

# Interacting with **external resources** using **runners** (aka **comodels**)

Danel Ahman

(joint work with Andrej Bauer)

University of Ljubljana, Slovenia

CHoCoLa meeting, Lyon, 17.10.2019

# Today's plan

- **Computational effects** and **external resources** in PL
- **Issues with standard approaches** to **external resources**
- **Runners** – a natural model for **top-level runtime**
- **T-runners** – for also modelling **non-top-level runtimes**
- Turning **T**-runners into a **useful programming construct**
- Demonstrate the use of runners through **programming examples**

**Computational effects**  
**and**  
**external resources**

# Computational effects in PL

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- Using **monads** (as in HASKELL)

```
type St a = String → (a,String)
```

```
f :: St a → St (a,a)
```

```
f c = c >>= (\ x → c >>= (\ y → return (x,y)))
```

- Using **alg. effects** and **handlers** (as in EFF, FRANK, KOKA)

```
effect Get : int
```

```
effect Put : int → unit
```

```
let g (c:Unit → a!{Get,Put}) : int → a * int =
```

```
  with st_handler handle (perform (Put 42); c ())
```

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- Both are good for **faking comp. effects** in a pure language!  
But what about effects that need access to the **external world**?

# External resources in PL

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- Declare a **signature of monads** or **algebraic effects**, e.g.,

```
(* System.IO *)  
type IO a  
openFile :: FilePath → IOMode → IO Handle
```

```
(* pervasives.eff *)  
effect RandomInt : int → int  
effect RandomFloat : float → float
```

- And then **treat them specially** in the compiler, e.g., in EFF

```
(* eff/src/backends/runtime/eval.ml *)  
let rec top_handle op =  
  match op with  
  | Value v → v  
  | Call (RandomInt, v, k) →  
    top_handle (k (Const.of_integer (Random.int (Value.to_int v))))  
  | ...
```



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but there are **some issues** with that approach ...

# First issue

- Difficult to cover all possible use cases
  - **external resources hard-coded** into the top-level runtime
  - **non-trivial to change** what's available and how it's implemented

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So here's the hack I added. We should do something a bit more principled

In `pervasives.eff`:

```
effect Write : (string*string) -> unit
```

in `eval.ml`, under `let rec top_handle op =` add the case:

```
| "Write" ->
  (match v with
  | V.Tuple vs ->
    let (file_name :: str :: _) = List.map V.to_str vs in
    let file_handle = open_out_gen
                        [Open_wronly
                         ;Open_append
                         ;Open_creat
                         ;Open_text
                        ] 0o666 file_name in
    Printf.fprintf file_handle "%s" str;
    close_out file_handle;
    top_handle (k V.unit_value)
  )
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```

**This work — a principled modular (co)algebraic approach!**

# Second issue

- **Lack of linearity** for external resources

```
let f (s:string) =  
  let fh = fopen "foo.txt" in  
  fwrite (fh,s^s);  
  fclose fh;  
  return fh
```

```
let g s =  
  let fh = f s in fread fh
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(\* fh not open any more ! \*)

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(\* fh not open any more ! \*)

- We shall address these kinds of issues **indirectly (!)**,
  - by **not** introducing a linear typing discipline
  - but instead we make it convenient to **hide external resources**  
(addressing stronger typing disciplines in the future)

# Third issue

- **Excessive generality** of effect handlers

```
let f (s:string) =  
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let h = handler { fwrite (fh,s) k → return () }
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```

- But misuse of external resources can also be **purely accidental**

```
let nd_handler =  
  handler { choose () k → return (k true ++ k false) }  
  
let g (s1 s2:string) =  
  let fh = fopen "foo.txt" in  
  let b = choose () in  
  if b then (fwrite (fh,s1^s2)) else (fwrite (fh,s2^s1));  
  fclose fh
```



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- We shall address these kinds of issues **directly (!!)**,
  - by proposing a **restricted form of handlers** for resources
  - that support **controlled initialisation** and **finalisation**,
  - (and limit how general handlers can be used)

# Runners

# A natural model of **top-level runtime**

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- Given a **signature**<sup>1</sup>  $\Sigma$  of operation symbols ( $A_{\text{op}}, B_{\text{op}}$  are sets)

$$\text{op} : A_{\text{op}} \rightsquigarrow B_{\text{op}}$$

a **runner**<sup>2</sup>  $\mathcal{R}$  for  $\Sigma$  is given by a carrier  $|\mathcal{R}|$  and co-operations

$$\left( \overline{\text{op}}_{\mathcal{R}} : A_{\text{op}} \times |\mathcal{R}| \longrightarrow B_{\text{op}} \times |\mathcal{R}| \right)_{\text{op} \in \Sigma}$$

where we think of  $|\mathcal{R}|$  as a set of **runtime configurations**

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- For example, a natural **runner  $\mathcal{R}$  for  $S$ -valued state** signature

$$\left\{ \text{get} : \mathbb{1} \rightsquigarrow S \quad , \quad \text{set} : S \rightsquigarrow \mathbb{1} \right\}$$

is given by

$$|\mathcal{R}| \stackrel{\text{def}}{=} S \qquad \overline{\text{get}}_{\mathcal{R}}(\star, s) \stackrel{\text{def}}{=} (s, s) \qquad \overline{\text{set}}_{\mathcal{R}}(s', s) \stackrel{\text{def}}{=} (\star, s')$$

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# A natural model of **top-level runtime** ctd.

- Runners/comodels have been used for
  - **operational semantics** using tensors of models and comodels  
[Plotkin and Power '08]
  - **top-level implementation of algebraic effects** in  $\text{EFF}$   
[Bauer and Pretnar '15]and
- **stateful running** of algebraic effects [Uustalu '15]
- **linear-use state-passing translation**  
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- **stateful running** of algebraic effects [Uustalu '15]
- **linear-use state-passing translation** [Møgelberg and Staton '11, '14]
- The latter explicitly rely on one-to-one correspondence between
  - **runners**  $\mathcal{R}$
  - **monad morphisms**<sup>3</sup>  $r : \mathbf{Free}_{\Sigma}(-) \longrightarrow \mathbf{St}_{|\mathcal{R}|}$

---

<sup>3</sup> $\mathbf{Free}_{\Sigma}(X)$  is the free monad ind. defined with leaves  $\text{val } x$  and nodes  $\text{op}(a, \kappa)$ .

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- But what if this runtime is not **\*\*the\*\*** runtime?
  - hardware vs OSs
  - OSs vs VMs
  - VMs vs sandboxes

but also

- browsers vs web pages
- ...

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- But what if this runtime is not **\*\*the\*\*** runtime?
  - hardware vs OSs
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  - ...
- Unfortunately, runners, as defined above, are **not readily able to**
  - use **external resources**
  - **signal failure** caused by unavoidable circumstances
- But is there a **useful generalisation** that would achieve this?

# Effectful runners for modular top-levels

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- Møgelberg and Staton usefully observed that a **runner**  $\mathcal{R}$  is equivalently simply a family of **generic effects** for  $\mathbf{St}_{|\mathcal{R}|}$ , i.e.,

$$\left( \overline{\text{op}}_{\mathcal{R}} : A_{\text{op}} \longrightarrow \mathbf{St}_{|\mathcal{R}|} B_{\text{op}} \right)_{\text{op} \in \Sigma}$$

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- Building on this, we define a **T-runner**  $\mathcal{R}$  for  $\Sigma$  to be given by

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- The one-to-one correspondence with **monad morphisms**

$$r : \mathbf{Free}_{\Sigma}(-) \longrightarrow \mathbf{T}$$

simply amounts to the **universal property of free models**, i.e.,

$$r_X(\text{val } x) = \eta_X x \qquad r_X(\text{op}(a, \kappa)) = \underbrace{(r_X \circ \kappa)^{\dagger}(\overline{\text{op}}_{\mathcal{R}} a)}_{\text{op}_{\mathcal{M}}(a, r_X \circ \kappa)}$$

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- Observe that  $\kappa$  appears in a **tail call position** on the right!

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- We want a runner to be a bit like a **kernel of an OS**, i.e., to
  - (i) provide management of **(internal) resources**
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- **Algebraically** (and pragmatically), this amounts to taking
  - (i)  $\text{getenv} : \mathbb{1} \rightsquigarrow C$     &     $\text{setenv} : C \rightsquigarrow \mathbb{1}$
  - (ii)  $\text{op} : A_{\text{op}} \rightsquigarrow B_{\text{op}}$     ( $\text{op} \in \Sigma'$ , for some external  $\Sigma'$ )
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- The **induced monad** is then isomorphic to

$$\mathbf{T} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma'}((X \times C) + S)$$

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- The corresponding **T-runners**  $\mathcal{R}$  for  $\Sigma$  are then of the form

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- With this, our **T-runners**  $\mathcal{R}$  for  $\Sigma$  are (with “primitive” excs.)

$$\left( \overline{\text{op}}_{\mathcal{R}} : A_{\text{op}} \longrightarrow \mathbf{K}_C^{\Sigma' ! E_{\text{op}} \not\leq S} B_{\text{op}} \right)_{\text{op} \in \Sigma}$$

where we call  $\mathbf{K}_C^{\Sigma' ! E_{\text{op}} \not\leq S}$  a **kernel monad**, given by

$$\mathbf{K}_C^{\Sigma' ! E_{\text{op}} \not\leq S} B_{\text{op}} \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma'}(((B_{\text{op}} + E_{\text{op}}) \times C) + S)$$



**T-runners as a programming construct**  
(towards a core calculus for runners)

# T-runners as a programming construct

- First, we include **T-runners** for  $\Sigma$

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in our language **as values**, and **co-ops. as kernel code**, i.e.,

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let R = runner { op1 x1 → K1 , ... , opn xn → Kn } @ C
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- For instance, we can implement a **write-only file handle** as

```
let RFH = runner {  
  write s → if (length s > maxSize)  
    then (raise WriteSizeExceeded)  
    else (let fh = getenv () in  
      if (isValid fh) then (fwrite (fh,s)) else (kill IOError))  
} @ FileHandle
```

where

$$\Sigma \stackrel{\text{def}}{=} \{ \text{write} : \text{String} \rightsquigarrow 1 ! E \cup \{ \text{WriteSizeExceeded} \} \}$$

$$(\text{fwrite} : \text{FileHandle} \times \text{String} \rightsquigarrow 1 ! E) \in \Sigma' \quad S = \{ \text{IOError} \}$$

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- We make use of it to enable programmers to **run user code**:

```
using R @ Minit
run M
finally {return x @ c → Mret , ... raise e @ c → Me ... , ... kill s → Ms ...}
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where

(a **user monad**)

- $M_s$  are **user code**, modelled using  $\mathbf{U}^{\Sigma!E} X \stackrel{\text{def}}{=} \mathbf{Free}_\Sigma(X + E)$

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- $M_{\text{init}}$  produces the **initial kernel state**
- $M$  is the user code being **run using the runner**  $R$
- $M_{\text{ret}}$ ,  $M_e$ ,  $M_s$  **finalise** for return values, exceptions, and signals

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where (a **user monad**)

- $M_s$  are **user code**, modelled using  $\mathbf{U}^{\Sigma!E} X \stackrel{\text{def}}{=} \mathbf{Free}_\Sigma(X + E)$
- $M_{\text{init}}$  produces the **initial kernel state**
- $M$  is the user code being **run using the runner**  $R$
- $M_{\text{ret}}$ ,  $M_e$ ,  $M_s$  **finalise** for return values, exceptions, and signals
- $M_{\text{ret}}$  and  $M_e$  **depend on the final state**  $c$ , but  $M_s$  **does not**



# Controlled **initialisation** and **finalisation** ctd.

- For instance, we can define a PYTHON-esque **with construct**

```
with fileName do M
=
using R_FH @ (fopen fileName)
run M
finally {
  return x @ fh → fclose fh; return x ,
  raise WriteSizeExceeded @ fh → fclose fh; return () ,
  raise e @ fh → fclose fh; raise e , (* other exceptions in E are re-raised *)
  kill IOError → ... }
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  kill IOError → ... }
```

- the **file handle is hidden** from M
- M **can only call** `write : String → 1 ! E ∪ {WriteSizeExceeded}`  
but **not** (the external operations) `fopen` , `fclose` , and `fwrite`
- `fopen` and `fclose` are **limited to initialisation-finalisation**
- M can itself also catch `WriteSizeExceeded` to **re-try writing**

**A core calculus for  
programming with runners**

# Core calculus (syntax)

# Core calculus (syntax)

- **Ground types** (types of ops. and kernel state)

$$A, B, C ::= B \mid 1 \mid 0 \mid A \times B \mid A + B$$

- **Types**

$$\begin{aligned} X, Y &::= B \mid 1 \mid 0 \mid X \times Y \mid X + Y \\ &\mid X \xrightarrow{\Sigma} Y ! E \\ &\mid X \xrightarrow{\Sigma} Y ! E \Downarrow S @ C \\ &\mid \Sigma \Rightarrow \Sigma' \Downarrow S @ C \end{aligned}$$

- **Values**

$$\Gamma \vdash V : X$$

- **User computations**

$$\Gamma \vDash M : X ! E$$

- **Kernel computations**

$$\Gamma \vDash K : X ! E \Downarrow S @ C$$

# Core calculus (user computations)

$M, N ::= \text{return } V$

value

|  $\text{try } M \text{ with } \{\text{return } x \mapsto N, (\text{raise } e \mapsto N_e)_{e \in E}\}$

exception handler

|  $V W$

application

|  $\text{match } V \text{ with } \{\langle x, y \rangle \mapsto M\}$

product elimination

|  $\text{match } V \text{ with } \{\} X$

empty elimination

|  $\text{match } V \text{ with } \{\text{inl } x \mapsto M, \text{inr } y \mapsto N\}$

sum elimination

|  $\text{op}_X(V, (x . M), (N_e)_{e \in E_{\text{op}}})$

operation call

|  $\text{raise}_X e$

raise exception

|  $\text{using } V @ W \text{ run } M \text{ finally } \{$

run

$\text{return } x @ c \mapsto N,$

$(\text{raise } e @ c \mapsto N_e)_{e \in E},$

$(\text{kill } s \mapsto N_s)_{s \in S}\}$

|  $\text{kernel } K @ V \text{ finally } \{$

switch to kernel mode

$\text{return } x @ c \mapsto N,$

$(\text{raise } e @ c \mapsto N_e)_{e \in E},$

$(\text{kill } s \mapsto N_s)_{s \in S}\}$

# Core calculus (kernel computations)

$K, L ::=$	$\text{return}_C V$	value
	$\text{try } K \text{ with } \{\text{return } x \mapsto L, (\text{raise } e \mapsto L_e)_{e \in E}\}$	exception handler
	$V W$	application
	$\text{match } V \text{ with } \{\langle x, y \rangle \mapsto K\}$	product elimination
	$\text{match } V \text{ with } \{\}_{X@C}$	empty elimination
	$\text{match } V \text{ with } \{\text{inl } x \mapsto K, \text{inr } y \mapsto L\}$	sum elimination
	$\text{op}_{X@C}(V, (x . K), (L_e)_{e \in E_{\text{op}}})$	operation call
	$\text{raise}_{X@C} e$	raise exception
	$\text{kill}_{X@C} s$	send signal
	$\text{getenv}_C(c . K)$	get state
	$\text{setenv}(V, K)$	set state
	$\text{user } M \text{ with } \{\text{return } x \mapsto K, (\text{raise } e \mapsto L_e)_{e \in E}\}$	switch to user mode

# Core calculus (type system and eq. theory)



# Core calculus (type system and eq. theory)

- For example, the **typing rule for running user comps.** is

$$\begin{array}{c}
 \Gamma \vdash V : \Sigma \Rightarrow \Sigma' \not\downarrow S @ C \quad \Gamma \vdash W : C \\
 \Gamma \Vdash M : X ! E \quad \Gamma, x : X, c : C \Vdash' N_{ret} : Y ! E' \\
 \frac{(\Gamma, c : C \Vdash' N_e : Y ! E')_{e \in E} \quad (\Gamma \Vdash' N_s : Y ! E')_{s \in S}}{\Gamma \Vdash' \text{using } V @ W \text{ run } M \text{ finally } \{ \text{return } x @ c \mapsto N_{ret} , \\
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 \end{array}$$

- and the **main  $\beta$ -equation for running user comps.** is

$$\begin{array}{l}
 \Gamma \Vdash' \text{using } R @ W \text{ run } (\text{op}_X(V, (y.M), (M_e)_{e \in E_{op}})) \text{ finally } F \\
 \equiv \text{kernel } K_{op}[V/x_{op}] @ W \text{ finally } \{ \\
 \quad \text{return } y @ c' \mapsto \text{using } R @ c' \text{ run } M \text{ finally } F , \\
 \quad (\text{raise } e @ c' \mapsto \text{using } R @ c' \text{ run } M_e \text{ finally } F)_{e \in E_{op}} , \\
 \quad (\text{kill } s \mapsto N_s)_{s \in S} \} : Y ! E'
 \end{array}$$

# Core calculus (type system and eq. theory)

- The calculus also includes **subtyping**, and **subsumption rules**

$$\frac{\Gamma \vdash V : A \quad A <: B}{\Gamma \vdash V : B}$$

$$\frac{\Gamma \Vdash M : A ! E \quad \Sigma \subseteq \Sigma' \quad A <: B \quad E \subseteq E'}{\Gamma \Vdash' M : B ! E'}$$

$$\frac{\begin{array}{cccc} \Gamma \Vdash K : A ! E \downarrow S @ C & \Sigma \subseteq \Sigma' & & \\ A <: B & E \subseteq E' & S \subseteq S' & C = C' \end{array}}{\Gamma \Vdash' K : B ! E' \downarrow S' @ C'}$$

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- We use  $C = C'$  to have (standard) **proof-irrelevant subtyping**
- Otherwise, instead of just  $C <: C'$ , we would need a **lens**  $C' \leftrightarrow C$

# Core calculus (semantics)

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- **Monadic semantics**, for concreteness in **Set**, using
  - **user monads**  $\mathbf{U}^{\Sigma!E} X \stackrel{\text{def}}{=} \mathbf{Free}_{\Sigma}(X + E)$
  - **kernel monads**  $\mathbf{K}_C^{\Sigma!E \not\vdash S} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}(((X + E) \times C) + S)$

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  - **kernel monads**  $\mathbf{K}_C^{\Sigma!E \not\downarrow S} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}(((X + E) \times C) + S)$
- (At a high level) the **judgements are interpreted** as

$$\llbracket \Gamma \vdash V : X \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \llbracket X \rrbracket$$

$$\llbracket \Gamma \models M : X ! E \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \mathbf{U}^{\Sigma!E} \llbracket X \rrbracket$$

$$\llbracket \Gamma \models K : X ! E \not\downarrow S @ C \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma!E \not\downarrow S} \llbracket X \rrbracket$$

## Core calculus (semantics ctd.)

- However, to prove **coherence** of the semantics (**subtyping!**), we actually give the semantics in the **subset fibration**



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- For instance, **kernel computations** are interpreted as

$$\begin{array}{ccc}
 \llbracket \Gamma \rrbracket & \xrightarrow{\llbracket \Gamma \models K : X ! E \not\leq S @ C \rrbracket} & \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma ! E \not\leq S} \llbracket X \rrbracket \\
 \downarrow \subseteq & & \downarrow \subseteq \\
 \llbracket \Gamma^s \rrbracket & \xrightarrow{\llbracket \Gamma^s \vdash K : X^s @ C \rrbracket} & \mathbf{K}_{\llbracket C \rrbracket}^{\mathcal{O} ! E \not\leq S + \{\perp\}} \llbracket X^s \rrbracket
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where  $\Gamma^s \vdash K : X^s @ C$  is a **skeletal kernel typing judgement**

- No essential obstacles to extending to **Sub(Cpo)** and beyond
- **Ground type restriction** on  $C$  needed to stay within **Sub(...)**
  - Otherwise, analogously to subtyping, we'd need **lenses** instead

## **Implementing runners**

**Experimenting with the theory in practice**

# Experimenting with the theory in practice

- A **small experimental language** COOP<sup>4</sup>
  - Implements the core calculus with few extras
  - The interpreter is directly based on the denotational semantics
  - Top-level containers for running external (OCaml) code

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- A **HASKELL library** HASKELL-COOP
  - A shallow-embedding of the core calculus in HASKELL
  - Uses one of the Freer monad implementations underneath
  - Again, the operational aspects implement the denot. semantics
  - Top-level containers for arbitrary HASKELL monads
  - Examples make use of HASKELL's features (GADTs, ...)

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  - Top-level containers for arbitrary HASKELL monads
  - Examples make use of HASKELL's features (GADTs, ...)
- Both still need some finishing touches, but will be public soon

---

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**Runners in action**

Runners can be **vertically nested**

# Runners can be vertically nested

- ```
using RFH @ (fopen fileName)
run (
  using RFC @ (return "")
  run M
  finally {
    return x @ str → write str; return x ,
    raise WriteSizeExceeded @ str → write str; raise WriteSizeExceeded }
)
finally {
  return x @ fh → ... , raise e @ fh → ... , kill IOError → ... }
```

where the **file contents runner** (with  $\Sigma' = \{\}$ ) is defined as

```
let RFC = runner {
  write str' → let str = getenv () in
    if (length (str^str') > max) then (raise WriteSizeExceeded)
    else (setenv (str^str'))
} @ String
```

# Vertical nesting for instrumentation

# Vertical nesting for instrumentation

- ```
using RSniffer @ (return 0)
run M
finally {
  return x @ c →
    let fh = fopen "nsa.txt" in fwrite (fh,toStr c); fclose fh; return x }
```

where the **instrumenting runner** is defined as

```
let RSniffer = runner {
  ... ,
  op a → let c = getenv () in
    setenv (c + 1);
    op a ,
  ...
} @ Nat
```

(\* forwards op outwards \*)

- The runner  $R_{\text{Sniffer}}$  implements the same sig.  $\Sigma$  that  $M$  is using
- As a result, the runner  $R_{\text{Sniffer}}$  is **invisible** from  $M$ 's viewpoint

# Vertical nesting for **active monitoring**

# Vertical nesting for **active monitoring**

- First, we define a runner for **integer-valued ML-style state** as

**type** IntHeap = (Nat  $\rightarrow$  (Int + 1))  $\times$  Nat

**type** Ref = Nat

```
let RIntState = runner {  
  alloc x  $\rightarrow$  let h = getenv () in                                     (* alloc : Int  $\rightsquigarrow$  Ref ! {} *)  
    let (r,h') = heapAlloc h x in  
    setenv h';  
    return r ,  
  
  deref r  $\rightarrow$  let h = getenv () in                                     (* deref : Ref  $\rightsquigarrow$  Int ! {} *)  
    match (heapSel h r) with  
    | inl x  $\rightarrow$  return x  
    | inr ()  $\rightarrow$  kill ReferenceDoesNotExist ,  
  
  assign r y  $\rightarrow$  let h = getenv () in                                 (* assign : Ref  $\times$  Int  $\rightsquigarrow$  1 ! {} *)  
    match (heapUpd h r y) with  
    | inl h'  $\rightarrow$  setenv h'  
    | inr ()  $\rightarrow$  kill ReferenceDoesNotExist  
}
```

@ IntHeap

## Vertical nesting for **active monitoring** ctd.

- Next we define a runner for **monotonicity layer** on top of  $R_{\text{IntState}}$



# Vertical nesting for **active monitoring** ctd.

- Next we define a runner for **monotonicity layer** on top of  $R_{\text{IntState}}$

**type** MonMemory = Ref  $\rightarrow ((\text{Int} \rightarrow \text{Int} \rightarrow \text{Bool}) + 1)$

```
let RMonState = runner {  
  mAlloc x rel  $\rightarrow$  let r = alloc x in (* : Int  $\times$  Ord  $\rightsquigarrow$  Ref ! {} *)  
    let m = getenv () in  
    setenv (memAdd m r rel);  
    return r,  
  
  mDeref r  $\rightarrow$  deref r , (* monDeref : Ref  $\rightsquigarrow$  Int ! {} *)  
  
  mAssign r y  $\rightarrow$  let x = deref r in (* : Ref  $\times$  Int  $\rightsquigarrow$  1 ! {MV} *)  
    let m = getenv () in  
    match (memSel m r) with  
    | inl rel  $\rightarrow$  if (rel x y)  
        then (assign r y)  
        else (raise MonotonicityViolation)  
    | inr  $\rightarrow$  kill PreorderDoesNotExist  
}  
@ MonMemory
```

## Vertical nesting for **active monitoring** ctd.

- We can then perform **runtime monotonicity verification** as

# Vertical nesting for **active monitoring** ctd.

- We can then perform **runtime monotonicity verification** as

```
using RIntState @ ((fun _ → inr ()), 0)      (* init empty ML-style heap *)
run (

  using RMonState @ (fun _ → inr ())          (* init empty preorders memory *)
  run (

    let r = mAlloc 0 (≤) in
    mAssign r 1;
    mAssign r 0;                               (* RMonState raises exception *)
    mAssign r 2

  )
  finally { ... }

)
finally { ... }
```

Runners can also be **horizontally paired**

# Runners can also be horizontally paired

- Given runners for  $\Sigma$  and  $\Sigma'$

```
let R1 = runner { ... , op1i x → K1i , ... } @ C1  
let R2 = runner { ... , op2j x → K2j , ... } @ C2
```

we can **pair them** to get a runner for  $\Sigma + \Sigma'$

```
let R = runner { ... ,  
  op1i x → let (c,c') = getenv () in  
    user (kernel (K1i x) @ c finally {  
      return y @ c'' → return (inl (inl y,c'')),  
      raise e @ c'' → return (inl (inr e,c'')), (* e ∈ Eop1i *)  
      kill s → return (inr s) } (* s ∈ S1 *)  
    finally {  
      return (inl (inl y,c'')) → setenv (c'',c'); return y,  
      return (inl (inr e,c'')) → setenv (c'',c'); raise e,  
      return (inr s) → kill s },  
  ... ,  
  op2j x → ..., (* analogously to above, just on 2nd comp. of state *)  
  ... } @ C1 × C2
```

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      return (inl (inr e,c'')) → setenv (c'',c'); raise e,  
      return (inr s) → kill s },  
  ... ,  
  op2j x → ..., (* analogously to above, just on 2nd comp. of state *)  
  ... } @ C1 × C2
```

- For instance, this way we can build a runner for **IO and state**

## Other examples (in `HASKELL`)

## Other examples (in HASKELL)

- More general forms of **(ML-style) state** (for general Ref A)
  - if the host language allows it, we use GADTs, etc for safety
  - some examples extract a footprint from a larger memory
- **Combinations** of different effects and runners
  - in particular the combination of IO and state
  - good use case for both vertical and horizontal composition
- KOKA-style **ambient values** and **ambient functions**
  - **ambient values** are essentially **mutable variables/parameters**
  - **ambient functions** are **applied in their lexical context**
  - a runner that treats **amb. fun. application as a co-operation**
  - amb. funs. are stored in a context-depth-sensitive heap
  - the appl. co-operation restores the heap to the lexical context



# Other examples (ambient functions)

```
module Control.Runner.Ambients

...

ambCoOps :: Amb a -> Kernel sig AmbHeap a
ambCoOps (Bind f) =
  do h <- getEnv;
    (f,h') <- return (ambHeapAlloc h f);
    setEnv h';
    return f
ambCoOps (Apply f x) =
  do h <- getEnv;
    (f,d) <- return (ambHeapSel h f (depth h));
    user
      (run
        ambRunner
          (return (h {depth = d}))
          (f x)
          ambFinaliser)
    return
ambCoOps (Rebind f g) =
  do h <- getEnv;
    setEnv (ambHeapUpd h f g)

ambRunner :: Runner '[Amb] sig AmbHeap
ambRunner = mkRunner ambCoOps
```

```
module AmbientsTests where

import Control.Runner
import Control.Runner.Ambients

ambFun :: AmbVal Int -> Int -> AmbEff Int
ambFun x y =
  do x <- getVal x;
    return (x + y)

test1 :: AmbEff Int
test1 =
  withAmbVal
    (4 :: Int)
    (\ x ->
      withAmbFun
        (ambFun x)
        (\ f ->
          do rebindVal x 2;
            applyFun f 1))

test2 = ambTopLevel test1
```

# Wrapping up

- **Runners** are a natural model of **top-level runtime**
- We propose **T-runners** to also model **non-top-level runtimes**
- We have turned **T**-runners into a **(practical ?) programming construct**, that supports controlled initialisation and finalisation
- I showed you some **combinators** and **programming examples**
- Two **implementations** in the works, COOP & HASKELL-COOP
- **Ongoing** and **future**: lenses in subtyping and semantics, cat. of runners, handlers, case studies, refinement typing, compilation, ...

This project has received funding from the European Union's Horizon 2020 research and innovation programme under the Marie Skłodowska-Curie grant agreement No 834146.



This material is based upon work supported by the Air Force Office of Scientific Research under award number FA9550-17-1-0326.

**Thank you!**





# Core calculus (semantics ctd.)

$$\begin{aligned} \llbracket \Gamma \models' \text{using } V @ W \text{ run } M \text{ finally } \{ \text{return } x @ c \mapsto N_{ret} , \\ (\text{raise } e @ c \mapsto N_e)_{e \in E} , \\ (\text{kill } s \mapsto N_s)_{s \in S} \} : Y ! E' \rrbracket_\gamma \stackrel{\text{def}}{=} \dots \end{aligned}$$

- $\llbracket V \rrbracket_\gamma = \mathcal{R} = \left( \overline{\text{op}}_{\mathcal{R}} : \llbracket A_{\text{op}} \rrbracket \longrightarrow \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma' ! E_{\text{op}} \not\leq S} \llbracket B_{\text{op}} \rrbracket \right)_{\text{op} \in \Sigma}$
- $\llbracket W \rrbracket_\gamma \in \llbracket C \rrbracket$
- $\llbracket M \rrbracket_\gamma \in \mathbf{U}^{\Sigma ! E} \llbracket A \rrbracket$
- $\llbracket \text{return } x @ c \mapsto N_{ret} \rrbracket_\gamma \in \llbracket A \rrbracket \times \llbracket C \rrbracket \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
- $\llbracket (\text{raise } e @ c \mapsto N_e)_{e \in E} \rrbracket_\gamma \in E \times \llbracket C \rrbracket \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
- $\llbracket (\text{kill } s \mapsto N_s)_{s \in S} \rrbracket_\gamma \in S \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
- allowing us to use the **free model property** to get

$$\mathbf{U}^{\Sigma ! E} \llbracket A \rrbracket \xrightarrow{r_{\llbracket A \rrbracket} + E} \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma' ! E \not\leq S} \llbracket A \rrbracket \xrightarrow{(\lambda \llbracket N_{ret} \rrbracket_\gamma)^\dagger} \llbracket C \rrbracket \Rightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$$

and then apply the resulting composite to  $\llbracket M \rrbracket_\gamma$  and  $\llbracket W \rrbracket_\gamma$