UMD DATA605 - Big Data Systems

8.2: Map Reduce

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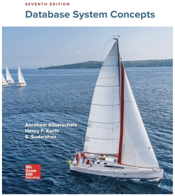
- **References**
  - Silbershatz: Chap 10
  - Ghemawat et al.: *The Google File System*, 2003
  - Dean et al.: *MapReduce: Simplified Data Processing on Large Clusters*, 2004



SCIENCE ACADEMY

MapReduce: Overview

- **MapReduce programming model**
  - Inspired by functional programming (e.g., Lisp)
  - Common pattern of parallel programming to process large number of records
- **Basic algorithm**
  - Apply `map()` to each record
  - Group results by key
  - Apply `reduce()` to results of `map()`
- **Example**
  - *Goal*: Sum length of all tuples in a document
    - E.g.,
      `[() (a,) (a, b) (a, b, c)]`
  - *map(function, set of values)*
    - Apply function to each value (e.g., len)
      `map(len, [(), (a), (a, b), (a, b, c))]) -> [0, 1, 2, 3]`
  - *reduce(function, set of values)*
    - Combine values using a binary function (e.g., add)
      `reduce(add, [0, 1, 2, 3]) -> 6`

MapReduce: Overview

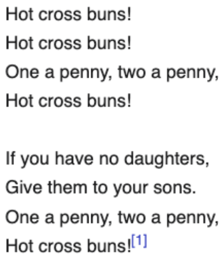
- **Structure of computation**
  - *Read input*
    - Sequentially or in parallel
  - *Map*
    - Extract / compute from records
  - *Group by key*
    - Sort and shuffle
  - *Reduce*
    - Aggregate, summarize, filter, transform
  - *Write result*
- **Division of responsibilities**
  - User specifies `map()` and `reduce()` functions to solve problem
  - MapReduce framework (e.g., Hadoop, Spark) implements algorithm

MapReduce: Word Count

- **Word Count**
  - “Hello world” of MapReduce
  - Huge text file (can’t fit in memory)
  - Count occurrences of each distinct word
- **Linux solution**

```
> more doc.txt
One a penny, two a penny, hot cross buns.
> words doc.txt | sort | uniq -c
a 2
buns 1
cross 1
...
```

  - `words` outputs words one per line
  - Unix pipeline is parallelizable in MapReduce sense
- **Sample application**
  - Analyze web server logs for popular URLs



Hot cross buns!
Hot cross buns!
One a penny, two a penny,
Hot cross buns!

If you have no daughters,
Give them to your sons.
One a penny, two a penny,
Hot cross buns![1]

MapReduce: Word Count

Action

Read input

Map:

- Invoke **map**() on each input record
- Emit 0 or more output data items

Group by key:

- Gather all outputs from **map**() stage
- Collect outputs by keys

Reduce:

- Combine the list of outputs with same keys

Python code

```
values = read(file_name)

def map(values):
  # values: words in document
  for word in values:
        emit(word, 1)

def reduce(key, values):
   # key: a word
   # value: a list of counts
   result = 0
   # result = sum(values)
   for count in values:
    result += count
   emit(key, result)
```

Example

"One a penny, two a penny, hot cross buns."

Map:

```
[("one", 1), ("a", 1),
("penny", 1),("two", 1),
("a", 1), ("penny", 1),
("hot", 1), ("cross", 1),
("buns", 1)]
```

Group by key:

```
[("a", [1, 1]),
("buns", [1]),
("cross", [1]),
("hot", [1]),
("one", [1]),
("penny", [1, 1]),
("two", [1])]
```

Reduce:

```
[("one", 1),
("a", 2),
("penny", 2),
("two", 1),
("hot", 1),
("cross", 1),
("buns", 1)]
```

MapReduce: Log Processing

- **Goal**:
  - Log file recording access to a website with format (date, hour, filename)
  - Find how many times each file is accessed during Feb 2013
- **Input**
  - Read file and split into lines
- **Map**
  - Parse each line into 3 fields
  - If date is in the required interval emit(dir_name, 1)
- **GroupBy**
  - Reduce key is the filename
  - Accumulate all (key, value) with the same filename
- **Reduce**
  - Add values for each list of (key, value) with the same filename
  - Output number of accesses to each file
- **Output**
  - Write results on disk separated by

After Input

```
2013/02/21 10:31:22.00EST  /slide-
2013/02/21 10:43:12.00EST  /slide-
2013/02/22 18:26:45.00EST  /slide-
2013/02/22 18:26:48.00EST  /exer-d
2013/02/22 18:26:54.00EST  /exer-d
2013/02/22 20:53:29.00EST  /slide-
```

After Map

```
['/slide-dir/11.ppt', 1), ...)]
```

After GroupBy

```
[('/slide_dir/11.ppt', 1), ...,
('/slide-dir/12.ppt', [1, 1]), ...
```

After Reduce

```
[('/slide_dir/11.ppt', 1), ...,
('/slide-dir/12.ppt', 2), ...]
```

Output

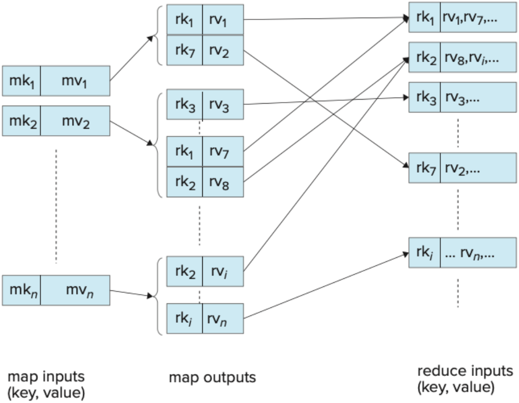
```
/slide_dir/11.ppt 1
...
/slide-dir/12.ppt 2
...
```

MapReduce: Interfaces

- **Input**: Read key-value pairs `List[Tuple[k, v]]`
- **Programmer** specifies two methods `map` and `reduce`
- **Map**
  - `Map(Tuple[k, v]) → List[Tuple[k, v]]`
  - Take a key-value pair and output a set of key-value pairs
    - E.g., key is a file, value is the number of occurrences
    - “One a penny” → `[("One", 1), ("a", 1), ("penny", 1)]`
  - There is one `Map` call for every `(k, v)` pair
- **GroupBy**
  - `GroupBy(List[Tuple[k, v]]) → List[Tuple[k, List[v]]]`
  - Group and optionally sort all the records with the reduce key
- **Reduce**
  - `Reduce(Tuple[k, List[v]]) → Tuple[k, v]`
  - All values `v'` with same key `k'` are reduced together
  - There is one `Reduce` call per unique key `*k'`
- **Output**: write key-value pairs `List[Tuple[k, v]]`

MapReduce: Data Flow

- Focusing on MapReduce flow of the data to expose the parallelism

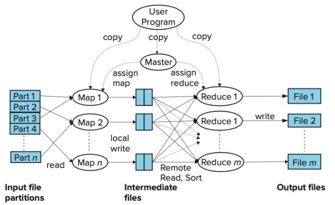


Input Map GroupBy Reduce

- **Input**
- **Map**
  - mk_i = map keys
  - mv_i = map values
- **GroupBy**
  - Shuffle / collect the data
- **Reduce**
  - rk_i = reduce keys
  - rv_i = reduce values
  - Reduce outputs are not shown

MapReduce: Parallel Data Flow

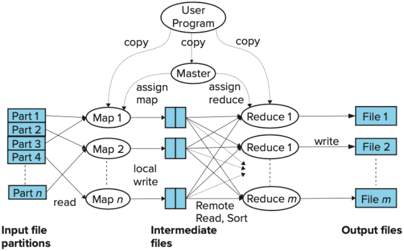
- **User program** specifies map/reduce code
  - *MasterNode* sends code to all computing nodes
  - *Machines* reused for multiple computations (`Map`, `Reduce`) at different times
  - All operations use HDFS as storage



- **Map**
  - *n* data chunks to process
  - Functions executed in parallel on *k* machines
  - Output data saved on disk
- **GroupBy / Sort**
  - Output data sorted and partitioned by reduce key
  - Files created for each `Reduce` task
- **Reduce**
  - Functions executed in parallel on multiple machines
  - Each works on part of the data
  - Output data saved on disk

MasterNode Responsibilities

- *MasterNode* **coordinates / schedule tasks**
  - Task status: idle, in-progress, completed
  - Schedule idle tasks as workers become available
  - `Map` task completion sends location and sizes of intermediate files to Master
  - Master informs `Reduce` tasks
  - Schedule idle `Reduce` tasks
- *MasterNode* **pings workers to detect failures**
  - Heartbeat



Dealing with Failures

- **Map worker failure**
  - Reset failed map tasks to idle
  - Notify reduce workers when task is rescheduled
- **Reduce worker failure**
  - Reset in-progress tasks to idle
  - Restart reduce task
- **Master failure**
  - Abort MapReduce task
  - Notify client

How Many `Map` and `Reduce` Jobs?

- Number of map tasks $= M$
- Number of reduce tasks $= R$
- Number of worker nodes $= N$
- Typically $M \gg N$
  - Pros:
    - Improve dynamic load balancing
    - Speed up recovery from worker failures
  - Cons:
    - More communication between *MasterNode* and *WorkerNodes*
    - Lots of smaller files
- Typically $R > N$
- Usually $R < M$, output is spread across fewer files

Refinements: Backup Tasks

- **Problem**
  - Slow workers significantly lengthen the job completion time
  - Slow workers due to:
    - Older processor
    - Not enough RAM
    - Other jobs on the machine
    - Bad disks
    - OS thrashing / virtual memory hell
- **Solution**
  - Near the end of `Map` / `Reduce` phase
    - Spawn backup copies of tasks
    - Whichever one finishes first “wins”
- **Result**
  - Shorten job completion time

Refinement: Combiners

- **Problem**
  - Often a `Map` task produces many pairs for the same key `k` `[(k1, v1), (k1, v2), ...]`
  - E.g., common words in the word count example
  - Increase complexity of the `GroupBy` stage
- **Solution**
  - Pre-aggregate values in the `Map` with a `Combine` `[k1, (v1, v2, ...), k2, ([...])]`
  - `Combine` is usually the same as the `Reduce` function
  - Works only if `Reduce` function is commutative and associative
- **Result**
  - Better data locality
  - Less shuffling and reordering
  - Less network / disk traffic

Refinement: Partition Function

- **Problem**
  - Users want to control key partitioning
  - Inputs to `Map` tasks created by contiguous input file splits
  - Default partition function: `hash(key) mod R`
  - Ensure records with the same intermediate key go to the same worker
- **Solution**
  - Override hash function:
  - E.g., `hash(hostname(URL)) mod R` ensures URLs from a host end up in the same output file

Implementations of MapReduce

- There are many implementations of map reduce
  - **Google**
    - Not available outside Google
  - **Hadoop**
    - Open-source in Java
    - Uses HDFS for storage
    - Hadoop Wiki: Intro, Getting Started, Map/Reduce Overview
  - **Amazon Elastic MapReduce (EMR)**
    - Hadoop MapReduce on Amazon EC2
    - Also runs Spark, HBase, Hive,
  - **Spark**
  - **Dask**