Model Learning and Model Checking of SSH Implementations



Introduction

- protocols: SSH, TLS, SMTP, FTP, TCP, UDP...
- many implementations per protocol
 - implementations MUST/SHOULD/MAY *adhere* to the specifications...





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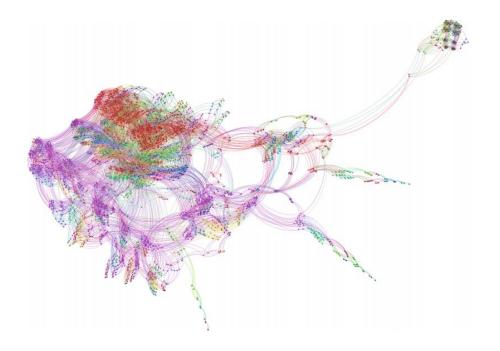
conformance testing





- → automatically infers models for concrete implementations
- → checking conformance of models may be difficult

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Model Checking

- → automatically checks conformance of models to specifications
- → requires models and formalized specifications

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Model Learning + Model Checking

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Application of ML+MC on SSH (a real world protocol):

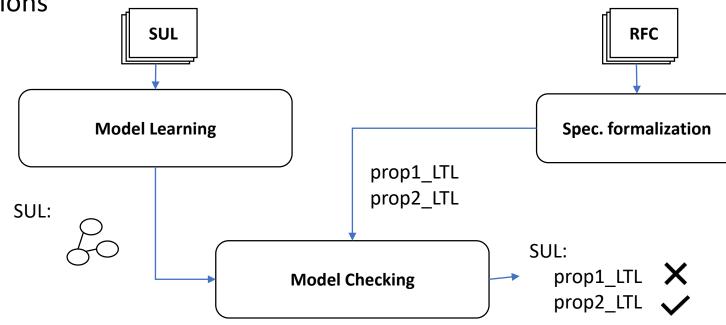
- 1. use Model Learning to infer models of 3 SSH server implementations
- 2. formalize specifications from the SSH RFC standards
- 3. use Model Checking to verify models against these specification

Model Learning + Model Checking

- → automatically infers models for concrete implementations
- → automatically checks conformance of models to specifications

→ requires formalized specifications

Schematic Overview



Model Learning + Model Checking

- → automatically infers models for concrete implementations
- → automatically checks conformance of models to specifications
- → requires formalized specifications

enough for a publication?

Model Learning + Model Checking

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Model Learning:

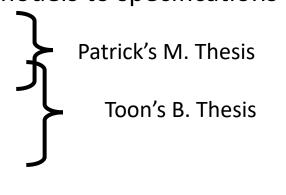
- → requires mapper component
- → requires thorough testing

Model Checking:

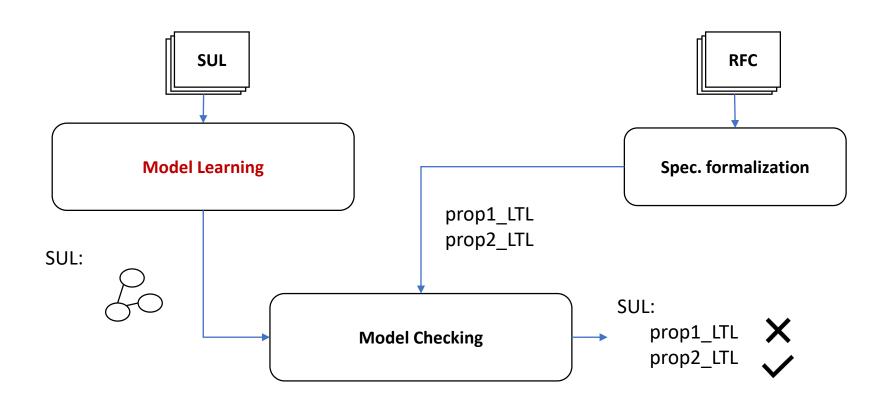
- → requires model transformation
- → requires counterexample validation

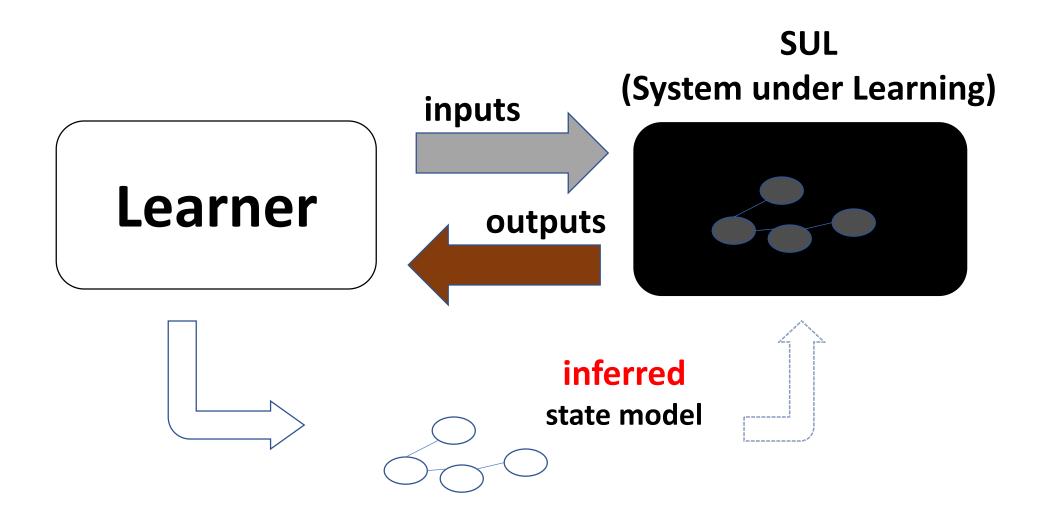
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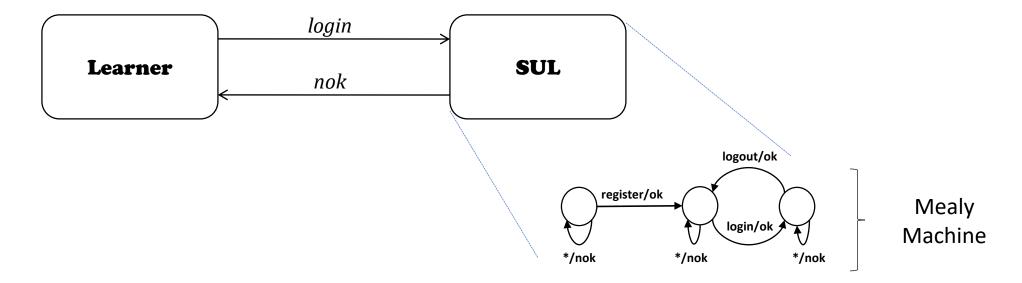


Publication





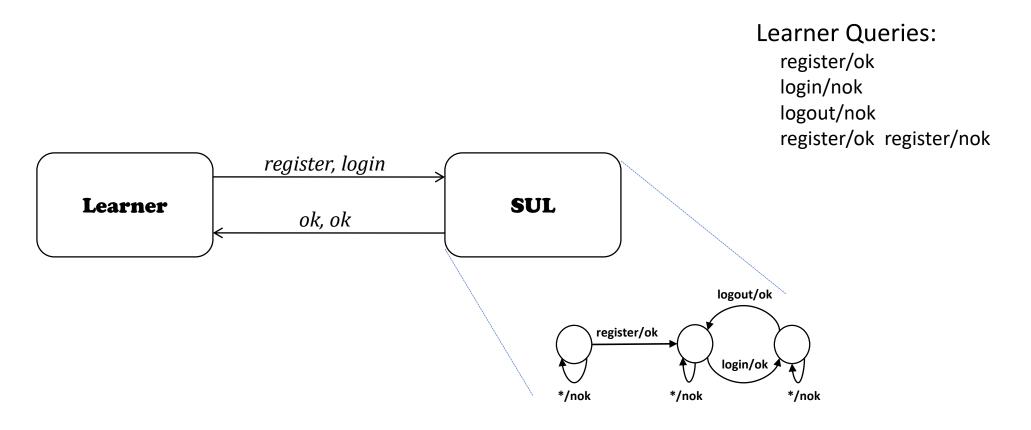
Learner Queries: register/ok



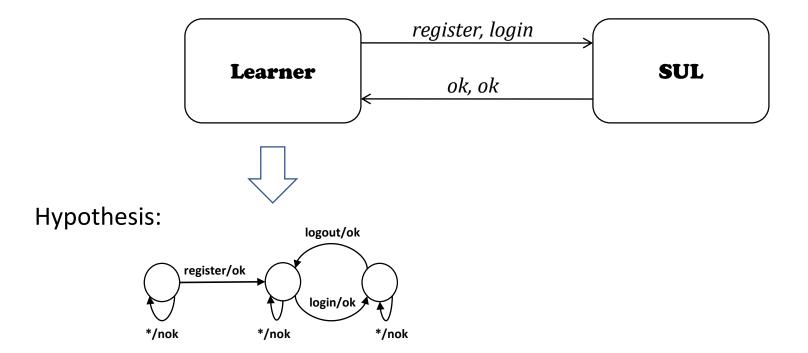
Input Alphabet:

[register, login, logout]

Output Alphabet: [ok, nok]



Input Alphabet: [register, login, logout]



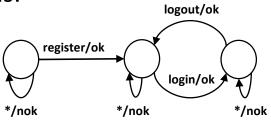
Learner Queries:

register/ok login/nok logout/nok register register/ ok nok

Learner

SUL

Hypothesis:





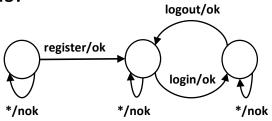
Tester

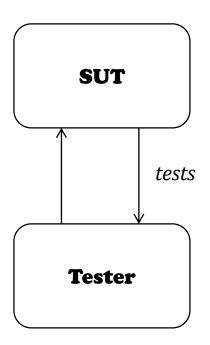
Test Queries:

register register login/ok nok nok



Hypothesis:

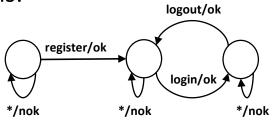


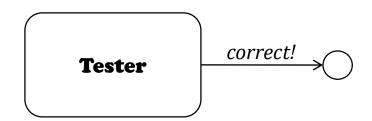


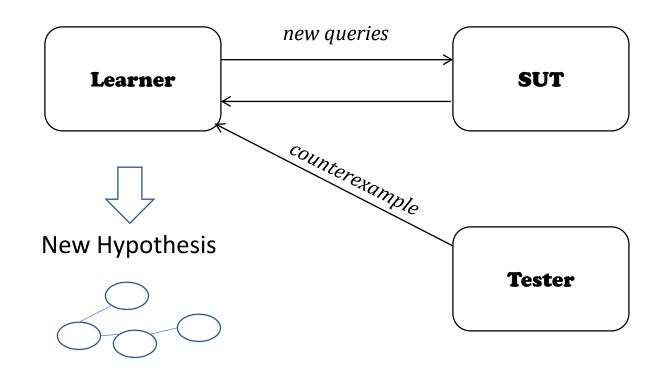
Learner

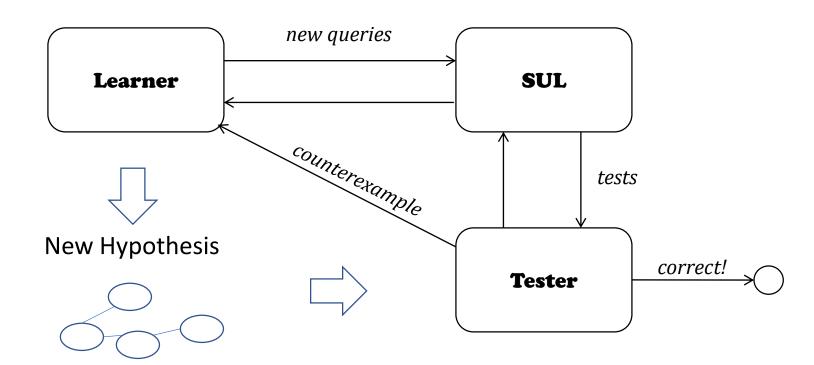
SUT

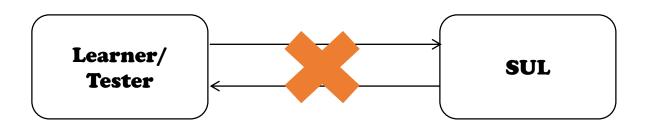
Hypothesis:

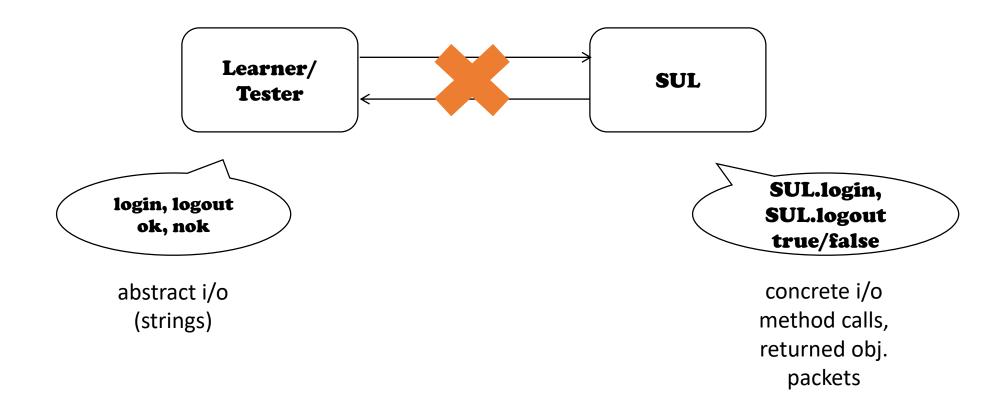


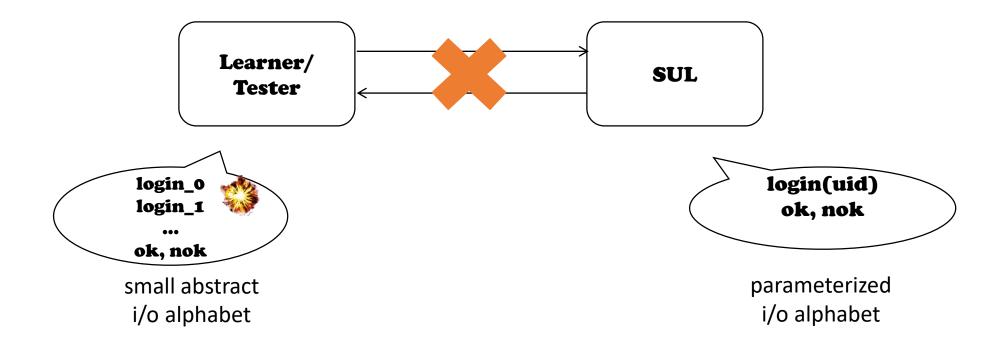




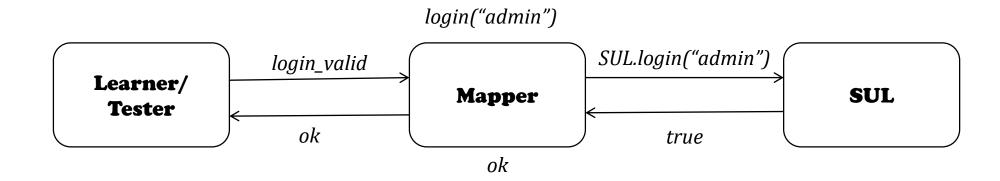






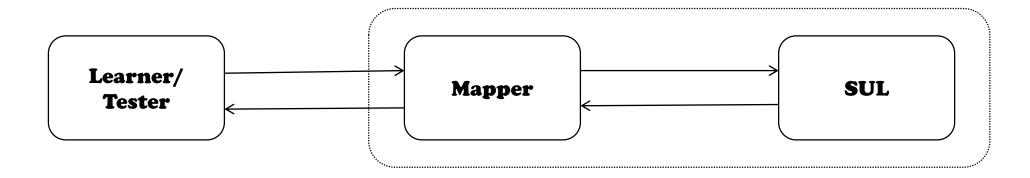


abstract i/o param i/o concrete i/o



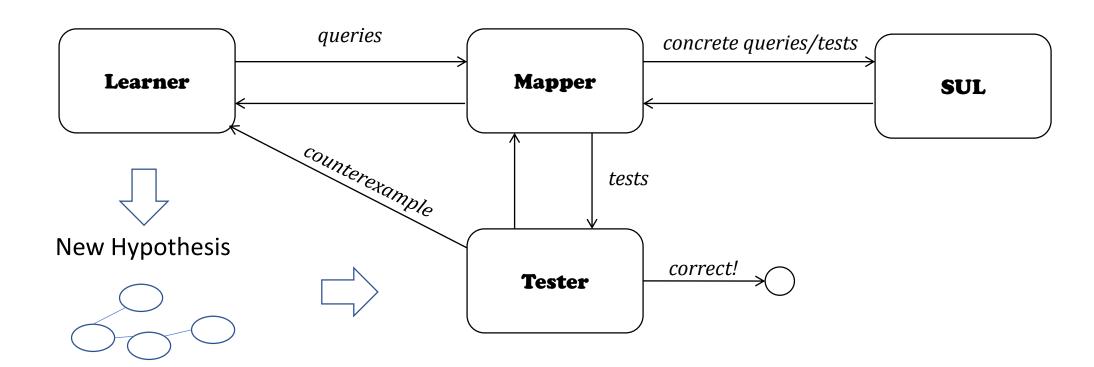
Mapper

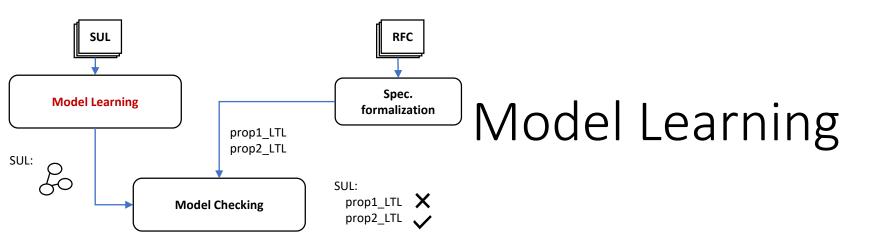
- 1. translates:
 - between abstract and param. i/o
 - between param. i/o and concrete i/o

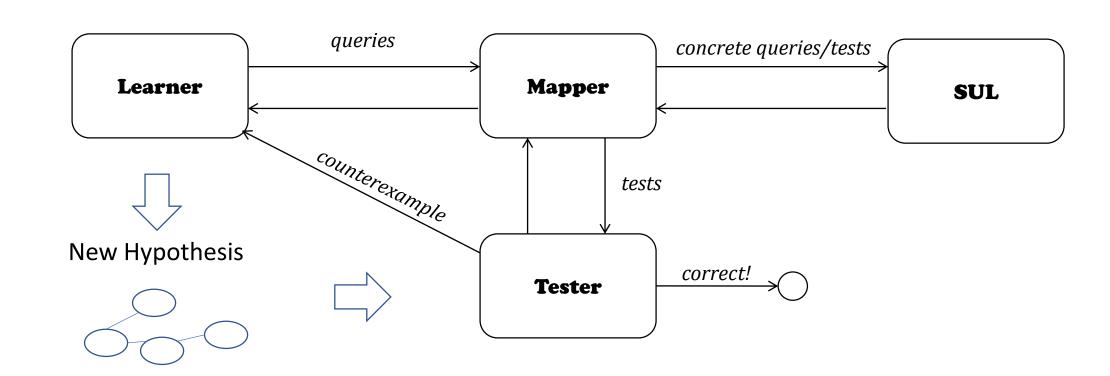


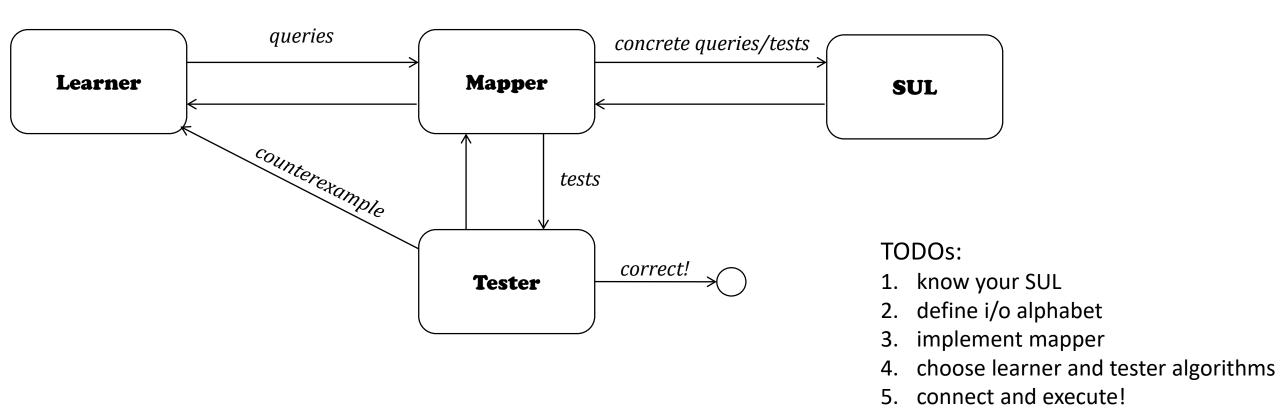
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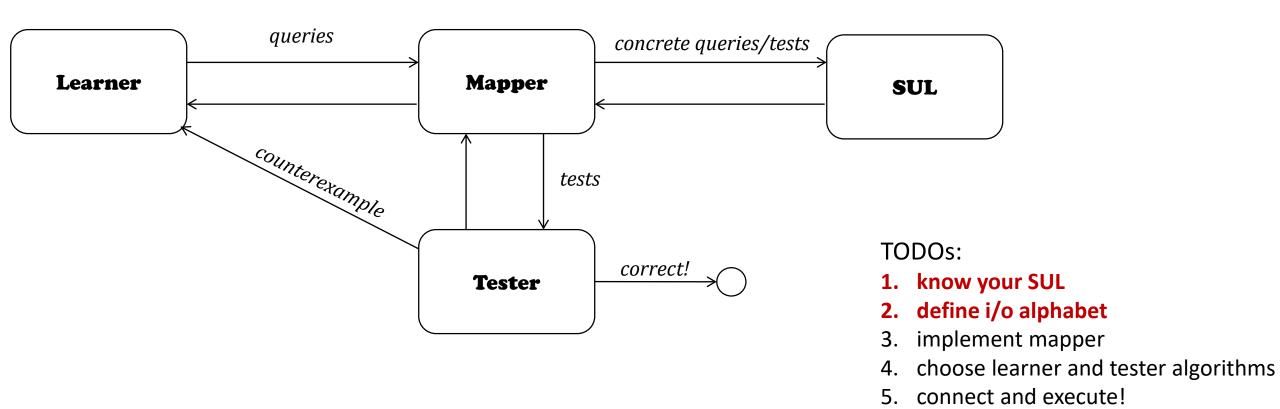
- 1. translates:
 - between abstract and param. i/o
 - > between param. i/o and concrete i/o
- 2. gives a (deterministic) Mealy Machine representation
 - removes time dependencies, non-determinism..











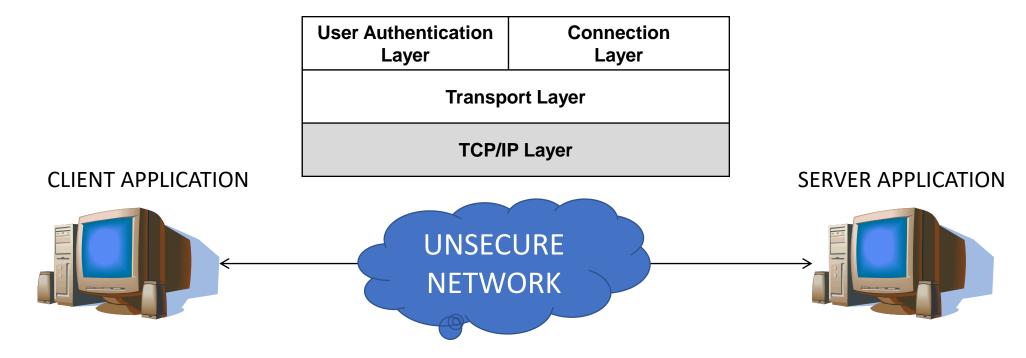
- protocol for operating network services (e.g. terminal) securely over an unsecured network
- client/server application layer protocol, runs on top of TCP



- > protocol for operating network services (e.g. terminal) securely over an unsecured network
- client/server application layer protocol, runs on top of TCP
- Learner + Mapper replaces the SSH CLIENT, goal learn the SSH Server!



- comprises three layers which interoperate (no encapsulation)
- each layer responsible for each of the 3 protocol steps,
- for each we define the happy flow at an abstract level



- > 3 steps
 - 1. establish a secure connection (by exchanging keys)

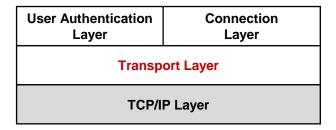
User Authentication Layer Connection Layer

Transport Layer

TCP/IP Layer



- > 3 steps
 - 1. establish a secure connection (by exchanging keys)





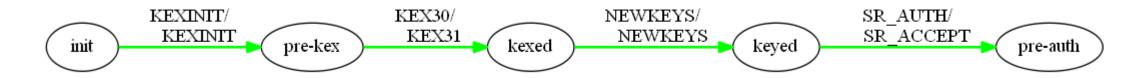
- > 3 steps
 - 1. establish a secure connection (by exchanging keys)
 - 1. exchange preferences (KEXINIT)
 - perform key exchange (KEXxx)
 - put new keys to use (NEWKEYS)
 - 4. engage the auth. service (SR_AUTH)

User Authentication Connection
Layer Layer

Transport Layer

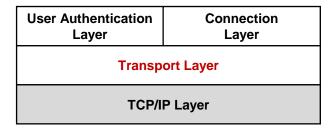
TCP/IP Layer

Happy flow:



Other inputs: DEBUG, IGNORE, DISCONNECT..
Other outputs: DEBUG, IGNORE, DISCONNECT..

- > 3 steps
 - 1. establish a secure connection (by exchanging keys)





- 3 steps
 - 1. establish a secure connection (by exchanging keys) key re-exchange (rekey): same procedure, old keys

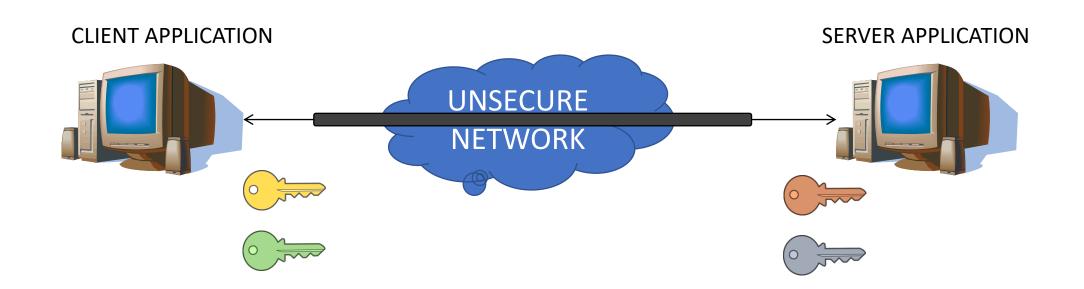
User Authentication Connection
Layer Layer

Transport Layer

TCP/IP Layer

are replaced by new ones

can happen any time after the initial key exchange protocol should not affect operation of higher layer protocols

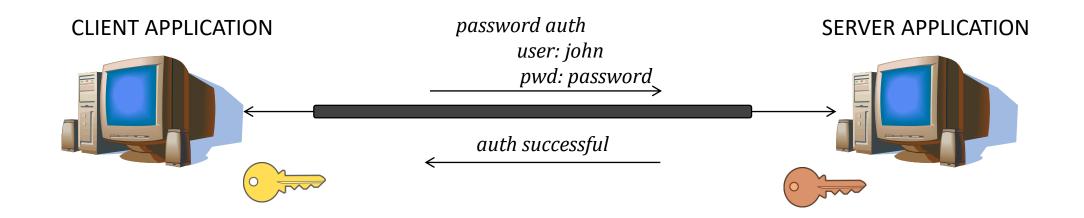


- 3 steps
 - 1. establish a secure connection (by exchanging keys)
 - 2. authentication with server

User Authentication Layer

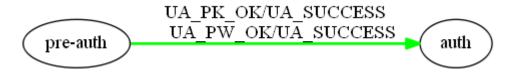
Transport Layer

TCP/IP Layer



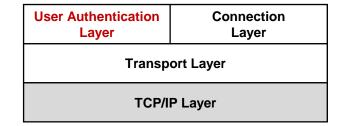
- > 3 steps
 - 1. establish a secure connection (by exchanging keys)
 - 2. authentication with server
 - user/public key auth. UA_PK_OK
 - user/password auth. UA_PW_OK
 - none auth. UA_NONE

Happy flow:



Other inputs: UA_NONE, UA_PK_NOK, UA_PW_NOK...

Other outputs: UA_FAILURE

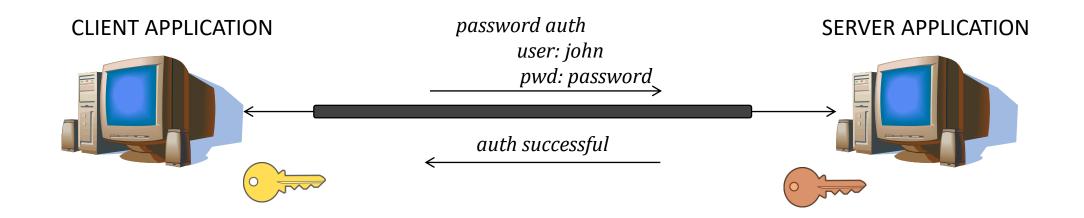


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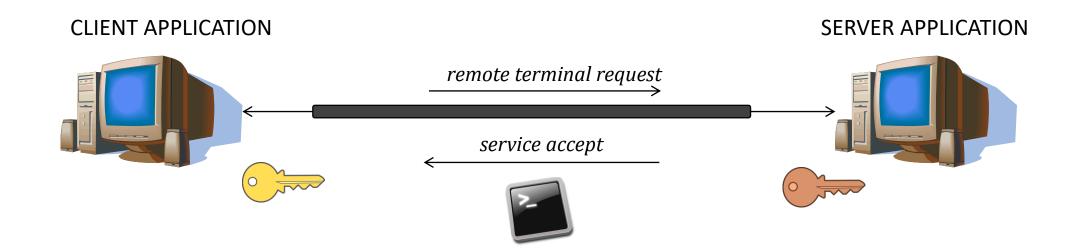


- 3 steps
 - 1. establish a secure connection (by exchanging keys)
 - 2. authentication with server
 - 3. access network services (say remote terminal)

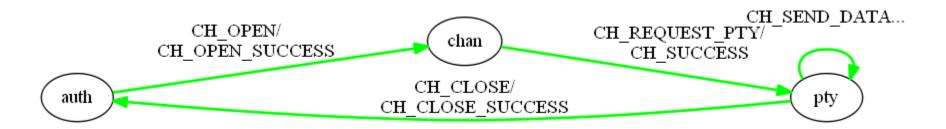
User Authentication Connection
Layer Layer

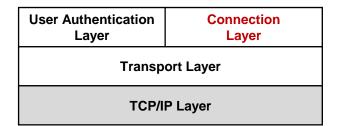
Transport Layer

TCP/IP Layer



- > 3 steps
 - 1. establish a secure connection (by exchanging keys)
 - 2. authentication with server
 - access network services (say remote terminal)
 - 1) open channel (CH_OPEN)
 - request term. service over channel (CH_REQUEST_PTY)
 - channel data management (CH_SEND_DATA)
 - 4) close channel (CH_CLOSE)



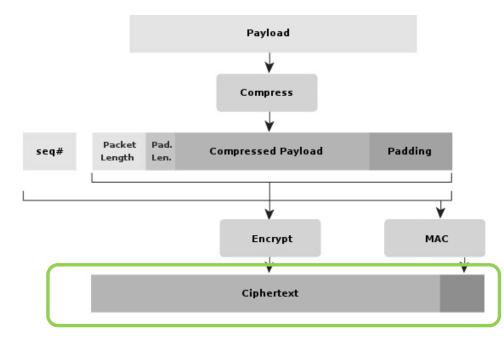


Mapper task

1. translate between abstract, parametrized and concrete i/o

AUTH_PW_OK AUTH_REQUEST("password", "john"...)

- 1. know your SUL
- 2. define i/o alphabet
- 3. implement mapper
- 4. choose learner and tester algorithms
- 5. connect and execute!



Mapper task

- 1. translate between abstract, parametrized and concrete i/o
 - > needs to be able to encrypt/decrypt compress/decompress
 - stores information in variables: encryption keys, session ID, sequence number...

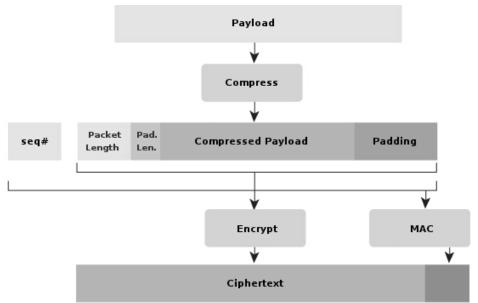
→ implemented by adapting an existing SSH suite implementation

(Paramiko)

AUTH_PW_OK

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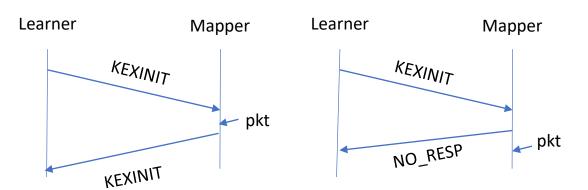


TODOs:

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Mapper task

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 - needs to be able to encrypt/decrypt compress/decompress
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 - → implemented by adapting an existing SSH suite implementation (Paramiko)
- 2. ensure deterministic Mealy Machine representation
 - reliable setting of timing parameters (e.g. NO_RESP timing parameter)



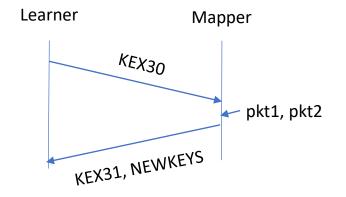
false NO_RESP, mapper should have waited longer

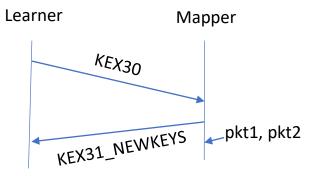
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 - enforce one output per input by concatenating ('_') multiple responses to an input into one output



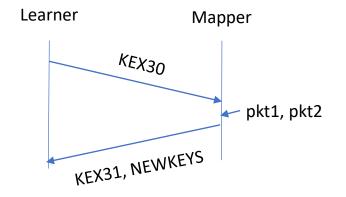


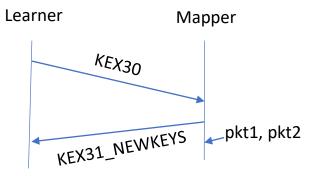
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Mapper task

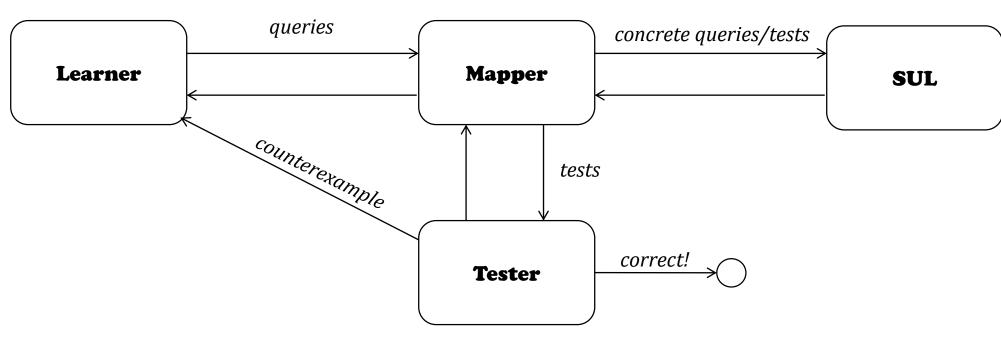
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LearnLib algorithms:

L* , Observation Pack



Tester Algorithms:

Random Walk, W Method,

Yannakakis (Random + Exhaustive)

- 1. know your SUL
- 2. define i/o alphabet
- 3. implement mapper
- 4. choose learner and tester algorithms
- 5. connect and execute!

Note on testing:

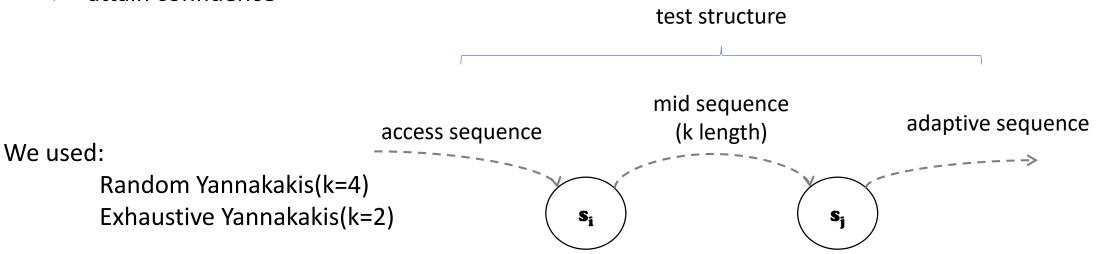
- > testing can never guarantee correctness
- > exhaustive test algs. ensure a well defined level of confidence
 - but lack penetration
- random test algs. have penetration \rightarrow more likely to find CEs
 - but give no formal confidence

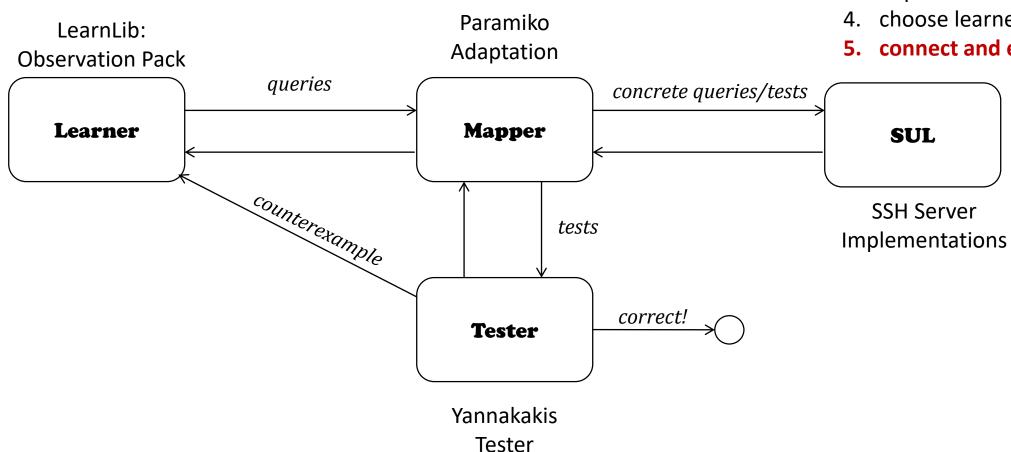
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Example Yannakakis

- →random:
 - > choose bigger k
 - > random mid sequences
- → exhaustive
 - > choose smaller k
 - generate for all mid-sequences
 - attain confidence

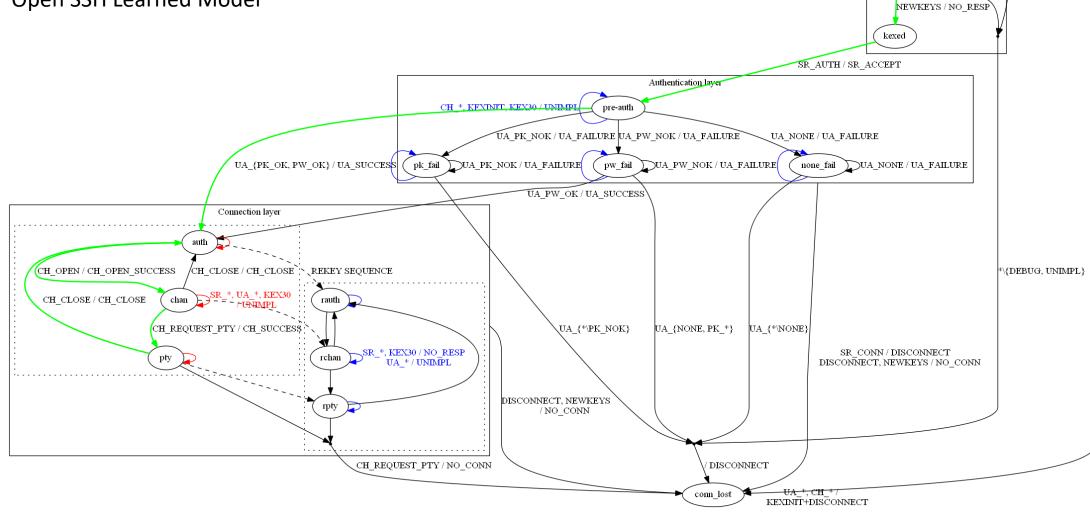
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- know your SUL
- define i/o alphabet
- implement mapper
- choose learner and tester algorithms
- connect and execute!

Open SSH Learned Model



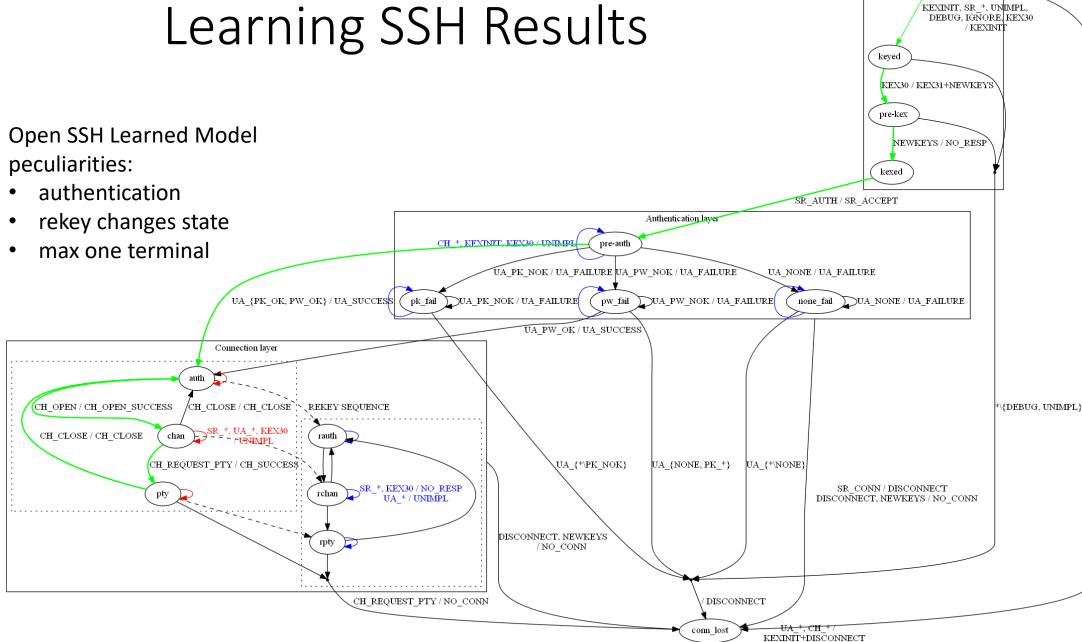
Transport layer

KEX30 / KEX31+NEWKEYS

keyed

pre-kex

KEXINIT, SR_*, UNIMPL, DEBUG, IGNORE, KEX30 / KEXINIT



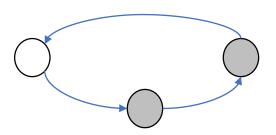
Transport layer

SUT	States		Hypotheses	Mem. Q.	Test Q.
OpenSSH 6.9p1-2	31		4	19836	76418
BitVise 7.23	65		15	24996	58423
DropBear v2014.65	29		8	8357	64478

- rekey (3 step sequence)
- buffering
- > mapper induced behavior

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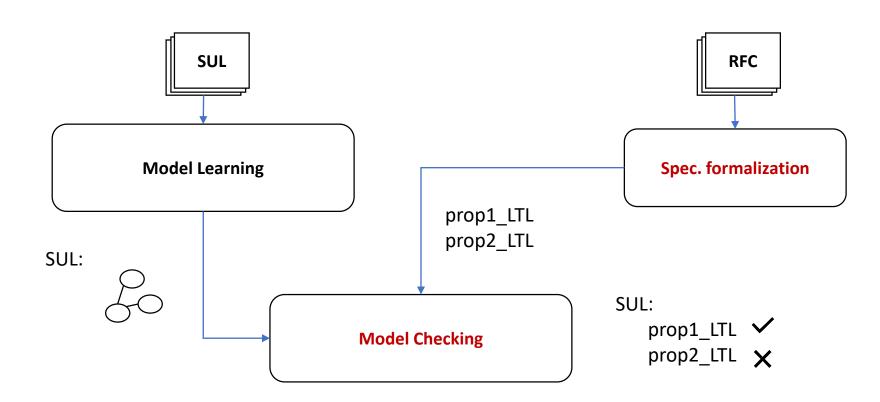


- state permitting rekey
- rekey state

SUT	States		Hypotheses	Mem. Q.	Test Q.
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- > rekey (3 step sequence)
- buffering
- > mapper induced behavior
 - remember, we learn SUL + mapper, not SUL alone

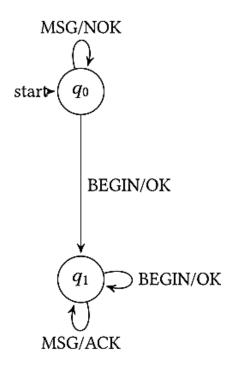
What was done



- > we used NuSMV:
 - > supports LTL, CTL and Real Time CTL specifications
 - > requires conversion to a .SMV model

- we used NuSMV:
 - > supports LTL, CTL and Real Time CTL specifications
 - requires conversion to a .SMV model
 - wrote script to automatically perform this conversion

Mealy Machine



NuSMV

- Model
- → kripke structure with:
 - > state function(next)
 - output function(out)

```
MODULE main
  VAR state : {q0, q1};
  inp : {BEGIN, MSG};
  out : {OK, NOK, ACK};
  ASSIGN
  init(state) := q0;
  next(state) := case
    state = q0 & inp = BEGIN: q1;
    state = q0 & inp = MSG: q0;
    state = q1 & inp = BEGIN: q1;
    state = q1 & inp = MSG: q1;
    esac;
    state = q0 & inp = BEGIN: OK;
    state = q0 & inp = MSG: NOK;
    state = q1 & inp = BEGIN: OK;
    state = q1 & inp = MSG: ACK;
  esac;
```

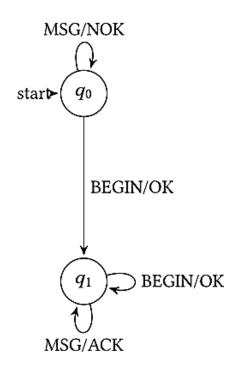
- we used NuSMV:
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MODULE main

VAR state : {q0, q1};

requires conversion to a .SMV model

Mealy Machine



NuSMV Model

- → kripke structure with:
 - state function(next)
 - output function(out)

```
inp : {BEGIN, MSG};
out : {OK, NOK, ACK};
ASSIGN
init(state) := q0;
next(state) := case
  state = q0 \& inp = BEGIN: q1;
                                    G (inp=BEGIN -> out=OK)
  state = q0 \& inp = MSG: q0;
  state = q1 & inp = BEGIN: q1;
  state = q1 & inp = MSG: q1;
                                    G (out=OK ->
  esac;
                                       X (inp=MSG -> out=ACK) )
  state = q0 & inp = BEGIN: OK;
  state = q0 & inp = MSG: NOK;
  state = q1 & inp = BEGIN: OK;
                                    G U (out=OK ->
  state = q1 & inp = MSG: ACK;
                                          X (inp=MSG -> out=NOK))
esac;
```

- we used NuSMV:
 - > supports LTL, CTL and Real Time CTL specifications

MODULE main

ASSIGN

esac;

esac;

VAR state : {q0, q1};

inp : {BEGIN, MSG};
out : {OK, NOK, ACK};

init(state) := q0;
next(state) := case

state = q0 & inp = BEGIN: q1;

state = q1 & inp = BEGIN: q1;
state = q1 & inp = MSG: q1;

state = q0 & inp = BEGIN: OK;

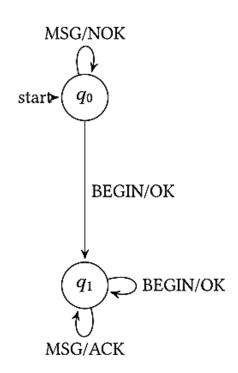
state = q0 & inp = MSG: NOK; state = q1 & inp = BEGIN: OK;

state = q1 & inp = MSG: ACK;

state = q0 & inp = MSG: q0;

requires conversion to a .SMV model

Mealy Machine



NuSMV Model

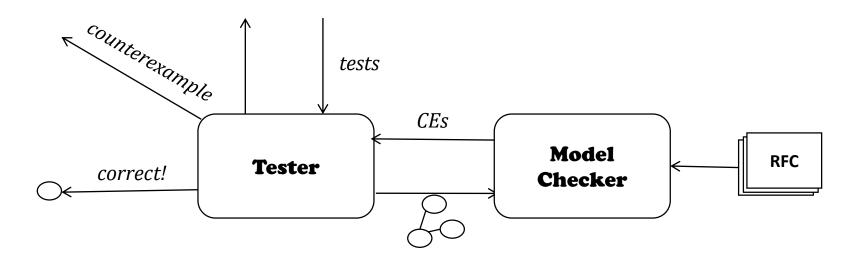
- → kripke structure with:
 - > state function(next)
 - output function(out)

```
G (inp=BEGIN -> out=OK) ✓
```

$$G (out=OK -> X (inp=MSG -> out=ACK))$$

$$G U (out=OK -> X (inp=MSG -> out=NOK))$$

- we used NuSMV:
 - > supports LTL, CTL and Real Time CTL specifications
 - requires conversion to a .SMV model
- > specification either holds or counterexample (CE) given
 - > CE may
 - ➤ agree with the SUL → non-conformance
 - → disagree with the SUL → a CE for the learner
 - > thus, all CEs must first be confirmed by running it on the system
 - > integrated model checker into testing s.t. all CEs are confirmed



- > LTL formulas with both forward and past modalities
- > checked on the mapper + SUL assembly (not only on the SUL itself), thus results not fully translatable
- > 4 types:
 - > basic properties: describe the SUL + mapper setup, all true
 - > security properties: define the overriding goal of each layer
 - rekey properties: is rekey allowed (does it not disconnect) does rekey preserve state?
 - > functional properties: are MUST/SHOULD statements met

Only one SSH connection is made and once it is gone, it is gone.

- basic properties
- > security properties
- > rekey properties
- functional properties

```
connection
is gone

G (out=NO CONN - >

G (out=NO CONN | out=CH MAX | out=CH NONE))
```

SUL no longer responds

We consider an transport layer state machine secure if there is:

no path from the initial state to the point where the authentication service is invoked without exchanging and employing cryptographic keys.

- basic properties
- security properties
- > rekey properties
- functional properties

```
G ( hasReqAuth - >
O ( ( inp=NEWKEYS & out=NO RESP ) &
O ( ( inp=KEX30 & out=KEX31_NEWKEYS) &
O ( out=KEXINIT ) ) )
```

```
SSH_MSG_CHANNEL_CLOSE
Upon receiving this message, a party MUST send back an SSH_MSG_CHANNEL_CLOSE unless it has already sent this message for the channel.
(RFC 4254, p 9)
```

- basic properties
- > security properties
- > rekey properties
- functional properties

```
SSH_MSG_CHANNEL_CLOSE
Upon receiving this message, a party MUST send back an SSH_MSG_CHANNEL_CLOSE unless it has already sent this message for the channel.
(RFC 4254, p 9)
```

- basic properties
- > security properties
- > rekey properties
- functional properties

→ in red, predicates not expressed in RFC statement, yet deducted from context
 → formalization forces clarification

```
SSH_MSG_USERAUTH_SUCCESS MUST be sent only once. (RFC 4252 p. 5)
```

G (out=UA SUCCESS - > X G out != UA SUCCESS)

- basic properties
- > security properties
- > rekey properties
- functional properties

```
SSH_MSG_USERAUTH_SUCCESS has been sent, any further authentication requests received after that SHOULD be silently ignored. (RFC 4252 p. 5)
```

```
G ( out=UA SUCCESS - > X ( ( authReq - > out=NO RESP ) W (connLost | kexStarted) ) )
```

key exchange does not affect the protocols that lie above the SSH transport layer.

(RFC 4253 p. 24)

- > state based property:
 - → cannot be efficiently formulated by LTL
 - → checked using script

- basic properties
- security properties
- rekey properties
- functional properties

Model Checking Results

	Property	Key word	OpenSSH	Bitvise	DropBear
Security	Trans.		√	✓	✓
	Auth.		√	√	√
Rekey	Pre-auth.		X	✓	√
	Auth.		✓	X	✓
Funct.	Prop. 6	MUST	√	✓	✓
	Prop. 7	MUST	✓	✓	✓
	Prop. 8	MUST	X*	X	√
	Prop. 9	MUST	√	✓	√
	Prop. 10	MUST	\checkmark	✓	✓ -
	Prop. 11	SHOULD	X*	X*	√
	Prop. 12	MUST	√	√	X

UA_SUCC once
 NO_RESP after UA_SUCC
 (*X sends UNIMPL)
 CH_CLOSE after CH_CLOSE

Conclusions and Future Work

Conformance checking of the SSH protocol,

using model learning & model checking

- ➤ inferred models for 3 SSH server implementations
- run extensive testing on models
- formalized and checked models against security properties, as well as server RFC MUST/SHOULD requirements
- > found inconsistencies with limited security impact

Future work:

- > formalize mapper so it is clear what it does (and a concretization can be made)
- ➤ make mapper abstraction less impactful → reduce num. of mapper induced states
- ➤ learn SSH client, model check assembly client/server
- > replace classical learner by a register automata learner, extract parameters and infer their related behavior