

Research vision

Danel Ahman

University of Ljubljana

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Programming Languages

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let r = alloc 0 in r := r + 1; r
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Typed Programming Languages

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let r = alloc 0 in r := r + 1; r : ref Nat
```

lightweight and modular

specification

verification

documentation

correct by construction

In today's world
software is
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Typed Programming Languages

But what about behaviour?

```
let r = alloc 0 in r := r + 1; r : ref Nat
```

r is fresh ?

!r > 0 ?

other
effects
like I/O ?

user-
defined
effects ?

...

State of affairs in type-based reasoning

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Values

Well-understood, uniform,
and thoroughly studied! :)

- Refinement types

`Odd \sqsubseteq Nat` `Even \sqsubseteq Nat`

`Vec A n = {l : List A | len l = n}`

- Dependent types

`Vec a h =`

`| nil : Vec a 0`

`| cons : ... -> Vec a (n+1)`

- Agda, Coq, F*, Idris, L.Haskell, ...

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$\text{Vec } A \ n = \{l : \text{List } A \mid \text{len } l = n\}$

● Dependent types

$\text{Vec } a \ h =$

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Effects and behaviour

Scattered landscape, effect-
specific, little uniformity! :(

● Hoare Type Theory (state)

$M : \Psi.X.\{P\}x:A\{Q\}$

● F* (state, exceptions, but no I/O)

$M : \text{ST } A \ \text{wp}_{\text{ST}}$

~~$M : \text{IO } A \ \text{wp}_{\text{IO}}$~~

● Session Types (I/O & channels)

$c : ?\text{Nat}.\text{!String}.\text{!Nat}.T$

● Graded monads, param. monads

Research vision: let's be uniform!

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- © **Goal:** a general, uniform framework for reasoning about **effects**
 - wide range of effects (state, I/O, exceptions, probability, ...)
 - primitive and user-defined effects
 - combinations of effects

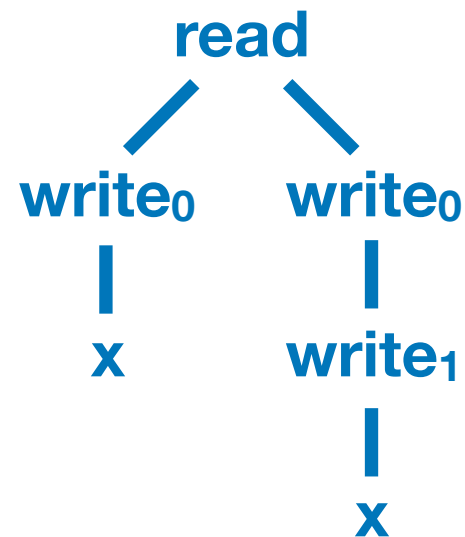
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- ◎ **Answer:** **algebraic effects** and **effect handlers**
 - given by operations and equations
 - reveal the **fundamental underlying tree-like structure**
 - handlers are homomorphic tree transformers

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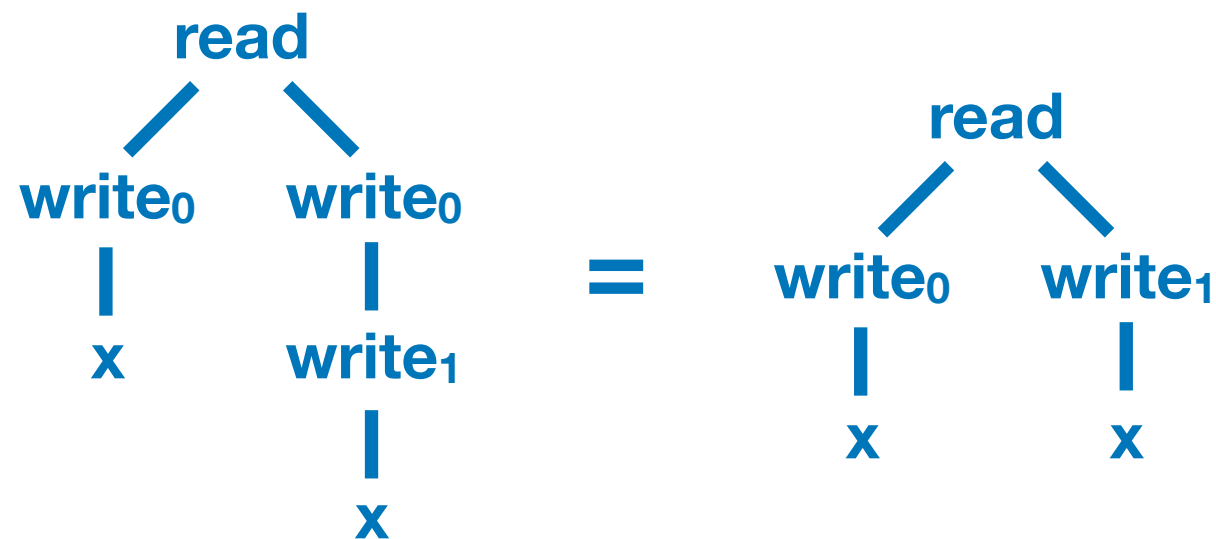
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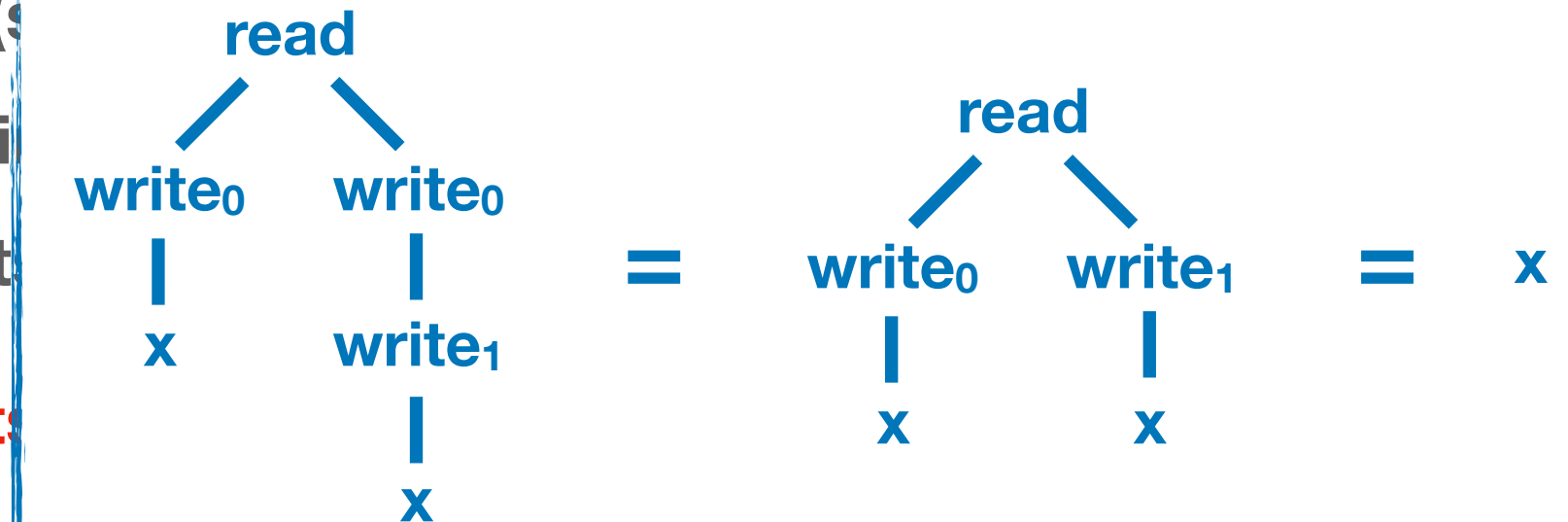
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- **State of the art:** very popular (!) but **effect systems** too coarse grained (!)
 - concurrency, delimited control, monadic reflection, ...
 - Multicore OCaml , Uber's Pyro tool , Eff , Koka , Frank , ...
 - $M : A ! \{ \text{read} , \text{write} , \text{throw} \}$

Specifics: modal logic based c. ref. types

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- ◎ **Simple idea:** exploit the **underlying tree-like structure** of **effects**!
 - comp. ref. types should denote sets of comp. trees
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$$\langle \text{op} \rangle(\psi_1, \dots, \psi_n) = \left\{ \begin{array}{c} \text{op} \\ \swarrow \quad \searrow \\ t_1 \quad \dots \quad t_n \end{array} \mid t_1 \in \psi_1 \wedge \dots \wedge t_n \in \psi_n \right\}$$

$$M : A ! \psi$$

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- uniform across **all algebraic effects**
- can encode Hoare Logic, Session Types, $\text{HL} \otimes \text{ST}$, ...

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- ◎ **Challenges:**

- operations taking arguments and binding values
- lifting non-linear effect equations to such refinement types
- effect instances, generativity, and locality (my current focus in LJ)

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- ◎ **Probabilistic programming**
 - algebraic effects and handlers as in Uber's Pyro tool

Temporal planning

◎ Year 1

- modal logic (design, model and proof theory)
- instances, generativity, locality

◎ Year 2

- declarative PL design (type-and-effect system)
- meta-theory (denotational and operational)
- encodings of existing specification styles

◎ Year 3

- algorithmic PL design (type-and-effect inference)
- implementation
- case studies and applications

Conclusions



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- **Software is everywhere!**
- **We had better know what it does!**
- **General and uniform frameworks already exist for values!**
- **But only scattered, effect-specific frameworks for behaviour!**
- **My research will seek to rectify this situation**
 - inspired by algebraic effects and effect handlers
 - exploiting modalities in computational refinement types
 - with a wide range of potential application areas