

Research vision

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Programming Languages

```
let r = alloc 0 in r := !r + 1; r
```



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Typed Programming Languages

```
let r = alloc 0 in r := !r + 1; r : ref Nat
```

lightweight and modular

specification

verification

documentation

correct by construction

In today's world
software is
everywhere!

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Typed Programming Languages

But what about behaviour?

```
let r = alloc 0 in r := !r + 1; r : ref Nat
```

r is fresh

$!r > 0$

other
effects
like I/O ?

user-
defined
effects ?

...

State of affairs in type-based reasoning

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Values

Well-understood, uniform,
and thoroughly studied! :)

● Refinement types

`Odd ⊆ Nat Even ⊆ Nat`

`Vec A n = { l : List A | len l = n }`

● Dependent types

`Vec a n = Σ l : List A. (len l = n)`

● Homogeneous implementations

- Agda, Coq, Idris, F*, L.Haskell, ...

State of affairs in **type-based reasoning**

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● **Refinement types**

$\text{Odd} \sqsubseteq \text{Nat}$ $\text{Even} \sqsubseteq \text{Nat}$

$\text{Vec } A \ n = \{ \ell : \text{List } A \mid \text{len } \ell = n \}$

● **Dependent types**

$\text{Vec } a \ n = \sum \ell : \text{List } A. (\text{len } \ell = n)$

● **Homogeneous implementations**

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Effects and behaviour

Scattered landscape, effect-specific, little uniformity! :(

● **Hoare Type Theory** (state)

$M : \Psi.X.\{P\}x:A\{Q\}$

● **F*** (state, exceptions, but no I/O)

$M : \text{ST } A \ \text{wp}_{\text{ST}}$

~~$M : \text{IO } A \ \text{wp}_{\text{IO}}$~~

● **Session Types** (I/O & channels)

$c : ?\text{Nat}.!\text{String}.!\text{Nat}.T$

● **Graded monads, param. monads**

My vision: no need for this non-uniformity!

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- © **Goal:** a general, uniform framework for reasoning about **effects**
 - wide range of effects (state, I/O, exceptions, probability, ...)
 - primitive and user-defined effects
 - combinations of effects

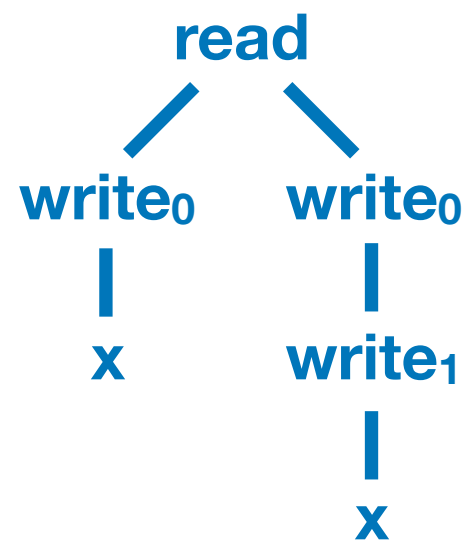
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- © **Answer:** **algebraic effects** and **effect handlers** (rather than just monads)
 - operations and equations
 - reveal the **fundamental underlying tree-like structure** of effects
 - **effect handlers** are algebras; **handling** is homomorphism application

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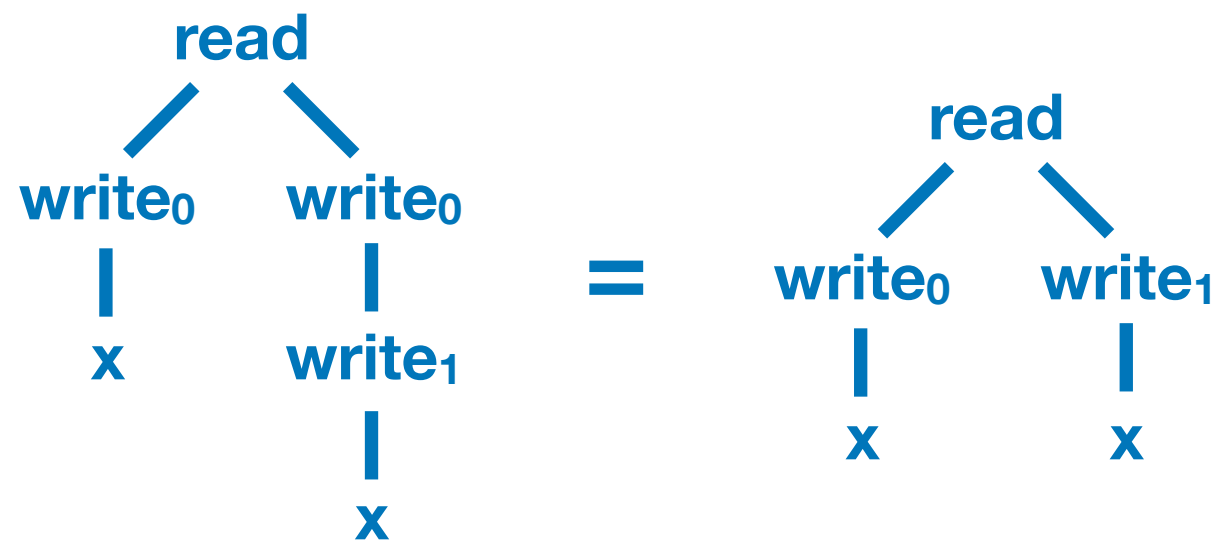
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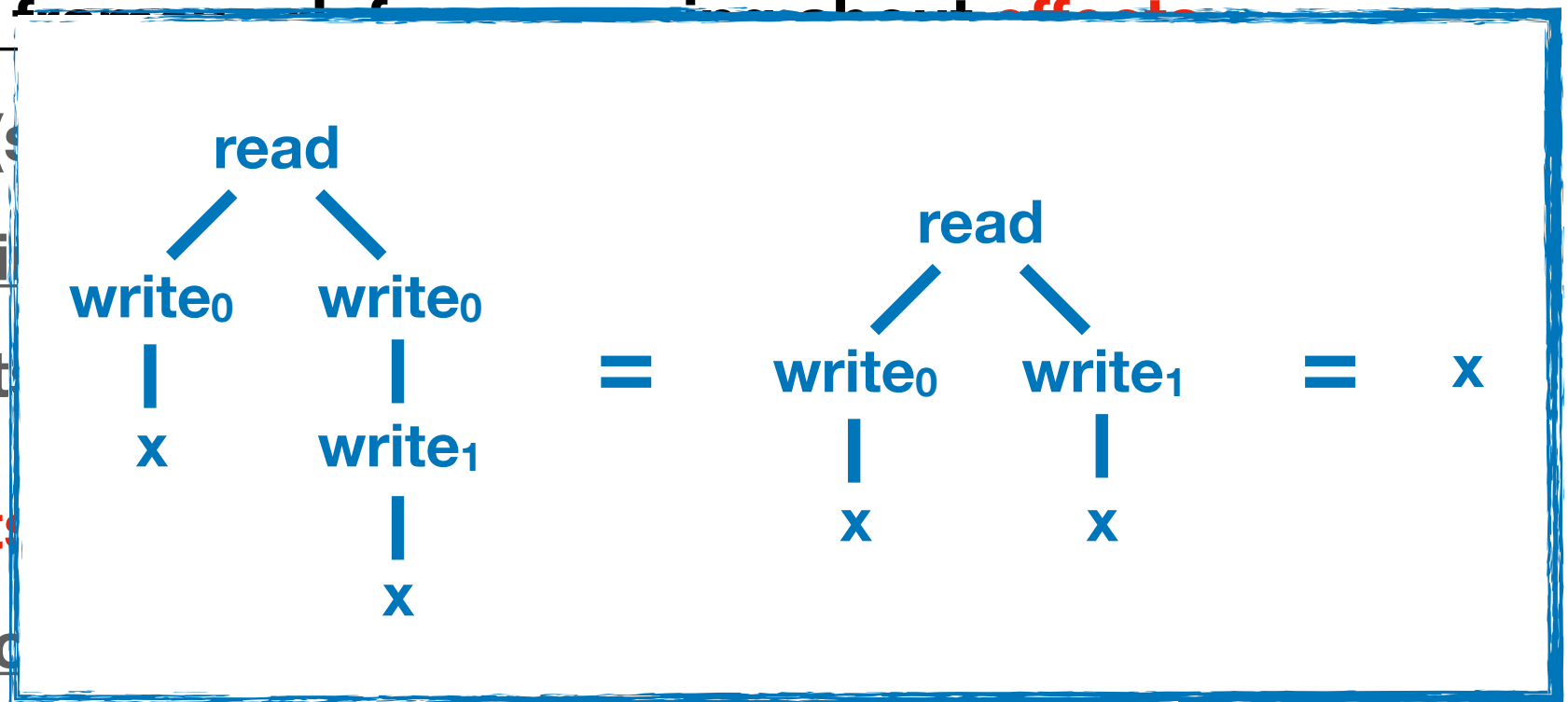
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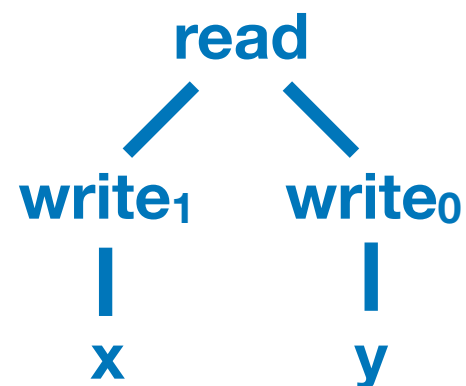
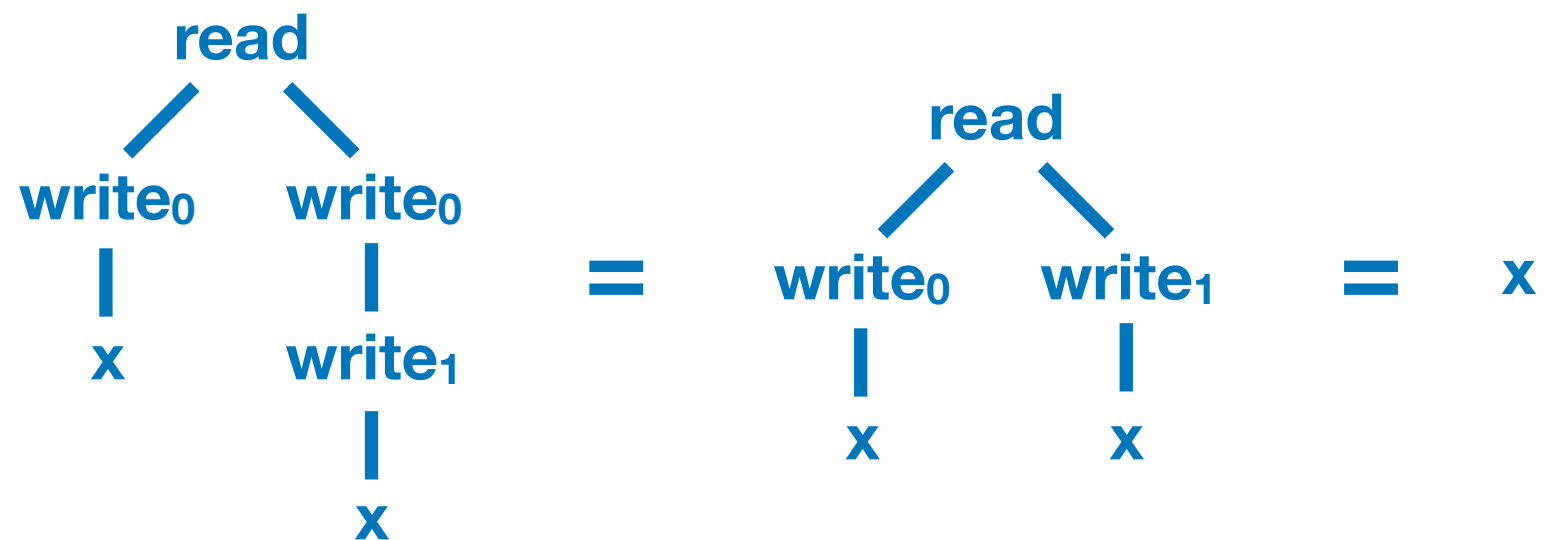
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Goal: a general, uniform framework for reasoning about effects

- wide range of effects (stateful computations)
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Answer: algebraic effects

- operations and equations that capture the essential properties of effects
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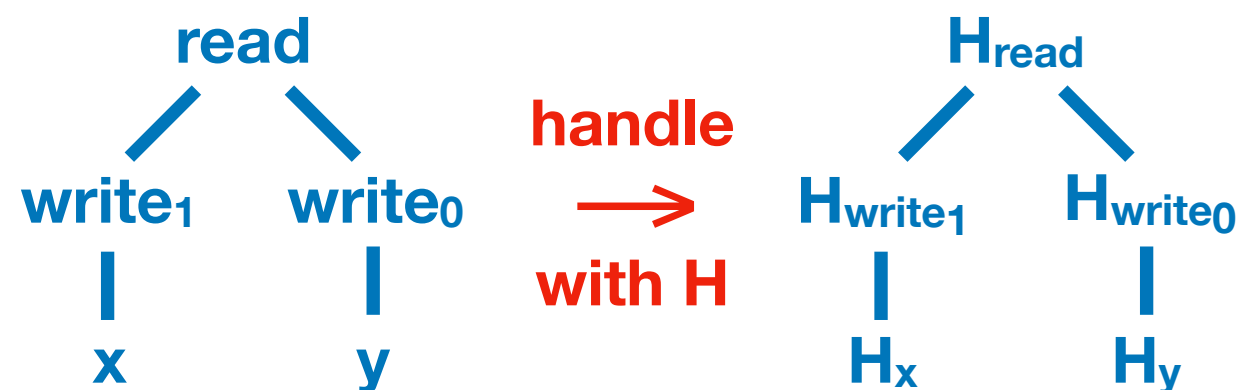
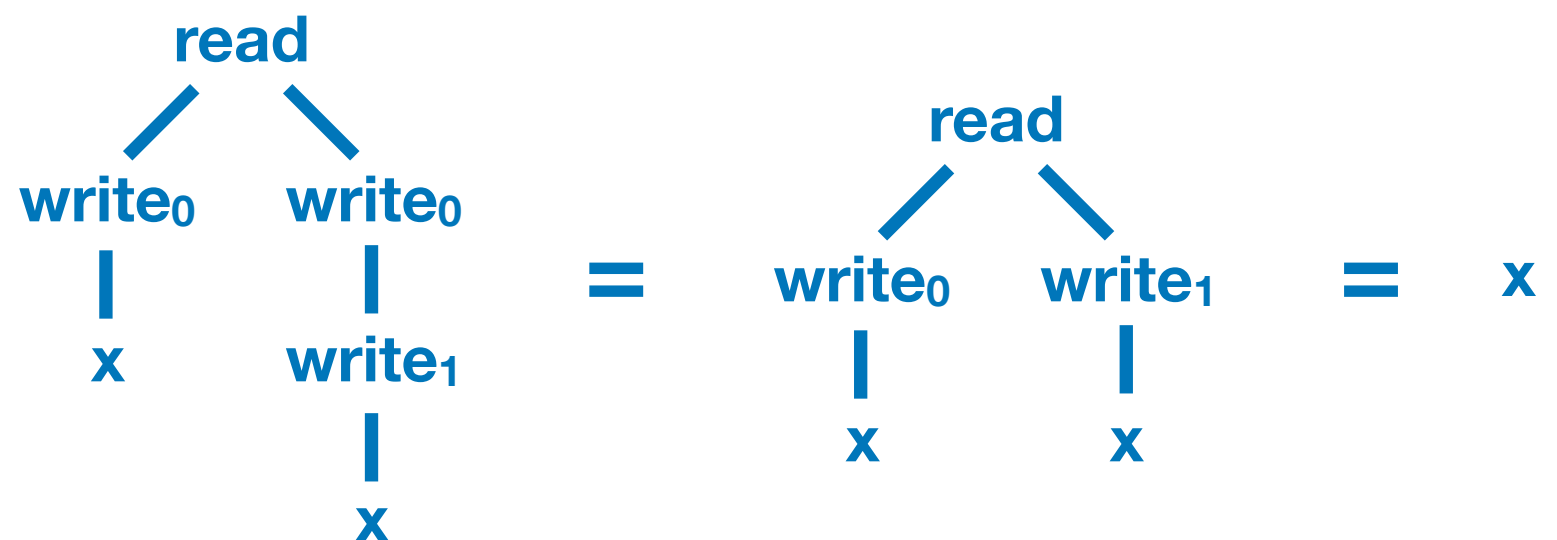
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- **State of the art:** very popular (!) but **effect systems** too coarse grained (!)
 - concurrency, probability, delimited control, monadic reflection, ...
 - Multicore OCaml, Uber's Pyro tool, Eff, Koka, Frank, ...
 - $M : A ! \{ \text{read}, \text{write}, \text{throw} \}$

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- ◎ **Simple idea:** exploit the **underlying tree-like structure** of **effects**!
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$$\{\{1\}\} A \{Q\} = A ! \bigvee_{i \in \{0,1\}, q \in Q} \langle \text{rd} \rangle (\langle \text{wr}_i \rangle (\text{ret}), \langle \text{wr}_q \rangle (\text{ret}))$$

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- ◎ **Challenges:**

- lifting non-linear effect equations (read $x \ x = x$)
- operations with value params. and variable binding
- effect instances, generativity, and locality (my current focus in LJ)
- dynamic nature of handlers

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 - partitioning, spatial layout, ...

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- (multi-)**handlers** based concurrency
- Scala's promises and futures as a **monotonic state effect**

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- ◎ Concurrency
 - (multi-)**handlers** based concurrency
 - Scala's promises and futures as a **monotonic state effect**
- ◎ Probabilistic programming
 - sample as an **algebraic effect**; condition as a **handler** (cf Pyro)

Temporal view

◎ Year 1

- modal logic (design, model and proof theory)
- instances, generativity, locality
- case studies and applications

◎ Year 2

- declarative PL design (type-and-effect system)
- meta-theory (denotational and operational)
- case studies and applications

◎ Year 3

- algorithmic PL design (type-and-effect inference)
- implementation
- case studies and applications

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◎ Year 3

- algorithmic
- implementation
- case studies

◎ In parallel

- continue collaborations with the F* team
- continue collaborations on container datatypes
- forge new collaborations in Oxford (and elsewhere)

Conclusions

- Software is everywhere!
- We had better know what it does!
- General and uniform frameworks already exist for values!
- But only scattered, effect-specific frameworks for behaviour!
- My research will seek to rectify this situation
 - a uniform and widely applicable approach
 - inspired by algebraic effects and effect handlers
 - both foundational theory and exciting applications