

# Algorithms for the Online Time Dependant Freeze-Tag Problem

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## Abstract

*The abstract lives here.*

## I. INTRODUCTION

The Freeze-Tag Problem (FTP) [1] is a problem in the the field robotics in which a strategy to awaken a swarm must be found. In the original problem, we are given a swarm sleeping robots and a single awake robot. Sleeping robots are awakened when touched by an awake robot.

In this paper we examine a variation of FTP [2] in which each robot has a release time associated with it. Awake robots do not become aware of a sleeping robot's existence until it's release time has been met. We call this variation the Online Time Dependant Freeze-Tag Problem ( OTDFTP).

**Related Work.** Hammer et al. [2] found the Offline Time Dependant Freeze Tag Problem has a lower bound of  $7/3 - \epsilon$ , for any  $\epsilon > 0$ .

**Preliminaries.** Let  $R = \{r_0, r_1, r_2, \dots, r_n\} \subset M$  be the set of  $n$  robots in some continuous metric space  $M$ .  $M$  is a  $d$ -dimensional Euclidean space with distances measured according to an  $L_p$  metric.

Unless otherwise stated, the robot  $r_0$  is the *source robot* and is the only robot that is not in sleep mode. Each robot has 3 modes. In awake mode, a robot is fully functional and free to move. In sleep mode a robot is completely inactive and unable to be activated even if touched by an awake robot. Once the release time of a robot has been reached it enters *wait mode* and awake robots can now sense it's position. In wait mode a robot is available for activation.

Let  $\sigma = (r, v)_1, (r, v)_2, \dots, (r, v)_m$  be a set of locations of robots ordered by their release time.

**Summary of Results.** We find that the competitive ratio of the Nearest Neighbor heuristic is  $5/2$ .

## II. WAKE UP STRATEGIES

### I. Nearest Neighbor Algorithm

To keep robots from traveling in a pack we utilize claims in which when a robot chooses a target, it will claim that target and no other robots will target that robot. [3]

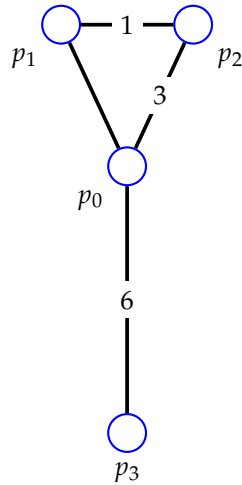
**Algorithm 1** Returns the nearest unclaimed sleeping robot.

**Precondition:**  $A$  is the set of sleeping robots and  $t$  be the robot's current target.

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function NEAREST NEIGHBOR( $A, t$ )
    if  $t \neq \text{NULL}$  then
        return  $t$ 
    if  $A.\text{size} = 0$  then
        return  $\text{NULL}$ 
     $m \leftarrow \text{NULL}$ 
    for  $\text{robot}$  in  $A$  do
        if  $\neg \text{robot.claimed}$  then
            if  $m = \text{NULL}$  then
                 $m \leftarrow \text{robot}$ 
            else
                 $a \leftarrow \text{dist}(\text{self}, \text{robot})$ 
                 $b \leftarrow \text{dist}(\text{self}, m)$ 
                if  $a < b$  then
                     $m \leftarrow \text{robot}$ 
    return  $m$ 
    
```

**Theorem 1.** For any  $\epsilon > 0$ , there exists an instance of the OTDFTP for which the Nearest Neighbor Algorithm results in a makespan less than  $\frac{5}{2} - \epsilon$  times optimal.



**Figure 1**

*Proof.* Let  $G = V, E$  be the graph in Figure 1. The source robot  $r_0$  activates at vertex  $v_0$ . At time  $t = 0$  robot  $r_1$  enters the sleeping state and is immediately awakened by  $r_0$ . At  $t = 3$   $r_2$

enters sleeping state at  $v_1$  and  $r_3$  enters sleeping state at  $v_2$ .  $r_0$  claims  $r_2$  and begins moving down edge  $(v_0, v_1)$ .  $r_1$  claims  $r_3$  and begins moving down edge  $(v_0, v_2)$ .  $r_0$  and  $r_1$  arrive and awaken their respective targets at  $t = 6$ . At the same time  $r_4$  enters sleeping state at  $v_3$  and is claimed by  $r_0$ .  $r_0$  travels down  $(v_1, v_0)$  to  $v_0$  then travels down  $(v_0, v_3)$ .  $r_0$  arrives at  $v_0$  at  $t = 15$ .

In the optimum solution, the source robot  $r_0$  still starts at  $v_0$  and awakens  $r_1$  at  $t = 0$ .  $r_0$  immediately moves down  $(v_0, v_3)$  and  $r_1$  moves down  $(v_0, v_1)$ . At  $t = 3$ ,  $r_1$  arrives at  $v_1$  as  $r_2$  and  $r_3$  are entering sleeping state.  $r_1$  immediately claims and awakens  $r_2$  then claims  $r_2$  and begins moving down  $(v_1, v_2)$ .  $r_1$  arrives at  $v_2$  and awakens  $r_3$  at  $t = 4$ . At  $t = 6$   $r_0$  arrives at  $v_3$  as  $r_4$  is entering sleep mode and immediately awakens it.

This gives us a makespan of 15 for the Nearest Neighbor algorithm and a makespan of 6 for the optimal solution to Figure 1. This gives us a competitive ratio of  $\frac{5}{2}$ .  $\square$

## II. Density Based Algorithm

## III. Sibling Based Algorithm

## REFERENCES

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