CMPE 282 Cloud Services MapReduce Design Patterns Join

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Join Patterns

What: Bringing data sets together

Why: I want to mash my different data sources together

- Reduce-side join
- Replicated join
- Composite join
- Cartesian product

RDBMS Join Refresher

Table 5-3. Inner Join of A + B on User ID

Table 5-1. Table A

User ID	Reputation	Location
3	3738	New York, NY
4	12946	New York, NY
5	17556	San Diego, CA
9	3443	Oakland, CA

Table 5-2. Table B

User ID	Post ID	Text
3	35314	Not sure why this is getting downvoted.
3	48002	Hehe, of course, it's all true!
5	44921	Please see my post below.
5	44920	Thank you very much for your reply.
8	48675	HTML is not a subset of XML!

A.User ID	A.Reputation	A.Location	B.User ID	B.Post ID	B.Text
3	3738	New York, NY	3	35314	Not sure why this is getting downvoted.
3	3738	New York, NY	3	48002	Hehe, of course, it's all true!
5	17556	San Diego, CA	5	44921	Please see my post below.
5	17556	San Diego, CA	5	44920	Thank you very much for your reply.

Table 5-4. Left Outer Join of A + B on User ID

A.User ID	A.Reputation	A.Location	B.User ID	B.Post ID	B.Text
3	3738	New York, NY	3	35314	Not sure why this is getting downvoted.
3	3738	New York, NY	3	48002	Hehe, of course, it's all true!
4	12946	New York, NY	null	null	null
5	17556	San Diego, CA	5	44921	Please see my post below.
5	17556	San Diego, CA	5	44920	Thank you very much for your reply.
9	3443	Oakland, CA	null	null	null

Table 5-5. Right Outer Join of A + B on User ID

A.User ID	A.Reputation	A.Location	B.User ID	B.Post ID	B.Text
3	3738	New York, NY	3	35314	Not sure why this is getting downvoted.
3	3738	New York, NY	3	48002	Hehe, of course, it's all true!
5	17556	San Diego, CA	5	44921	Please see my post below.
5	17556	San Diego, CA	5	44920	Thank you very much for your reply.
null	null	null	8	48675	HTML is not a subset of XML!

Reduce Side Join 1/3

Intent

Join large multiple data sets together by some foreign key

Motivation

- Simple to implement in Reducers
- Supports all the different join operations (inner, outer)
- No limitation on the size of your data sets

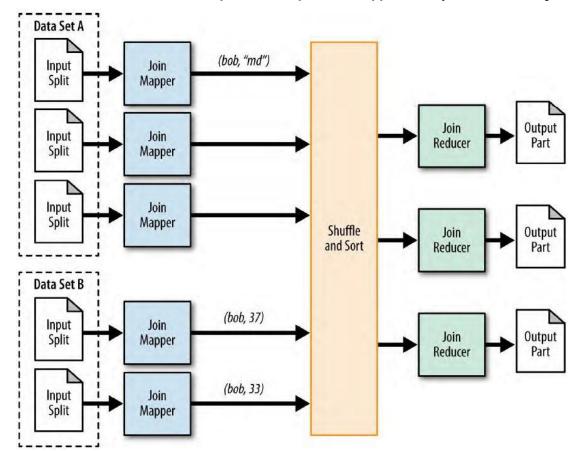
Applicability

- Multiple large data sets are being joined by a foreign key
- You want the flexibility of being able to execute any join operation
- A large amount of network bandwidth

Reduce Side Join 2/3

Structure

- Mapper: prepares [(foreign key, record)]
- Reducer: receives (FK, list(values)) and performs join operation



Reduce Side Join 3/3

SQL

```
SELECT T1.ID, T1.A, T2.C
FROM T1 [INNER|LEFT|RIGHT] JOIN T2
ON T1.ID=T2.ID
```

- Performance analysis
 - # of part files == # of reduce tasks
 - Heavy network traffic
 - Utilize relatively more reducers than your analytic
- Ex: ReduceSideJoinDriver.java
 - In-1: Users.xml
 - In-2: Comments.xml



Replicated Join 1/3

Intent

Eliminates the need to shuffle any data to the reduce phase

Motivation

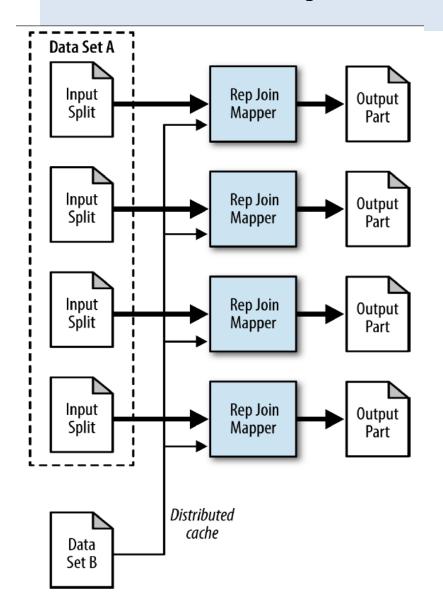
 All the data sets except the very large one are essentially read into memory during the setup phase of each map task, which is limited by the JVM heap

Applicability

 All of the data sets, except for the large one, can be fit into main memory of each map task

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Replicated Join 2/3



Structure

- Map-only pattern
- Read all files from the distributed cache and store them into in-memory lookup tables

Replicated Join 3/3

- Consequences
 - # of part files == # of map tasks
 - The part files contain the full set of joined records
- Performance analysis
 - A replicated join can be the fastest type of join because there is no reducer required
 - The amount of data that can be stored safely inside JVM
- Ex: ReplicatedJoinDriver.java

Cartesian Product 1/3

Intent

 Pair up and compare every single record with every other record in a data set

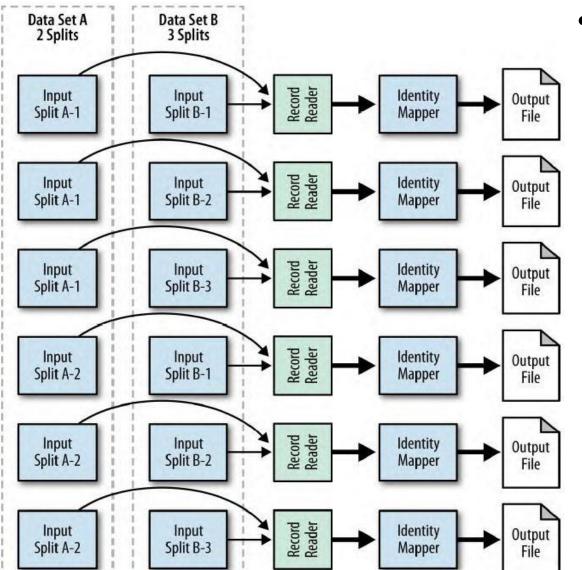
Motivation

- Simply pairs every record of a data set with every record of all the other data sets
- To analyze relationships between one or more data sets

Applicability

- You want to analyze relationships between all pairs of individual records
- You've exhausted all other means to solve this problem.
- You have no time constraints on execution time

Cartesian Product 2/3



Structure

- Map-only
- RecordReader job

Cartesian Product 3/3

Consequences

- The final data set is made up of tuples equivalent to the number of input data sets
- Every possible tuple combination from the input records is represented in the final output
- SQL

SELECT * FROM tableA, tableB;

- Performance Analysis
 - A massive explosion in data size O(n^2)
 - If a single input split contains a thousand records, the right input split needs to be read a thousand times before the task can finish
 - If a single task fails for an odd reason, the whole thing needs to be restarted
- Ex: CartesianProduct.java

References

- Donald Miner and Adam Shook, MapReduce Design Patterns.
 - http://oreil.ly/mapreduce-design-patterns
 - https://github.com/adamjshook/mapreducepatterns