A fibrational view on computational effects

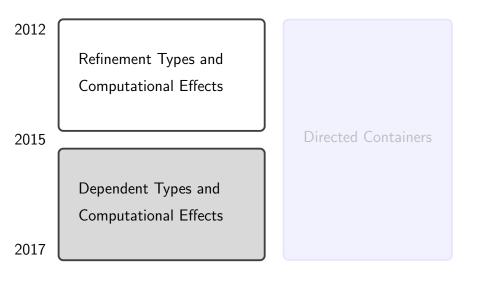
(or some things I did during my PhD)

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Prosecco Team, Inria Paris

Edinburgh, 28 November 2017

Outline – what I did during my PhD



Outline – dependent types

The Curry-Howard correspondence:

```
\begin{array}{lll} \text{Simple Types} & \sim & \text{Propositional Logic} & & (\text{Nat}, \text{String}, \ldots) \\ \\ \text{Dependent Types} & \sim & \text{Predicate Logic} & & (\Sigma, \Pi, =, \ldots) \end{array}
```

A tiny example: we can use dep. types to express sorted lists

$$\ell \colon (\mathsf{List} \; \mathsf{Nat}) \vdash \mathsf{Sorted}(\ell) \;\; \stackrel{\ \, \mathsf{def}}{=} \;\; \stackrel{\forall}{\mathsf{\Pi}} \; i \colon \mathsf{Nat} \; . \; (0 < i < \mathsf{len} \; \ell) \; \rightarrow \; (\ell[i \text{-}1] \leq \ell[i])$$

which in turn could be used to type a sorting function

$$\forall$$
 sort : $\Pi \ell$: (List Nat) . $\Sigma \ell'$: (List Nat) . (Sorted $(\ell) \times \dots$)

Large examples: CompCert (Coq), miTLS and HACL* (F*), ...

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\ell: (List Nat) \vdash Sorted(\ell) \stackrel{\text{def}}{=} \Pi i: Nat. (0 < i < \text{len } \ell) \rightarrow (\ell[i-1] \le \ell[i]) which in turn could be used to type a sorting function
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\forall sort : \Pi \ell: (List Nat). \Sigma \ell': (List Nat). \left( \text{Sorted}(\ell) \times \dots \right)
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Outline – computational effects

Examples:

- state
- exceptions
- nondeterminism
- I/O
- . . .

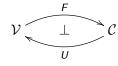
Meta-languages and models:

• based on monads (T, η, μ)

(Moggi)

(Levy)

based on adjunctions



based on algebraic presentations

(Plotkin and Power)

get : $1 \rightarrow S$ put : $S \rightarrow 1$ + equations

We investigate the combination of

```
• dependent types  (\Pi, \Sigma, V =_{\mathcal{A}} W, ...)
```

• computational effects (state, nondeterminism, I/O, ...)

Two guiding problems

- effectful programs in types (e.g., get and put in types)
- types of effectful programs (e.g., of sequential composition)

- tell a mathematically natural story
- use established math. techniques
- cover a wide range of comp. effects
- discover smth. interesting

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- cover a wide range of comp. effects (alg. effects, continuations)
- discover smth. interesting (using handlers to reason about effects)

(type-dependency in the presence of effects)

Q: Should we allow situations such as Sorted[receive(y.M)/ ℓ]?

A1: In this talk, we say not directly

- types should only depend on static information about effects
- we allow dependency on effectful comps. via analysing thunks

A2: But we are also looking into the direct case

- type-dependency needs to be "homomorphic", but not only so
- intuitively, lift Sorted(ℓ) to Sorted[†](c), where c: T(List Chr)

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Aim: Types should only depend on static info about effects

Solution: CBPV/EEC style distinction between vals. and comps.

- value types $\Gamma \vdash A$ (MLTT + thunks + ...)
- computation types $\Gamma \vdash \underline{C}$ (dep. typed CBPV/EEC
- where Γ contains only value variables $x_1: A_1, \ldots, x_n: A_n$

Could have also considered Moggi's λ_{ML} and Levy's FGCBV

- building on CBPV/EEC gives a more general story
- especially for the treatment of sequential composition
- and also for integrating dependent- and effect-typing (ongoing)

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(e.g., sequential composition)

The problem: The standard typing rule for seq. composition

$$\frac{\Gamma \vdash_{\overline{c}} M : FA \qquad \Gamma, x : A \vdash_{\overline{c}} N : \underline{C}}{\Gamma \vdash_{\overline{c}} M \text{ to } x : A \text{ in } N : \underline{C}}$$

is not correct any more because x can appear free in the type

(

in the conclusion

Aim: To fix the typing rule of sequential composition

Option 1: We could restrict the free variables in \underline{C} : [Levy'04] $\underline{\Gamma \vDash M : FA \qquad \Gamma \vdash \underline{C} \qquad \Gamma, x : A \vDash N : \underline{C}}$

But: sometimes it is useful if \underline{C} can depend on x!

if we consider

fopen (return true, return false) to x: Bool in N

• then it would be natural to let \underline{C} depend on x, e.g.,

 $x: Bool \vdash \underline{C}(x) \stackrel{\text{def}}{=} \text{if } x \text{ then "allow fread, fwrite, and fclose"}$ else "allow fopen"

(needs more expressive comp. types than we consider here)

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But: then comp. types would be singleton-like!?!

However, smth. like this is probably needed for the direct case.

Option 3: In the monadic metalanguage λ_{ML} , one could try

$$\frac{\Gamma \vdash M : TA \qquad \Gamma, x : A \vdash N : TB(x)}{\Gamma \vdash M \text{ to } x : A \text{ in } N : T(\Sigma x : A.B)}$$

But: what makes this a principled solution? Why is it correct?

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Option 4: We draw inspiration from algebraic effects
and combine it with restricting <u>C</u> in seq. comp. (Option 1)

E.g., consider the non-det. program (for $x : \text{Nat } \vdash N : \underline{C}(x)$) $M \stackrel{\text{def}}{=} \text{choose} (\text{return 4}, \text{return 2}) \text{ to } x : \text{Nat in } N$

After tossing the coin, this program evaluates as either N[4/x] : $\underline{C}[4/x]$ or N[2/x] : $\underline{C}[2/x]$

Idea: M denotes an element of the coproduct of algebras $\underline{C}[4/x] + \underline{C}[2/x] \stackrel{\text{def}}{=} F\left(U\left(\underline{C}[4/x]\right) + U\left(\underline{C}[2/x]\right)\right)_{/\equiv}$

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Putting these ideas together

(eMLTT: a core dep.-typed language with comp. effects)

eMLTT – value and comp. types

Value types: MLTT + thunks + ...

$$A, B ::=$$
Nat $\mid 1 \mid 0 \mid \Pi x : A.B \mid \Sigma x : A.B \mid V =_A W \mid U \subseteq | \dots |$

• $U \subseteq C$ is the type of thunked (i.e., suspended) computations

Computation types: dep.-typed version of EEC's comp. types

$$\underline{C}, \underline{D} ::= FA \mid \Pi x : A . \underline{C} \mid \Sigma x : A . \underline{C}$$

- F A is the type of computations returning values of type A
- Πx: A.C is the type of dependent effectful functions
 - generalises CBPV/EEC's comp. types $A \to \underline{C}$ and $\underline{C} \times \underline{D}$
- $\Sigma x: A.\underline{C}$ is the type of dep. pairs of values and effectful comps.
 - captures the intuition about seq. comp. and coprods. of algebras
 - generalises EEC's comp. types $!A \otimes C$ and $C \oplus D$

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eMLTT – value and comp. terms

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Value terms: MLTT + thunks + ... V, W ::= x \mid zero \mid succ V \mid ... \mid thunk M \mid ...
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equational theory based on intensional MLTT

Comp. terms: dep.-typed version of CBPV/EEC's comp. terms

```
\begin{array}{lll} M,N ::= & \operatorname{force} V \\ & | & \operatorname{return} V \\ & | & M \operatorname{to} x : A \operatorname{in} N \\ & | & \lambda x : A . M \\ & | & MV \\ & | & \langle V,M \rangle & (\operatorname{comp.} \Sigma \operatorname{intro.}) \\ & | & M \operatorname{to} \langle x : A,z : \underline{C} \rangle \operatorname{in} K & (\operatorname{comp.} \Sigma \operatorname{elim.}) \end{array}
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But: Value and comp. terms alone do not suffice, as in EEC!

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eMLTT - homomorphism terms

Note: We need to define K in such a way that the intended left-to-right evaluation order is preserved, e.g., consider

$$\Gamma \vdash_{\overline{c}} \langle V, M \rangle \text{ to } \langle x : A, \underline{z} : \underline{C} \rangle \text{ in } \underline{K} = \underline{K}[V/x, M/\underline{z}] : \underline{D}$$

Homomorphism terms: dep.-typed version of EEC's linear terms

$$K, L := z$$
 (linear comp. vars.)
 $\mid K \text{ to } x : A \text{ in } M$
 $\mid \lambda x : A, K$
 $\mid KV$
 $\mid \langle V, K \rangle$ (comp. $\Sigma \text{ intro.}$)
 $\mid K \text{ to } \langle x : A, z : C \rangle \text{ in } L$ (comp. $\Sigma \text{ elim.}$)

Typing judgments:

- Γ ⋈ V : A
- Γ la M : C
- $\Gamma \mid z : \underline{C} \mid_{\overline{h}} K : \underline{D}$ (linear in z; comp. bound to z happens first

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 to $\langle x : A, z : \underline{C} \rangle$ in $K = K[V/x, M/z] : \underline{D}$

Homomorphism terms: dep.-typed version of EEC's linear terms

```
\begin{array}{lll} \textit{K}, \textit{L} ::= & \textit{z} & \text{(linear comp. vars.)} \\ & | & \textit{K} \text{ to } x : \textit{A} \text{ in } \textit{M} \\ & | & \lambda x : \textit{A} . \textit{K} \\ & | & \textit{K} \textit{V} \\ & | & \langle \textit{V}, \textit{K} \rangle & \text{(comp. } \Sigma \text{ intro.)} \\ & | & \textit{K} \text{ to } \langle x : \textit{A}, \textit{z} : \underline{\textit{C}} \rangle \text{ in } \textit{L} & \text{(comp. } \Sigma \text{ elim.)} \end{array}
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eMLTT – typing sequential composition

We can then account for type-dependency in seq. comp. as

$$\frac{\Gamma, x : A \vDash N : \underline{C}(x)}{\Gamma \vDash M : FA} \frac{\Gamma, x : A \vDash N : \underline{C}(x)}{\Gamma, x : A \vDash \langle x, N \rangle : \Sigma x : A . \underline{C}(x)}$$
$$\Gamma \vDash M \text{ to } x : A \text{ in } \langle x, N \rangle : \Sigma x : A . \underline{C}(x)$$

The seq. comp. rule for $\lambda_{\rm ML}$ is justified by the type isomorphism

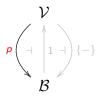
$$\frac{\Gamma \vdash A \qquad \Gamma, x : A \vdash B(x)}{\Gamma \vdash U(\Sigma x : A . F(B)) \cong UF(\Sigma x : A . B) = T(\Sigma x : A . B)}$$

Categorical semantics of eMLTT

(fibrations + adjunctions)

Fibred adjunction models – value part

Given by a split closed comprehension category p, as in

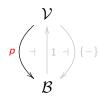


such that

- p has split fibred strong colimits of shape 0 and 2
 - (in thesis, also Jacobs-style axiomatisation for arbitrary shapes)
 - (all one needs are cocones and fully-faithfulness of induced func.)
- p has weak split fibred strong natural numbers
- p has split intensional propositional equality

Fibred adjunction models – value part

Given by a split closed comprehension category p, as in

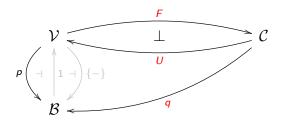


allowing us to define a partial interpretation fun. [-], that maps:

- a context Γ to and object $\llbracket \Gamma \rrbracket$ in \mathcal{B} , with
 - $\bullet \quad \llbracket \diamond \rrbracket \stackrel{\mathsf{def}}{=} 1$
 - $\llbracket \Gamma, x : A \rrbracket \stackrel{\text{def}}{=} \{ \llbracket \Gamma; A \rrbracket \}$ (if $x \notin Vars(\Gamma)$ and $\llbracket \Gamma; A \rrbracket$ is defined)
- a context Γ and a value type A to an object $[\![\Gamma;A]\!]$ in $\mathcal{V}_{[\![\Gamma]\!]}$
- a context Γ and a value term V to $[\![\Gamma;V]\!]:1_{[\![\Gamma]\!]}\longrightarrow A$ in $\mathcal{V}_{[\![\Gamma]\!]}$

Fibred adjunction models – effects part

Given by a split fibration q and a split fib. adjunction $F \dashv U$, as in



such that

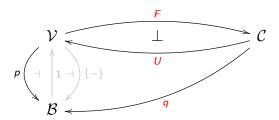
- q has split dependent p-products (comp. ∏-type; r. adj. to wk.)
- q has split dependent p-coproducts (comp. Σ -type; I. adj. to wk.)

and to account for the full calculus presented in the thesis,

• q admits split fibred pre-enrichment in p (hom. function type $-\circ$)

Fibred adjunction models - effects part

Given by a split fibration q and a split fib. adjunction $F \dashv U$, as in



we extend the partial interpretation fun. [-] so that it maps:

- a ctx. Γ and a comp. type \underline{C} to an object $\llbracket \Gamma ; \underline{C} \rrbracket$ in $\mathcal{C}_{\llbracket \Gamma \rrbracket}$
- a ctx. Γ and a comp. term M to $\llbracket \Gamma; M \rrbracket : 1_{\llbracket \Gamma \rrbracket} \longrightarrow U(\underline{C})$ in $\mathcal{V}_{\llbracket \Gamma \rrbracket}$
- a ctx. Γ , a comp. var. z, a comp. type \underline{C} , and a hom. term K to $[\![\Gamma;z:\underline{C};K]\!]:[\![\Gamma;\underline{C}]\!]\longrightarrow \underline{D}$ in $\mathcal{C}_{[\![\Gamma]\!]}$

Fibred adjunction models – correctness

Theorem (Soundness):

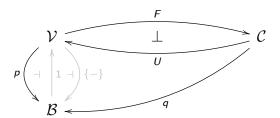
- If $\Gamma \vdash \underline{C}$, then $[\![\Gamma;\underline{C}]\!] \in \mathcal{C}_{[\![\Gamma]\!]}$
- $\bullet \ \, \text{If} \,\, \Gamma \models M : \underline{C}, \, \text{then} \,\, [\![\Gamma;M]\!] : 1_{[\![\Gamma]\!]} \longrightarrow \mathit{U}([\![\Gamma;\underline{C}]\!])$
- If $\Gamma \mid z : \underline{C} \models K : \underline{D}$, then $\llbracket \Gamma; z : \underline{C}; K \rrbracket : \llbracket \Gamma; \underline{C} \rrbracket \longrightarrow \llbracket \Gamma; \underline{D} \rrbracket$
- $\bullet \ \ \mathsf{If} \ \Gamma \vdash \underline{\mathcal{C}} = \underline{\mathcal{D}}, \ \mathsf{then} \ \llbracket \Gamma ; \underline{\mathcal{C}} \rrbracket = \llbracket \Gamma ; \underline{\mathcal{D}} \rrbracket \in \mathcal{C}_{\llbracket \Gamma \rrbracket}$
- ...

Theorem (Classifying model):

• The well-formed syntax of eMLTT forms a fib. adjunction model.

Theorem (Completeness):

• If two types or terms are equal in all fibred adjunction models, then they are also equal in the equational theory of eMLTT.



Example 1 (identity adjunctions):

• sound as long as we haven't included any actual comp. effects

Example 2 (simple fibrations from enriched adj. models of EEC):

• doesn't support any real type dependency (constant families

Example 3 (families fibrations and lifting of adjunctions):

- $\bullet \ (\llbracket \Gamma \rrbracket, \llbracket A \rrbracket) \in \mathsf{Fam}(\mathsf{Set}) \qquad \qquad (\mathsf{where} \ \llbracket A \rrbracket \in \llbracket \Gamma \rrbracket \longrightarrow \mathsf{Set})$
- $\bullet \ \, \big(\big[\![\Gamma]\!], \big[\![\underline{C}]\!] \big) \in \mathsf{Fam}(\mathcal{D}) \qquad \qquad (\mathsf{where} \ \big[\![\underline{C}]\!] \in \big[\![\Gamma]\!] \longrightarrow \mathcal{D})$

Example 4 (continuous families and CPO-enriched monads):

- $(\llbracket \Gamma \rrbracket, \llbracket A \rrbracket) \in \mathsf{CFam}(\mathsf{CPO})$ (where $\llbracket A \rrbracket \in \llbracket \Gamma \rrbracket \longrightarrow \mathsf{CPO}^{\mathit{EP}}$)
- $(\llbracket \Gamma \rrbracket, \llbracket \underline{C} \rrbracket) \in \mathsf{CFam}(\mathsf{CPO}^\mathsf{T})$ (where $\llbracket \underline{C} \rrbracket \in \llbracket \Gamma \rrbracket \longrightarrow (\mathsf{CPO}^\mathsf{T})^{EP} \rrbracket$
- Theorem: cod_{CPO} not suitable because CPO not a LCCC.

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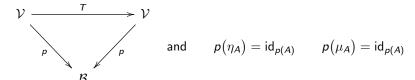
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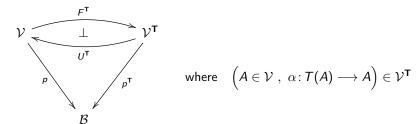
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Example 5 (EM-resolutions of split fibred monads):

• given a split fibred monad $\mathbf{T} = (T, \eta, \mu)$ on p, i.e.,



we consider models based on the EM-resolution of T



and show that three familiar results hold for this situation

Example 5 (EM-resolutions of split fibred monads):

• Theorem 1: If p supports Π -types, then p^{T} also supports Π -types

$$\Pi_A^{\mathsf{T}}(B,\beta) \ \stackrel{\scriptscriptstyle\mathsf{def}}{=} \ \left(\Pi_A(B),\beta_{\Pi_A^{\mathsf{T}}}\right)$$

Prop.: Every T on a split closed comp. cat. has a dep. strength

$$\sigma_A: \Sigma_A \circ T \longrightarrow T \circ \Sigma_A \qquad (A \in \mathcal{V})$$

• Theorem 2: If σ_A are natural isos., then p^T supports Σ -types

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Theorem 3: If p supports Σ-types and p^T has split fibred reflexive coequalizers, then p^T also supports Σ-types

$$\Sigma_A^{\mathsf{T}}(B,\beta) \stackrel{\text{def}}{=} F^{\mathsf{T}}(\Sigma_A(B))_{/\equiv}$$

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Algebraic effects

Fibred effect theories \mathcal{T}_{eff} :

signatures of dep. typed operation symbols

$$\frac{\cdot \vdash I \qquad x_i \colon I \vdash O \qquad I \text{ and } O \text{ are pure value types}}{\text{op} \colon (x_i \colon I) \longrightarrow O}$$

equipped with equations on derivable effect terms

In eMLTT:

$$M ::= \dots \mid \operatorname{op}_{V}^{C}(x.M)$$

General algebraicity equations (in addition to eff. th. eqs.):

$$\frac{\Gamma \trianglerighteq V: I - \Gamma, x: O[V/x_i] \trianglerighteq M: \underline{C} - \Gamma | z: \underline{C} \trianglerighteq_{\mathbb{F}} K: \underline{D}}{\Gamma \trianglerighteq K[\operatorname{op}_V^{\underline{C}}(x.M)/z] = \operatorname{op}_V^{\underline{D}}(x.K[M/z]): \underline{D}} (\operatorname{op}: (x_i: I) \longrightarrow O)$$

Sound semantics: based on

•
$$p : \mathsf{Fam}(\mathsf{Set}) \longrightarrow \mathsf{Set}$$
 and $q : \mathsf{Fam}(\mathsf{Mod}(\mathcal{L}_{\mathcal{T}_{\mathsf{eff}}}, \mathsf{Set})) \longrightarrow \mathsf{Set}$

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Algebraic effects – examples

Example 1 (interactive I/O):

- ullet read : $1 \longrightarrow \mathsf{Chr}$ $(\mathsf{Chr} \stackrel{\mathsf{def}}{=} 1 + \ldots + 1)$ write : $\mathsf{Chr} \longrightarrow 1$
- no equations

Example 2 (global state with location-dependent store type):

```
• \diamond \vdash \mathsf{Loc}

\ell : \mathsf{Loc} \vdash \mathsf{Val}

\diamond \vdash \mathsf{v} \mathsf{isDec}_{\mathsf{Loc}} : \Pi\ell : \mathsf{Loc} . \Pi\ell' : \mathsf{Loc} . (\ell =_{\mathsf{Loc}} \ell') + (\ell =_{\mathsf{Loc}} \ell' \to 0)
```

- get : $(\ell : \mathsf{Loc}) \longrightarrow \mathsf{Val}$ put : $(\Sigma \ell : \mathsf{Loc.Val}) \longrightarrow 1$
- five equations (two of them branching on isDecLoc

Example 3 (dep. typed update monads $TX \stackrel{\text{def}}{=} \Pi_{s:S}$. $Ps \times X$)

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Handlers of algebraic effects

(for programming and extrinsic reasoning)

Handlers of alg. effects – for programming

Idea: Generalisation of exception handlers [Plotkin,Pretnar'09]

Handler = Algebra and Handling = Homomorphism

Usual term-level presentation:

Fig. M handled with $\{\operatorname{op}_{\mathsf{x}_{\mathsf{v}}}(\mathsf{x}_k)\mapsto \mathsf{N}_{\operatorname{op}}\}_{\operatorname{op}\in\mathcal{T}_{\operatorname{eff}}}$ to y:A in C \mathbb{C} satisfying

(return V) handled with $\{...\}_{\mathsf{op}\in\mathcal{T}_{\mathsf{eff}}}$ to $y\!:\!A$ in $N_{\mathsf{ret}} = N_{\mathsf{ret}}[V/x]$ (op $\frac{C}{V}(x.M)$) handled with $\{...\}_{\mathsf{op}\in\mathcal{T}_{\mathsf{eff}}}$ to $y\!:\!A$ in $N_{\mathsf{ret}} = N_{\mathsf{op}}[V/x_V][.../x_k]$

Typical use case for programming:

- write your programs using alg. ops. (e.g., get and put)
- use handlers to provide fit-for-purpose impl. (e.g., $S \to X \times S$)

Handlers of alg. effects – for programming

Usual term-level presentation:

```
\Gamma \vdash M \text{ handled with } \{ \operatorname{op}_{\mathsf{x}_{\mathsf{v}}}(\mathsf{x}_{\mathsf{k}}) \mapsto \mathsf{N}_{\operatorname{op}} \}_{\operatorname{op} \in \mathcal{T}_{\operatorname{eff}}} \text{ to } y : A \operatorname{in}_{\underline{C}} \mathsf{N}_{\operatorname{ret}} : \underline{C} satisfying
```

```
\begin{split} & (\text{return } V) \text{ handled with } \{...\}_{\mathsf{op} \in \mathcal{T}_{\mathsf{eff}}} \text{ to } y \colon\! A \text{ in } N_{\mathsf{ret}} = N_{\mathsf{ret}}[V/x] \\ & (\mathsf{op}_V^{\underline{C}}(x.M)) \text{ handled with } \{...\}_{\mathsf{op} \in \mathcal{T}_{\mathsf{eff}}} \text{ to } y \colon\! A \text{ in } N_{\mathsf{ret}} = N_{\mathsf{op}}[V/x_v][.../x_k] \end{split}
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satisfying  (\text{return } V) \text{ handled with } \{...\}_{\mathsf{op} \in \mathcal{T}_{\mathsf{eff}}} \text{ to } y \colon A \text{ in } N_{\mathsf{ret}} = N_{\mathsf{ret}}[V/x]
```

 $\Gamma \vdash M$ handled with $\{ op_{x_k}(x_k) \mapsto N_{op} \}_{op \in \mathcal{T}_{eff}}$ to $y : A \text{ in }_C N_{ret} : \underline{C}$

 $(\operatorname{op}_{V}^{\underline{C}}(x.M)) \text{ handled with } \{...\}_{\operatorname{op}\in\mathcal{T}_{\operatorname{eff}}} \text{ to } y:A \text{ in } N_{\operatorname{ret}} = N_{\operatorname{op}}[V/x_{v}][.../x_{k}]$

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- write your programs using alg. ops. (e.g., get and put)
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Idea: Using a derived handle-into-values handling construct

M handled with $\{\operatorname{op}_{x_v}(x_k)\mapsto V_{\operatorname{op}}\}_{\operatorname{op}\in\mathcal{T}_{\operatorname{eff}}}$ to $y\colon A$ in B V_{ret} can define natural predicates (essentially, dependent types)

$$\Gamma \vdash P : UFA \rightarrow \mathcal{U}$$

by

- ullet equipping a universe ${\cal U}$ with an algebra for $\mathcal{T}_{\sf eff}$, and
- using the above handle-into-values construct to define P

Note 1: P(thunk M) computes a proof obligation for M

- a universe $\mathcal U$ closed under Nat, 1, 0, +, Σ , and Π
- a type-based treatment of handlers $\underline{C} ::= \ldots \mid \langle A; \overrightarrow{V_{\sf op}}; \overrightarrow{W_{\sf eq}} \rangle$
- function extensionality (actually, it's a bit more extensional)

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Example 1 (Evaluation Logic style modalities):

- Given a predicate $P:A\to \mathcal{U}$ on return values, we define a predicate $\Diamond P:UFA\to \mathcal{U}$ on I/O-computations as
- $\Diamond P \stackrel{\text{def}}{=} \lambda x : UFA . (\text{force } x) \text{ handled with } \{...\}_{\text{op} \in \mathcal{T}_{\text{IO}}} \text{ to } y : A \text{ in}_{\mathcal{U}} P y$ using the handler given by

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$$\Diamond P\left(\texttt{thunk}\left(\texttt{read}(x\,.\,\texttt{write}_{'e'}(\texttt{return}\,V))\right)\right) = \widehat{\Sigma}\,x\,: \mathsf{El}(\widehat{\mathsf{Chr}})\,.\,P\,\,V$$

• To get the necessity modality $\Box P$, use $\widehat{\Pi} x : El(\widehat{\mathsf{Chr}})$ in V_{read}

Example 1 (Evaluation Logic style modalities):

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Example 2 (Dijkstra's weakest precondition semantics):

Given a postcondition on return values and final states

$$Q: A \to S \to \mathcal{U}$$
 $(S \stackrel{\text{def}}{=} \Pi x: \text{Loc. Val})$

we define a precondition for stateful comps. on initial states

$$\mathsf{wp}_{\mathcal{Q}}: \mathit{UFA} o \mathit{S} o \mathcal{U}$$

by

i) handling the given comp. into a state-passing function using

$$V_{
m get},\,V_{
m put}$$
 on $S o (\mathcal{U} imes S)$ and $V_{
m ret}$ " $=$ " Q

- ii) feeding in the initial state; and iii) projecting out ${\cal U}$
- Theorem: wp_Q satisfies expected properties of WPs, e.g., $\operatorname{wp}_Q\left(\operatorname{thunk}\left(\operatorname{return}V\right)\right) = \lambda x_S : S \cdot Q \cdot V \cdot x_S$

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$$wp_Q (thunk (return V)) = \lambda x_S : S . Q V x_S$$

$$wp_Q (thunk (put_{\langle \ell, V \rangle}(M))) = \lambda x_S : S . wp_Q (thunk M) (x_S[\ell \mapsto V])$$

Example 3 (Patterns of allowed I/O-effects):

Assuming an inductive type Protocol, given by

e: Protocol
$$\mathbf{r}: (\mathsf{Chr} \to \mathsf{Protocol}) \to \mathsf{Protocol}$$
 $\mathsf{w}: (\mathsf{Chr} \to \mathcal{U}) \to \mathsf{Protocol} \to \mathsf{Protocol}$

• Then, we define the predicate (rel. between comps. and protoco

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Conclusion

In work we told a mathematically natural story of combining

dependent types and computational effects

In particular, we saw

- a clean core language of dependent types and comp. effects
- a natural category-theoretic semantics
- alg. effects and handlers, in particular, for reasoning using
 - Evaluation Logic style modalities
 - Dijkstra's weakest precondition semantics of state
 - patterns of allowed (I/O)-effects

Things to look at:

- type-dependency on computations (e.g., in seq. composition)
- more expressive comp. types (par. adjunctions, Dijkstra monads)

Thank you!

D. Ahman.

Fibred Computational Effects. (PhD Thesis, 2017)

D. Ahman, N. Ghani, G. Plotkin.

Dependent Types and Fibred Computational Effects. (FoSSaCS'16)

D. Ahman.

Handling Fibred Computational Effects. (POPL'18)

Digression: dep. elimination of 0 and +

The coproduct type A + B:

[Jacobs'99]

- require $p: \mathcal{V} \longrightarrow \mathcal{B}$ to have split fibred coproducts $A +_X B$, and
- $\langle \{\mathsf{inl}_A\}^*, \{\mathsf{inr}_B\}^* \rangle : \mathcal{V}_{\{A+_XB\}} \longrightarrow \mathcal{V}_{\{A\}} \times \mathcal{V}_{\{B\}}$ to be fully-faith.
- allows one to interpret dependent case analysis, i.e.,

$$\mathcal{V}_{\{A\}}\Big(1_{\{A\}}, \{\operatorname{inl}_A\}^*(C)\Big) \times \mathcal{V}_{\{B\}}\Big(1_{\{B\}}, \{\operatorname{inr}_B\}^*(C)\Big) \\ \cong \\ \mathcal{V}_{\{A+_{X}B\}}\Big(1_{\{A+_{X}B\}}, C\Big)$$

provides semantics for

$$\frac{\Gamma, y_1 : A \trianglerighteq W_1 : C[\operatorname{inl}_A y_1/x] \quad \Gamma, y_2 : B \trianglerighteq W_2 : C[\operatorname{inr}_B y_2/x]}{\Gamma, x : A + B \trianglerighteq \operatorname{case} x \text{ of } (\operatorname{inl}(y_1) \mapsto W_1, \operatorname{inr}(y_2) \mapsto W_2) : C[\operatorname{inr}_B y_2/x]}$$

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Digression: dep. elimination of colimits

A generalisation:

[Ahman'17]

- Idea: fully-faith. for cocones $A \longrightarrow A \circledast_X B \longleftarrow B$ is enough, and we can generalise this to all split fibred colimits
- Theorem:
 - if for every object $X \in \mathcal{B}$ and diagram $J : \mathcal{D} \longrightarrow \mathcal{V}_X$ there exists a cocone $\underline{\operatorname{in}}^J : J \longrightarrow \Delta(\underline{\operatorname{colim}}(J))$ in \mathcal{V}_X ,
 - such that f*(in^J_D) = in^{f*oJ}_D, for any f : X → Y, and such that the unique mediating functor

$$\begin{split} & \langle \{\underline{\operatorname{in}}_D^J\}_{D \in \mathcal{D}}^* \rangle : \mathcal{V}_{\{\underline{\operatorname{colim}}(J)\}} \longrightarrow \operatorname{lim}(\widehat{J}) \\ & \text{s fully-faithful (for } \widehat{J} : \mathcal{D}^{op} \longrightarrow \operatorname{Cat, where } \widehat{J}(D) = \mathcal{V}_{\{J(D)\}}) \end{split}$$

then p has split fibred colimits of shape D, and
 p supports dependent elimination for them (analogously to +x)

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