Handling Fibred Algebraic Effects

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INRIA Paris

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Dependent Types

and

Logical Reasoning

Algebraic Effects

and

Effect Handlers

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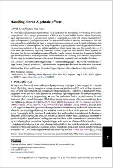

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1 INTRODUCTION

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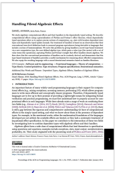
Handling Fibred Algebraic Effects

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Algebraic Effects

and

Effect Handlers

What can we do?

How to do it?

Outline

- Setting the scene
  - **Algebraic effects** and their **handlers**
  - An effectful dependently typed **core calculus** (FoSSaCS'16)

    [A., Ghani, Plotkin'16]
- What can we gain from handlers + dependent types?
  - Modular programming with handlers + expressiveness of d. types
  - **Extrinsic reasoning** about effectful computations
- Extending the FoSSaCS'16 calculus with alg. effects and handlers
  - Take 1: The common **term-level def.** of handlers (has issues)
  - Take 2: A new **type-level treatment** of handlers

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Algebraic effects

- Moggi taught us to model comp. effects using **monads** $(T, \eta, (-)^\dagger)$

$$\eta_A : A \to TA \qquad (f : A \to TB)^\dagger_{A,B} : TA \to TB$$

- Plotkin and Power showed that most of these monads arise from
  - **operation symbols** – representing the **sources** of effects

$$\mathsf{raise} : \mathsf{Exc} \longrightarrow 0 \qquad \mathsf{get} : \mathsf{Loc} \longrightarrow \mathsf{Val} \qquad \mathsf{put} : \mathsf{Loc} \times \mathsf{Val} \longrightarrow 1$$

  - **equations** – describing the computational **behaviour**

$$\ell{:}\mathsf{Loc} \mid w{:}1 \vdash \mathtt{get}_\ell(x.\mathtt{put}_{\langle \ell, x\rangle}(w(\star))) = w(\star)$$

- The algebraic approach significantly simplifies
  - **choosing** a monad/adjunction to model a given language
  - modelling **combinations** of two or more comp. effects
  - **generic** effectful programming (via **handlers**)

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Handlers of algebraic effects

- Plotkin and Pretnar's **handlers** of algebraic effects
  - generalisation of exception handlers
  - given by **redefining** the given ops. (handlers denote **algebras**)
  - many uses – stream redirection, state, rollbacks, concurrency, ...

- Usually included in languages using the **handling** construct

$$M \ \texttt{handled with} \ \{\mathrm{op}_{x_v}(x_k) \mapsto N_{\mathrm{op}}\}_{\mathrm{op} \in \mathcal{S}_{\mathrm{eff}}} \ \texttt{to} \ y\!:\!A \ \texttt{in}_{\underline{C}} \ N_{\mathrm{ret}}$$

  interpreted using the **homomorphism** $FA \longrightarrow \langle U\underline{C}, \overrightarrow{N_{\mathrm{op}}}\rangle$, i.e.,

$$(\mathrm{op}_V(y.M)) \ \texttt{handled with} \ \{\ldots\}_{\mathrm{op} \in \mathcal{S}_{\mathrm{eff}}} \ \texttt{to} \ y\!:\!A \ \texttt{in}_{\underline{C}} \ N_{\mathrm{ret}}$$
$$=$$
$$N_{\mathrm{op}}[V/x_v][\lambda\, y\!:\!O \,.\, \texttt{thunk}\,(M \ \texttt{handled with} \ \ldots)/x_k]$$

  and

$$(\texttt{return}\ V) \ \texttt{handled with} \ \{\ldots\}_{\mathrm{op} \in \mathcal{S}_{\mathrm{eff}}} \ \texttt{to} \ y\!:\!A \ \texttt{in}_{\underline{C}} \ N_{\mathrm{ret}} \ = \ N_{\mathrm{ret}}[V/y]$$

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A core dependently typed effectful calculus

- Natural extension of Martin Löf's (intensional) type theory
  - clear distinction between **values** and **computations** (CBPV, EEC)

- **Value types** ($\Gamma \vdash A$) and **computation types** ($\Gamma \vdash \underline{C}$)

$$A, B ::= \ldots \mid U\underline{C} \qquad \underline{C}, \underline{D} ::= FA \mid \Pi x{:}A \,.\, \underline{C} \mid \boxed{\Sigma x{:}A \,.\, \underline{C}}$$

- **Value terms** ($\Gamma \vdash V : A$)

$$V, W ::= \ldots \mid \mathtt{thunk}\, M$$

- **Computation terms** ($\Gamma \vdash M : \underline{C}$)

$$M, N ::= \mathtt{return}\, V \mid M \;\mathtt{to}\; x{:}A \;\mathtt{in}_{\underline{C}}\; N \mid \lambda x{:}A \,.\, M \mid M\, V \mid \langle V, M \rangle \mid \boxed{M \;\mathtt{to}\; (x{:}A, z{:}\underline{C}) \;\mathtt{in}_{\underline{D}}\; K} \mid \mathtt{force}_{\underline{C}}\, V$$

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The calculus we work in this paper

- We work in an extension of the FoSSaCS'16 calculus, with
  - a Tarski-style **value universe** $\mathcal{U}$
    - with **codes** written as $\widehat{\Pi}$, $\widehat{\Sigma}$, $\widehat{0}$, $\widehat{1}$, ...
    - but thinking of them as $\forall$, $\exists$, $\bot$, $\top$, ...
  - fibred **algebraic effects**
    - dep. typed **operation symbols** $\mathsf{op} : (x_v : I) \longrightarrow O$
    - ops. determine **computation terms** $\mathsf{op}^{\underline{C}}_V(y : O[V/x_v] \,.\, M)$
    - effect equations determine **definitional equations**
  - a **derivable** "into-comps." variant of **handlers** and **handling**

$$M \;\texttt{handled with}\; \{\mathsf{op}_{x_v}(x_k) \mapsto N_{\mathsf{op}}; \overrightarrow{W_{\mathsf{eq}}}\}_{\mathsf{op} \in \mathcal{S}_{\mathsf{eff}}} \;\texttt{to}\; y : A \;\texttt{in}_{\underline{C}}\; N_{\mathsf{ret}}$$

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Reasoning about effectful computations

- Handlers are useful for **extrinsic reasoning**!

- They help us to reason about effectful computations $M : FA$
  - Can be used to define **predicates** $P : UFA \to \mathcal{U}$ by
    1) equipping $\mathcal{U}$ (or a resp. type) with an **algebra** structure
    2) **handling** the given computation using that algebra
  - Intuitively, $P\ (\texttt{thunk}\, M)$ computes a **proof obligation** for M
  - We discuss **three examples** of such predicates

- Also, can be an alternative to mon. reification for **rel. reasoning**
  - E.g., relating **stateful comps.** $M, N : FA$ as **functions** $S \to A \times S$
  - Not touched upon in this paper
  - See [Grimm et al.'18] for **reification-based** relational reasoning

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- Handlers are useful for **extrinsic reasoning**!

- They help us to reason about effectful computations $M : FA$
  - Can be used to define **predicates** $P : UFA \to \mathcal{U}$ by
    1) equipping $\mathcal{U}$ (or a resp. type) with an **algebra** structure
    2) **handling** the given computation using that algebra
  - Intuitively, $P\ (\mathtt{thunk}\, M)$ computes a **proof obligation** for M
  - We discuss **three examples** of such predicates

- Also, can be an alternative to mon. reification for **rel. reasoning**
  - E.g., relating **stateful comps.** $M, N : FA$ as **functions** $S \to A \times S$
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Ex1: Lifting predicates to effectful comps.

- Given a predicate $P : A \to \mathcal{U}$ on **return values**,

  we define a predicate $\Box P : UFA \to \mathcal{U}$ on **(I/O)-comps.** as

$$\Box P \stackrel{\text{def}}{=} \lambda\, y : UFA \,.\, (\texttt{force}\, y)\ \texttt{handled with}\ \{\ldots\}_{\mathsf{op} \in \mathcal{S}_{\mathsf{I/O}}}\ \texttt{to}\ y' : A\ \texttt{in}_{\mathcal{U}}\ P\, y'$$

  using the **handler** given by

$$\begin{array}{lll} \mathsf{read}(x_k) & \mapsto & \widehat{\Pi}\, y : \mathsf{El}(\widehat{\mathsf{Chr}}) \,.\, x_k\ y & \qquad (\text{where } x_k : \mathsf{Chr} \to \mathcal{U}) \\ \mathsf{write}_{x_v}(x_k) & \mapsto & x_k\ \star & \qquad (\text{where } x_v : \mathsf{Chr},\ x_k : 1 \to \mathcal{U}) \end{array}$$

- $\Box P$ is similar to the **necessity modality** from Evaluation Logic

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Ex2: Dijkstra's weakest precondition sem.

- Given a postcondition on **return values** and **final states**

$$Q : A \to S \to \mathcal{U} \qquad\qquad (S \stackrel{\text{def}}{=} \Pi\,\ell\,{:}\,\mathsf{Loc}\,.\,\mathsf{Val}(\ell))$$

  we define a precondition for **stateful comps.** on **initial states**

$$\mathsf{wp}_Q : UFA \to S \to \mathcal{U}$$

  by

  1) handling the given comp. into a **state-passing function** using

$$V_{\mathsf{get}}\,, V_{\mathsf{put}} \quad \text{on} \quad S \to \mathcal{U} \times S \qquad\qquad \text{and} \qquad\qquad V_{\mathsf{ret}}\ \text{“}{=}\text{”}\ Q$$

  2) feeding in the **initial state**; and **3)** projecting out the **value of** $\mathcal{U}$

- Then, wp_Q satisfies the **expected properties**, such as

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Ex3: Allowed patterns of (I/O)-effects

- Assuming an inductive type of **I/O-protocols**, given by

$$\mathtt{e} : \text{Protocol} \qquad \mathtt{r} : (\text{Chr} \to \text{Protocol}) \to \text{Protocol}$$

$$\mathtt{w} : (\text{Chr} \to \mathcal{U}) \times \text{Protocol} \to \text{Protocol}$$

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Outline

- Setting the scene
  - **Algebraic effects** and their **handlers**
  - An effectful dependently typed **core calculus** (FoSSaCS'16)

    [A., Ghani, Plotkin'16]
- What can we gain from handlers + dependent types?
  - Modular programming with handlers + expressiveness of d. types
  - **Extrinsic reasoning** about effectful computations
- Extending the FoSSaCS'16 calculus with alg. effects and handlers
  - Take 1: The common **term-level def.** of handlers (has issues)
  - Take 2: A new **type-level treatment** of handlers

Extending the FoSSaCS'16 calculus

- We assume given a **fibred effect theory** $\mathcal{T} = (\mathcal{S}, \mathcal{E})$
- First, we extend the calculus with **algebraic effects** as follows:
  - we extend the **computation terms** with

$$M, N \ ::= \ \ldots \ \mid \ \mathsf{op}^{\underline{C}}_{V}(y \,{:}\, O[V/x_v] \,.\, M) \qquad (\mathsf{op} : (x_v \,{:}\, I) \longrightarrow O \in \mathcal{S})$$

  - we extend the **equational theory** with equations given in $\mathcal{E}$
  - we capture the **interaction** of comp. terms and ops. with the eq.

$$\frac{\Gamma \vdash V : I \quad \Gamma, x \,{:}\, O[V/x_v] \vdash M : \underline{C} \quad \Gamma \mid z \,{:}\, \underline{C} \vdash K : \underline{D}}{\Gamma \vdash K[\mathsf{op}^{\underline{C}}_{V}(x.M)/z] = \mathsf{op}^{\underline{D}}_{V}(x.K[M/z]) : \underline{D}} \ (\mathsf{op} : (x_v \,{:}\, I) \longrightarrow O \in \mathcal{S})$$

- Second, we extend the calculus with a support for **handlers** ...

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$$\Gamma \vdash \mathtt{write}_{\mathsf{a}}(\mathtt{return}\,\star) = \mathtt{write}_{\mathsf{z}}(\mathtt{return}\,\star) : F1$$

- At a very high-level, the problem is (see the paper for details)
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How to proceed?

- Possible ways to solve this unsoundness problem
  - **Option 1: Change** the FoSSaCS'16 calculus
    - change the equational theory of homomorphism terms
    - hom. terms would not denote homomorphisms any more
    - investigated for exceptions in CBPV with stacks by [Levy'06]
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    - extend it so that handling for comp. terms is derivable
    - while making sure that the calculus remains sound
    - **key idea:** comp. types and handlers both denote **algebras**
    - extended calculus admits a natural denotational semantics

Take 2: A type-level treatment of handlers

- Instead, we extend the FoSSaCS'16 **computation types** with
  - a **user-defined algebra type**

    $$\underline{C}\,,\underline{D} \;::=\; \ldots \;\mid\; \langle A; \overrightarrow{V_{\mathrm{op}}}; \overrightarrow{W_{\mathrm{eq}}} \rangle$$

    where

    - A is the **carrier** value type
    - $\overrightarrow{V_{\mathrm{op}}}$ is a set of user-defined **operations**
    - $\overrightarrow{W_{\mathrm{eq}}}$ is a set of **witnesses** of equational proof obligations

- As a result, we can derive the **handing construct** as

$$M \texttt{ handled with } \{\mathtt{op}_{x_v}(x_k) \mapsto N_{\mathrm{op}}; \overrightarrow{W_{\mathrm{eq}}}\}_{\mathtt{op} \in \mathcal{S}_{\mathrm{eff}}} \texttt{ to } y\!:\!A \texttt{ in}_{\underline{C}}\ N_{\mathrm{ret}}$$

$$\stackrel{\mathrm{def}}{=}$$

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  and similarly for the "**into-values**" variant of it

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Conclusion

- In conclusion
  - handlers are natural for defining **predicates on computations**
    - lifting predicates from return values to computations
    - Dijkstra's weakest precondition semantics of state
    - specifying patterns of allowed (I/O)-effects
  - they admit a principled **type-based treatment**

- See the paper for
  - **formal details** of what I have shown you today
  - families fibrations based **denotational semantics**
  - discussion about the calculus's inherent **extensional nature**
  - **Agda code** for the example predicates $P : UFA \to \mathcal{U}$

Thank you!

Questions?