实验报告

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1	实验题目 - 网络路由实验二	

2 实验内容

- 基于已有框架代码,实现计算路由器路由表项的相关操作
 - Hello/LSU 消息的相关操作
- 运行实验
 - 运行网络拓扑 (topo/mospf_topo.py)
 - 在各个路由器节点上执行 disable_arp.sh, disable_icmp.sh, disable_ip_forward.sh), 禁止 协议栈的相应功能
 - 运行./mospfd, 使得各个节点生成一致的链路状态数据库
 - 等待一段时间后,每个节点生成完整的路由表项
 - 在节点 h1 上 ping/traceroute 节点 h2

3 实验流程

3.1 文件列表

```
lab11
homework.pdf
11-mospf
mospf_daemon.c
mospf_database.c
```

3.2 相关代码

• 补充通知没有 hello 的邻居节点

```
list_for_each_entry(iface, &(instance->iface_list), list){
     if (list_empty(&(iface->nbr_list))){
       lsa [number].subnet = htonl(iface->ip);
       lsa [number].mask = htonl(iface->mask);
       lsa[number].rid = htonl(0);
       number++;
    }else{
       list_for_each_entry(nbr, &(iface->nbr_list), list){
         lsa[number].subnet = htonl(nbr->nbr_ip);
         lsa [number].mask = htonl(nbr->nbr_mask);
10
         lsa [number].rid = htonl(nbr->nbr id);
11
12
         number++;
13
14
    }
15 }
• 广度优先搜索
void update_rtable()
2 {
       clear_rtable();
3
       init rtable();
4
       struct queue_t *head, *current, *prev;
       head = (struct queue_t *)malloc(sizeof(struct queue_t));
       init_list_head((struct list_head *)head);
       char *visited = (char *) malloc(MAX_NODE_NUM * sizeof(char));
       char *rid_visited = (char *) malloc(MAX_NODE_NUM * sizeof(char));
10
       memset(visited, 0, MAX NODE NUM * sizeof(char));
11
       memset(rid_visited, 0, MAX_NODE_NUM * sizeof(char));
13
       iface_info_t *iface;
14
       mospf\_nbr\_t *nbr;
15
       list_for_each_entry(iface, &instance->iface_list, list){
16
           list_for_each_entry(nbr, &iface->nbr_list, list){
               prev = (struct queue_t *)malloc(sizeof(struct queue_t));
18
               prev \rightarrow rid = nbr \rightarrow nbr_id;
               prev \rightarrow gw = nbr \rightarrow nbr _ip;
20
               prev->iface = iface;
21
               list_add_tail((struct list_head *)prev, (struct list_head *)head);
23
24
           visited [hash(iface->mask & iface->ip)] = 1;
25
26
       struct mospf_lsa *lsa;
       mospf\_db\_entry\_t \ *db;
27
       current = (struct queue_t *)((struct list_head *)head)->next;
28
       while (current != head) {
           list for each entry(db, &mospf db, list){
30
                if (db->rid == current->rid){
31
                    lsa = db->array;
32
                    for (int i = 0; i < db \rightarrow nadv; i++, lsa++){
33
                        if (rid_visited[hash(lsa->rid)] == 0){
                            prev = (struct queue_t *)malloc(sizeof(struct queue_t));
35
                            prev \rightarrow rid = lsa \rightarrow rid;
37
                            prev->gw = current->gw;
                            prev->iface = current->iface;
38
                            list_add_tail((struct list_head *)prev, (struct list_head *)
39
       head);
                            rid visited [hash(lsa\rightarrowrid)] = 1;
41
                        if (visited [hash(lsa->subnet)] == 0){
42
                            rt_entry_t *rt_entry;
43
                            rt_entry = (rt_entry_t *) malloc(sizeof(rt_entry_t));
44
                            init_list_head(&rt_entry->list);
                            rt_entry->dest = lsa->subnet;
46
```

```
rt_entry->gw = current->gw;
47
                            rt_entry->mask = lsa->mask;
48
                            rt_entry->iface = current->iface;
                            rt_{entry}->flags = 0;
50
                            memcpy(rt_entry->if_name, current->iface->name, 16);
51
                            add_rt_entry(rt_entry);
52
                            visited [hash(lsa->subnet)] = 1;
53
                        }
55
56
                   break;
               }
57
58
           prev = current;
59
           current = (struct queue_t *)((struct list_head *)current)=>next;
60
61
           free (prev);
62
       free (head);
63
       free(visited);
64
       free(rid_visited);
65
66
       print_rtable();
67 }
```

4 实验结果

⊗⊖ □ '	'Node: r1"				⊗ ● □ "N	lode: h1"				
destination 10.0.1.0 10.0.2.0 10.0.3.0 10.0.4.4 10.0.5.4 10.0.6.4	gateway 0x00000000 0x00000000 0x00000000 0x0a000202 0x0a000303 0x0a000202	netmask 0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00	flags 1 1 1 0 0	r1-eth0 r1-eth1 r1-eth2 r1-eth1	rtt min/avg/max/mdev = 0.108/0.285/0.365/0.105 ms root@12-ubuntu:"/Desktop/10-mospf# ping 10.0.6.22 -c 4 PING 10.0.6.22 (10.0.6.22) 56(84) bytes of data. 64 bytes from 10.0.6.22: icmp_seq=1 ttl=61 time=0.177 ms 64 bytes from 10.0.6.22: icmp_seq=2 ttl=61 time=0.122 ms 64 bytes from 10.0.6.22: icmp_seq=3 ttl=61 time=0.122 ms 64 bytes from 10.0.6.22: icmp_seq=3 ttl=61 time=0.292 ms 64 bytes from 10.0.6.22: icmp_seq=4 ttl=61 time=0.290 ms 10.0.6.22 ping statistics 4 packets transmitted, 4 received, 0% packet loss, time 3050ms rtt min/avg/max/mdev = 0.122/0.197/0.292/0.063 ms rtt min/avg/max/mdev = 0.122/0.197/0.292/0.063 ms root@12-ubuntu:"/Desktop/10-mospf# traceroute 10.0.6.22 traceroute to 10.0.6.22 (10.0.6.22), 30 hops max, 60 byte packe 1 10.0.1.1 (10.0.1.1) 0.042 ms 0.010 ms 0.007 ms 2 10.0.2.2 (10.0.2.2) 0.026 ms 0.014 ms 0.014 ms 3 10.0.4.4 (10.0.4.4) 0.033 ms 0.022 ms 0.023 ms 4 * * * 5 * * * *					
destination 	9ateway 0x00000000 0x0000000 0x0a000201 0x0a000201 0x0a000404 0x0a000404	netmask 0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00	flags 1 1 0 0 0	if_name r2-eth0 r2-eth1 r2-eth0 r2-eth1 r2-eth1						
⊗ ⊕ ⊕ "1	● ■ "Node: r3"					⊗ 🖨 🎟 "Node: r4"				
destination	gateway 	netmask 	flags	if_name	destination	gateway	netmask	flags	if_name	
10.0.3.0 10.0.5.0 10.0.1.1 10.0.2.2 10.0.4.2 10.0.6.4	0x00000000 0x00000000 0x0a000301 0x0a000301 0x0a000504 0x0a000504	0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00	1 1 0 0 0	r3-eth0 r3-eth1 r3-eth0 r3-eth0 r3-eth1 r3-eth1	10.0.4.0 10.0.5.0 10.0.6.0 10.0.2.1 10.0.3.1	0x00000000 0x00000000 0x00000000 0x0a000402 0x0a000503 0x0a000402	0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00 0xffffff00	1 1 1 0 0	r4-eth0 r4-eth1 r4-eth2 r4-eth0 r4-eth1 r4-eth0	

5 结果分析

• 数据库建立拓扑后,用广度优先搜索获得最短路径