

Interacting with the **external world** using **comodels** (aka **runners**)

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Gallinette seminar, Nantes, 14.10.2019

The plan

- **Computational effects** and **external resources** in PL
- **Runners** – a natural model for **top-level runtime**
- **T-runners** – for also modelling **non-top-level runtimes**
- Turning **T**-runners into a **useful programming construct**
- Some **programming examples**
- Some **implementation details**

Computational effects
and
external resources

Computational effects in PL

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- Using **monads** (as in HASKELL)

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type St a = String → (a,String)
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```
f :: St a → St (a,a)
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f c = c >>= (\x → c >>= (\y → return (x,y)))
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- Using **alg. effects** and **handlers** (as in EFF, FRANK, KOKA)

```
effect Get : int
```

```
effect Put : int → unit
```

```
let g (c:Unit → a!{Get,Put}) =
```

```
  with st_handler handle (perform (Put 42); c ()) (* : int → a * int *)
```

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```
  with st_handler handle (perform (Put 42); c ()) (* : int → a * int *)
```

- Both are good for **faking comp. effects** in a pure language!
But what about effects that need access to the **external world**?

External resources in PL

External resources in PL

- Declare a **signature of monads** or **algebraic effects**, e.g.,

```
(* System.IO *)  
type IO a  
openFile :: FilePath → IOMode → IO Handle
```

```
(* pervasives.eff *)  
effect RandomInt : int → int  
effect RandomFloat : float → float
```

- And then **treat them specially** in the compiler, e.g.,

```
(* eff/src/backends/eval.ml *)  
let rec top_handle op =  
  match op with  
  | ...
```

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but there are some issues with that approach ...

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- Difficult to cover all possible use cases
 - **external resources hard-coded** into the top-level runtime
 - **non-trivial to change** what's available and how it's implemented

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So here's the hack I added. We should do something a bit more principled

In `pervasives.eff`:

```
effect Write : (string*string) -> unit
```

in `eval.ml`, under `let rec top_handle op =` add the case:

```
| "Write" ->
  (match v with
  | V.Tuple vs ->
    let (file_name :: str :: _) = List.map V.to_str vs in
    let file_handle = open_out_gen
                        [Open_wronly
                         ;Open_append
                         ;Open_creat
                         ;Open_text
                        ] 0o666 file_name in
    Printf.fprintf file_handle "%s" str;
    close_out file_handle;
    top_handle (k V.unit_value)
  )
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This talk — a principled modular (co)algebraic approach!

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- **Lack of linearity** for external resources

```
let f (s:string) =  
  let fh = fopen "foo.txt" in  
  fwrite (fh,s^s);  
  fclose fh;  
  return fh
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let g s =  
  let fh = f s in fread fh
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- We shall address these kinds of issues **indirectly**,
 - by **not** introducing a linear typing discipline
 - but instead make it convenient to **hide** external resources

Third issue

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- **Excessive generality** of effect handlers

```
let f (s:string) =  
  let fh = fopen "foo.txt" in  
  fwrite (fh,s^s);  
  fclose fh  
  
let h = handler { fwrite (fh,s) k → return () }  
  
let f' s = handle (f "bar") with h
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where misuse of external resources can also be **purely accidental**

```
let g (s:string) =  
  let fh = fopen "foo.txt" in  
  let b = choose () in  
  if b then (fwrite (fh,s)) else (fwrite (fh,s^s));  
  fclose fh  
  
let nd_handler =  
  handler { choose () k → return (k true ++ k false) }
```

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- We shall address these kinds of issues **directly**,
 - by proposing a **restricted form** of handlers for resources
 - that support **controlled initialisation** and **finalisation**,
 - and **limit** how general handlers can be used

Runners enter the spotlight

A natural model of **top-level runtime**

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- Given a **signature**¹ Σ of operation symbols ($A_{\text{op}}, B_{\text{op}}$ are sets)

$$\text{op} : A_{\text{op}} \rightsquigarrow B_{\text{op}}$$

a **runner**² \mathcal{R} for Σ is given by a carrier $|\mathcal{R}|$ and co-operations

$$\left(\overline{\text{op}}_{\mathcal{R}} : A_{\text{op}} \times |\mathcal{R}| \longrightarrow B_{\text{op}} \times |\mathcal{R}| \right)_{\text{op} \in \Sigma}$$

¹We consider runners for signatures, but the work generalises to alg. theories.

²In the literature also known as **comodels** for Σ (or for an algebraic theory).

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- For example, a natural runner \mathcal{R} for **S-valued state**

$$\text{get} : \mathbb{1} \rightsquigarrow S \quad \text{set} : S \rightsquigarrow \mathbb{1}$$

is given by

$$|\mathcal{R}| \stackrel{\text{def}}{=} S \quad \overline{\text{get}}_{\mathcal{R}}(\star, s) \stackrel{\text{def}}{=} (s, s) \quad \overline{\text{set}}_{\mathcal{R}}(s, s) \stackrel{\text{def}}{=} (\star, s)$$

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A natural model of **top-level runtime** ctd.

- Runners/comodels have been used for
 - **operational semantics** using tensors of models and comodels
[Plotkin and Power '08]
and
 - **stateful running** of algebraic effects [Uustalu '15]
 - **linear-use state-passing translation**
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and
 - **stateful running** of algebraic effects [Uustalu '15]
 - **linear-use state-passing translation** [Møgelberg and Staton '11, '14]
- The latter explicitly rely on one-to-one correspondence between
 - **runners** \mathcal{R}
 - **monad morphisms**³ $r : \mathbf{Free}_{\Sigma}(-) \longrightarrow \mathbf{St}_{|\mathcal{R}|}$

where

$$\mathbf{St}_C X \stackrel{\text{def}}{=} C \Rightarrow X \times C$$

³ $\mathbf{Free}_{\Sigma}(X)$ is the free monad ind. defined with leaves $\text{val } x$ and nodes $\text{op}(a, \kappa)$.

A natural model of **top-level runtime** ctd.

- For our purposes, we see runners

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as describing how operations affect **runtime configurations** $|\mathcal{R}|$

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- But what if this runtime is not **the** runtime?
 - hardware vs OS
 - OS vs VMs
 - VMs vs sandboxes

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- Unfortunately, runners, as defined above, are **not readily able to**
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 - **signal failure** caused by unavoidable circumstances

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- Unfortunately, runners, as defined above, are **not readily able to**
 - use **external resources**
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- But is there a **useful generalisation** that would achieve this?

Effectful runners for modular top-levels

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- Møgelberg and Staton usefully observed that a **runner** \mathcal{R} is equivalently simply a family of **generic effects** for $\mathbf{St}_{|\mathcal{R}|}$, i.e.,

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- The one-to-one correspondence with **monad morphisms**

$$r : \mathbf{Free}_{\Sigma}(-) \longrightarrow \mathbf{T}$$

now simply amounts to the **univ. property of free models**, e.g.,

$$r_X(\text{val } x) = \eta_X x \qquad r_X(\text{op}(a, \kappa)) = (r_X \circ \kappa)^{\dagger}(\overline{\text{op}}_{\mathcal{R}} a)$$

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- Observe that κ appears in a **tail call position** on the right!

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- What would be a **useful class of monads** \mathbf{T} to use?

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- What would be a **useful class of monads \mathbf{T}** to use?
- We want a runner to be a bit like a **kernel** of an OS, i.e., to
 - (i) provide management of **(internal) resources**
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- **Algebraically** (and pragmatically), this amounts to taking
 - (i) $\text{getenv} : \mathbb{1} \rightsquigarrow C$, $\text{setenv} : C \rightsquigarrow \mathbb{1}$
 - (ii) $\text{op} : A_{\text{op}} \rightsquigarrow B_{\text{op}}$ ($\text{op} \in \Sigma'$, for some external Σ')
 - (iii) $\text{kill} : S \rightsquigarrow \mathbb{0}$s.t., (i) satisfy state equations; and (i) commute with (ii) and (iii)

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- The **induced monad** is then isomorphic to

$$\mathbf{T} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma'}((X \times C) + S)$$

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- The corresponding **T-runners** \mathcal{R} for Σ are then of the form

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- With this, our **T-runners** \mathcal{R} for Σ are (with “primitive” excs.)

$$\left(\overline{\text{op}}_{\mathcal{R}} : A_{\text{op}} \longrightarrow \mathbf{K}_C^{\Sigma'!E_{\text{op}} \not\vdash S} B_{\text{op}} \right)_{\text{op} \in \Sigma}$$

where we call $\mathbf{K}_C^{\Sigma'!E \not\vdash S}$ a **kernel monad**, given by

$$\mathbf{K}_C^{\Sigma'!E \not\vdash S} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}(((X + E) \times C) + S)$$

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we can easily accommodate co-operations as **kernel code**

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let R = runner { op1 x1 → K1 , ... , opn xn → Kn } @ C
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- For instance, we can implement a **write-only file handle** as

```
let RFH = runner {  
  write s → if (length s > max)  
    then (raise WriteSizeExceeded)  
    else (let fh = getenv () in  
          if (isValid fh) then (fwrite (fh,s)) else (kill IOError))  
} @ FileHandle
```

where

$$(\text{fwrite} : \text{FileHandle} \times \text{String} \rightsquigarrow 1 + E) \in \Sigma'$$

$$\Sigma \stackrel{\text{def}}{=} \{ \text{write} : \text{String} \rightsquigarrow 1 + E \cup \{\text{WriteSizeExceeded}\} \} \quad \text{IOError} \in S$$

Controlled **initialisation** and **finalisation**

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- We make use of it to enable one to **run user code**:

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using R @ Minit  
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finally {return x @ c → Mret , ... raise e @ c → Me ... , ... kill s → Ms ...}
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(a **user monad**)

- M_s are **user code**, modelled using $\mathbf{U}^{\Sigma!E} X \stackrel{\text{def}}{=} \mathbf{Free}_\Sigma(X + E)$

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- M is the user code being **run using the runner** R
- M_{ret} , M_e , M_s **finalise** for return values, exceptions, and signals

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- M is the user code being **run using the runner** R
- M_{ret} , M_e , M_s **finalise** for return values, exceptions, and signals
- M_{ret} and M_e **depend on the final state** c , but M_s **does not**

Controlled **initialisation** and **finalisation** ctd.

- For instance, we can define a PYTHON-esque **with construct**

```
with fileName do M
=
using R_FH @ (fopen fileName)
run M
finally {
  return x @ fh → fclose fh; return x ,
  raise e @ fh → fclose fh; raise e ,
  kill s → return () }
```

- Importantly, here
 - the **file handle is hidden** from `M`
 - `M` can only use `write` but not `fopen` and `fclose`
 - `write` : $\text{String} \rightsquigarrow 1 + E \cup \{\text{WriteSizeExceeded}\}$
 - `fopen` and `fclose` are limited to **initialisation-finalisation**

**A core calculus for
programming with runners**

Core calculus (syntax)

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- **Ground types** (types of ops. and kernel state)

$$A, B, C ::= B \mid 1 \mid 0 \mid A \times B \mid A + B$$

- **Types**

$$\begin{aligned} X, Y &::= B \mid 1 \mid 0 \mid X \times Y \mid X + Y \\ &\mid X \xrightarrow{\Sigma} Y ! E \\ &\mid X \xrightarrow{\Sigma} Y ! E \Downarrow S @ C \\ &\mid \Sigma \Rightarrow \Sigma' \Downarrow S @ C \end{aligned}$$

- **Values**

$$\Gamma \vdash V : X$$

- **User computations**

$$\Gamma \Vdash M : X ! E$$

- **Kernel computations**

$$\Gamma \Vdash K : X ! E \Downarrow S @ C$$

Core calculus (user computations)

$$\begin{aligned} M, N ::= & \text{ return } V \\ & | \text{ try } M \text{ with } \{ \text{ return } x \mapsto M_{\text{return}}, \\ & \quad (\text{ raise } e \mapsto N_e)_{e \in E} \} \\ & | V \ W \\ & | \text{ match } V \text{ with } \{ \langle x, y \rangle \mapsto M \} \\ & | \text{ match } V \text{ with } \{ \} _X \\ & | \text{ match } V \text{ with } \{ \text{ inl } x \mapsto M, \text{ inr } y \mapsto N \} \\ & | \text{ op}_X(V, (x . M), (N_e)_{e \in E}) \\ & | \text{ raise}_X e \\ & | \text{ using } V @ W \text{ run } M \text{ finally } \{ \text{ return } x @ c \mapsto M, \\ & \quad (\text{ raise } e @ c \mapsto N_e)_{e \in E} \\ & \quad (\text{ kill } s \mapsto N_s)_{s \in S} \} \\ & | \text{ kernel } K @ V \text{ finally } \{ \text{ return } x @ c \mapsto M, \\ & \quad (\text{ raise } e @ c \mapsto N_e)_{e \in E} \\ & \quad (\text{ kill } s \mapsto N_s)_{s \in S} \} \end{aligned}$$

Core calculus (kernel computations)

$$\begin{aligned} K, L ::= & \text{return}_C V \\ & | \text{try } K \text{ with } \{ \text{return } x \mapsto K_{\text{return}}, \\ & \quad (\text{raise } e \mapsto L_e)_{e \in E} \} \\ & | V W \\ & | \text{match } V \text{ with } \{ \langle x, y \rangle \mapsto K \} \\ & | \text{match } V \text{ with } \{ \}_{X@C} \\ & | \text{match } V \text{ with } \{ \text{inl } x \mapsto K, \text{inr } y \mapsto L \} \\ & | \text{op}_{X@C}(V, (x . K), (L_e)_{e \in E}) \\ & | \text{raise}_{X@C} e \\ & | \text{kill}_{X@C} s \\ & | \text{getenv}_C(c . K) \\ & | \text{setenv}(V, K) \\ & | \text{user } M \text{ with } \{ \text{return } x \mapsto K_{\text{return}}, \\ & \quad (\text{raise } e \mapsto L_e)_{e \in E} \} \end{aligned}$$

Core calculus (type system and eq. theory)

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- For example, the **typing rule for running user comps.** is

$$\begin{array}{c}
 \Gamma \vdash V : \Sigma \Rightarrow \Sigma' \not\downarrow S @ C \quad \Gamma \vdash W : C \\
 \Gamma \Vdash M : X ! E \quad \Gamma, x : X, c : C \Vdash' N_{ret} : Y ! E' \\
 \hline
 \frac{(\Gamma, c : C \Vdash' N_e : Y ! E')_{e \in E} \quad (\Gamma \Vdash' N_s : Y ! E')_{s \in S}}{\Gamma \Vdash' \text{using } V @ W \text{ run } M \text{ finally } \{ \text{return } x @ c \mapsto N_{ret} , \\
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 \text{(kill } s \mapsto N_s)_{s \in S} \} : Y ! E'}
 \end{array}$$

- and the **main β -equation for running user comps.** is

$$\begin{aligned}
 &\Gamma \preceq' \text{using } R_C @ W \text{ run } (\text{op}_X(V, (x.M), (M_e)_{e \in E_{\text{op}}})) \text{ finally } F \\
 &\equiv \text{kernel } R_{\text{op}}[V] @ W \text{ finally } \{ \\
 &\quad \text{return } x @ c' \mapsto \text{using } R_C @ c' \text{ run } M \text{ finally } F , \\
 &\quad (\text{raise } e @ c' \mapsto \text{using } R_C @ c' \text{ run } M_e \text{ finally } F)_{e \in E_{\text{op}}} , \\
 &\quad (\text{kill } s \mapsto N_s)_{s \in S} \} : Y ! E'
 \end{aligned}$$

Core calculus (type system and eq. theory)

- The calculus also includes **subtyping**, and **subsumption rules**

$$\frac{\Gamma \vdash V : A \quad A <: B}{\Gamma \vdash V : B}$$

$$\frac{\Gamma \Vdash M : A ! E \quad \Sigma \subseteq \Sigma' \quad A <: B \quad E \subseteq E'}{\Gamma \Vdash' M : B ! E'}$$

$$\frac{\begin{array}{cccc} \Gamma \Vdash K : A ! E \downarrow S @ C & \Sigma \subseteq \Sigma' & & \\ A <: B & E \subseteq E' & S \subseteq S' & C = C' \end{array}}{\Gamma \Vdash' K : B ! E' \downarrow S' @ C'}$$

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- We use $C = C'$ to have (standard) **proof-irrelevant subtyping**
- Otherwise, instead of just $C <: C'$, we would need a **lens** $C' \leftrightarrow C$

Core calculus (semantics)

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- **Monadic semantics**, for simplicity in **Set**, using
 - **user monads** $\mathbf{U}^{\Sigma!E} X \stackrel{\text{def}}{=} \mathbf{Free}_{\Sigma}(X + E)$
 - **kernel monads** $\mathbf{K}_C^{\Sigma!E \not\vdash S} X \stackrel{\text{def}}{=} C \Rightarrow \mathbf{Free}_{\Sigma}(((X + E) \times C) + S)$

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- (At a high level) the **judgements are interpreted** as

$$\llbracket \Gamma \vdash V : X \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \llbracket X \rrbracket$$

$$\llbracket \Gamma \models M : X ! E \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \mathbf{U}^{\Sigma!E} \llbracket X \rrbracket$$

$$\llbracket \Gamma \models K : X ! E \not\vdash S @ C \rrbracket : \llbracket \Gamma \rrbracket \longrightarrow \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma!E \not\vdash S} \llbracket X \rrbracket$$

Core calculus (semantics ctd.)

- However, to prove **coherence** of the semantics (**subtyping!**), we actually give the semantics in the **subset fibration**

Core calculus (semantics ctd.)

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- For instance, **kernel computations** are interpreted as

$$\begin{array}{ccc}
 \llbracket \Gamma \rrbracket & \xrightarrow{\llbracket \Gamma \models K : X ! E \not\leq S @ C \rrbracket} & \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma ! E \not\leq S} \llbracket X \rrbracket \\
 \downarrow \subseteq & & \downarrow \subseteq \\
 \llbracket \Gamma^s \rrbracket & \xrightarrow{\llbracket \Gamma^s \vdash K : X^s @ C \rrbracket} & \mathbf{K}_{\llbracket C \rrbracket}^{\mathcal{O} ! E \not\leq S + \{\perp\}} \llbracket X^s \rrbracket
 \end{array}$$

where $\Gamma^s \vdash K : X^s @ C$ is a **skeletal kernel typing judgement**

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 \end{array}$$

where $\Gamma^s \vdash K : X^s @ C$ is a **skeletal kernel typing judgement**

- No essential obstacles to extending to **Sub(Cpo)** and beyond
- **Ground type restriction** on C needed to stay within **Sub(—)**
 - Otherwise, analogously to subtyping, we'd need **lenses** instead

Core calculus (semantics ctd.)

$$\llbracket \Gamma \models' \text{using } V @ W \text{ run } M \text{ finally } \{ \text{return } x @ c \mapsto N_{ret} , \\ (\text{raise } e @ c \mapsto N_e)_{e \in E} , \\ (\text{kill } s \mapsto N_s)_{s \in S} \} : Y ! E' \rrbracket_\gamma \stackrel{\text{def}}{=} \dots$$

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- $\llbracket V \rrbracket_\gamma = \mathcal{R} = \left(\overline{\text{op}}_{\mathcal{R}} : \llbracket A_{\text{op}} \rrbracket \longrightarrow \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma' ! E_{\text{op}} \not\vdash S} \llbracket B_{\text{op}} \rrbracket \right)_{\text{op} \in \Sigma}$
- $\llbracket W \rrbracket_\gamma \in \llbracket C \rrbracket$
- $\llbracket M \rrbracket_\gamma \in \mathbf{U}^{\Sigma' ! E} \llbracket A \rrbracket$
- $\llbracket \text{return } x @ c \mapsto N_{ret} \rrbracket_\gamma \in \llbracket A \rrbracket \times \llbracket C \rrbracket \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
- $\llbracket (\text{raise } e @ c \mapsto N_e)_{e \in E} \rrbracket_\gamma \in E \times \llbracket C \rrbracket \longrightarrow \mathbf{U}^{\Sigma' ! E'} \llbracket B \rrbracket$
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Core calculus (semantics ctd.)

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- $\llbracket (\text{kill } s \mapsto N_s)_{s \in S} \rrbracket_\gamma \in S \longrightarrow \mathbf{U}^{\Sigma'!E'} \llbracket B \rrbracket$
- allowing us to use the **free model property** to get

$$\mathbf{U}^{\Sigma!E} \llbracket A \rrbracket \xrightarrow{r_{\llbracket A \rrbracket} + E} \mathbf{K}_{\llbracket C \rrbracket}^{\Sigma'!E \not\leq S} \llbracket A \rrbracket \xrightarrow{(\lambda \llbracket N_{ret} \rrbracket_\gamma)^\dagger} \llbracket C \rrbracket \Rightarrow \mathbf{U}^{\Sigma'!E'} \llbracket B \rrbracket$$

and then apply the resulting composite to $\llbracket M \rrbracket_\gamma$ and $\llbracket W \rrbracket_\gamma$

Runners in action

Runners can be **vertically nested**

Runners can be **vertically nested**

- ```
using RFH @ (fopen fileName)
run (
 using RFC @ (return "")
 run M
 finally {
 return x @ str → write str; return x ,
 raise e @ str → write str; raise e }
)
finally {
 return x @ fh → fclose fh; return x ,
 raise e @ fh → fclose fh; raise e , kill IOError → return ()}
```

where the **file contents runner** (with  $\Sigma' = \emptyset$ ) is defined as

```
let RFC = runner {
 write str → let str' = getenv () in
 if (length (str^str') > max) then (raise WriteSizeExceeded)
 else (setenv (str^str'))
} @ String
```

# Vertical nesting for instrumentation

# Vertical nesting for instrumentation

- ```
using RSniffer @ (return 0)
run M
finally {
  return x @ c →
    let fh = fopen "nsa.txt" in fwrite (fh, nat_to_str c); fclose fh }
```

where the **instrumenting runner** is defined as

```
let RSniffer = runner {
  ... ,
  op a → let c = getenv () in
    setenv (c + 1);
    op a ,
  ...
} @ Nat
```

(* forwards op outwards *)

- The runner R_{Sniffer} implements the same sig. Σ that M is using
- As a result, the runner R_{Sniffer} is **invisible** from M 's viewpoint

Vertical nesting for **active monitoring**

Vertical nesting for **active monitoring**

- First, we define a runner for **integer-valued ML-style state** as

type IntHeap = (Nat \rightarrow (Int + 1)) \times Nat

type Ref = Nat

```
let RIntState = runner {  
  alloc x  $\rightarrow$  let h = getenv () in  
    let (r, h') = heapAlloc h x in  
    setenv h';  
    return r ,  
  
  deref r  $\rightarrow$  let h = getenv () in  
    match (heapSel h r) with  
    | inl x  $\rightarrow$  return x  
    | inr ()  $\rightarrow$  kill ReferenceDoesNotExist ,  
  
  assign r y  $\rightarrow$  let h = getenv () in  
    match (heapUpd h r y) with  
    | inl h'  $\rightarrow$  setenv h'  
    | inr ()  $\rightarrow$  kill ReferenceDoesNotExist  
}
```

 @ IntHeap

Vertical nesting for **active monitoring** ctd.

- Next we define a runner for **monotonicity layer** on top of R_{IntState}

Vertical nesting for **active monitoring** ctd.

- Next we define a runner for **monotonicity layer** on top of R_{IntState}

```
type MonMemory = Ref  $\rightarrow ((\text{Int} \rightarrow \text{Int} \rightarrow \text{Bool}) + 1)$ 
```

```
let RMonState = runner {  
  monAlloc x rel  $\rightarrow$  let r = alloc x in  
    let m = getenv () in  
    setenv (memAdd m r rel);  
    return r,  
  
  monDeref r  $\rightarrow$  deref r ,  
  
  monAssign r y  $\rightarrow$  let x = deref r in  
    let m = getenv () in  
    match (memSel m r) with  
    | inl rel  $\rightarrow$  if (rel x y)  
      then (assign r y)  
      else (raise MonotonicityViolation)  
    | inr  $\rightarrow$  kill PreorderDoesNotExist  
}
```

@ IntHeap

Vertical nesting for **active monitoring** ctd.

- We can then perform **runtime monotonicity verification** as

Vertical nesting for **active monitoring** ctd.

- We can then perform **runtime monotonicity verification** as

```
using RIntState @ ((fun _ → inr ()), 0)           (* empty ML-style heap *)

run (
  using RMonState @ (fun _ → inr ())              (* empty preorders memory *)

  run (
    let r = monAlloc 0 (≤) in
    monAssign r 1;
    monAssign r 0;      (* RMonState raises MonotonicityViolation exception *)
    monAssign r 2 )

  finally {return x @ _ → return x ,
           raise MonotonicityViolation @ _ → ... ,
           kill PreorderDoesNotExist → ... })

finally {return x @ _ → return x ,
         kill ReferenceDoesNotExist → ... }
```

Runners can also be **horizontally paired**

Runners can also be horizontally paired

- Given a runner for Σ

```
let R1 = runner { ... , op1i x → k1i , ... } @ C1
```

and a runner for Σ'

```
let R2 = runner { ... , op2j x → k2j , ... } @ C2
```

we can **pair them** to get a runner for $\Sigma \cup \Sigma'$

```
let R = runner {  
  ... ,  
  op1i x → let (c,c') = getenv () in  
             let (x,c'') = k1i x in  
             setenv (c'',c');  
             return x ,  
  ... ,  
  op2j x → ... (* analogously to above *) ,  
  ...  
} @ C1 × C2
```


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  ...  
} @ C1 × C2
```

- For instance, this way we can build a **runner for IO and state**

Other examples

Other examples

- More general forms of **(ML-style) state** (for general Ref A)
 - if the host language allows it, we use GADTs, etc for safety
 - some examples extract a footprint from a larger memory
- **Combinations** of different effects and runners
 - in particular the combination of IO and state
 - good use case for both vertical and horizontal composition
- KOKA-style **ambient values** and **ambient functions**
 - ambient values are essentially **mutable variables/parameters**
 - ambient functions are **applied in their lexical context**
 - a runner that treats **amb. fun. application as a co-operation**
 - amb. funs. are stored in a context-depth-sensitive heap
 - the appl. co-operation restores the heap to the lexical context

Implementing runners

Experimenting with the **theory in practice**

Experimenting with the theory in practice

- A **small experimental language** COOP⁴
 - Implements the core calculus with few extras
 - The interpreter is directly based on the denotational semantics
 - Top-level containers for running external (OCaml) code

⁴coop [/ku:p/] – a cage where small animals are kept, especially chickens

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- A **HASKELL library** HASKELL-COOP
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 - Again, the operational aspects implement the denot. semantics
 - Top-level containers for arbitrary HASKELL monads
 - Examples make use of HASKELL's features (GADTs, ...)

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 - Examples make use of HASKELL's features (GADTs, ...)
- Both still need some finishing touches, but will be public soon

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Experimenting with the theory in practice

```
module AmbientsTests where

import Control.Runner
import Control.Runner.Ambients

ambFun :: AmbVal Int -> Int -> AmbEff Int
ambFun x y =
  do x <- getVal x;
  return (x + y)

test1 :: AmbEff Int
test1 =
  withAmbVal
    (4 :: Int)
    (\ x ->
      withAmbFun
        (ambFun x)
        (\ f ->
          do rebindVal x 2;
          applyFun f 1))

test2 = ambToplevel test1
```

Wrapping up

- **Runners** are a natural model of **top-level runtime**
- We propose **T-runners** to also model **non-top-level runtimes**
- We have turned **T-runners** into a **(practical ?) programming construct**, that supports controlled initialisation and finalisation
- I showed you some **combinators** and **programming examples**
- Two **implementations** in the works, COOP & HASKELL-COOP
- **Future:** lenses in subtyping and semantics, category of runners, handlers, larger case studies, refinement typing, compilation, ...

This project has received funding from the European Union's Horizon 2020 research and innovation programme under the Marie Skłodowska-Curie grant agreement No 834146.



This material is based upon work supported by the Air Force Office of Scientific Research under award number FA9550-17-1-0326.

Thank you!

