

# 人工智能基础 最佳 split 点的寻找

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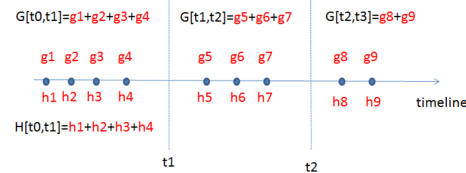
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## 1 问题

“Introduction to Boosted Trees” 这个 slides 的 39 页上的问题，  
提交伪代码，并分析时间复杂度

### Questions to check if you really get it

- Time series problem



- All that is important is the structure score of the splits

$$Obj = -\frac{1}{2} \sum_{j=1}^T \frac{G_j^2}{H_j + \lambda} + \gamma T$$

- Top-down greedy, same as trees
- Bottom-up greedy, start from individual points as each group, greedily merge neighbors
- Dynamic programming, can find optimal solution for this case

ORZ，写完以后想起来要用动态规划写，但好像我用的是分治 >.<。但讲道理差别不大嘛，动态规划用上一次的东西往下传，在这个题里面，分治也可以把这个信息往下传嘛！差别不是很大嘛！懒得改了 =.=

## 2 伪码

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### Algorithm 1 Find The Split

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```

procedure MAIN(g,h,n,a)
    k = 1
    FindSplit(g,h,n,a)
    return Split
end procedure
procedure GETSCORE(G,H,T)
    Obj =  $-\frac{1}{2} \sum_j \frac{G[j]^2}{H[j]+\lambda} + 3\gamma$ 
    return Obj ▷ 计算分数
end procedure
procedure GETGAIN(GL,GR,HL,HR)
    Gain =  $\frac{1}{2} [\frac{GL^2}{HL+\lambda} + \frac{GR^2}{HR+\lambda} - \frac{(GL+GR)^2}{(HL+HR)+\lambda}] - \gamma$ 
    return Gain ▷ 计算增益
end procedure
procedure FINDSPLIT(g,h,n,a) ▷ g,h 为 n 个梯度数据, a 为增益阈值
    G[1] =  $\sum g[i]$ 
    h[1] =  $\sum h[i]$ 
    Obj[0] = GetScore(G, H, 1) ▷ 用 0 位存储不分割的分数
    Split[0] = 0 ▷ 用 0 位存储分割点, 0 是不分割
    Gain[0] = a ▷ 用 0 位暂时存储比较值, a 是阈值
    for i = 1 to n do
        GR =  $\sum_{j=1}^i g[j]$ 
        GR =  $\sum_{j=i+1}^n g[j]$ 
        HL =  $\sum_{j=1}^i h[j]$ 
        HR =  $\sum_{j=i+1}^n h[j]$ 
        Gain[i] = GetGain(GL, GR, HL, HR)
        if Gain[i] > Gain[0] then ▷ 贪心, 找出最大的增益
            Gain[0] = Gain[i]
            Split[0] = i
        end if
    end for

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**Algorithm 2** Find The Split

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```
if Split != 0 then      ▷ 如果找到了分割点，保存数据，分治继续寻找
    Obj[k] = Obj[k-1] - Gain[0]
    Split[k] = Split[0]
    k = k + 1
    for i = 1 to Split do
        gl[i] = g[i]
        hl[i] = h[i]
    end for
    for i = Split + 1 to n do
        gr[i - Split] = g[i]
        hr[i - Split] = h[i]
    end for
    FindSplit(gl, hl, Split, a)
    FindSplit(gr, hr, n - Split, a)
end if
end procedure
```

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### 3 分析

单层贪心，遍历  $O(n)$ ；分治处理多层， $O(\log n)$ 。  
综上，时间复杂度是  $O(n \log n)$ 。

### 4 参考

毕竟英文渣，看了网上翻译版的 =.=

<http://www.52cs.org/?p=429>

[http://blog.sina.com.cn/s/blog\\_7103b28a0102w6qa.html](http://blog.sina.com.cn/s/blog_7103b28a0102w6qa.html)