Interacting with the external world using comodels (aka runners)

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The plan

- **Computational effects** and **external resources** in PL
- **Runners** – a natural model for **top-level runtime**
- **T-runners** – for also modelling **non-top-level runtimes**
- Turning **T**-runners into a **useful programming construct**
- Some **programming examples**
- Some **implementation details**

Computational effects

and

external resources

Computational effects in PL

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- Using **monads** (as in Haskell)

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type St a = String → (a,String)

f :: St a → St (a,a)
f c = c  >>=  (\x → c  >>=  (\y → return (x,y)))
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```

- Using **alg. effects** and **handlers** (as in Eff, Frank, Koka)

```
effect Get : int
effect Put : int → unit

let g (c:Unit → a!{Get,Put}) =
  with statehandler handle (perform (Put 42); c ())  (* : int → a * int *)
```

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let g (c:Unit → a!{Get,Put}) =
  with statehandler handle (perform (Put 42); c ())  (* : int → a * int *)
```

- Both are good for **faking comp. effects** in a pure language!

  But what about effects that need access to the **external world**?

External resources in PL

External resources in PL

- Declare a **signature of monads** **or** **algebraic effects**, e.g.,

```
(* System.IO *)
type IO a
openFile :: FilePath → IOMode → IO Handle
```

```
(* pervasives.eff *)
effect RandomInt : int → int
effect RandomFloat : float → float
```

- And then **treat them specially** in the compiler, e.g.,

```
(* eff/src/backends/eval.ml *)
let rec tophandle op =
  match op with
  | ...
```

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but there are some issues with that approach ...

First issue

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  - **non-trivial to change** what's available and how it's implemented

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So here's the hack I added. We should do something a bit more principled

In `pervasives.eff`:

```
effect Write : (string*string) -> unit
```

in `eval.ml`, under `let rec top_handle op =` add the case:

```
    | "Write" ->
      (match v with
       | V.Tuple vs ->
          let (file_name :: str :: _)  = List.map V.to_str vs in
          let file_handle = open_out_gen
                              [Open_wronly
                              ;Open_append
                              ;Open_creat
                              ;Open_text
                              ] 0o666 file_name in
          Printf.fprintf file_handle "%s" str;
          close_out file_handle;
          top_handle (k V.unit_value)
      )
```

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This talk — a principled modular (co)algebraic approach!

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- **Lack of linearity** for external resources

```
let f (s:string) =
  let fh = fopen "foo.txt" in
  fwrite (fh,s^s);
  fclose fh;
  return fh

let g s =
  let fh = f s in fread fh
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- We shall address these kinds of issues **indirectly**,
  - by **not** introducing a linear typing discipline
  - but instead make it convenient to **hide** external resources

Third issue

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- **Excessive generality** of effect handlers

```
let f (s:string) =
  let fh = fopen "foo.txt" in
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let h = handler { fwrite (fh,s) k → return () }

let f' s = handle (f  "bar") with h
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where misuse of external resources can also be **purely accidental**

```
let g (s:string) =
  let fh = fopen "foo.txt" in
  let b = choose () in
  if b then (fwrite (fh,s)) else (fwrite (fh,s^s));
  fclose fh

let nondeterminismhandler =
  handler { choose () k → return (k true ++ k false) }
```

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- We shall address these kinds of issues **directly**,
  - by proposing a **restricted form** of handlers for resources
  - that support **controlled initialisation** and **finalisation**,
  - and **limit** how general handlers can be used

Runners enter the spotlight

A natural model of top-level runtime

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- Given a **signature**[1] Σ of operation symbols ($A_{\mathsf{op}}, B_{\mathsf{op}}$ are sets)

$$\mathsf{op} : A_{\mathsf{op}} \rightsquigarrow B_{\mathsf{op}}$$

  a **runner**[2] $\mathcal{R}$ for Σ is given by a carrier $|\mathcal{R}|$ and co-operations

$$\Big(\overline{\mathsf{op}}_{\mathcal{R}} : A_{\mathsf{op}} \times |\mathcal{R}| \longrightarrow B_{\mathsf{op}} \times |\mathcal{R}|\Big)_{\mathsf{op} \in \Sigma}$$

[1] We consider runners for signatures, but the work generalises to alg. theories.
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- For example, a natural runner $\mathcal{R}$ for **S-valued state**

$$\mathsf{get} : \mathbb{1} \rightsquigarrow S \qquad \mathsf{set} : S \rightsquigarrow \mathbb{1}$$

  is given by

$$|\mathcal{R}| \stackrel{\text{def}}{=} S \qquad \overline{\mathsf{get}}_{\mathcal{R}}\,(\star, s) \stackrel{\text{def}}{=} (s, s) \qquad \overline{\mathsf{set}}_{\mathcal{R}}\,(s, s) \stackrel{\text{def}}{=} (\star, s)$$

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A natural model of top-level runtime ctd.

- Runners/comodels have been used for
  - **operational semantics** using tensors of models and comodels
    [Plotkin and Power '08]

    and
  - **stateful running** of algebraic effects [Uustalu '15]
  - **linear-use state-passing translation**
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  - **stateful running** of algebraic effects [Uustalu '15]
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- The latter explicitly rely on one-to-one correspondence between
  - **runners** $\mathcal{R}$
  - **monad morphisms**[3] $\mathsf{r} : \mathbf{Free}_\Sigma(-) \longrightarrow \mathbf{St}_{|\mathcal{R}|}$

  where

$$\mathbf{St}_C\, X \overset{\text{def}}{=} C \Rightarrow X \times C$$

[3] $\mathbf{Free}_\Sigma(X)$ is the free monad ind. defined with leaves $\mathsf{val}\, x$ and nodes $\mathsf{op}(a, \kappa)$.

A natural model of top-level runtime ctd.

- For our purposes, we see runners

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  as describing how operations affect **runtime configurations** $|\mathcal{R}|$

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- But what if this runtime is not **the** runtime?
  - hardware vs OS
  - OS vs VMs
  - VMs vs sandboxes

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  - use **external resources**
  - **signal failure** caused by unavoidable circumstances

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- But is there a **useful generalisation** that would achieve this?

Effectful runners for modular top-levels

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- Møgelberg and Staton usefully observed that a **runner** $\mathcal{R}$ is equivalently simply a family of **generic effects** for $\mathbf{St}_{|\mathcal{R}|}$, i.e.,

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- The one-to-one correspondence with **monad morphisms**

$$\mathsf{r} : \mathbf{Free}_{\Sigma}(-) \longrightarrow \mathbf{T}$$

now simply amounts to the **univ. property of free models**, e.g.,

$$\mathsf{r}_X\,(\mathsf{val}\, x) = \eta_X\, x \qquad\qquad \mathsf{r}_X\,(\mathsf{op}(a, \kappa)) = (\mathsf{r}_X \circ \kappa)^{\dagger}(\overline{\mathsf{op}}_{\mathcal{R}}\, a)$$

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- Observe that κ appears in a **tail call position** on the right!

Effectful runners for modular top-levels ctd.

- What would be a **useful class of monads** $\mathbf{T}$ to use?

Effectful runners **for** **modular top-levels** **ctd.**

- What would be a **useful class of monads** $\mathbf{T}$ to use?
- We want a runner to be a bit like a **kernel** of an OS, i.e., to
  - **(i)** provide management of **(internal) resources**
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- **Algebraically** (and pragmatically), this amounts to taking
  - **(i)** $\mathsf{getenv} : \mathbb{1} \rightsquigarrow C$, $\mathsf{setenv} : C \rightsquigarrow \mathbb{1}$
  - **(ii)** $\mathsf{op} : A_{\mathsf{op}} \rightsquigarrow B_{\mathsf{op}}$ ($\mathsf{op} \in \Sigma'$, for some external Σ')
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  s.t., **(i)** satisfy state equations; and **(i)** commute with **(ii)** and **(iii)**

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- The **induced monad** is then isomorphic to

$$\mathbf{T}\,X \quad \stackrel{\text{def}}{=} \quad C \Rightarrow \mathbf{Free}_{\Sigma'}\big((X \times C) + S\big)$$

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- The corresponding **T-runners** $\mathcal{R}$ for Σ are then of the form

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- With this, our **T-runners** $\mathcal{R}$ for Σ are (with "primitive" excs.)

$$\left(\overline{\mathsf{op}}_{\mathcal{R}} : A_{\mathsf{op}} \longrightarrow \mathbf{K}_C^{\Sigma' ! E_{\mathsf{op}} \lightning S}\, B_{\mathsf{op}}\right)_{\mathsf{op} \in \Sigma}$$

  where we call $\mathbf{K}_C^{\Sigma ! E \lightning S}$ a **kernel monad**, given by

$$\mathbf{K}_C^{\Sigma ! E \lightning S}\, X \quad \stackrel{\text{def}}{=} \quad C \Rightarrow \mathbf{Free}_{\Sigma}\big(((X + E) \times C) + S\big)$$

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  we can easily accommodate co-operations as **kernel code**

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let R = runner {  op₁ x₁ → K₁  ,  ...  ,  opₙ xₙ → Kₙ  } @ C
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let R = runner { op₁ x₁ → K₁ , ... , opₙ xₙ → Kₙ } @ C
```

- For instance, we can implement a **write-only file handle** as

```
let R_FH = runner {
  write s → if (length s > max)
              then (raise WriteSizeExceeded)
              else (let fh = getenv () in
                      if (isValid fh) then (fwrite (fh,s)) else (kill IOError))
} @ FileHandle
```

where
$$\left(\mathrm{fwrite} : \mathrm{FileHandle} \times \mathrm{String} \rightsquigarrow 1 + E\right) \in \Sigma'$$
$$\Sigma \stackrel{\mathrm{def}}{=} \{\ \mathrm{write} : \mathrm{String} \rightsquigarrow 1 + E \cup \{\mathrm{WriteSizeExceeded}\}\ \} \qquad \mathrm{IOError} \in S$$

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- Recall that the components r_X of the monad morphism

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using R **@** $\mathsf{M}_{\mathsf{init}}$
run M
finally $\{$**return** x **@** $\mathsf{c} \to \mathsf{M}_{\mathsf{ret}}$, ... **raise** e **@** $\mathsf{c} \to \mathsf{M}_{\mathsf{e}}$ ... , ... **kill** $\mathsf{s} \to \mathsf{M}_{\mathsf{s}}$ ...$\}$

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where

(a **user monad**)

- Ms are **user code**, modelled using $\mathbf{U}^{\Sigma!E}\,X \stackrel{\text{def}}{=} \mathbf{Free}_\Sigma(X + E)$

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- $\mathsf{M}_{\mathsf{init}}$ produces the **initial kernel state**
- M is the user code being **run using the runner** R
- $\mathsf{M}_{\mathsf{ret}}$, M_{e}, M_{s} **finalise** for return values, exceptions, and signals

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  - $\mathsf{M}_{\mathsf{ret}}$, M_{e}, M_{s} **finalise** for return values, exceptions, and signals
  - $\mathsf{M}_{\mathsf{ret}}$ and M_{e} **depend on the final state** c, but M_{s} **does not**

Controlled **initialisation** and **finalisation** ctd.

- For instance, we can define a PYTHON-esque **with construct**

```
with fileName do M
=
using R_FH @ (fopen fileName)
run M
finally {
  return x @ fh → fclose fh; return x ,
  raise  e @ fh → fclose fh; raise e   ,
  kill   s      → return ()   }
```

- Importantly, here
  - the **file handle is hidden** from `M`
  - `M` can only use `write` but not `fopen` and `fclose`
    - write : String $\rightsquigarrow$ $1 + E \cup \{\text{WriteSizeExceeded}\}$
  - `fopen` and `fclose` are limited to **initialisation-finalisation**

A core calculus for programming with runners

Core calculus (syntax)

Core calculus (syntax)

- **Ground types** (types of ops. and kernel state)

$$A, B, C \quad ::= \quad \mathsf{B} \mid 1 \mid 0 \mid A \times B \mid A + B$$

- **Types**

$$\begin{array}{rcl} X, Y & ::= & \mathsf{B} \mid 1 \mid 0 \mid X \times Y \mid X + Y \\ & \mid & X \xrightarrow{\Sigma} Y \mathrel{!} E \\ & \mid & X \xrightarrow{\Sigma} Y \mathrel{!} E \lightning S \mathbin{@} C \\ & \mid & \Sigma \Rightarrow \Sigma' \lightning S \mathbin{@} C \end{array}$$

- **Values**

$$\Gamma \vdash V : X$$

- **User computations**

$$\Gamma \mathrel{\vdash^{\Sigma}} M : X \mathrel{!} E$$

- **Kernel computations**

$$\Gamma \mathrel{\vdash^{\Sigma}} K : X \mathrel{!} E \lightning S \mathbin{@} C$$

$$
\begin{aligned}
M ::= {} & \mathbf{return}\, V \;\mid\; \mathbf{try}\; M \;\mathbf{with}\; \{\; \mathbf{return}\; x \mapsto N_{val} \;,\; \big(\mathbf{raise}\; e \mapsto N_e\big)_{e \in E} \;\} \\
\mid {} & V\, W \;\mid\; \mathbf{match}\; V \;\mathbf{with}\; \{\; \langle x_1, x_2 \rangle \mapsto N \;\} \\
\mid {} & \mathbf{match}\; V \;\mathbf{with}\; \{\}_X \;\mid\; \mathbf{match}\; V \;\mathbf{with}\; \{\; \mathbf{inl}\; x_1 \mapsto N_1 \;,\; \mathbf{inr}\; x_2 \mapsto N_2 \;\} \\
\mid {} & \mathsf{op}_X\, V\, (x.M)\, (N_e)_{e \in E_{\mathsf{op}}} \;\mid\; \mathbf{raise}_X\, e \\
\mid {} & \mathbf{using}\; V \;@\; W \;\mathbf{run}\; M \;\mathbf{finally}\; \{\; \mathbf{return}\; x \;@\; c \mapsto N_{val} \;, \\
& \qquad \big(\mathbf{raise}\; e \;@\; c \mapsto N_e\big)_{e \in E} \;, \\
& \qquad \big(\mathbf{kill}\; s \mapsto N_s\big)_{s \in S} \;\} \\
\mid {} & \mathbf{exec}\; K \;@\; W \;\mathbf{finally}\; \{\; \mathbf{return}\; x \;@\; c \mapsto N_{val} \;, \\
& \qquad \big(\mathbf{raise}\; e \;@\; c \mapsto N_e\big)_{e \in E} \;, \\
& \qquad \big(\mathbf{kill}\; s \mapsto N_s\big)_{s \in S} \;\}
\end{aligned}
$$

$$
\begin{aligned}
K ::= {} & \mathbf{return}_C\, V \;\mid\; \mathbf{try}\; K \;\mathbf{with}\; \{\; \mathbf{return}\; x \mapsto L_{val} \;,\; \big(\mathbf{raise}\; e \mapsto L_e\big)_{e \in E} \;\} \\
\mid {} & V\, W \;\mid\; \mathbf{match}\; V \;\mathbf{with}\; \{\; \langle x_1, x_2 \rangle \mapsto L \;\} \\
\mid {} & \mathbf{match}\; V \;\mathbf{with}\; \{\}_{X @ C} \;\mid\; \mathbf{match}\; V \;\mathbf{with}\; \{\; \mathbf{inl}\; x_1 \mapsto L_1 \;,\; \mathbf{inr}\; x_2 \mapsto L_2 \;\} \\
\mid {} & \mathsf{op}_{X @ C}\, V\, (x.K)\, (L_e)_{e \in E_{\mathsf{op}}} \;\mid\; \mathbf{raise}_{X @ C}\, e \;\mid\; \mathbf{kill}_{X @ C}\, s \\
\mid {} & \mathbf{getenv}_C\, (c.K) \;\mid\; \mathbf{setenv}\, V\, K \\
\mid {} & \mathbf{exec}\; M \;\mathbf{finally}\; \{\; \mathbf{return}\; x \mapsto L_{val} \;,\; \big(\mathbf{raise}\; e \mapsto L_e\big)_{e \in E} \;\}
\end{aligned}
$$

Fig. 1. Syntax of user and kernel computations

Core calculus (type system and eq. theory)

Core calculus (type system and eq. theory)

- For example, the **typing rule for running user comps.** is

$$
\frac{\begin{array}{c}
\Gamma \vdash V : \Sigma \Rightarrow \Sigma' \lightning S \mathbin{@} C \qquad \Gamma \vdash W : C \\
\Gamma \mathrel{\vdash^{\Sigma}} M : X \mathbin{!} E \qquad \Gamma, x \mathord{:} X, c \mathord{:} C \mathrel{\vdash^{\Sigma'}} N_{ret} : Y \mathbin{!} E' \\
\big(\Gamma, c \mathord{:} C \mathrel{\vdash^{\Sigma'}} N_e : Y \mathbin{!} E'\big)_{e \in E} \qquad \big(\Gamma \mathrel{\vdash^{\Sigma'}} N_s : Y \mathbin{!} E'\big)_{s \in S}
\end{array}}{\begin{array}{l}
\Gamma \mathrel{\vdash^{\Sigma'}} \mathbf{using}\ V \mathbin{@} W\ \mathbf{run}\ M\ \mathbf{finally}\ \{\ \mathbf{return}\ x \mathbin{@} c \mapsto N_{ret}\ , \\
\qquad\qquad\qquad\qquad\qquad\qquad \big(\mathbf{raise}\ e \mathbin{@} c \mapsto N_e\big)_{e \in E}\ , \\
\qquad\qquad\qquad\qquad\qquad\qquad \big(\mathbf{kill}\ s \mapsto N_s\big)_{s \in S}\ \} : Y \mathbin{!} E'
\end{array}}
$$

Core calculus (type system and eq. theory)

- For example, the **typing rule for running user comps.** is

$$\frac{\Gamma \vdash V : \Sigma \Rightarrow \Sigma' \not\downarrow S \mathbin{@} C \qquad \Gamma \vdash W : C \\ \Gamma \models^{\Sigma} M : X \mathrel{!} E \qquad \Gamma, x\!:\!X, c\!:\!C \models^{\Sigma'} N_{ret} : Y \mathrel{!} E' \\ \big(\Gamma, c\!:\!C \models^{\Sigma'} N_e : Y \mathrel{!} E'\big)_{e\in E} \qquad \big(\Gamma \models^{\Sigma'} N_s : Y \mathrel{!} E'\big)_{s\in S}}{\begin{array}{l}\Gamma \models^{\Sigma'} \textbf{using } V \mathbin{@} W \textbf{ run } M \textbf{ finally } \{ \textbf{return } x \mathbin{@} c \mapsto N_{ret}\ , \\ \qquad\qquad\qquad\qquad\qquad\qquad\qquad \big(\textbf{raise } e \mathbin{@} c \mapsto N_e\big)_{e\in E}\ , \\ \qquad\qquad\qquad\qquad\qquad\qquad\qquad \big(\textbf{kill } s \mapsto N_s\big)_{s\in S}\ \} : Y \mathrel{!} E'\end{array}}$$

- and the **main β-equation for running user comps.** is

$$\begin{array}{l}\Gamma \models^{\Sigma'} \textbf{using } R_C \mathbin{@} W \textbf{ run } (\mathsf{op}_X\, V\, (x.M)\, (M_e)_{e\in E_{\mathsf{op}}}) \textbf{ finally } F \\ \quad \equiv \textbf{exec } R_{op}[V] \mathbin{@} W \textbf{ finally } \{ \\ \qquad\qquad \textbf{return } x \mathbin{@} c' \mapsto \textbf{using } R_C \mathbin{@} c' \textbf{ run } M \textbf{ finally } F\ , \\ \qquad\qquad \big(\textbf{raise } e \mathbin{@} c' \mapsto \textbf{using } R_C \mathbin{@} c' \textbf{ run } M_e \textbf{ finally } F\big)_{e\in E_{\mathsf{op}}}\ , \\ \qquad\qquad \big(\textbf{kill } s \mapsto N_s\big)_{s\in S}\ \} : Y \mathrel{!} E'\end{array}$$

Core calculus (type system and eq. theory)

- The calculus also includes **subtyping**, and **subsumption rules**

$$\frac{\Gamma \vdash V : A \qquad A <: B}{\Gamma \vdash V : B}$$

$$\frac{\Gamma \vdash^{\Sigma} M : A \,!\, E \qquad \Sigma \subseteq \Sigma' \qquad A <: B \qquad E \subseteq E'}{\Gamma \vdash^{\Sigma'} M : B \,!\, E'}$$

$$\frac{\begin{array}{c}\Gamma \vdash^{\Sigma} K : A \,!\, E \lightning S \,@\, C \qquad \Sigma \subseteq \Sigma' \\ A <: B \qquad E \subseteq E' \qquad S \subseteq S' \qquad C = C'\end{array}}{\Gamma \vdash^{\Sigma'} K : B \,!\, E' \lightning S' \,@\, C'}$$

Core calculus (type system and eq. theory)

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- $C = C'$ to have (standard) **proof-irrelevant subtyping**
- Otherwise, instead of just $C <: C'$, we would need a **lens** $C' \leftrightarrow C$

Core calculus (semantics)

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- **Monadic semantics**, for simplicity in **Set**, using
  - **user monads** $\mathbf{U}^{\Sigma ! E}\, X \stackrel{\text{def}}{=} \mathbf{Free}_{\Sigma}(X + E)$
  - **kernel monads** $\mathbf{K}_{C}^{\Sigma ! E \lightning S}\, X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}\big(((X + E) \times C) + S\big)$

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  - **kernel monads** $\mathbf{K}_{C}^{\Sigma ! E \lightning S}\, X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}\big(((X + E) \times C) + S\big)$

- (At a high level) the **judgements are interpreted** as

$$[\![\Gamma \vdash V : X]\!] : [\![\Gamma]\!] \longrightarrow [\![X]\!]$$

$$[\![\Gamma \overset{\Sigma}{\vdash} M : X \,!\, E]\!] : [\![\Gamma]\!] \longrightarrow \mathbf{U}^{\Sigma ! E}[\![X]\!]$$

$$[\![\Gamma \overset{\Sigma}{\vdash} K : X \,!\, E \lightning S \mathbin{@} C]\!] : [\![\Gamma]\!] \longrightarrow \mathbf{K}_{[\![C]\!]}^{\Sigma ! E \lightning S}[\![X]\!]$$

Core calculus (semantics ctd.)

- However, to prove **coherence** of the semantics (**subtyping!**), we actually give the semantics in the **subset fibration**

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- For instance, **kernel computations** are interpreted as

$$
\begin{tikzcd}
\llbracket \Gamma \rrbracket \arrow[rrr, "\llbracket \Gamma \vdash^{\Sigma} K \,:\, X \,!\, E \lightning S \,@\, C \rrbracket"] \arrow[d, "\subseteq"'] & & & \mathbf{K}^{\Sigma ! E \lightning S}_{\llbracket C \rrbracket} \llbracket X \rrbracket \arrow[d, "\subseteq"] \\
[\![\![\Gamma^s]\!]\!] \arrow[rrr, "{[\![\![\Gamma^s \vdash K \,:\, X^s \,@\, C]\!]\!]}"'] & & & \mathbf{K}^{\mathcal{O} ! \mathcal{E} \lightning \mathcal{S} + \{\dagger\}}_{[\![\![C]\!]\!]} [\![\![X^s]\!]\!]
\end{tikzcd}
$$

  where $\Gamma^s \vdash K : X^s \;@\; C$ is a **skeletal kernel typing judgement**

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[\![\![\Gamma^s]\!]\!] \arrow[r, "{[\![\![\Gamma^s \vdash K \,:\, X^s \,@\, C]\!]\!]}"'] & \mathbf{K}^{\mathcal{O} ! \mathcal{E} \lightning \mathcal{S} + \{\lightning\}}_{[\![\![C]\!]\!]} [\![\![X^s]\!]\!]
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$$

  where $\Gamma^s \vdash K : X^s \;@\; C$ is a **skeletal kernel typing judgement**

- No essential obstacles to extending to $\mathbf{Sub}(\mathbf{Cpo})$ and beyond

- **Ground type restriction** on C needed to stay within $\mathbf{Sub}(-)$
  - Otherwise, analogously to subtyping, we'd need **lenses** instead

Core calculus (semantics ctd.)

$$
\begin{aligned}
\llbracket \Gamma \vdash^{\Sigma'} \textbf{using } V \mathbin{@} W \textbf{ run } M \textbf{ finally } \{\ & \textbf{return } x \mathbin{@} c \mapsto N_{ret}\ , \\
& \left(\textbf{raise } e \mathbin{@} c \mapsto N_e\right)_{e \in E}\ , \\
& \left(\textbf{kill } s \mapsto N_s\right)_{s \in S}\ \} : Y \ !\ E' \rrbracket_\gamma \ \stackrel{\text{def}}{=}\ \ldots
\end{aligned}
$$

Core calculus (semantics ctd.)

$$\begin{aligned}\llbracket \Gamma \vdash^{\Sigma'} \textbf{using } V \mathbin{@} W \textbf{ run } M \textbf{ finally } \{\, & \textbf{return } x \mathbin{@} c \mapsto N_{ret} \,, \\ & \left(\textbf{raise } e \mathbin{@} c \mapsto N_e\right)_{e \in E} \,, \\ & \left(\textbf{kill } s \mapsto N_s\right)_{s \in S} \,\} : Y \,!\, E' \rrbracket_\gamma \stackrel{\text{def}}{=} \ \ldots\end{aligned}$$

- $\llbracket V \rrbracket_\gamma = \mathcal{R} = \left(\overline{\mathsf{op}}_{\mathcal{R}} : \llbracket A_{\mathsf{op}} \rrbracket \longrightarrow \mathbf{K}^{\Sigma' ! E_{\mathsf{op}} \lightning S}_{\llbracket C \rrbracket} \llbracket B_{\mathsf{op}} \rrbracket\right)_{\mathsf{op} \in \Sigma}$
- $\llbracket W \rrbracket_\gamma \in \llbracket C \rrbracket$
- $\llbracket M \rrbracket_\gamma \in \mathbf{U}^{\Sigma ! E} \llbracket A \rrbracket$
- $\llbracket \textbf{return } \mathsf{x} \mathbin{@} \mathsf{c} \to N_{ret} \rrbracket_\gamma \in \llbracket A \rrbracket \times \llbracket C \rrbracket \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
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Core calculus (semantics ctd.)

$$\llbracket \Gamma \vdash^{\Sigma'} \textbf{using } V \mathbin{@} W \textbf{ run } M \textbf{ finally } \{ \textbf{return } x \mathbin{@} c \mapsto N_{ret} \,, \;\left(\textbf{raise } e \mathbin{@} c \mapsto N_e\right)_{e \in E} \,, \;\left(\textbf{kill } s \mapsto N_s\right)_{s \in S} \} : Y \,!\, E' \rrbracket_\gamma \stackrel{\text{def}}{=} \;\ldots$$

- $\llbracket V \rrbracket_\gamma = \mathcal{R} = \left(\overline{\mathrm{op}}_{\mathcal{R}} : \llbracket A_{\mathrm{op}} \rrbracket \longrightarrow \mathbf{K}^{\Sigma' ! E_{\mathrm{op}} \lightning S}_{\llbracket C \rrbracket} \llbracket B_{\mathrm{op}} \rrbracket\right)_{\mathrm{op} \in \Sigma}$
- $\llbracket W \rrbracket_\gamma \in \llbracket C \rrbracket$
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- $\llbracket (\textbf{kill } \mathsf{s} \to N_s)_{s \in S} \rrbracket_\gamma \in S \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
- allowing us to use the **free model property** to get

$$\mathbf{U}^{\Sigma ! E} \llbracket A \rrbracket \xrightarrow{\mathsf{r}_{\llbracket A \rrbracket + E}} \mathbf{K}^{\Sigma' ! E \lightning S}_{\llbracket C \rrbracket} \llbracket A \rrbracket \xrightarrow{(\lambda \llbracket N_{ret} \rrbracket_\gamma)^{\ddagger}} \llbracket C \rrbracket \Rightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$$

  and then apply the resulting composite to $\llbracket M \rrbracket_\gamma$ and $\llbracket W \rrbracket_\gamma$

Runners in action

Runners can be vertically nested

Runners can be vertically nested

- ```
  using R_FH @ (fopen fileName)
  run (
    using R_FC @ (return "")
    run M
    finally {
      return x @ str → write str; return x ,
      raise   e @ str → write str; raise e   }
  )
  finally {
    return x @ fh → fclose fh; return x ,
    raise   e @ fh → fclose fh; raise e   , kill IOError → return ()}
  ```

  where the **file contents runner** (with $\Sigma' = \mathbb{0}$) is defined as

  ```
  let R_FC = runner {
    write str → let str' = getenv () in
                if (length (str^str') > max) then (raise WriteSizeExceeded)
                                                 else  (setenv (str^str'))
  } @ String
  ```
  ```

Vertical nesting for instrumentation

Vertical nesting for **instrumentation**

- ```
  using R_Sniffer @ (return 0)
  run M
  finally {
    return x @ c →
        let fh = fopen "nsa.txt" in fwrite (fh,nat_to_str c); fclose fh }
  ```

  where the **instrumenting runner** is defined as

  ```
  let R_Sniffer = runner {
    ... ,
    op a → let c = getenv () in
             setenv (c + 1);
             op a ,                        (* forwards op outwards *)
    ...
  } @ Nat
  ```

- The runner $\mathsf{R}_{\mathsf{Sniffer}}$ implements the same sig. Σ that `M` is using
- As a result, the runner $\mathsf{R}_{\mathsf{Sniffer}}$ is **invisible** from `M`'s viewpoint
  ```

Vertical nesting for active monitoring

Vertical nesting for **active monitoring**

- First, we define a runner for **integer-valued ML-style state** as

```
type IntHeap = (Nat → (Int + 1)) × Nat                    type Ref = Nat
```

```
let R_IntState = runner {
  alloc x → let h = getenv () in
            let (r,h') = heapAlloc h x in
            setenv h';
            return r ,

  deref r → let h = getenv () in
            match (heapSel h r) with
            | inl x → return x
            | inr () → kill ReferenceDoesNotExist ,

  assign r y → let h = getenv () in
               match (heapUpd h r y) with
               | inl h' → setenv h'
               | inr () → kill ReferenceDoesNotExist
} @ IntHeap
```

Vertical nesting for **active monitoring** ctd.

- Next we define a runner for **monotonicity layer** on top of $\mathsf{R}_{\mathsf{IntState}}$

Vertical nesting for active monitoring ctd.

- Next we define a runner for **monotonicity layer** on top of $\mathsf{R}_{\mathsf{IntState}}$

```
type MonMemory = Ref → ((Int → Int → Bool) + 1)
```

```
let R_MonState = runner {
  monAlloc x rel → let r = alloc x in
                   let m = getenv () in
                   setenv (memAdd m r rel);
                   return r,

  monDeref r → deref r ,

  monAssign r y → let x = deref r in
                  let m = getenv () in
                  match (memSel m r) with
                  | inl rel → if (rel x y)
                              then (assign r y)
                              else (raise MonotonicityViolation)
                  | inr → kill PreorderDoesNotExist
} @ IntHeap
```

Vertical nesting for **active monitoring** ctd.

- We can then perform **runtime monotonicity verification** as

Vertical nesting for active monitoring ctd.

- We can then perform **runtime monotonicity verification** as

```
using R_IntState @ ((fun _ → inr ()) , 0)          (* empty ML−style heap *)

run (
 using R_MonState @ (fun _ → inr ())          (* empty preorders memory *)

 run (
  let r = monAlloc 0 (⩽) in
  monAssign r 1;
  monAssign r 0;     (* R_MonState raises MonotonicityViolation exception *)
  monAssign r 2 )

 finally {return x @ _ → return x ,
          raise MonotonicityViolation @ _ → ... ,
          kill   PreorderDoesNotExist       → ... })

finally {return x @ _ → return x ,
         kill ReferenceDoesNotExist → ... }
```

Runners can also be horizontally paired

Runners can also be **horizontally paired**

- Given a runner for Σ

```
let R₁ = runner { ... , op₁ᵢ x → k₁ᵢ , ... } @ C₁
```

and a runner for Σ'

```
let R₂ = runner { ... , op₂ⱼ x → k₂ⱼ , ... } @ C₂
```

we can **pair them** to get a runner for $\Sigma \uplus \Sigma'$

```
let R = runner {
  ... ,
  op₁ᵢ x → let (c,c') = getenv () in
           let (x,c'') = k₁ᵢ x in
           setenv (c'',c');
           return x ,
  ... ,
  op₂ⱼ x → ... (* analogously to above *) ,
  ...
} @ C₁ × C₂
```

Runners can also be horizontally paired

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           return x ,
  ... ,
  op₂ⱼ x → ... (* analogously to above *) ,
  ...
} @ C₁ × C₂
```

- For instance, this way we can build a **runner for IO and state**

Other examples

Other examples

- More general forms of **(ML-style) state** (for general `Ref A`)
  - if the host language allows it, we use GADTs, etc for safety
  - some examples extract a footprint from a larger memory

- **Combinations** of different effects and runners
  - in particular the combination of IO and state
  - good use case for both vertical and horizontal composition

- KOKA-style **ambient values** and **ambient functions**
  - ambient values are essentially **mutable variables/parameters**
  - ambient functions are **applied in their lexical context**
  - a runner that treats **amb. fun. application as a co-operation**
  - amb. funs. are stored in a context-depth-sensitive heap
  - the appl. co-operation restores the heap to the lexical context

Implementing runners

Experimenting with the theory in practice

Experimenting with the theory in practice

- A **small experimental language** Coop[4]
  - Implements the core calculus with few extras
  - The interpreter is directly based on the denotational semantics
  - Top-level containers for running external (OCaml) code

[4] coop [/ku:p/] – a cage where small animals are kept, especially chickens

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  - Top-level containers for running external (OCaml) code
- A HASKELL **library** HASKELL-COOP
  - A shallow-embedding of the core calculus in HASKELL
  - Uses one of the Freer monad implementations underneath
  - Again, the operational aspects implement the denot. semantics
  - Top-level containers for arbitrary HASKELL monads
  - Examples make use of HASKELL's features (GADTs, ...)

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  - Top-level containers for running external (OCaml) code
- A Haskell **library** Haskell-Coop
  - A shallow-embedding of the core calculus in Haskell
  - Uses one of the Freer monad implementations underneath
  - Again, the operational aspects implement the denot. semantics
  - Top-level containers for arbitrary Haskell monads
  - Examples make use of Haskell's features (GADTs, ...)
- Both still need some finishing touches, but will be public soon

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Experimenting with the theory in practice

```haskell
module AmbientsTests where

import Control.Runner
import Control.Runner.Ambients

ambFun :: AmbVal Int -> Int -> AmbEff Int
ambFun x y =
  do x <- getVal x;
     return (x + y)

test1 :: AmbEff Int
test1 =
  withAmbVal
    (4 :: Int)
    (\ x ->
      withAmbFun
        (ambFun x)
        (\ f ->
          do rebindVal x 2;
             applyFun f 1))

test2 = ambTopLevel test1
```

Wrapping up

- **Runners** are a natural model of **top-level runtime**
- We propose **T-runners** to also model **non-top-level runtimes**
- We have turned **T**-runners into a **(practical ?) programming construct**, that supports controlled initialisation and finalisation
- I showed you some **combinators** and **programming examples**
- Two **implementations** in the works, COOP & HASKELL-COOP
- **Future:** lenses in subtyping and semantics, category of runners, handlers, bigger case studies, refinement typing, compilation, ...

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