

# Research vision

Danel Ahman

University of Ljubljana

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# Programming Languages

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let r = alloc 0 in r := !r + 1; r
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# Typed Programming Languages

```
let r = alloc 0 in r := !r + 1; r : ref Nat
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lightweight and modular

specification

verification

documentation

correct by construction

In today's world  
software is  
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# Typed Programming Languages

But what about behaviour?

```
let r = alloc 0 in r := !r + 1; r : ref Nat
```

r is fresh

$!r > 0$

other  
effects  
like I/O ?

user-  
defined  
effects ?

...

# State of affairs in type-based reasoning

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## **Values**

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**Well-understood, uniform,  
and thoroughly studied! :)**

### ● **Refinement types**

`Odd ⊆ Nat      Even ⊆ Nat`

`Vec A n = { l : List A | len l = n }`

### ● **Dependent types**

`Vec a n = Σ l : List A. (len l = n)`

### ● **Homogeneous implementations**

- Agda, Coq, Idris, F\*, L.Haskell, ...



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### ● Refinement types

$\text{Odd} \sqsubseteq \text{Nat}$        $\text{Even} \sqsubseteq \text{Nat}$

$\text{Vec } A \ n = \{ \ell : \text{List } A \mid \text{len } \ell = n \}$

### ● Dependent types

$\text{Vec } a \ n = \sum \ell : \text{List } A. (\text{len } \ell = n)$

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## Effects and behaviour

Scattered landscape, effect-  
specific, little uniformity! :(

### ● Hoare Type Theory (state)

$M : \Psi.X.\{P\}x:A\{Q\}$

### ● F\* (state, exceptions, but no I/O)

$M : \text{ST } A \ \text{wp}_{\text{ST}}$

~~$M : \text{IO } A \ \text{wp}_{\text{IO}}$~~

### ● Session Types (I/O & channels)

$c : ?\text{Nat}.!\text{String}.!\text{Nat}.T$

### ● Graded monads, param. monads

**My vision: no need for this non-uniformity!**

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- © **Goal:** a general, uniform framework for reasoning about **effects**
  - wide range of effects (state, I/O, exceptions, probability, ...)
  - primitive and user-defined effects
  - combinations of effects

# My vision: **no need for this non-uniformity!**

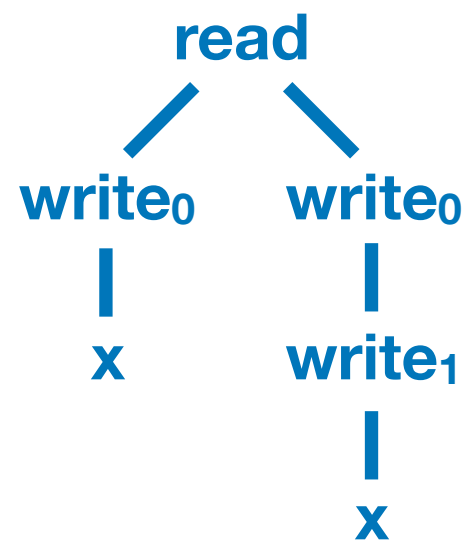
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  - operations and equations
  - reveal the **fundamental underlying tree-like structure** of effects
  - **effect handlers** are algebras; **handling** is homomorphism application

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## Goal: a general, uniform framework for reasoning about effects

- wide range of effects (stateful computations)
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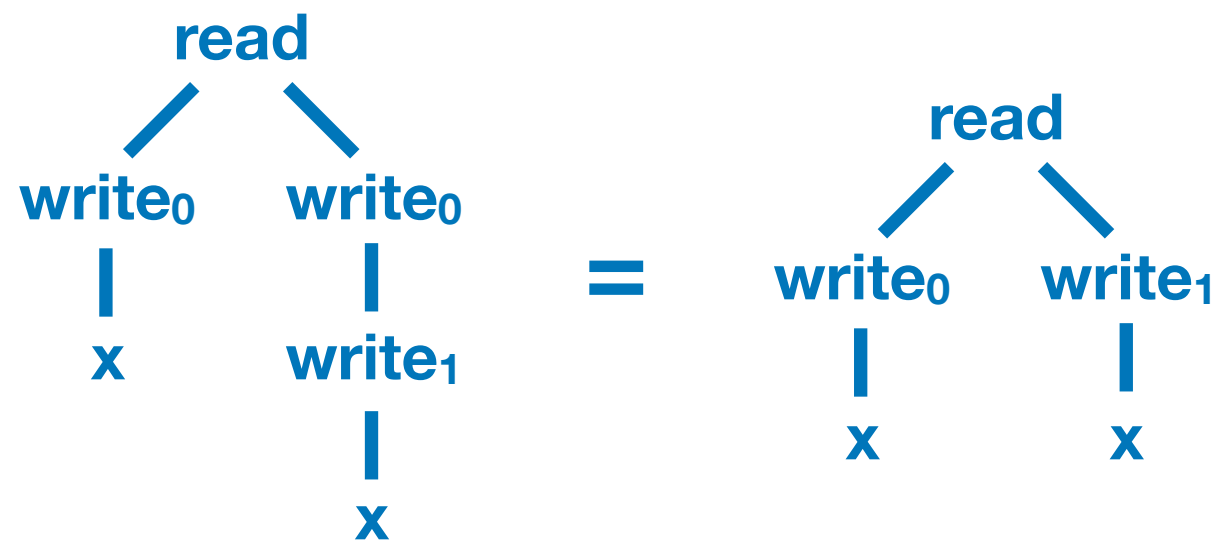
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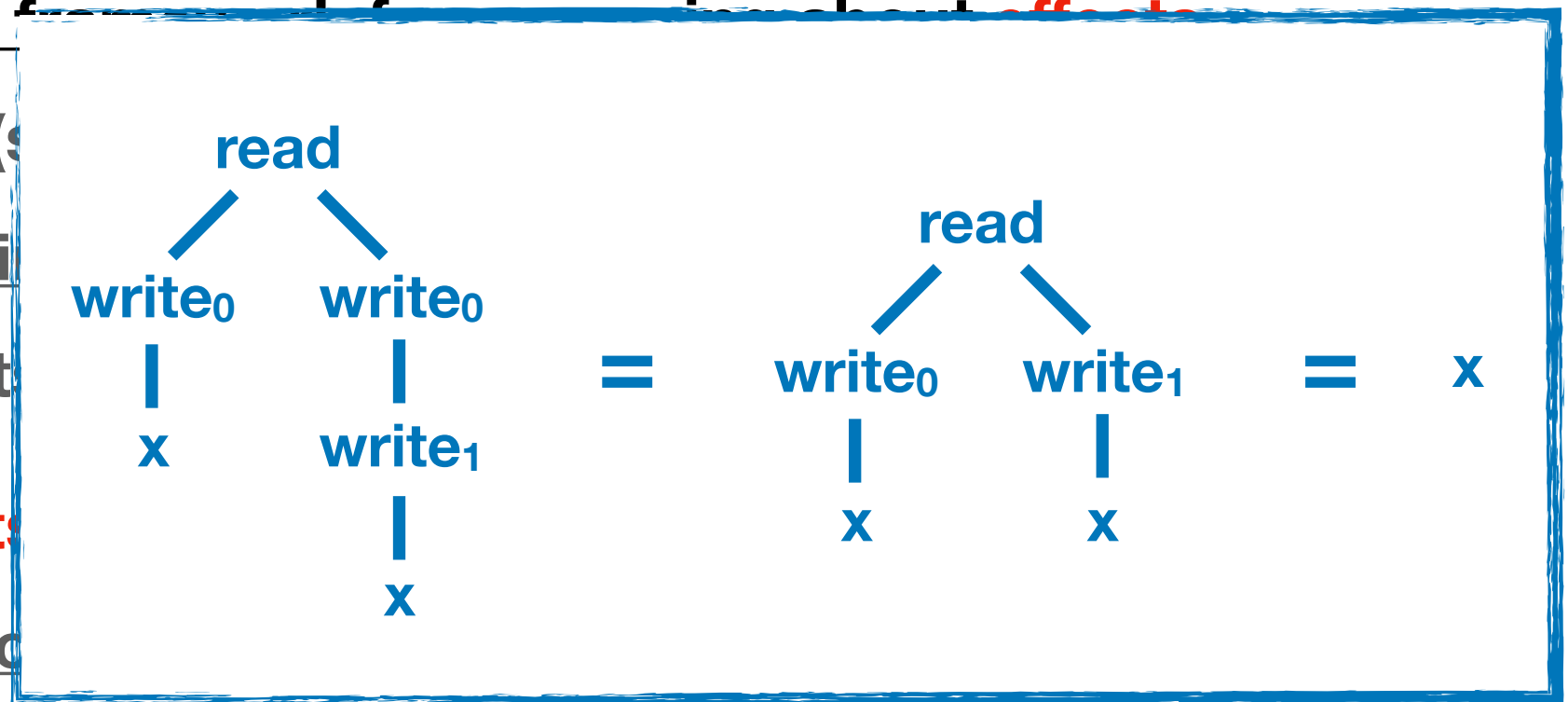
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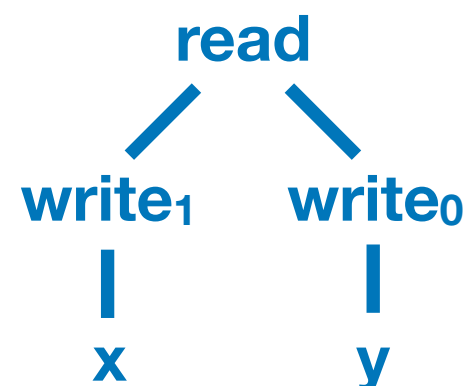
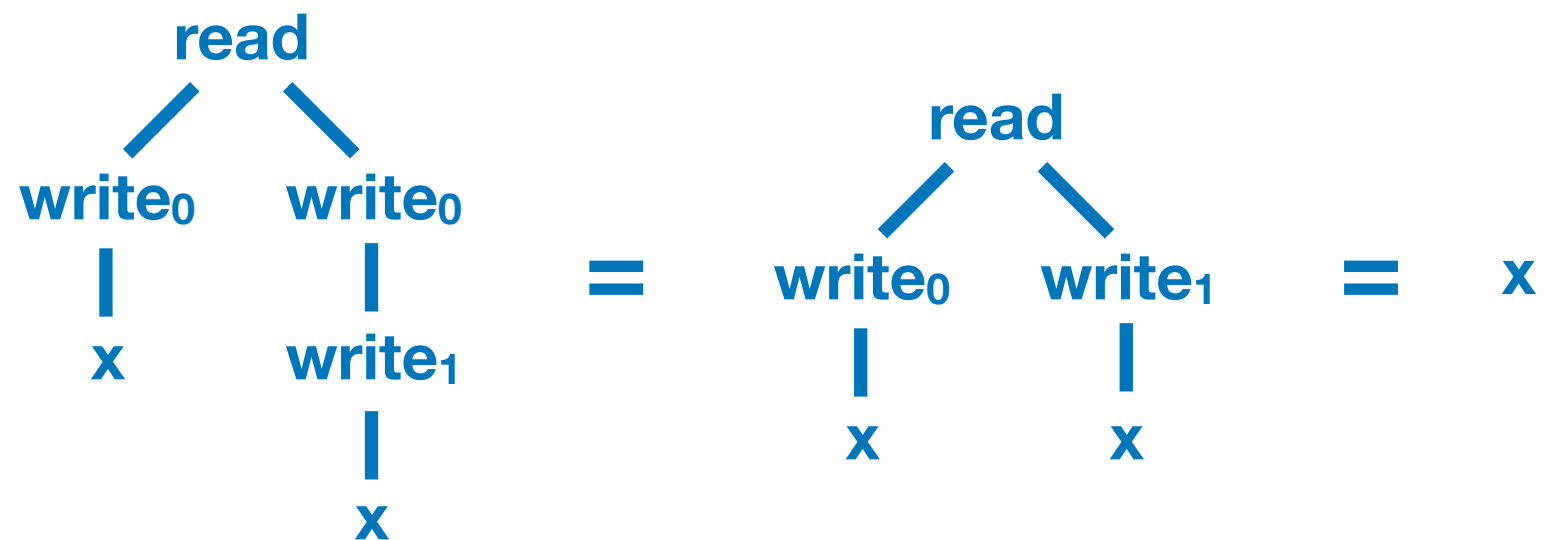
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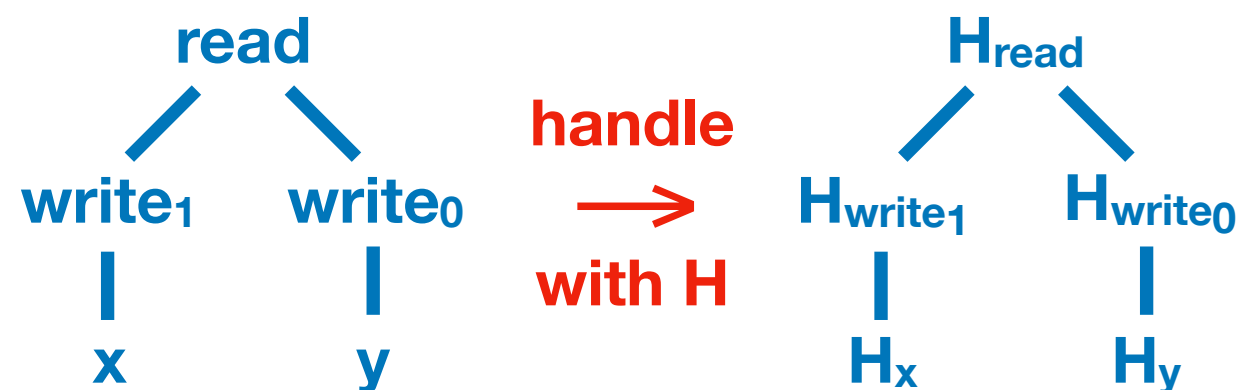
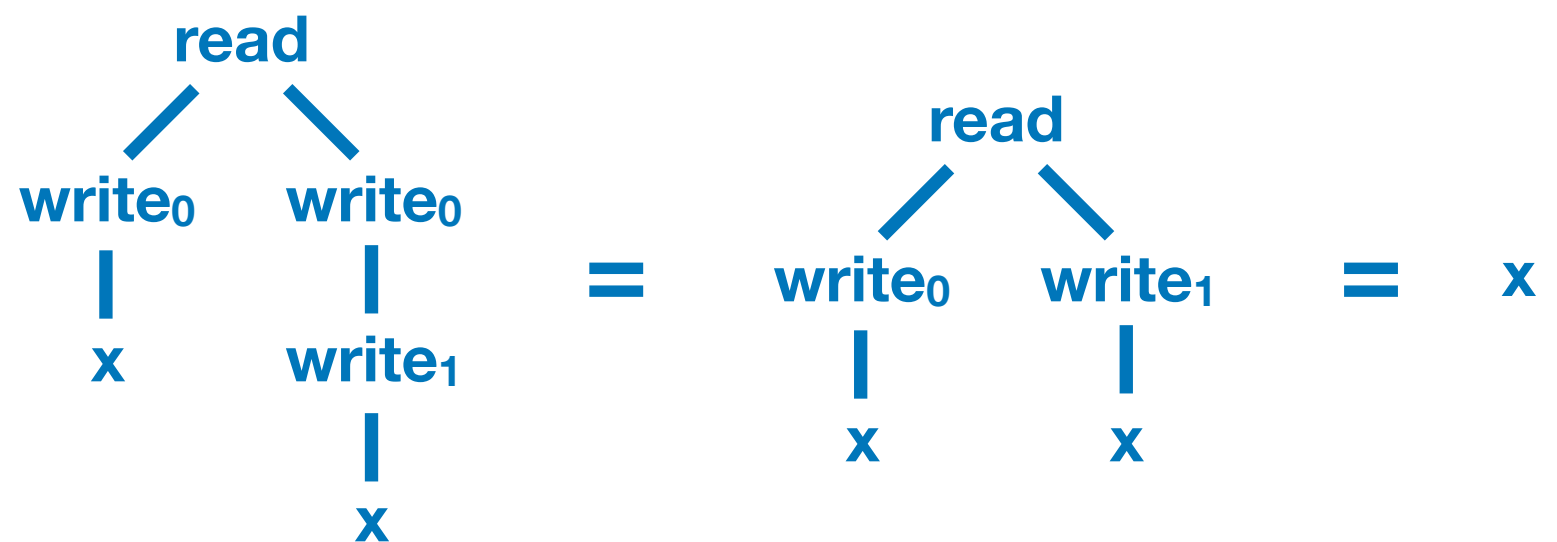
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- **State of the art:** very popular (!) but **effect systems** too coarse grained (!)
  - concurrency, probability, delimited control, monadic reflection, ...
  - Multicore OCaml, Uber's Pyro tool, Eff, Koka, Frank, ...
  - $M : A ! \{ \text{read}, \text{write}, \text{throw} \}$

The plan: modal logic based c. ref. types

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- ◎ **Simple idea:** exploit the **underlying tree-like structure** of **effects**!
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$$\{\{1\}\} A \{Q\} = A ! \bigvee_{i \in \{0,1\}, q \in Q} \langle \text{rd} \rangle (\langle \text{wr}_i \rangle (\text{ret}), \langle \text{wr}_q \rangle (\text{ret}))$$

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- ◎ **Challenges:**

- non-linear effect equations (read  $x \ x = x$ )
- operations with value params. and variable binding
- effect instances, generativity, and locality (my current focus in LJ)
- dynamic nature of handlers

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  - partitioning, spatial layout, ...

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- ◎ Concurrency
  - (multi-)**handlers** based concurrency
  - Scala's promises and futures as a **monotonic state effect**
- ◎ Probabilistic programming
  - sample as an **algebraic effect**; condition as a **handler** (cf Pyro)



# Temporal view

## ◎ Year 1

- modal logic (design, model and proof theory)
- instances, generativity, locality
- case studies and applications

## ◎ Year 2

- declarative PL design (type-and-effect system)
- meta-theory (denotational and operational)
- case studies and applications

## ◎ Year 3

- algorithmic PL design (type-and-effect inference)
- implementation
- case studies and applications

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- algorithmic
- implementation
- case studies

## ◎ In parallel

- continue collaborations with the F\* team
- continue collaborations on container datatypes
- forge new collaborations in Oxford (and elsewhere)

# Conclusions

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- Software is everywhere!
- We had better know what it does!
- General and uniform frameworks already exist for values!
- But only scattered, effect-specific frameworks for behaviour!
- My research will seek to rectify this situation
  - a uniform and widely applicable approach
  - inspired by algebraic effects and effect handlers
  - both foundational theory and exciting applications