

# DUALITY IN MONOIDAL CATEGORIES

A tale of how details are important sometimes

Based on `arXiv:2301.03545` with Sebastian Halbig.



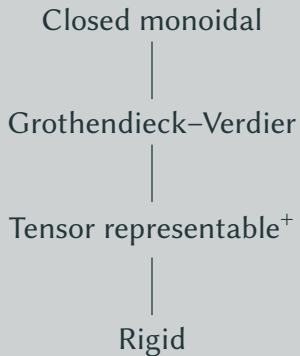
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# The Goal: establishing connections



# Notions of duality

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# Closed monoidal categories

## Definition

A (strict) monoidal category  $(\mathcal{C}, \otimes, 1)$  is called *(left) closed*, if for every  $x \in \mathcal{C}$  there exists an adjunction

$$- \otimes x : \mathcal{C} \rightleftarrows \mathcal{C} : [x, -]$$

with unit  $\eta^x : - \Longrightarrow [x, - \otimes x]$  and counit  $\varepsilon^x : [x, -] \otimes x \Longrightarrow -$ .

## Example

$(\mathbf{Cat}, \times)$ ,  $(\mathbf{k}\text{-Mod}, \otimes_{\mathbf{k}})$ ,  $([\mathcal{C}^{\text{op}}, \mathbf{Set}], \star)$ ,  $(\mathbb{S}\text{-Mod}, \circ)$ .

# (Co)evaluation morphisms in closed categories

Every object  $x$  in a closed category  $\mathcal{C}$  is equipped with natural *(co)evaluation* morphisms

$$\mathrm{coev}_x := \eta_1^x: 1 \longrightarrow [x, x], \qquad \mathrm{ev}_x := \varepsilon_1^x: [x, 1] \otimes x \longrightarrow 1$$

satisfying e.g.,

$$\begin{array}{ccc} x & \xrightarrow{\eta_1^x \otimes x} & [x, x] \otimes x \\ & \searrow \mathrm{id}_x & \downarrow \varepsilon_x^x \\ & & x \end{array}$$

We need a coherent isomorphism  $\phi_x: [x, x] \longrightarrow x \otimes [x, 1]$ .

# Rigid monoidal categories

## Definition

A closed monoidal category  $\mathcal{C}$  is *(left) rigid monoidal* if for every  $x \in \mathcal{C}$  there exists a natural isomorphism  $\phi_y^x: [x, y] \xrightarrow{\sim} y \otimes [x, 1]$ , compatible with the evaluation and coevaluation; e.g.,

$$\begin{array}{ccc}
 [x, x] \otimes x & & \\
 \phi_x^x \otimes x \downarrow & \searrow \varepsilon_x^x & \\
 x \otimes [x, 1] \otimes x & \xrightarrow{x \otimes \varepsilon_1^x} & x
 \end{array}
 \quad \rightsquigarrow \quad
 \begin{array}{ccc}
 x & \xrightarrow{\text{coev}_x \otimes x} & x \otimes x \otimes x \\
 \searrow \text{id}_x & & \downarrow x \otimes \text{ev}_x \\
 & & x
 \end{array}$$

## Example

$$(\text{vect}_k, \otimes_k), \text{Ad}_C^\infty \leq ([C, C], \circ).$$

# Grothendieck–Verdier categories

## Definition

A *Grothendieck–Verdier* (also called *\*-autonomous*) category comprises a monoidal category  $\mathcal{C}$  and an object  $d \in \mathcal{C}$ , such that there exists an equivalence  $D: \mathcal{C}^{\text{op}} \rightarrow \mathcal{C}$ , and a natural isomorphism

$$C(x \otimes y, d) \cong C(x, Dy). \qquad \frac{\text{id}: Dy \rightarrow Dy}{\text{ev}_y: Dy \otimes y \rightarrow d}$$

## Example

$(\Lambda M, \cap, 0)$ ,  $(\text{Chu}(\mathcal{C}, d), \otimes_{\mathcal{C}}, d)$ ,  $(eCe, \otimes_{\mathcal{C}}, De)$ ,  $(\text{ban}_{\mathbf{R}}, \otimes_{\mathbf{R}}, \mathbf{R})$ .

# Tensor representability

Every rigid category yields adjunctions

$$- \otimes x : \mathcal{C} \rightleftarrows \mathcal{C} : - \otimes {}^{\vee}x.$$

what can we say about the category  $\mathcal{C}$ ?

## Definition

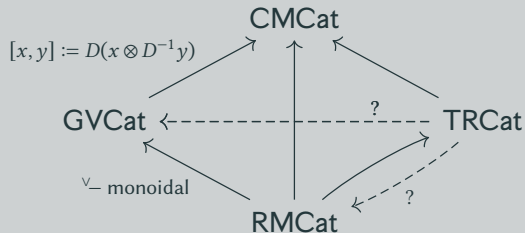
A monoidal category is *tensor representable*, if for every  $x \in \mathcal{C}$  there is an adjunction as in Equation (1).



# Tensor representability in context

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# Relationships among different types of duality



# Tensor representable categories are Grothendieck–Verdier

## Theorem ([HZ23, Theorem 3.2])

*If for every  $x \in \mathcal{C}$ , there exist adjunctions*

$$- \otimes Lx \dashv - \otimes x \dashv - \otimes Rx,$$

*then  $\mathcal{C}$  is Grothendieck–Verdier.*

## Proof.

The assumption implies that  $L$  is a quasi-inverse of  $R$ ; for example,

$$\mathcal{C}(LRx, y) \cong \mathcal{C}(1, y \otimes Rx) \cong \mathcal{C}(x, y).$$

Set  $D := R$ . In order to show that  $\mathcal{C}$  is a Grothendieck–Verdier category, we need to find some  $d \in \mathcal{C}$  such that

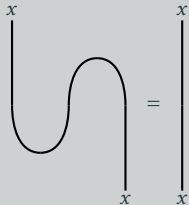
$$\mathcal{C}(y \otimes x, d) \cong \mathcal{C}(y, Dx).$$



**Is tensor representability and  
rigidity the same thing?**

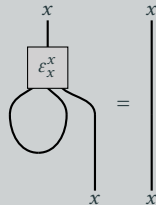
# The problem in string diagrams

The rigid case



$$(x \otimes \text{ev}_x) \circ (\text{coev}_x \otimes x) = \text{id}_x$$

The tensor representable case

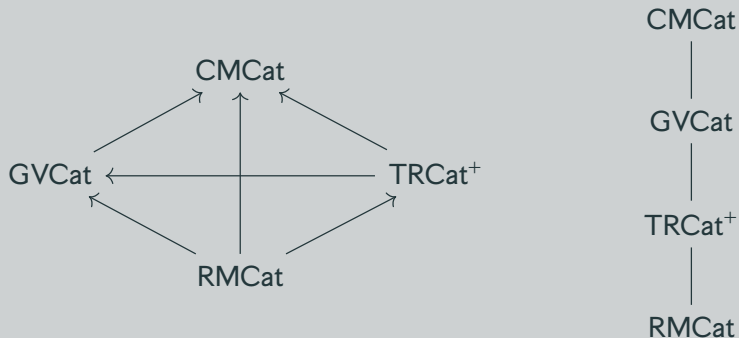


$$\varepsilon_x^x \circ (\eta_1^x \otimes x) = \text{id}_x$$

# A negative result

**Theorem ([HZ23, Theorems 2.5 and 4.15])**

*There exist categories that are tensor representable, but not rigid.*



# Mackey functors

## Definition

Let  $k$  be a field. A *finite-dimensional Mackey functor*  $M$  is a coproduct preserving functor  $M: \mathbf{Sp}_G \rightarrow \mathbf{vect}_k$  such that  $Mx$  is finite-dimensional for every  $x \in \mathbf{Sp}_G$ .

## Theorem

*The category  $\mathbf{mky}$  of finite-dimensional Mackey functors is closed monoidal:*

$$(M \star N)x := \int^{a,b} k\mathbf{Sp}_G(a \otimes b, x) \otimes_k Ma \otimes_k Nb,$$

$$[M, N]x := \int_{a,b} \mathbf{vect}_k(k\mathbf{Sp}_G(x \otimes a, b), \mathbf{vect}_k(Ma, Nb)) \cong \mathbf{mky}(M(x^* \otimes -), N).$$

# Mackey functors in positive characteristic

## Theorem ([HZ23, Theorem 4.15])

*The category  $\text{mky}$  is tensor representable, but not rigid.*

### Proof idea.

1. In  $\text{vect}_k$ :  $\int_y \int^x \cong \int^x \int_y$ .
2. By [Day06]:  $\int^x \cong \int_y$ .
3. By [PS07, Section 9]:  $\text{mky}$  is  $^*$ -autonomous.
4. Calculation: it is even tensor representable.
5. By [Bou05, Lemma 2.2]:  $\text{mky}$  is rigid iff all short exact sequences split.
6. If the order of  $G$  divides the characteristic of  $k$ , then  $\text{mky}$  is not rigid.





# Thank you for your attention!

arXiv:2301.03545

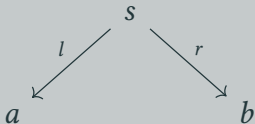


- [Bou05] Serge Bouc. “**The Burnside dimension of projective Mackey functors**”. English. In: *RIMS Kôkyûroku* 1440 (2005), pp. 107–120. ISSN: 1880-2818.
- [Day06] Brian J. Day. “**Compact convolution**”. In: *arXiv e-prints* (2006). arXiv: math/0605463 [math.CT].
- [HZ23] Sebastian Halbig and Tony Zorman. “**Duality in Monoidal Categories**”. In: *arXiv e-prints* (2023). arXiv: 2301.03545 [math.CT].
- [PS07] Elango Panchadcharam and Ross Street. “**Mackey functors on compact closed categories**”. English. In: *J. Homotopy Relat. Struct.* 2.2 (2007), pp. 261–293. ISSN: 2193-8407.

# Spans

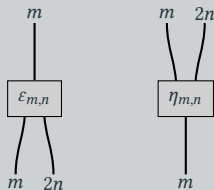
## Definition

Let  $G$  be a finite group, and write  $G\text{-set}$  for the category of finite  $G$ -sets. The category  $\text{Sp}_G$  of *spans* of finite  $G$ -sets has, as objects, finite  $G$ -sets, and as morphisms *isomorphism classes of spans*



# The explicit construction

Define a category  $\mathcal{D}$  whose objects are natural numbers, with morphisms generated by compositions and tensor products of the identities, as well as



subject to appropriate naturality conditions. Define a *monoidal localisation*  $\mathcal{C}$  of  $\mathcal{D}$ , in which adjunction properties for  $\varepsilon$  and  $\eta$  hold. By studying strong monoidal functors  $\mathcal{C} \rightarrow \mathbf{vect}_k$ , we find that the snake equation and the identity must lie in different equivalence classes.