Project 6 File System 设计文档

中国科学院大学 [王苑铮] [2018.1.20]

1. 文件系统初始化设计

请至少包含以下内容

(1) 请用图表示你设计的文件系统对磁盘的布局(布局上可以不考虑 boot block 的大小,直接从逻辑地址第 0 块开始),并说明各部分占用的磁盘空间大小,例如 superblock,inode 的元数据等

Super block 备份 Super block

Inode Bit map Data block Bit map

inode

Data block

布局,以及每部分的大小和换算关系。(单位为块数。4KB 每块)。TOTAL_DATA_SIZE 为最大数据空间,此处为 4G。

(2) 你如何实现 superblock 的备份?如何判断 superblock 损坏,以及当有一个 superblock 损坏时你的文件系统如何正常启动?

备份:另有一个备份的 superblock 块每次更新 superblock 时也一起更新备份的 backup_superblock。判断 superblock 损坏:读出 superblock 后里面的 magic_number 和我定义的 magic_number 不同。当有一个 superblock 损坏时,先读出备份的 backup_ssuperblock,如果备份块里的 magic_number 正确,则认为只是 superblock 坏了,把备份块写给 superblock 就可以了。如果备份块里面的 magicnumber 也不对,则认为是整个文件系统都损坏了,或者新的磁盘尚未初始化文件系统,此时就要 mkfs(),初始化 superblock,inode 数组,bitmap,根目录等数据结构,写进磁盘里。之后 mount()即可启动。

(3) 请列出你设计的 superblock 和 inode 数据结构,并阐明各项含义。请说明你设计的文件系统能支持的最大文件大小,最多文件数目,以及单个目录下能支持的最多文件/子目录数目。

```
typedef enum {
    DIRECTORY,
    SOFT_LINK,
    NORMAL_FILE
} type;

//size:4+4+4+4+4+4*DIRECT_DATA_BLOCKS_NUM+4+4+8*3+4 32+24=56

*typedef struct inode_disk_t{
    // complete it
    int file_size;
    int file_indirect_blocks_num;
    int dentry_num;

//如果是目录,则代表目录项的数量。如果是文件,这一项为6

type file_type;
    int blocks_number;
    int blocks_number;
    int direct_block[DIRECT_DATA_BLOCKS_NUM];
    // 文件类型。支持目录文件、软连接文件、普通文件
    int indirect_block[DIRECT_DATA_BLOCKS_NUM];
    // 文件类型。发持目录文件、软连接文件、普通文件
    int indirect_block;
    int indirect_block;
    int links;

//连接数量

uint64_t last_access_time;
    uint64_t last_access_time;
    uint64_t last_modified_time;
    uint64_t reate_time;
    // L次到达的时间(读、修改)
    uint64_t reate_time;
    // L次到达的时间
    uint64_t reate_time;
    // L次到达的时间
    int indef_t create_time;
    // Mile principle of the complete it
    char file_name[MAX_FILE_NAME_LEN];
    int inode;
}dentry;
```

最多文件数目=datablock 数目=4G/4K=1M 个文件

限定一个目录只能使用一个块。目录项 dentry 中,单个文件名字最长不能超过 60char,加上一个 int 索引 inode 好,一个目录项为 64B。1 块 4K,可以放 64 放文件/子目录。

(4) 请说明你设计的文件系统的块分配策略,按需分配还是有设计其他分配策略?

按需分配,分配的块数为 ceil(文件大小/4K)个块。

(5) 如果完成了 bonus,请通过举例说明你设计的面向 SSD 的文件系统采用异地 更新模式后文件数据块在磁盘上的布局,例如新写一个文件后和更新同一个 文件后,该文件数据块在磁盘上的布局

没做

(6)设计或实现过程中遇到的问题和得到的经验(如果有的话可以写下来,不是必需项)

2. 文件操作设计

请至少包含以下内容

(1) 请说明 link 和 unlink 的操作流程

```
int p6fs_unlink(const char *path)
{
    //解析路径
    int second_last_dir_inode_index,last_file_inode_index;
    char file_name[MAX_FILE_NAME_LEN];
    path_analizing_result result = path_analize(path,file_name,&second_last_dir_inode_index,&last_file_inode_index);
    //未投到
    if(result == mid_not_find || result == second_last_find__last_not_find)
        return -ENOENT;
    //如果ink的是个目录
    //目录不允许被link、因为可能link到自己产生死循环
    if(inode_disk_table[last_file_inode_index].file_type == DIRECTORY)
        return -EISDIR;

    //删除目录项
    remove_dentry_in_directory_by_inode(second_last_dir_inode_index,file_name);

    //如果对应的文件inode没有连接则回收inode和空间
    --inode_disk_table[last_file_inode_index].links;
    if(inode_disk_table[last_file_inode_index].links == 0){
        recycle_inode_and_space(last_file_inode_index);
    }

    //写入disk
    write_inode_disk_table();
    return 0;
}
```

见上面代码的注释

(2) 请说明 rename 涉及的操作流程

```
int p6fs_rename(const char *path, const char *newpath)

{
    //解析原路径
    int second_last_dir_inode_index,last_file_inode_index;
    char file_name[MAX_FILE_NAME_LEN];
    path_analizing_result result = path_analize(path,file_name,&second_last_dir_inode_index,&last_file_inode_index);
    //未表到
    if(result == mid_not_find || result == second_last_find__last_not_find)
        return -teNcNT;
    //如果要重命名的是根目录
    if( strcmp(path,'/')==0 )
        return -tePERM; /* Operation not permitted */

    //解析部路径
    int new_second_last_dir_inode_index,new_last_file_inode_index;
    char new_file_name[MAX_FILE_NAME_LEN];
    path_analizing_result new_result = path_analize(newpath,new_file_name,&new_second_last_dir_inode_index,&new_last_file_inode_index);
    //未表到
    if(new_result == mid_not_find)
    neturn -teNcENT;
    //新日录中间名文件已存在
    if(new_result == second_last_find__last_find)
    return -teExIST;

//在原路径倒数第二级目录里删除dentry
    remove_dentry_in_directory_by_inode(second_last_dir_inode_index,new_file_name);

//在斯路径倒数第二级目录序加斯dentry
    add_dentry_in_directory_by_inode(new_second_last_dir_inode_index,new_file_name,last_file_inode_index);

return 0;
```

见以上代码的注释

(3)设计或实现过程中遇到的问题和得到的经验(如果有的话可以写下来,不是必需项)

3. 目录操作设计

请至少包含以下内容

(1) 请说明 rmdir 的操作流程?

```
int p6fs_rmdir(const char *path)
{
    int second_last_dir_inode_index,last_dir_inode_index;
    char file_name[MAX_FILE_NAME_LEN];
    path_analizing_result result = path_analize(path,file_name,&second_last_dir_inode_index,&last_dir_inode_index);

    //未找到
    if(result == mid_not_find || result == second_last_find__last_not_find)
        return -ENOENT;

    //找到的不是不自录
    if(inode_disk_table[last_dir_inode_index].file_type != DIRECTORY)
        return -ENOTDIR;

    //找到的是根目录
    if( strcmp(path,"/")==0 )
        return -FPERM; /* Operation not permitted */

        //要删除的目录非空
    if(inode_disk_table[last_dir_inode_index].dentry_num > 2 )
        return -ENOTEMPTY;/* Directory not empty */

        //回收空间
        recycle_inode_and_space(last_dir_inode_index);
        //在父目录中取得dentry
        remove_dentry_in_directory_by_inode(second_last_dir_inode_index,file_name);
        return 0;
```

见以上代码的注释

(2)设计或实现过程中遇到的问题和得到的经验(如果有的话可以写下来,不是必需项)

4. 关键函数功能

请列出上述各项功能设计里,你觉得关键的函数或代码块,及其作用

1.绝对路径解析函数。

```
strcpy(path_cpy1,path);
strcpy(path_cpy2,path);
457
458
            if(strcmp(path,"/") == 0){
   *second_last_inode_index = 0;
   *last_inode_index = 0;
   strcpy(last_file_name,"/");
   return second_last_find__last_find;
460
461
465
466
467
469
            int max_deepth=0;
            char* p=strtok(path_cpy1,"/");
while(p != NULL){
471 ▼
                  ++max_deepth;
strcpy(last_file_name,p);
p = strtok(NULL,"/");
472
473
475
476
479
480
           int i,current_inode_index;
           inclinet_inde_index=0,p=strtok(path_cpy2,"/"); p!=NULL && i<max_deepth; ++i){
    if(lookup_directory(current_inode_index,p) == -1 && i<max_deepth-1){</pre>
482
483
484
                     *second_last_inode_index = -1;
*last_inode_index = -1;
485
486
487
                             n mid_not_find;
                   se if(lookup_directory(current_inode_index,p) == -1 && i==max_deepth-1){
488
489
490
                      *second_last_inode_index = current_inode_index;
                      *last_inode_index = -1;
return second_last_find__last_not_find;
491
493
494
495
496
497
498
499
500
501
                }
else if(lookup_directory(current_inode_index,p) != -1 && i==max_deepth-1){
                      *second_last_inode_index = current_inode_index;
                      *last_inode_index = lookup_directory(current_inode_index,p);
return second_last_find__last_find;
                     current_inode_index = lookup_directory(current_inode_index,p);
p = strtok(NULL,"/");
502
503
```

2.当文件需要的 size 大于文件当前的块数的 size 时,申请新的 datablock 对文件扩容

```
int file_increase_block_num(int file_inode_index,int new_size){
898
899
900
901
902
          inode_disk_t* inode = &inode_disk_table[file_inode_index];
          if(new size <= inode->blocks number * BLOCK SIZE)
               return 0;
904
905
997
          //直接娱! 在difect_block[]至國家引的【数据块】 //difect_datablock
//间接数据块: 通过二级间接【索引块】索引到的【数据块】 //indirect_datablock
//一级间接索引块: indirect_block索引到的一个【索引块】 //first_indirect_block
//二级级间接索引块: 用过一级间接【索引块】索引到的一个【索引块】//second_indirect_block
908
909
911
912
          //四种块新的数量
//旧的【数据块】总数量(非索引块)
916
          int old_size = inode->file_size;
920
          int old_direct_datablock_num;
          int old_indirect_datablock_num;
921
922
          int old_first_indirect_block_num;
          int old_second_indirect_block_num;
924
          int new_direct_datablock_num;
927
          int new_indirect_datablock_num;
          int new_first_indirect_block_num;
928
          int new_second_indirect_block_num;
930
931
932
          int old_datablock_num = ceil_division(inode->file_size,BLOCK_SIZE);
          int new_datablock_num = ceil_division(new_size
                                                                           ,BLOCK SIZE);
```

```
if(old_datablock_num < DIRECT_DATA_BLOCKS_NUM && new_datablock_num <= DIRECT_DATA_BLOCKS_NUM){</pre>
                                           = old_datablock_num;
             old_direct_datablock_num
             old_indirect_datablock_num
             old_first_indirect_block_num = 0;
             old_second_indirect_block_num = 0;
             new_direct_datablock_num
                                           = new_datablock_num;
             new_indirect_datablock_num
             new_first_indirect_block_num = 0;
943
944
             new_second_indirect_block_num = 0;
         }
else if(old_datablock_num <= DIRECT_DATA_BLOCKS_NUM && new_datablock_num > DIRECT_DATA_BLOCKS_NUM){
             old_direct_datablock_num
                                           = old_datablock_num;
             old_indirect_datablock_num
948
             old_first_indirect_block_num = 0;
old_second_indirect_block_num = 0;
949
950
                                           = DIRECT_DATA_BLOCKS_NUM;
             new_direct_datablock_num
             new_indirect_datablock_num
                                          = new_datablock_num - DIRECT_DATA_BLOCKS_NUM;
             new first indirect block num = 1;
             new_second_indirect_block_num = ceil_division(new_indirect_datablock_num, BLOCK_SIZE/sizeof(int));
         .
else if(old_datablock_num > DIRECT_DATA_BLOCKS_NUM && new_datablock_num > DIRECT_DATA_BLOCKS_NUM){
             old_direct_datablock_num
                                           = DIRECT_DATA_BLOCKS_NUM;
             old_indirect_datablock_num
                                            = old_datablock_num - DIRECT_DATA_BLOCKS_NUM;
             old_first_indirect_block_num = 1;
             old_second_indirect_block_num = ceil_division(old_indirect_datablock_num, BLOCK_SIZE/sizeof(int));
                                           = DIRECT_DATA_BLOCKS_NUM;
             new_direct_datablock_num
                                           = new_datablock_num - DIRECT_DATA_BLOCKS_NUM;
             new_indirect_datablock_num
             new_first_indirect_block_num = 1;
             new_second_indirect_block_num = ceil_division(new_indirect_datablock_num, BLOCK_SIZE/sizeof(int));
```

```
if( count_bitmap(datablock_bitmap) + add_all > DATA_BLOCKS_NUM )
            980
            int new_datablock_index[add_datablock_num];
            for(i=0; i<add_datablock_num ; ++i){
                 new_datablock_index[i] = apply_available_bit(datablock_bitmap);
990
                 if(new_datablock_index[i] == -1){
                      return -ENOSPC; // 可用空间不足
                 set_bitmap(datablock_bitmap, new_datablock_index[i] ,USED);
            int new_first_indirect_block_index;
            for(i=0 ; i<add_first_indirect_block_num ; ++i ){</pre>
                 new_first_indirect_block_index = apply_available_bit(datablock_bitmap);
                 if(new_first_indirect_block_index == -1){
000
                      return -ENOSPC;
                 set_bitmap(datablock_bitmap, new_first_indirect_block_index ,USED);
003
L004
            int new_second_indirect_block_index[add_second_indirect_block_num];
            for(i=0; i<add_second_indirect_block_num; ++i){</pre>
                 new_second_indirect_block_index[i] = apply_available_bit(datablock_bitmap);
L008
                 if(new_second_indirect_block_index[i] == -1){
                      return -ENOSPC;
1011
                 set_bitmap(datablock_bitmap, new_second_indirect_block_index[i] ,USED);
1014
1015
1017
1018
1019
1020
         indirect_block_t *first_indirect_block =(indirect_block_t*)malloc(sizeof(indirect_block_t));
if(first_indirect_block == NULL)
1024
1025
         return -ENOMEN; //内存不足 indirect_block_t*)malloc(sizeof(indirect_block_t)); if(second_indirect_block == NULL)
1026
1027
1029
1030
                   - ENOMEM: /
1031
         int new_datablock_i=0,new_second_indirect_block_i=0; //新申请的数据块的index
int j,k; //i:在直接块数组里的偏移, j: 二级间接索引块在一级间接索引块里的偏移, k: 间接数据块在二级间接索引块里的偏移
         int j,k;
         for(i=old_datablock_num; i<new_datablock_num && i<DIRECT_DATA_BLOCKS_NUM && new_datablock_i<add_datablock_num; ++i,++new_datablock_i){
   inode->direct_block[i] = new_datablock_index[new_datablock_i];
1034
1036
         if(add_first_indirect_block_num > 0)
  inode->indirect_block = new_first_indirect_block_index;
1039
1041
         if(add_first_indirect_block_num > 0){
1043
1044
             device_clear_sector(DATA_BLOCKS_LOC + inode->indirect_block);
     //read_debug_directory();
//read_debug_indirect_block();
1046
1047
1048
         device_read_sector(first_indirect_block,DATA_BLOCKS_LOC + inode->indirect_block);
1049
          for(j=old_second_indirect_block_num; j<new_second_indirect_block_num; ++j,++new_second_indirect_block_i){
    first_indirect_block->indirect_block_table[j] = new_second_indirect_block_index[new_second_indirect_block_i];
1051
1053
1054
      //read_debug_directory()
//read_debug_indirect_blo
```

int add_all = add_datablock_num + add_first_indirect_block_num + add_second_indirect_block_num;

```
device_write_sector(first_indirect_block,DATA_BLOCKS_LOC + inode->indirect_block);
              //光東允上一个改集兩的二級問接索可與
int old_down_offset_in_first_indirect_block = old_indirect_datablock_num / (BLOCK_SIZE/sizeof(int));
int old_offset_in_second_indirect_block = old_indirect_datablock_num % (BLOCK_SIZE/sizeof(int));
device_read_sector(second_indirect_block,DATA_BLOCKS_LOC + first_indirect_block->indirect_block_table[ old_down_offset_in_first_indirect_block
for(k=old_offset_in_second_indirect_block; k<BLOCK_SIZE/sizeof(int) && new_datablock_i<add_datablock_num; ++k,++new_datablock_i){
    second_indirect_block->indirect_block_table[k] = new_datablock_index[new_datablock_i];
              device write sector(second indirect_block,DATA_BLOCKS_LOC + first_indirect_block->indirect_block_table[ old_down_offset_in_first_indirect_block_table]
        for(j=old_second_indirect_block_num; j<new_second_indirect_block_num && new_datablock_i<add_datablock_num;++j){
//read_debug_directory();
//read_debug_indirect_block();</pre>
                   device_clear_sector(DATA_BLOCKS_LOC + first_indirect_block->indirect_block_table[j]);
 1078
1079
                   device_read_sector(second_indirect_block,DATA_BLOCKS_LOC + first_indirect_block->indirect_block_table[j]);
for(k=0; kdBLOCK_SIZE/sizeof(int) && new_datablock_i<add_datablock_num; ++k,++new_datablock_i){
    second_indirect_block->indirect_block_table[k] = new_datablock_index[new_datablock_i];
1085
1086
                   device_write_sector(second_indirect_block,DATA_BLOCKS_LOC + first_indirect_block->indirect_block_table[j]);
                 inode->file_size = new_size;
1091
                 inode->blocks_number = new_datablock_num;
inode->last_modified_time = time(NULL);
1092
1093
1094
1095
1096 //read_debug_directory();
1097 //read_debug_indirect_block();
             write_inode_disk_table();
1098
1099 //read_debug_directory();
1100 //read_debug_indirect_block();
1101
                 write_bitmap(datablock_bitmap,"datablock");
1102 //read_debug_directory();
1103 //read_debug_indirect_block();
1104
1105
1106
                 int new_datablock_index_array[new_datablock_num];
1107
                 get_file_datablock_index_array(new_datablock_index_array,file_inode_index);
1108
                  int m;
                 for(m=old_datablock_num ; m<new_datablock_num ; ++m){</pre>
1110 //read_debug_directory
1111 //read_debug_indirect_
1112
                      device_clear_sector(DATA_BLOCKS_LOC + new_datablock_index_array[m]);
1113 //read_debug_dire
1114 //read_debug_indirect_block();
1115
1116
                 free(first indirect block);
1117
                 free(second_indirect_block);
1118
1119
```

3.在目录中添加一个目录项

return 0;

1121 }

```
add_dentry_in_directory_by_inode(int directory_inode_index,char file_name[],int file_inode_index){
directory_t* directory=(directory_t*)malloc(file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_disk_table[directory_inode_index].file_size_to_blocks_num_one_more(inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inode_inod
                                          if(directory == NULL){
    return -ENOMEM;
607
                                          read_file(directory,directory_inode_index);
610
612
                                          if(MAX_FILE_SIZE - inode_disk_table[directory_inode_index].file_size < sizeof(dentry)){</pre>
                                                            free(directory);
return -ENOSPC;
614
615
616
617
618
619
                                        //处理当前目录
//如果需要增加新的dabatblock:
int new_datablock_index=-2;
int increase = ((inode_disk_table[directory_inode_index].file_size % BLOCK_SIZE) == 0);
if(increase){
621
622
623
624
625
626
                                                          new_datablock_index = apply_available_bit(datablock_bitmap);
if(new_datablock_index == -1){
627
628
629
630
                                                                              free(directory);
return -ENOSPC;
                                                                                set_bitmap(datablock_bitmap,new_datablock_index,USED);
                                         }
if(increase){
//如果需要增加新的dabatblock: 暂不处理
635
```

4.将一个由 inode 号索引的文件整个读进内存

```
int read_file(void* file_buf,int inode_index){
    inode_mem_t* inode_mem = inode_mem_index_to_pointer(inode_index);

if(file_buf == NULL){
    return -ENOBUFS;
}

if(inode_mem->inode_disk->file_size >= MAX_FILE_SIZE){
    printf("read file more than max size\n");
    return -EFBIG;
}

int i;

int i;

int i;

int blocks_number = inode_disk_table[inode_index].blocks_number;
    int datablock_index_array[blocks_number];
    get_file_datablock_index_array(datablock_index_array,inode_index);
    for(i=0; i<blocks_number; ++i){
        device_read_sector(file_buf+i*BLOCK_SIZE,DATA_BLOCKS_LOC+datablock_index_array[i]);
    return 0;
}
</pre>
```

参考文献

[1] [单击此处键入参考文献内容]