

Planar Maximal Covering with Ellipses

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- Covering problems
 - Set Cover Problem
 - Maximal Covering Problem
- Maximal Covering Location Problem (MCLP)
- Planar Maximal Covering Location Problem (PMCLP)
 - One disk: $\mathcal{O}(n^2)$ and $\mathcal{O}(n^2 \log n)$ algorithms
 - m disks: $\mathcal{O}(n^{2m-1} \log n)$ algorithm
- Goals
 - Develop a $\mathcal{O}(n^2 \log n)$ algorithm for the one disk case
 - Adapt it for the m ellipses case

Norms

Let $u \in \mathbb{R}^2$ and Q a 2 by 2 positive definite matrix

- Euclidean

$$\|u\|_2 = \sqrt{u^T u}$$

- Elliptical

$$\|u\|_Q = \sqrt{u^T Q u}$$

Ellipse

Given a center $c \in \mathbb{R}^2$ and a p.d. matrix Q , an ellipse is the set of points that satisfy

$$\|u - c\|_Q = 1$$

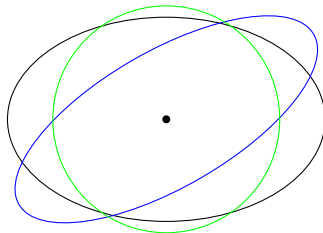
with \leq representing the set of covered points

Axis-parallel ellipse

Any 2 by 2 diagonal d.p. matrix determines an axis-parallel ellipse, which can also be described by

$$\frac{(x - c_x)^2}{a^2} + \frac{(y - c_y)^2}{b^2} = 1$$

Figure: Three curves representing the points that have distance equal to one.



Maximal Covering by Disks

One disk

$MCD(\mathcal{P}, 1)$ is the problem of placing one disk on the plane to cover a subset of a demand set \mathcal{P} maximizing the weights of the covered points.

$$\max_q w(\mathcal{P} \cap D(q)),$$

- $\mathcal{P} = \{p_1, \dots, p_n\}$ is the demand set with weights $w(p_i) > 0$
- $w(A)$, with $A \subset \mathcal{P}$ is the sum of weights of the points in A
- $D(q)$ is a disk with radius one with center at point q
- We will introduce an equivalent problem...

Maximum Weight Clique Problem

Let $\mathcal{D} = \{D_1, \dots, D_n\}$ be a set of n unit disks with weights $w_i > 0$. The maximum weight clique is defined as

$$\max_q \sum_{D_k \cap q \neq \emptyset} w_k,$$

- The disks are fixed at $\mathcal{P} = \{p_1, \dots, p_n\}$ with $w_k = w(p_k)$
- A clique is a non-empty intersection area of a subset of disks
- An optimal solution for the maximum weight clique is an optimal solution for $MCD(\mathcal{P}, 1)$.

Cobertura Maximal por Ellipses

content... trabalhos passados

Cobertura Maximal por Ellipses

m elipses

content... algoritmo

Trabalhos Futuros