

Research vision

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Programming Languages

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let r = alloc 0 in r := !r + 1; r
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Typed Programming Languages

```
let r = alloc 0 in r := !r + 1; r : ref Nat
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lightweight and modular

specification

verification

documentation

correct by construction

In today's world
software is
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Typed Programming Languages

But what about behaviour?

```
let r = alloc 0 in r := !r + 1; r : ref Nat
```

r is fresh

$!r > 0$

other
effects
like I/O ?

user-
defined
effects ?

...

State of affairs in type-based reasoning

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Values

Well-understood, uniform,
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- Refinement types

$\text{Odd} \sqsubseteq \text{Nat}$ $\text{Even} \sqsubseteq \text{Nat}$

$\text{Vec } A \ n = \{ \ell : \text{List } A \mid \text{len } \ell = n \}$

- Dependent types

$\text{Vec } a \ h =$

| $\text{nil} : \text{Vec } a \ 0$

| $\text{cons} : \dots \rightarrow \text{Vec } a \ (n+1)$

- Agda, Coq, F*, Idris, L.Haskell, ...

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Effects and behaviour

Scattered landscape, effect-
specific, little uniformity! :(

● Hoare Type Theory (state)

$M : \Psi.X.\{P\}x:A\{Q\}$

● F* (state, exceptions, but no I/O)

$M : \text{ST } A \ \text{wp}_{\text{ST}}$

~~$M : \text{IO } A \ \text{wp}_{\text{IO}}$~~

● Session Types (I/O & channels)

$c : ?\text{Nat}.!\text{String}.!\text{Nat}.T$

● Graded monads, param. monads

My vision: no need for this non-uniformity!

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- © **Goal:** a general, uniform framework for reasoning about **effects**
 - wide range of effects (state, I/O, exceptions, probability, ...)
 - primitive and user-defined effects
 - combinations of effects

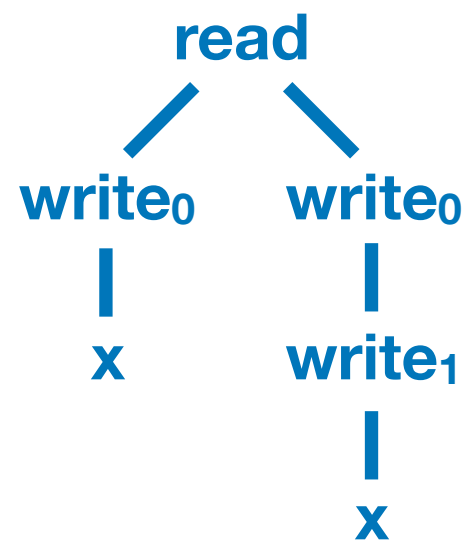
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- **Answer:** **algebraic effects** and **effect handlers** (rather than just monads)
 - operations and equations
 - reveal the **fundamental underlying tree-like structure** of effects
 - **effect handlers** are algebras; **handling** is homomorphism application

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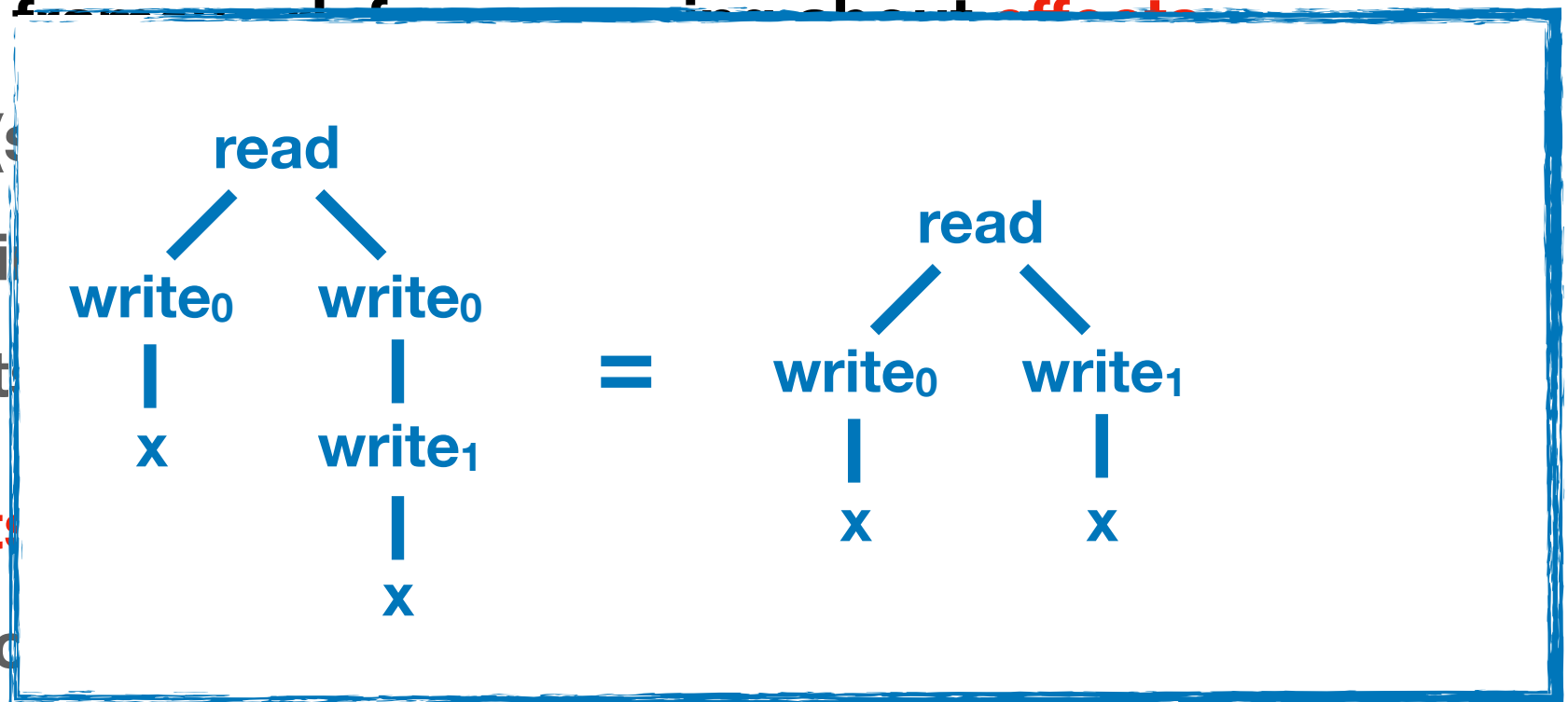
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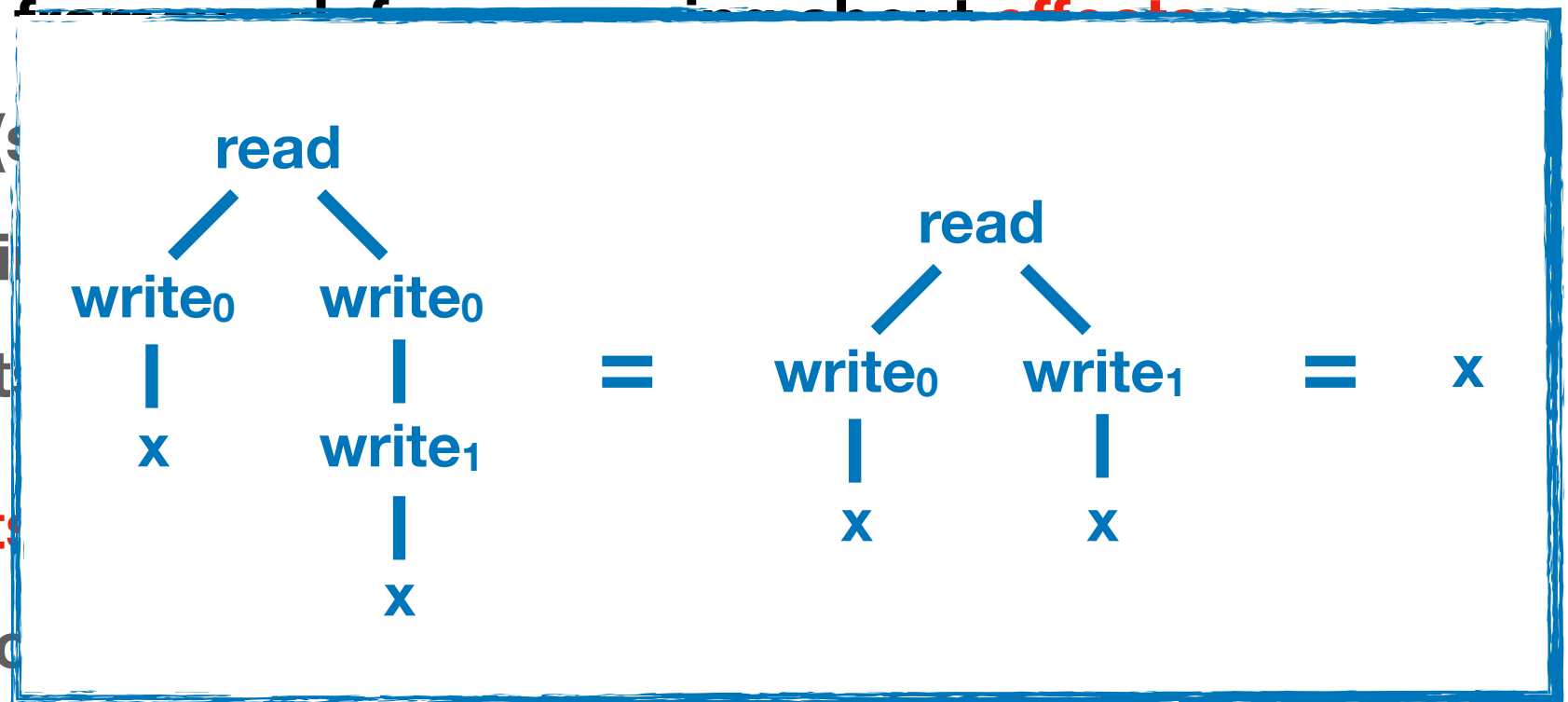
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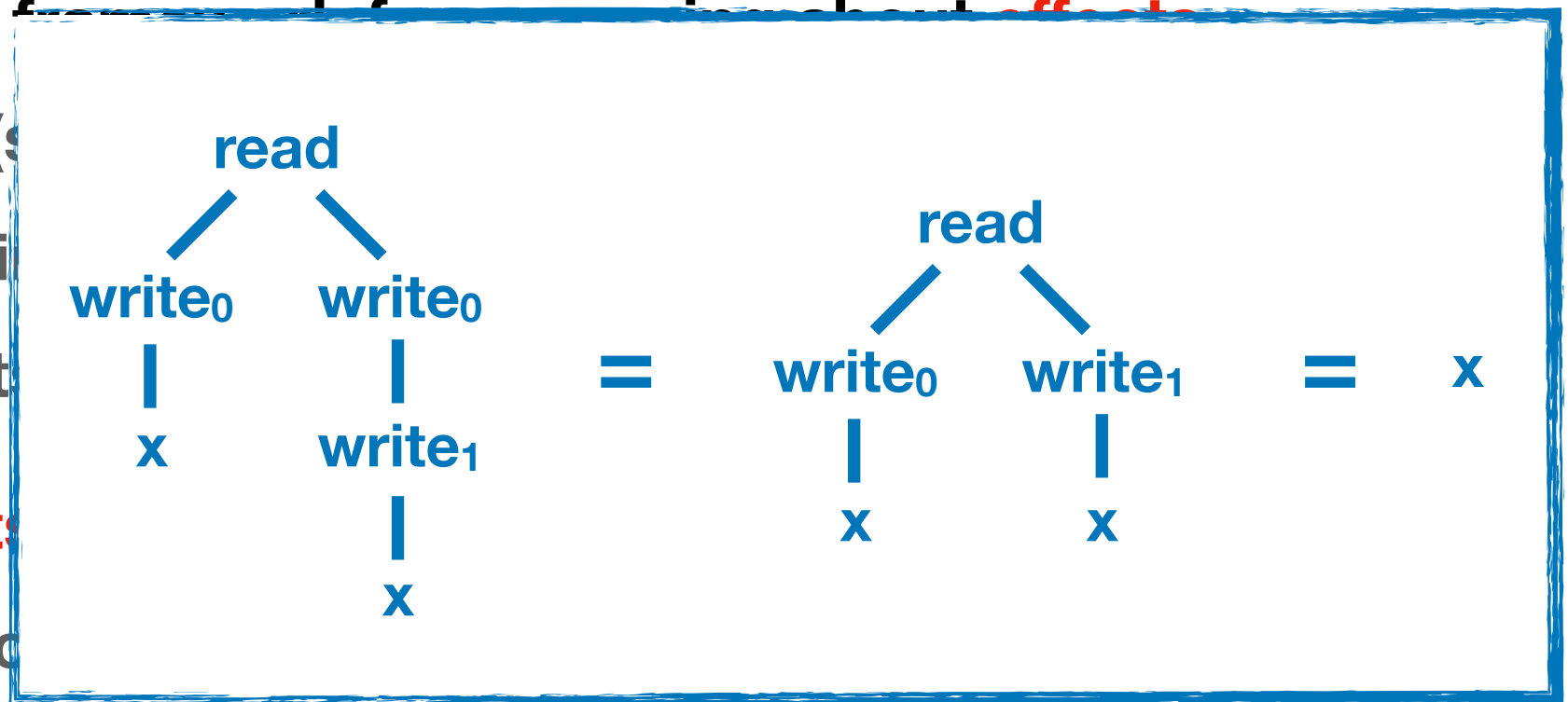
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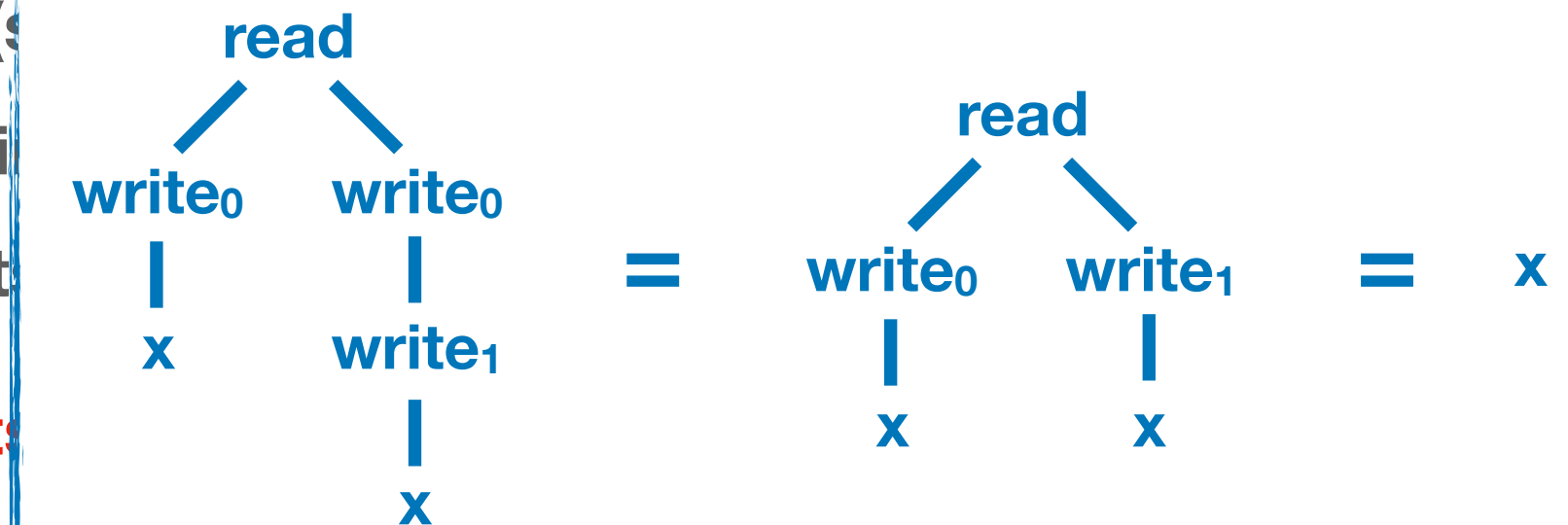
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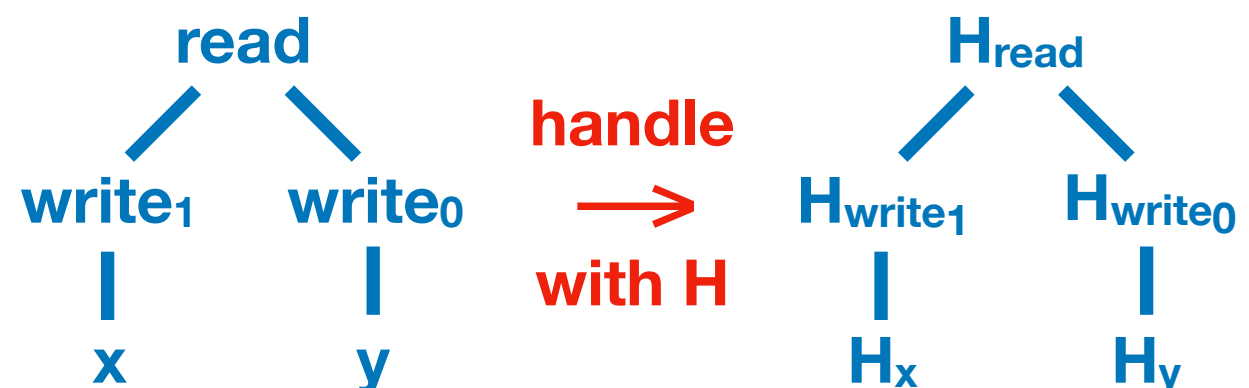
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- **State of the art:** very popular (!) but **effect systems** too coarse grained (!)
 - concurrency, probability, delimited control, monadic reflection, ...
 - Multicore OCaml, Uber's Pyro tool, Eff, Koka, Frank, ...
 - $M : A ! \{ \text{read}, \text{write}, \text{throw} \}$

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- ◎ **Simple idea:** exploit the **underlying tree-like structure** of **effects**!
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$$\{\{1\}\} A \{Q\} = A ! \bigvee_{i \in \{0,1\}, q \in Q} \langle \text{rd} \rangle (\langle \text{wr}_i \rangle (\text{ret}), \langle \text{wr}_q \rangle (\text{ret}))$$

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- ◎ **Challenges:**
 - non-linear effect equations
 - operations with value params. and variable binding
 - effect instances, generativity, and locality (my current focus in LJ)
 - dynamic nature of handlers

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 - commutative monoid structure (**an algebraic effect**)
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- ◎ Probabilistic programming
 - sample as an **algebraic effect**; condition as a **handler** (cf Pyro)

Temporal planning

◎ Year 1

- modal logic (design, model and proof theory)
- instances, generativity, locality

◎ Year 2

- declarative PL design (type-and-effect system)
- meta-theory (denotational and operational)
- encodings of existing specification styles

◎ Year 3

- algorithmic PL design (type-and-effect inference)
- implementation
- case studies and applications

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- algorithmic
- implementation
- case studies

◎ In parallel

- continue collaborations with the F* team
- continue collaborations on container datatypes
- forge new collaborations in Oxford (and elsewhere)

Conclusions

- Software is everywhere!
- We had better know what it does!
- General and uniform frameworks already exist for values!
- But only scattered, effect-specific frameworks for behaviour!
- My research will seek to rectify this situation
 - inspired by algebraic effects and effect handlers
 - exploiting modalities in computational refinement types
 - both foundational theory and exciting applications