_{0.001} (solid, square)
Random_{0.01} (dashed, circle)
Random_{0.1} (dotted, triangle)
Band (dot-dash, diamond)
- d small, similar time
 d large, Nystrom is 10–15%

APPROXIMATIONS TO \mathbf{U}



- y-axis, performance relative to \mathbf{U}_{cs}
- x-axis, approximation parameter q
- \mathbf{U}_{nys} , $\hat{\mathbf{U}}_{nys}$, $\hat{\mathbf{U}}_{cs}$, $\hat{\mathbf{U}}$
- Random_{0.001}
Random_{0.01}
Random_{0.1}
Band
- $d = 10$

CONCLUSIONS (SO FAR)

APPROXIMATE PCA

- For computing \mathbf{V} , CS beats Nyström in terms of accuracy, but is much slower for similar choices of the approximation parameter and d large.
- For computing \mathbf{U} , the naïve methods are bad, better to approximate \mathbf{V} and multiply, so see above
- $\hat{\mathbf{U}}$ really stinks. But it's most obvious. This also gets used frequently
- For more info, see the paper. Contains boring theory, more extensive simulations, and application to Enron email dataset.

EVEN BETTER APPROXIMATE PCA

- We really want the population versions: top d eigenvectors of $\Sigma \dots$

Atypical (better) result: Compressed regression

ESTIMATION RISK

Form a **loss function** $\ell : \mathbb{R}^p \times \mathbb{R}^p \rightarrow \mathbb{R}^+$

(For this talk, let's just specify $\ell(\hat{\theta}, \theta) = \|\hat{\theta} - \theta\|_2^2$)

The quality of an estimator is given by its **risk**

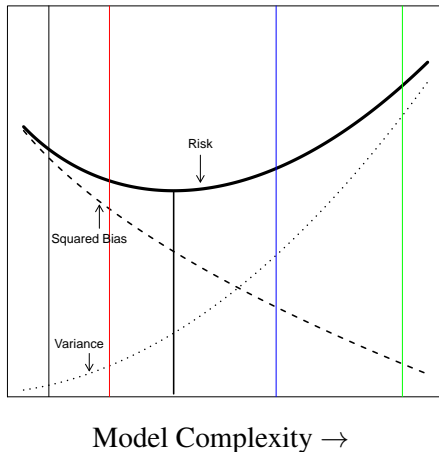
$$R(\hat{\theta}) = \mathbb{E} \|\hat{\theta} - \theta\|_2^2$$

This can be decomposed as:

$$\begin{aligned} R(\hat{\theta}) &= \|\mathbb{E}[\hat{\theta}] - \theta\|^2 + \mathbb{E} \|\hat{\theta} - \mathbb{E}[\hat{\theta}]\|^2 \\ &= \text{Bias}^2 + \text{Variance} \end{aligned}$$

There is a natural conservation between these quantities...

BIAS-VARIANCE TRADEOFF



Typical result examines the bias, atypical result also considers the variance

FULLY COMPRESSED REGRESSION

Let $Q \in \mathbb{R}^{q \times n}$

Let's solve the **fully compressed** least squares problem

$$\hat{\beta}_{FC} = \underset{\beta}{\operatorname{argmin}} \|Q\mathbf{X}\beta - QY\|_2^2$$

→ A common way to solve least squares problems that are:

- Very large or
- Poorly conditioned

The numerical/theoretical properties generally depend on Q , q

(Multiply quickly? Reduce dimension? Reduce Storage?)

THREE APPROXIMATIONS

Note:

$$\|Q(\mathbf{X}\beta - Y)\|_2^2 \propto \beta^\top \mathbf{X}^\top Q^\top Q \mathbf{X} \beta - 2\beta^\top \mathbf{X}^\top Q^\top Q Y$$

Full compression:

$$\hat{\beta}_{FC} = (\mathbf{X}^\top Q^\top Q \mathbf{X})^{-1} \mathbf{X}^\top Q^\top Q Y$$

Partial compression:

$$\hat{\beta}_{PC} = (\mathbf{X}^\top Q^\top Q \mathbf{X})^{-1} \mathbf{X}^\top Y$$

Convex combination compression:¹

$$\hat{\beta}_C = W \hat{\alpha} \quad W = [\hat{\beta}_{FC}, \hat{\beta}_{PC}] \quad \hat{\alpha} = \underset{\alpha}{\operatorname{argmin}} \|W\alpha - Y\|_2^2$$

¹ see the work of Stephen Becker CU Boulder Applied Math

WHY THESE?

- Turns out that FC is unbiased, and therefore worse than OLS (has high variance)
- On the other hand, PC is biased and empirics demonstrate low variance
- Combination should give better statistical properties

COMPRESSED REGRESSION

With this Q , $\hat{\beta}_C$ “works” in practice:

- Computational savings: $O\left(\frac{qnp}{s} + qp^2\right)$
- Approximately the same estimation risk as OLS

This is good, but we had a realization:

IF RIDGE REGRESSION IS BETTER THAN OLS, WHY NOT “POINT” THE APPROXIMATION AT RIDGE?

COMPRESSED RIDGE REGRESSION

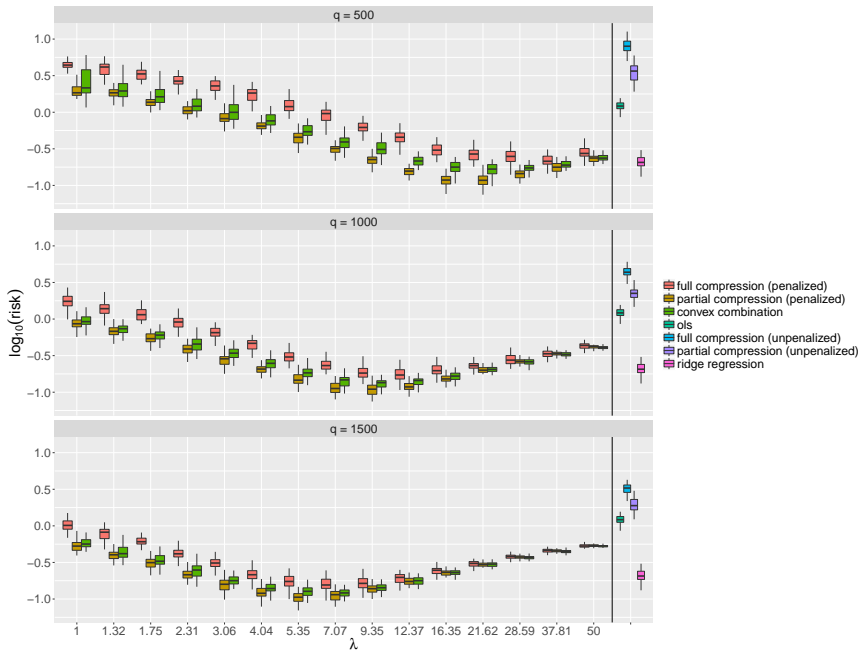
This means introducing a tuning parameter λ and defining:

$$\begin{aligned}\hat{\beta}_{PC}(\lambda) &= (\mathbf{X}^\top \mathbf{Q}^\top \mathbf{Q} \mathbf{X} + \lambda I)^{-1} \mathbf{X}^\top \mathbf{Y} \\ \hat{\beta}_{FC}(\lambda) &= (\mathbf{X}^\top \mathbf{Q}^\top \mathbf{Q} \mathbf{X} + \lambda I)^{-1} \mathbf{X}^\top \mathbf{Q}^\top \mathbf{Q} \mathbf{Y}\end{aligned}$$

(Everything else about the procedure is the same)

This has the same computational complexity, but has **much** lower risk

Let's look at an (a)typical result...



CONCLUSIONS

COMPRESSED REGRESSION

- In this case, approximation actually **improves** existing approaches
- We have ways of choosing the tuning parameter using the data
- Much more elaborate simulations and theory (current work)

GENERAL PHILOSOPHY

- Approximation is not necessarily a bad thing
- Don't just minimize it, that only considers the bias
- Examine statistical properties to calibrate
- In some cases, we can get the best of both worlds:
 - **easier computations**
 - **better statistical properties**