Computer & Information Security (372-1-460-1)

Access Control

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Access Control

- ITU-T (International Telecommunication Union)
- Recommendation X.800 defines access control as follows:

"The prevention of unauthorized use of a resource, including the prevention of use of a resource in an unauthorized manner."

Access Control Principles

Authentication:

 Verification that the claimed identity of a user or other system entity are valid.

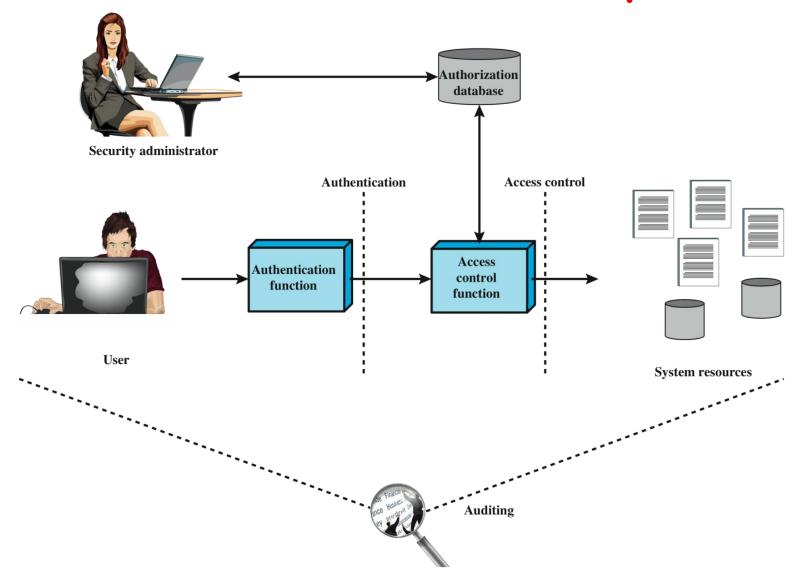
Authorization:

- The granting of a right or permission to a system entity to access a system resource.
- This function determines who is trusted for a given purpose.

Audit:

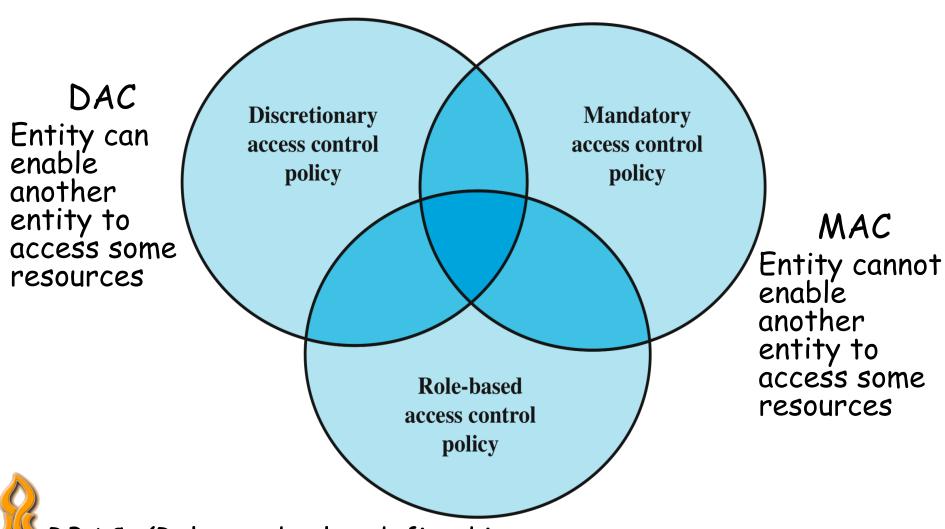
- An independent review and examination of system records and activities
- test for adequacy of system controls, to ensure compliance with established policy and operational procedures,
- _
 - detect breaches in security
 - recommend any indicated changes in control, policy and procedures.

Access Control Principles





Multiple Access Control Policies



RBAC -(Roles and rules defined in the system controls the access

Access Control Requirements

- Reliable input (authenticated user, IP address...)
- Support for proper specifications
- Least privilege
- Separation of duty
- Open and closed policies
- · Policy combinations and conflict resolution
- Administrative policies
- Dual control
- Authorization creep
- Security domain all resources and users that are working under the same security policy

Single-Sign-On (550)

Access Control Basic Elements

- subject entity capable of accessing objects
 - User or application gains access by process
 - typically held accountable for the actions they initiate
 - often have three classes:
 - owner (e.g files, system administrator, project leader),
 - Group (user may belong to authorization group who share same privileges)
 - World (the least amount of authorizations not included in Owner, Group)
- object resource to which access is controlled
 - entity used to contain and/or receive information
 - protection depends on the environment in which access control operates
- access right describes the way in which a subject may access an object
 - e.g. read, write, execute, delete, create, search

Discretionary Access Control (DAC)

- Scheme in which an entity may enable another entity to access some resource
- Often provided using an access matrix
 - one dimension consists of identified subjects that may attempt data access to the resources
 - the other dimension lists the objects that may be accessed
- Each entry in the matrix indicates the access rights of a particular subject for a particular object



Access Matrix

OBJECTS

	File 1	File 2	File 3	File 4
User A	Own Read Write		Own Read Write	
User B	Read	Own Read Write	Write	Read
User C	Read Write	Read		Own Read Write

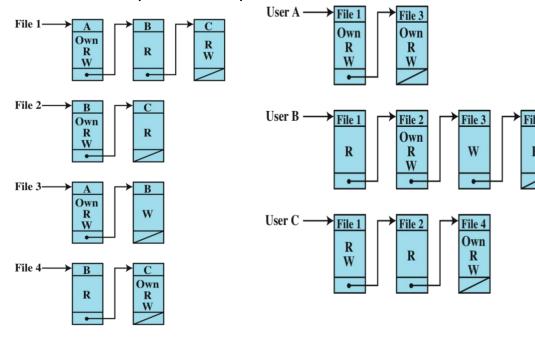
(a) Access matrix



SUBJECTS

Example of Access Control Structures

- (b) ACLs = Access Control Lists
 - Per columns = per object
- (c) Capability ticket
 - Per row= per subject (user\system)
 - Can also contain default\public entry





(b) Access control lists for files of part (a)

(c) Capability lists for files of part (a)

Example of Access Control Structures - Authorization Table

Subject	Access Mode	Object
A	Own	File 1
A	Read	File 1
A	Write	File 1
A	Own	File 3
A	Read	File 3
A	Write	File 3
В	Read	File 1
В	Own	File 2
В	Read	File 2
В	Write	File 2
В	Write	File 3
В	Read	File 4
С	Read	File 1
С	Write	File 1
С	Read	File 2
С	Own	File 4
С	Read	File 4
С	Write	File 4



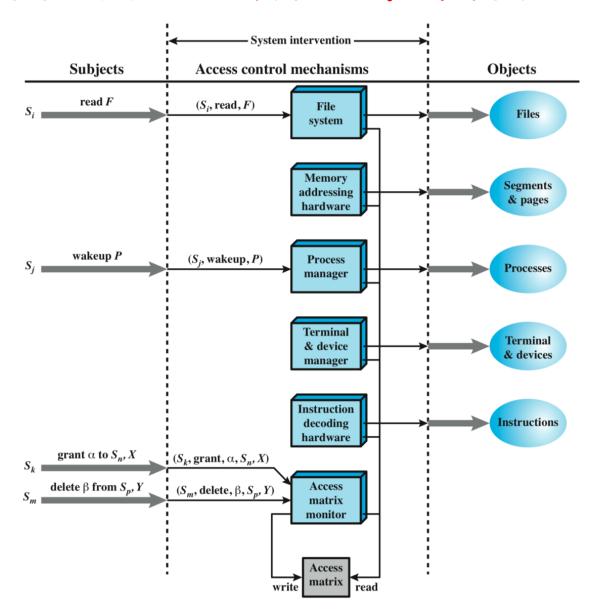
Extended Access Control Matrix

OBJECTS subjects files disk drives processes S_1 S_2 S_3 $\mathbf{F_1}$ $\mathbf{F_2}$ P_1 $\mathbf{P_2}$ $\mathbf{D_1}$ $\mathbf{D_2}$ read owner read * control owner wakeup wakeup seek owner control owner control **SUBJECTS** S_2 write * execute seek * owner S_3 control write stop

* - copy flag set



Access Control Function





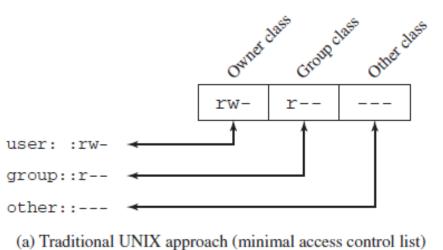
Access Control System Commands

Rule	Command (by S _o)	Authorization	Operation	
R1	transfer $\begin{cases} \alpha^* \\ \alpha \end{cases}$ to S, X	$'\alpha^{*'}$ in $A[S_0, X]$	store $\begin{cases} \alpha * \\ \alpha \end{cases}$ in $A[S, X]$	
R2	grant $\begin{cases} \alpha^* \\ \alpha \end{cases}$ to S, X	'owner' in $A[S_0, X]$	store $\begin{cases} \alpha * \\ \alpha \end{cases}$ in $A[S, X]$	
R3	delete α from S, X	'control' in $A[S_o, S]$ or 'owner' in $A[S_o, X]$	delete α from $A[S, X]$	
R4	$w \leftarrow \mathbf{read} \ S, X$	'control' in $A[S_o, S]$ or 'owner' in $A[S_o, X]$	copy $A[S, X]$ into w	
R5	create object X	None	add column for X to A ; store 'owner' in $A[S_0, X]$	
R6	destroy object X	'owner' in $A[S_0, X]$	delete column for X from A	
R7	create subject S	none	add row for <i>S</i> to <i>A</i> ; execute create object <i>S</i> ; store 'control' in <i>A</i> [<i>S</i> , <i>S</i>]	
R8	destroy subject S	'owner' in $A[S_0, S]$	delete row for <i>S</i> from <i>A</i> ; execute destroy object <i>S</i>	



Example: UNIX File Access Control

 Inode (Index node) = control structure needed for OS for particular file. Contains:



- Each user has unique user identification number (user ID)
- Each user is a member of a primary group identified by a group ID
- 12 protection bits
 - specify read, write, and execute permission for the owner of the file, members of the group and all other users
- First 9 bits include authorization to USER,GROUP and other.
- the owner ID, group ID, and protection bits are part of the file's inode

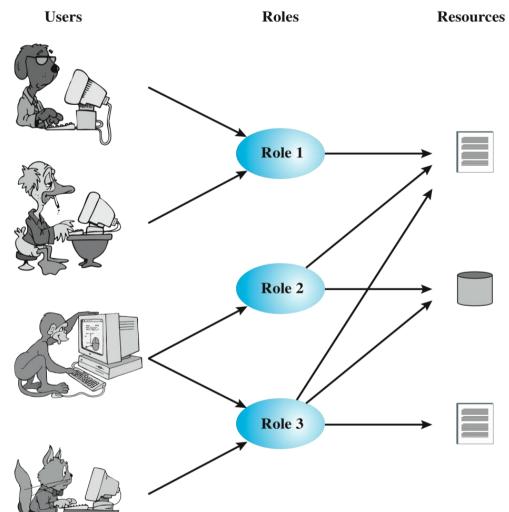


Traditional UNIX File Access Control

- Last 3 bits, specify additional behaviors:
- "set user ID"(SetUID)
- "set group ID"(SetGID)
 - system temporarily uses rights of the file owner / group in addition to the real user's rights when making access control decisions
 - enables privileged programs to access files / resources not generally accessible
- sticky bit
 - when applied to a directory it specifies that only the owner of any file in the directory can rename, move, or delete that file
- superuser
 - is exempt from usual access control restrictions
 - has system-wide access

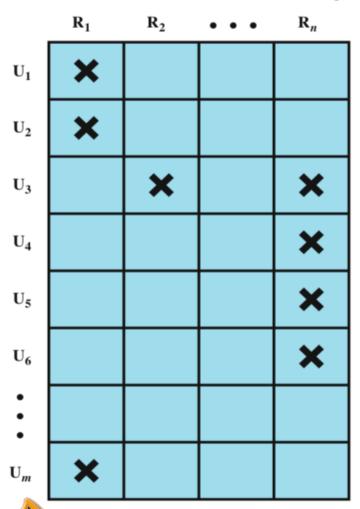


Role-Based Access Control (RBAC)





RBAC - Access Control Matrix



Assigning Users to Roles:

- More Users than Roles
- User Can be assigned to multiple roles
- Roles can be assigned to multiple users

Assigning Roles to Objects:

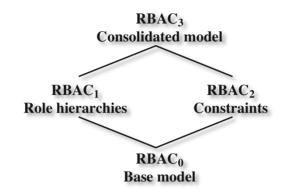
- More Objects than Roles
- Role can be treated as an object for role hierarchies definition.

		OBJECTS								
		$\mathbf{R_1}$	$\mathbf{R_2}$	\mathbf{R}_n	$\mathbf{F_1}$	$\mathbf{F_1}$	$\mathbf{P_1}$	$\mathbf{P_2}$	$\mathbf{D_1}$	$\mathbf{D_2}$
	R_1	control	owner	owner control	read *	read owner	wakeup	wakeup	seek	owner
ES	R_2		control		write *	execute			owner	seek *
ROLES	•									
	\mathbf{R}_n			control		write	stop			

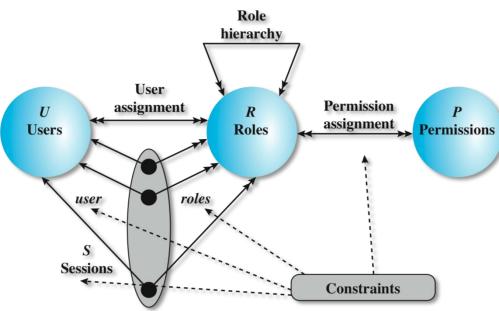
Role-Based Access Control Models

Models	Hierarchies	Constraints
$RBAC_0$	No	No
RBAC ₁	Yes	No
$RBAC_2$	No	Yes
RBAC ₃	Yes	Yes

Minimum functionality of RBAC system



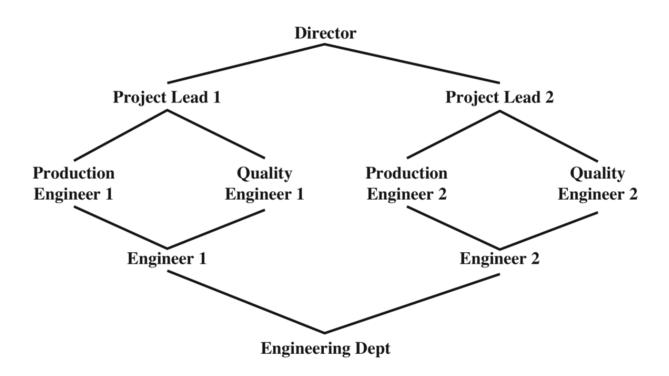
(a) Relationship among RBAC models



(b) RBAC models

Figure 4.9 A Family of Role-Based Access Control Models. $RBAC_0$ is the minimum requirement for an RBAC system. $RBAC_1$ adds role hierarchies and $RBAC_2$ adds constraints. $RBAC_3$ includes $RBAC_1$ and $RBAC_2$. [SAND96]

Example of Role Hierarchy - RBAC1





Constraints - RBAC2

- Provide a means of adapting RBAC to the specifics of administrative and security policies of an organization
- A defined relationship among roles or a condition related to roles
- Mutually exclusive roles
 - a user can only be assigned to one role in the set (either during a session or statically)
 - any permission (access right) can be granted to only one role in the set
- Cardinality
 - setting a maximum number with respect to roles
- Prerequisite roles
 - dictates that a user can only be assigned to a particular role if it is already assigned to some other specified role



Other Access Control Models

Context-bases access control (location)

Rule-based access control (calendar)

 Content-based access control (confidentiality)



Information classification Why is it important?

- Examples
 - unclassified, confidential, secret, top secret (used by US government)
 - public, internal use, confidential (used in the private sector)
- Everyone should know how sensitive or critical is the information in order to handle it appropriately
- To be able to provide proper security level and control measures (e.g., avoid unauthorized disclosure)
- To be able to comply with the organization management policy, regulations
- Follow best practices:
 - least privilege principle,
 - compartmentalization,
 - need-to-know,
 - separation of duties



Multilateral Security

- Multilevel Security (MLS) enforces access control up and down
- Simple hierarchy of security labels may not be flexible enough
- Multilateral security enforces access control across by creating compartments
- Suppose TOP SECRET divided into TOP SECRET {CAT} and TOP SECRET {DOG}
- Both are TOP SECRET but information flow restricted across the TOP SECRET level



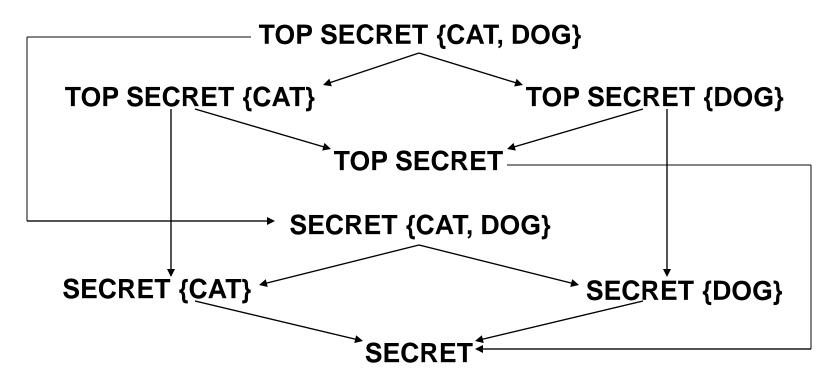
Multilateral Security

- Why compartments?
 - Why not create a new classification level?
- May not want either of
 - TOP SECRET {CAT} ≥ TOP SECRET {DOG}
 - TOP SECRET {DOG} ≥ TOP SECRET {CAT}
- Compartments allow us to enforce the need to know principle
 - Regardless of your clearance, you only have access to info that you need to know



Multilateral Security

Arrows indicate "≥" relationship







Information classification Process

- Define/develop functional policies and procedures
 - Identify the appropriate number of classification levels and their definitions
 - Define the classification process
 - Define data/information owners
 - Declassification
 - How information in each class should be handled/protected (e.g., labeling magnetic media and hard copy, encryption)



Roles and responsibilities

- Data\system owners
 - Classify data and systems
 - Set user access to data and systems
 - Decide on business continuity priorities
- Custodians
 - Handle data and systems that do not belong to them
 - Ensure the security of data and systems according to the owners' specification



Bell-LaPadula (BLP) Confidentiality Model (1973)

- By David Bell and Leonard LaPadula
- Hierarchical,
- Addresses the confidentiality of information in a system
- Mandatory access control
- Formalize the U.S. Department of Defense (DoD)
 multilevel security (MLS) policy; the basis for the Orange Book
- Simple security property read capability only
 - A subject can read information at or below his level of secrecy, but cannot read any information above his level of secrecy (no read up)

Higher secrecy level

- Star property write capability only
 - A subject can write information at or above his level of secrecy, but cannot write any information above his level of secrecy (no write down)

Assigned level

Lower secrecy level

Strong Star property – read and write capability

35 A subject can read and write information only at his level of secrecy



- MLS designed to restrict legitimate channels of communication
- May be other ways for information to flow
- For example, resources shared at different levels may signal information
- Covert channel: "communication path not intended as such by system's designers"



Covert Channel Example

- Alice has TOP SECRET clearance, Bob has CONFIDENTIAL clearance
- Suppose the file space shared by all users
- Alice creates file FileXYzW to signal "1" to Bob, and removes file to signal "0"
- Once each minute Bob lists the files
 - If file FileXYzW does not exist, Alice sent 0
 - If file FileXYzW exists, Alice sent 1
- Alice can leak TOP SECRET info to Bob!



Covert Channel Example

Alice: Create file Delete file Create file Delete file

Bob: Check file Check file Check file Check file

Data: 1 0 1 1 0

Time:



- Other examples of covert channels
 - Print queue
 - ACK messages
 - Network traffic, etc., etc., etc.
- When does a covert channel exist?
 - 1. Sender and receiver have a shared resource
 - 2. Sender able to vary property of resource that receiver can observe
 - 3. Communication between sender and receiver can be synchronized



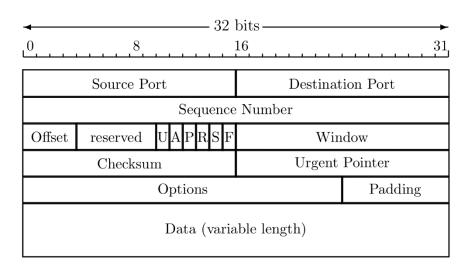
- · Covert channels exist almost everywhere
- Easy to eliminate covert channels...
 - Provided you eliminate all shared resources and all communication
- Virtually impossible to eliminate all covert channels in any useful system
 - DoD guidelines: goal is to reduce covert channel capacity to no more than 1 bit/second
 - Implication is that DoD has given up trying to eliminate covert channels!



- Consider 100MB TOP SECRET file
 - Plaintext version stored in TOP SECRET place
 - Encrypted with AES using 256-bit key,
 ciphertext stored in UNCLASSIFIED location
- Suppose we reduce covert channel capacity to 1 bit per second
- It would take more than 25 years to leak entire document thru a covert channel
- But it would take less than 5 minutes to leak 256-bit AES key thru covert channel!



Real-World Covert Channel



- · Hide data in TCP header "reserved" field
- Or use covert_TCP, tool to hide data in
 - Sequence number
 - ACK number



Real-World Covert Channel

- Hide data in TCP sequence numbers
- Tool: covert_TCP
- Sequence number X contains covert info

