Planar Maximal Covering with Ellipses

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ABSTRACT

Planar maximal covering with ellipses is an optimization problem where one wants to place ellipses on the plane to cover demand points, such that a function depending on the value of the covered points and on the cost of the ellipses that have been used is maximized. Initially, we developed an algorithm for the version of the problem where the ellipses are parallel to the coordinate axis. For the future, we intend to adapt an approximation algorithm developed for the planar maximal covering by disks and develop a method for the variant of the problem where the ellipses can be freely rotated.

Keywords: Optimization, Planar Maximal Covering Location Problem, maximal covering of points using ellipses.

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INTRODUCTION

The Minimum Cover Problem –also known as just the Set Cover Problem–, and the Maximal Covering Problem are the two main types of optimal covering problems found in the literature (KARATAS; RAZI; TOZAN, 2016).

One of the 21 Karp's NP-Complete problems[1] (KARP, 1972), the Minimum Cover Problem is very well explored and considered to be a classic. Given a demand set along with a collection of subsets of the demand set, the problem is to determine the minimum number of elements from the collection of subsets needed to cover the whole demand set. One of its most famous examples is the Minimum Vertex Cover defined over graphs, where the vertex set has to be covered by a subset of edges.

The second type of covering problems arose from the fact that covering almost all the demand set can be a lot cheaper than having to cover it all (QUILES; MARíN, 2015). This second type is known as Maximal Covering Location Problem (MCLP) and was introduced in Church and Velle (1974). In this first study, the author defined the problem on graphs, and the objective was to maximize the coverage of a demand set, which was a subset of the graph's vertices, by choosing the location of a facility set, which covered any vertex within a given coverage radius.

Just like the Minimum Cover problem, MCLP is NP-Hard (HATTA *et al.*, 2013) and both deterministic, using integer programming, and heuristic methods have been proposed to solve it. A very complete survey of developments in this area can be found in ReVelle, Eiselt and Daskin (2008).

In Church (1984) a new kind of MCLP named Planar Maximal Covering Location Problem (PMCLP) was introduced. Unlike its predecessor, this version of the problem was not defined on graphs. Instead, on PMCLP the demand set and the facilities are located in $\mathbb{R}^2$ and a facility's coverage area is determined by a distance function. Initially, the Euclidean distance was

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[1]   The decision version, which asks if there is a cover of size k, is NP-Complete.

considered and the idea behind the method proposed in Church (1984) was to convert PMCLP into an instance of MCLP and then utilize any of the previous developed exact methods to obtain a solution for PMCLP. This reduction was done by identifying a Candidate Locations Set (CLS) which represented the possible locations that needed to be evaluated for every facility, such that the optimal solution could be found. From the CLS, a network was built on which MCLP could be applied. Generating the CLS specifically for the case of Euclidean distance will be described here on Chapter 3.

Furthermore, some variations of PMCLP can also be found in the literature: in Younies and Wesolowsky (2007) PMCLP was studied under the block norm distance, in Craparo *et al.* (2019) a mean-shift algorithm for large scale[2] PMCLP was proposed, and in Bansal and Kianfar (2017) a version with partial coverage and rectangular demand and facility zones was introduced.

PMCLP under Euclidean distance is also found in the literature as the Maximum Covering by Disks (MCD) problem. Early works only tackled the one-disk version of it. In Chazelle and Lee (1986) a $\mathcal{O}(n^2)$ algorithm, which still stands as the best in terms of run-time complexity, was proposed beating the prior $\mathcal{O}(n^2 \log n)$ algorithm created by Drezner (1981). The m disks version of MCD was studied in Berg, Cabello and Har-Peled (2006) which had as its most important result a $(1-\varepsilon)$-approximation algorithm which runs in $\mathcal{O}(n \log n)$. To achieve its main goal, however, they developed a deterministic $\mathcal{O}(n^{2m-1} \log n)$ algorithm which gets employed into their approximation scheme. Additionally, in Aronov and Har-Peled (2008) one-disk maximum covering is proven to be 3SUM-HARD. This means that maximizing the number of points covered by a disk is as hard as finding three real numbers that sum to zero among n given real numbers.

Two versions of PMCLP are studied in this work, both of them related to ellipses. The first one was introduced in Canbolat and Massow (2009) and will be referred to here as Planar Maximal Covering by Ellipses (PMCE). This version does not allow the ellipses to rotate, so basically it only differs from its disks counterpart in the shape of the facility's coverage area. The second version was introduced in Andretta and Birgin (2013) and will be referred to here as Planar Maximal Covering by Ellipses with Rotation (PMCER). This version of the problem had no restriction on the rotation of ellipses–note that this is a slightly bigger modification of PMCLP compared to the first one, as not only the location for the facilities must be determined, but also their angle of rotation. The main motivation to study these two versions of PMCLP is that cellphone towers can have elliptically shaped coverage area, therefore, to determine what are the best locations to place m cellphone towers to maximize the amount of the population covered by their signal, an elliptical PMCLP is better-suited (CANBOLAT; MASSOW, 2009).

Only two articles have been found published in the literature that study this problem. In Canbolat and Massow (2009), a mixed-integer non-linear programming method was proposed as a first approach to the problem. For some instances, the method took too long and did not

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2   Numerical experiments were done for up to 3000 points.

find an optimal solution. For this reason, a heuristic method was developed using a technique called Simulated Annealing. Solutions for the instances that timed-out with the first method were then obtained. The problem was further explored in Andretta and Birgin (2013) which proposed a deterministic method that showed better performance obtaining optimal solutions for the instances, which the first method could not. Also, in Andretta and Birgin (2013), the version of the problem where the ellipses could be freely rotated was introduced and an exact method, which could not find optimal solutions for large instances; and a heuristic method were proposed for it. Despite the similarities, none of the works cited above base their development on the maximum covering by disks algorithms found in the literature.

The main results of this work are found in Chapter 5 where a new algorithm for the version of the elliptical PMCLP where the ellipses can be freely rotated is developed. The proposed algorithm has a runtime complexity of $\mathcal{O}(n^{3m})$ and one of its most important steps is finding every solution of a not-previously-known problem. In Chapter 6, this problem is introduced and an algorithm for it is presented. The rest of the work is structured in the following way: Chapter 2 introduces some definitions and results that are used throughout the next chapters; in Chapter 3, the maximum covering by disks problem is studied and a $\mathcal{O}(n^{2m})$ algorithm is proposed; in Chapter 4, the maximum covering by ellipses is introduced and the algorithm for the disks case is adapted for it; finally, Chapter 7 presents a conclusion as well as what is left as future work. Also, Appendix A determines with detail the intersection of two ellipses, which is used in the algorithm developed in Chapter 4.

NOTATION AND PRELIMINARIES

Some definitions and results that are used throughout the text are given in this chapter.

2.1 Elliptical and Euclidean norm functions

A norm in $\mathbb{R}^2$ is a function that maps every vector onto a non-negative real number satisfying homogeneity and the triangle inequality.

Let $u \in \mathbb{R}^2$ be a vector, the Euclidean norm of u is defined as

$$||u||_2 = \sqrt{u^T u}. \tag{2.1}$$

The elliptical norm, also known as weighted norm, takes a 2 by 2 positive definite matrix as its parameter. This matrix can be seen as a linear transformation of the Euclidean norm. The elliptical norm of $u \in \mathbb{R}^2$ is defined as

$$||u||_Q = \sqrt{u^T Q u}, \tag{2.2}$$

where Q is a 2 by 2 positive definite matrix. Note that the elliptical norm, when taking Q to be the identity matrix, becomes the Euclidean norm.

Determining the distance between two points, given a norm function, is done by calculating the norm of the vector defined by the difference between the two points. For example, the elliptical distance between the points $p, q \in \mathbb{R}^2$ is given by $||p - q||_Q$.

2.2 Disk

A circle (or circumference) is a set of points in $\mathbb{R}^2$ that have the same Euclidean distance, also known as radius, to another point, also referred to as the center of the circle. A unit circle is

a circle with radius equal to 1.

A disk is the set of points bounded by a circle. In other words let $c \in \mathbb{R}^2$. A unit disk with center c is the set of every point $p \in \mathbb{R}^2$ which satisfies

$$||p - c||_2^2 \leq 1. \tag{2.3}$$

2.3   Ellipse

An ellipse is a curve which is categorized, along with the parabola and the hyperbola, as a conic section. They get this name because conic sections are curves resulted from the intersection of a right circular cone in $\mathbb{R}^3$ with a plane (BRANNAN; ESPLEN; GRAY, 1999). From that definition, an equation which describes any conic section is given as follows

$$Ax^2 + Bxy + Cy^2 + Dx + Ey + F = 0, \tag{2.4}$$

where $A, B, C, D, E, F \in \mathbb{R}$ are fixed and $x, y \in \mathbb{R}$. Distinguishing an ellipse from the other conic sections ca be done using the condition given by Equation 2.5. More details about conic sections can be found in Ayoub (1993).

$$4AC - B^2 > 0. \tag{2.5}$$

Assuming the center of an ellipse is $c \in \mathbb{R}^2$, then Equation 2.4 can be rewritten as a quadratic form as follows

$$(p - c)^T Q(p - c) = 1, \tag{2.6}$$

with $p \in \mathbb{R}^2$ and Q being a 2 by 2 positive definite matrix which carries the parameters of the ellipse. From Equation 2.4, Q can be defined as follows

$$Q = \begin{pmatrix} A & \frac{B}{2} \\ \frac{B}{2} & C \end{pmatrix}.$$

Note that asking Q to be positive definite is the same as asking $4AC - B^2$ to be positive. This makes us arrive at the following definition of the ellipse.

Definition 1. Let $c \in \mathbb{R}^2$ be the center of an ellipse and Q be a 2 by 2 positive definite matrix. An ellipse is the set of every point $p \in \mathbb{R}^2$ such that $||p - c||_Q^2 = (p - c)^T Q(p - c) = 1$. Also, a point p is considered covered by an ellipse if $||p - c||_Q^2 = (p - c)^T Q(p - c) \leq 1$.

An alternative way to define an ellipse, which can be seen as just a property derived from the definition above, is to begin its construction with two points called foci and a constant $R \in \mathbb{R}$, with R being greater than the Euclidean distance between the two foci points (see Figure 1). The ellipse is, then, defined as the set of points whose distance to the foci is equal to R. In other words, let $f_1, f_2 \in \mathbb{R}^2$ be the two foci points, the ellipse is the set of every point $p \in \mathbb{R}^2$, such

that $||p - f_1||_2 + ||p - f_2||_2 = R$. It can be shown that this definition is equivalent to Definition 1, with the coverage of a point p being equivalent to $||p - f_1||_2 + ||p - f_2||_2 \leq R$.

Figure 1 – A non-axis-parallel ellipse and its foci points.

Source: Elaborated by the author.

Also, in Figure 1, the distance $a = ||p_a - c||_2$, where p_a is one of the intersection points of the ellipse with the major axis, is called the semi-major, and the distance $b = ||p_b - c||_2$, where p_b is one of the intersection points of the ellipse with the minor axis, is called the semi-minor. These two values are also referred to as the shape parameters of an ellipse. Let $d = ||c - f_1||_2$, then it is easy to see that $a = R - d$ and $b = \sqrt{\frac{R^2}{4} - d^2}$.

Finally, an ellipse is said to be axis-parallel if its major-axis (see Figure 1), which is the line that passes through its two foci points, is parallel to the x-axis.

2.3.1 Axis-parallel

An axis-parallel ellipse centered at $c = (c_x, c_y)$ can be described using Definition 1 with Q being a diagonal matrix [1]. This can be understood as a scaling transformation applied to the Euclidean norm.

Defining the matrix Q as

$$Q = \begin{pmatrix} \frac{1}{a^2} & 0 \\ 0 & \frac{1}{b^2} \end{pmatrix},$$

then, starting from Definition 1, we can obtain the following equation

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[1] The only non-zero terms are in the main diagonal.

$$(p-c)^T Q(p-c) = 1 \Rightarrow$$

$$(\frac{p_x - c_x}{a^2}, \frac{p_y - c_y}{b^2})^T (p_x - c_x, p_y - c_y) = 1 \Rightarrow$$

$$\frac{(p_x - c_x)^2}{a^2} + \frac{(p_y - c_y)^2}{b^2} = 1, \tag{2.7}$$

where a and b are the semi-major and semi-minor shape parameters, respectively.

Also, the coverage region is determined by just changing the equality to an inequality as follows

$$\frac{(p_x - c_x)^2}{a^2} + \frac{(p_y - c_y)^2}{b^2} \leq 1. \tag{2.8}$$

Another way to represent ellipses, which will be useful in some occasions, is through writing it as a curve, function of the angle with its major-axis (see Figure 2).

Figure 2 – The ellipse as a parametric curve.

Source: Elaborated by the author.

Let $c \in \mathbb{R}^2$ be the center of an ellipse with shape parameters $(a,b) \in \mathbb{R}^2_{>0}$. Then $\gamma : [0, 2\pi] \mapsto \mathbb{R}^2$ defines a curve which maps every angle onto a point on the ellipse and it is defined as follows

$$\gamma(t) = \begin{cases} x(t) = a\cos t + c_x, \\ y(t) = b\sin t + c_y. \end{cases} \tag{2.9}$$

2.3.2   Non-axis-parallel

A non-axis-parallel ellipse centered at $(c_x, c_y) \in \mathbb{R}^2$ can also be described by Definition 1. Nonetheless, a simpler equation is going to be used here. Besides the center, and the shape parameters $(a,b) \in \mathbb{R}^2_{>0}$, with $a > b$, an angle of rotation $\theta \in \mathbb{R}$ is given. This new parameter represents the angle between the x-axis and the major-axis of the ellipse, this can be seen on

Figure 3 where the dashed lines represent the ellipse's axes and the angle between the major-axis and the *x*-axis is displayed.

A rotated ellipse by θ can be transformed into a axis-parallel centered at the origin by applying two linear transformations: translation, so its center is at $(0,0)$, then rotation so its major-axis is parallel to the *x*-axis. Reversing these transformations produces the following equation for a non-axis-parallel ellipse

$$\frac{((x-c_x)\cos\theta+(y-c_y)\sin\theta)^2}{a^2}+\frac{((x-c_x)\sin\theta-(y-c_y)\cos\theta)^2}{b^2}=1. \qquad (2.10)$$

The coverage region of that same ellipse is given by every point $(x,y)\in\mathbb{R}^2$ that satisfies the following equation, which is the same as Equation 2.10 with the equality sign $(=)$ replaced by $(\le)$

$$\frac{((x-c_x)\cos\theta+(y-c_y)\sin\theta)^2}{a^2}+\frac{((x-c_x)\sin\theta-(y-c_y)\cos\theta)^2}{b^2}\le 1. \qquad (2.11)$$

Another important property of ellipses is shown on Figure 3. The two angles of rotation between the major-axis and the *x*-axis (θ and $\theta+\pi$) are actually equivalent–they produce the same ellipse. This symmetry is true for any angle of rotation, which means that θ is equivalent to $\theta+k\pi$, $k\in\mathbb{Z}$. Therefore, it is possible to conclude that to represent any ellipse, it is enough to specify the domain of θ to be $[0,\pi)$.

Figure 3 – The rotated ellipse.

Source: Elaborated by the author.

2.3.3 Notation

As stated by Definition 1, the word ellipse is used to refer to the set of points that satisfies the equality equation, which can be seen as the border of an ellipse's coverage area. For this

work, however, it is more convenient to refer directly to the coverage area of an ellipse and add a notation to express its border. For example, let E be an ellipse's coverage area, and $\mathcal{P} \subset \mathbb{R}^2$ a set of points, then $E \cap \mathcal{P}$ denotes the set of points from $\mathcal{P}$ inside the coverage area of that ellipse. When we need to refer specifically to the border of E, we use the boundary operator: ∂E.

2.4 Complex numbers

The set of complex numbers $\mathbb{C}$ is an extension of the set of real numbers $\mathbb{R}$ and can be very useful depending on the problem at hand. A thorough introduction on this topic is out of the scope and we just go through some basic properties that are going to be used later on Chapter 6.

A complex number is composed of a real part added to an imaginary part which is a multiple of the imaginary unit $i = \sqrt{-1}$. This is expressed as $a + bi$, with $a, b \in \mathbb{R}$. Because z is composed of two real numbers, the complex number system can be visualized on $\mathbb{R}^2$ as shown on Figure 4. This is also a good way to visualize Euler's Formula, as it can be also seen on Figure 4, which says that any complex number can be written in terms of its radius and polar angle as follows:

$$z = re^{i\theta} = r(\cos\theta + i\sin\theta), \tag{2.12}$$

with r being the length of the vector determined by the point z on the complex plane and $\theta = angle(z)$ being its polar angle. Note that *angle* is a function from $\mathbb{C}$ to $[0, 2\pi)$.
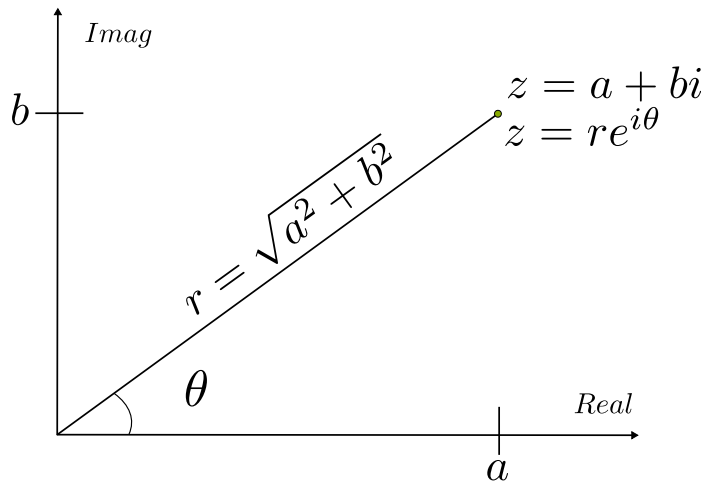

Figure 4 – The representation of a complex number on two dimensions.

The complex conjugate is another important operator that is utilized later. Let $z = a + bi \in \mathbb{C}$, then we refer to $\bar{z}$ as the complex conjugate of z and it is defined as:

$$\bar{z} = a - bi. \tag{2.13}$$

Finally, we need to state two observations about the angles of $-z$ and $\bar{z}$ with respect to the angle of z, this will be useful later on Chapter 6.

$$angle(\bar{z}) = 2\pi - angle(z) \tag{2.14}$$

$$angle(-z) = \pi + angle(z), \tag{2.15}$$

checking the validity of these two equalities can be done by just observing the symmetry between the points defined by $z, \bar{z}$, and $-z$ on the plane.

2.5 Polynomials and their roots

In this work, we are mostly interested in univariate polynomials defined over the complex numbers. A function $p_n : \mathbb{C} \mapsto \mathbb{C}$ is a n-degree polynomial if it can be written as

$$p_n(z) = \sum_{k=0}^{n} a_k z^k, \tag{2.16}$$

with $a_0, \ldots, a_n \in \mathbb{C}$. In this work, when a polynomial is written in the form of Equation 2.16 we say that it is in the power form or in the monomial form.

The famous Abel-Ruffini Theorem (a proof can be seen in Skopenkov (2015)) states that for polynomials of degree higher than four, there is no closed formula[2] to determine their roots. Fortunately, a numerical approach exists for higher-degree polynomials which works really well in practice.

In Horn and Johnson (1986, p. 195) a theorem is presented which says that for every univariate polynomial of degree n, there exists a companion matrix which is a $n \times n$ matrix, such that its eigenvalues are the zeros of that polynomial. For example, the companion matrix of a degree-5 polynomial written as Equation 2.16 is given by

$$\begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ -\dfrac{a_0}{a_5} & -\dfrac{a_1}{a_5} & -\dfrac{a_2}{a_5} & -\dfrac{a_3}{a_5} & -\dfrac{a_4}{a_5} \end{bmatrix}. \tag{2.17}$$

Finding every eigenvalue of a matrix can be done using the QR algorithm, which runs in $\mathcal{O}(n^3)$ and uses $\mathcal{O}(n^2)$ memory. A very complete introduction to this algorithm can be found on Watkins (2008). The idea is to convert the input matrix to the Hessenberg form, after that, the matrix preserves its form through the iterations and under some assumptions, convergence to the eigenvalues is achieved. Also, in Barel *et al.* (2010) it is pointed that companion matrices are already in the Hessenberg form, making the Hessenberg conversion step unnecessary. This and

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[2] A finite number of $+, -, \times, \div, \sqrt{}$.

other properties are used in Barel *et al.* (2010) on the development of a $\mathcal{O}(n^2)$ algorithm for the specific case of companion matrices.

In practice, LAPACK's ZGEEV routine is utilized–the user guide can be found in Anderson *et al.* (1999)–which is an implementation of the QR algorithm that returns every eigenvalue of a complex matrix.

2.6   Real trigonometric polynomial

The same definition found in Powell (1981, p. 150) for real trigonometric polynomials is given here. They are also referred to as truncated Fourier Series in Boyd (2006) and are given by

$$p_n(\theta) = \sum_{k=0}^{n} a_k \cos(k\theta) + \sum_{k=1}^{n} b_k \sin(k\theta). \tag{2.18}$$

We say that $p_n : \mathbb{R} \mapsto \mathbb{R}$ as defined by Equation 2.18 is a n-degree real trigonometric polynomial. An important property is stated in Powell (1981, p. 150), it says that a n-degree polynomial can have up to $2n$ distinct roots on the interval $[0, 2\pi)$. It also says that a function written in the format

$$\cos^j \theta \sin^k \theta \qquad j, k \in \mathbb{Z}_+ \ ,$$

can be transformed into a real trigonometric polynomial of degree $j + k$. Therefore, the following expression also represents a n-degree real trigonometric polynomial

$$\sum_{0 \le j+k \le n} c_{j,k} \cos^j \theta \sin^k \theta, \tag{2.19}$$

for some $\{c_{j,k} \in \mathbb{R} : 0 \le j + k \le n\}$.

MAXIMUM COVERING BY DISKS

In this chapter, we introduce a version of the classical Euclidean norm PMCLP where each facility has a given coverage radius. We refer to this problem as Maximum Covering by Disks (MCD). We also propose a method for it intending to adapt it for the elliptical PMCLP that will be introduced in the next chapter.

3.1 Definition

An instance of MCD is given by a set of n demand points $\mathcal{P} := \{p_1, \ldots, p_n\}$, with $p_j \in \mathbb{R}^2$; and a set of m disks $\mathcal{D} := \{D_1, \ldots, D_m\}$, which represents the facilities and are used to cover the demand set. Additionally, for every point in $\mathcal{P}$ a weight is given, and for every disk in $\mathcal{D}$ a coverage radius is given. The weight of the j-th point is denoted by $w_j \in \mathbb{R}_{>0}$, and the coverage radius of the j-th facility is denoted by $r_j \in \mathbb{R}_{>0}$.

A solution for an instance of MCD is determined by $Q := (q_1, \ldots, q_m) \in \mathbb{R}^{2m}$ which specifies the center of every disk in $\mathcal{D}$. The coverage area of j-th disk centered at q_j is denoted by $D_j(q_j)$. Finally, let $w : 2^{\mathcal{P}} \mapsto \mathbb{R}_{>0}$ be the function

$$w(A) = \sum_{j:p_j \in A} w_j, \tag{3.1}$$

which takes a subset of $\mathcal{P}$ and returns the sum of the weights of every point in that subset; then an optimal solution of MCD is given by the optmization problem:

$$\max_Q w\left(\bigcup_{j=1}^m \mathcal{P} \cap D_j(q_j)\right). \tag{3.2}$$

It is important to note that MCD is a slightly different PMCLP than the one introduced in the first study on the subject in Church (1984). There, a coverage radius is given for each

demand, rather than for each facility, and a demand is only considered covered if a facility is located within its radius.

3.1.1   CLS and CIPS

The method proposed in Church (1984) involves the creation of a Candidate List Set (CLS) which is a finite set of possible solutions for each facility. After a CLS is constructed, a complete search can be executed to find the best solution.

As this work is developed towards an exact method for MCD, the focus is on constructing a CLS that contains an optimal solution. This is done in Church (1984) where a circle is defined for every demand point, and a circle intersection point set (CIPS) is constructed containing the demand points and every pairwise circle intersection. We introduce a similar approach based on works found on the case for only one disk and describe a $\mathcal{O}(n^2 \lg n)$ algorithm for it.

3.2   One disk version

This version of the problem will be refered to as Maximum Covering by One Unit Disk (MCD-1) and it is just a specific case of MCD with one unit disk ($m = 1$ and $r_1 = 1$). We later use the algorithm for MCD-1 to construct a CLS which is guaranteed to contain an optimal solution for MCD.

Two exact methods for it have been found in the literature. A $\mathcal{O}(n^2)$ algorithm is proposed by Chazelle and Lee (1986) which improved the previously $\mathcal{O}(n^2 \log n)$ one proposed by Drezner (1981). As it has been mentioned, MCD-1 is a 3SUM-HARD problem, which means that it is as hard as the 3SUM problem (the problem of finding 3 real numbers that sum to 0, given n real numbers). Initially the lower bound of the 3SUM problem was conjectured to be $\Omega(n^2)$, matching the best algorithm for MCD-1, which meant that no better time-complexity could be achieved. Since then, however, better algorithms for 3SUM have been developed with a $\mathcal{O}(\frac{n^2}{poly(n)})$ run time complexity (KOPELOWITZ; PETTIE; PORAT, 2014).

In Drezner (1981), the main idea used to develop the $\mathcal{O}(n^2 \log n)$ algorithm is that, even though there are infinitely many points where the disk could be placed, only a few of them, a finite amount of $\mathcal{O}(n^2)$, needs to be considered for the method to find an optimal one. The algorithm, for every point, sorts the other points with respect to the angle they form with the first one. After that, the first point is placed on the border of the disk and, going through the sorted list, the algorithm inserts and removes points from the disk coverage. Also, when inserting and removing a point from the coverage, it only checks the disk centers that make the entering/leaving point to be on the border. Because the algorithm only checks the centers that make the disk have two points on its border, the number of centers it goes through is bounded by the number of pairs of points, which is $\binom{n}{2} = \mathcal{O}(n^2)$.

In Chazelle and Lee (1986) and Berg, Cabello and Har-Peled (2006), on the other hand, instead of working directly with MCD-1, an equivalent problem called Maximum Weight Clique (MWC) was introduced. The algorithm for MCD that is developed in this chapter also uses that equivalence. Because of that, MWC is introduced in the next section.

3.3 Maximum Weight Clique Problem

An instance of MWC is given by a set of points $\mathcal{P} := \{p_1, \ldots, p_n\}$, with $p_i \in \mathbb{R}^2$; and a set of unit disks $\mathcal{D} := \{D_1, \ldots, D_n\}$, such that D_i is centered at p_i, for $i = 1, \ldots, n$, and with the i-th disk having a weight $w_i \in \mathbb{R}_{>0}$ assigned to it.

A clique, in this context, is a non-empty intersection region of one or more disks, and its weight is the sum of the weights of those disks in the intersection. Following this, a solution for MWC can be defined as just a point $q \in \cup_{j=1}^n D_j$ which is inside any of the given disks in $\mathcal{D}$. From a solution q, the corresponding clique S can be obtained by intersecting every disk that contains q:

$$S = \bigcap_{j: q \in D_j} D_j. \tag{3.3}$$

With a geometric observation, though, the number of possible values for the solution can be reduced. Let $\partial \mathcal{D} = \{\partial D_1, \ldots, \partial D_n\}$ be the set of circles corresponding to each disk in $\mathcal{D}$, it is possible to see that any clique has at least one point on its border that is in the intersection of two circles in ∂D or, in the case where the clique is composed of only one disk, it contains obviously the center of its only disk. Therefore, q can be actually limited to the set of pairwise intersections of $\partial \mathcal{D}$ as well as the set of centers of each disk $\mathcal{P}$. Finally, an optimal solution of MWC is given by the optimization problem

$$\max_q \sum_{D_k \cap q \neq \emptyset} w_k, \tag{3.4}$$

with $q \in \{\partial D_i \cap \partial D_j : 1 \leq i < j \leq n\} \cup \mathcal{P}$. With this new observation, it can be shown that the number of possible solutions that need to be considered, so an optimal solution is found for an instance of MWC is bounded by $\binom{n}{2} + n = \mathcal{O}(n^2)$.

It is worth pointing out that MWC is a different problem than the maximum clique on a intersection graph (a graph where the vertices are the disks and an edge exists if there is an intersection between two disks). As shown in Figure 5, three disks could have non-empty pairwise intersection and still have an empty intersection of all of them together. That is why MWC is also referred to as the Maximum Geometric Clique Problem and the other version, when there is only the pairwise intersection condition, is referred to as the Maximum Graphical Clique Problem (DE; NANDY; ROY, 2014).

In Chazelle and Lee (1986), the method for MWC consists on building a planar graph on which the vertices were the $\binom{n}{2}$ pairwise intersection of the circumferences and the edges were the arcs of the circumferences connecting the intersections. With the graph constructed, a traversal is done to obtain the answer, thus the time complexity of $\mathcal{O}(n^2)$.

Figure 5 – Three disks that have non-empty pairwise intersection among them, but no common intersection.
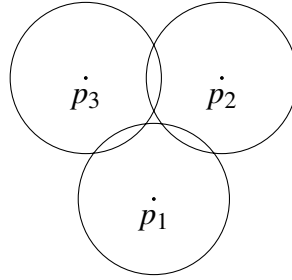
Source: Elaborated by the author.

As it has been mentioned, with the equivalence of the two problems, an optimal solution of the Maximum Weight Clique Problem is also an optimal solution of MCD-1, which means that a disk centered at q^* which is an optimal solution of MWC, will have a maximal weight covering of the demand set $\mathcal{P}$.

Given an instance of MCD-1, the equivalent MWC intance is obtained by defining the set $\mathcal{D}$ to contain the disks centered at $\mathcal{P}$ and setting the weight of every disk to be the weight of its corresponding point in $\mathcal{P}$. A disk D_i will represent the area where a disk can be placed in order to cover p_i. This means that a intersection between some disks is a region where a disk could be placed to cover the corresponding points.

In Figure 5, it can be seen that there is no point where a disk could be placed such that it would cover p_1, p_2 and p_3, nonetheless, in any of the pairwise intersections, a disk could be placed to cover the two corresponding points.

Formally, in MWC, if a point q lies inside $\bigcap_{k \in I} D_k$, with $I \subset \{1, \ldots, n\}$, then a disk centered at q will cover the points p_k, with $k \in I$ in the equivalent MCD-1 instance. Conversely, the same applies for a disk placed at q that covers points p_k, with $k \in I$ in the MCD-1 instance. It means that q will lie inside region $\bigcap_{k \in I} D_k$ in MWC.

3.3.1 An algorithm for the Maximum Weight Clique Problem

The algorithm described here is based on the one in Drezner (1981), also with some ideas from De, Nandy and Roy (2014) and Berg, Cabello and Har-Peled (2006). It has a run time complexity of $\mathcal{O}(n^2 \log n)$ and uses $\mathcal{O}(n)$ of extra space. It is worth noting, however, that a $\mathcal{O}((n+K) \log n)$ run time, with K being the number of intersections, can be obtained by

using the algorithm in Bentley and Ottmann (1979) to find all the intersections among the *n* circumferences.

In De, Nandy and Roy (2014) an important observation is made about the intersection regions of disks. Given an instance of MWC, any clique formed by a subset of $\mathcal{D}$ is bounded by the arcs of circles that intersect with it. Also, those arcs have the intersection of circles as its end-points. This can be seen on Figure 6 where the cliques that D_1 is part of are bounded by D_1's arcs which have its intersections with the other circles as end-points. Following this, a definition is presented to characterize the end-points of an arc bordering a clique.

Definition 2. Let D_i and D_j be two unit disks with non-empty intersection, and $(\theta_1, \theta_2) \in [0, 2\pi)^2$ be the two angles that ∂D_i and ∂D_j intersect, with the condition that (θ_1, θ_2) defines an arc (counter-clockwise order) of D_i that is the border of $D_i \cap D_j$. If D_i is tangent to D_j, then $\theta_1 = \theta_2$. Then, define $\Gamma_+(i, j) = \theta_1$ and $\Gamma_-(i, j) = \theta_2$, also we refer to them as opening and closing intersection angles respectively.

Figure 6 – Three disks and their intersection points.
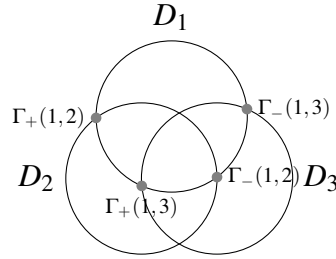
Source: Elaborated by the author.

In Figure 6, it is shown all the intersection points between D_1 with D_2 and D_3. Also, they are labeled according to Definition 2. Note that $\Gamma_+(1, 3) > \Gamma_-(1, 3)$ (the angles should be in the $[0, 2\pi)$ interval).

With Definition 2 in hand, we can establish the basis of the algorithm for MWC. The idea is to find for every disk, the best clique it is part of, this way an optimal solution is surely found.

Finding the best clique that D_i participates is done by traversing every intersection of ∂D_i with every circle from $\partial \mathcal{D}$ in counter-clockwise order while keeping a set of active disks. In other words, the set

$$A_i = \{\Gamma_+(i, j) : 1 \le j \le n\} \cup \{\Gamma_-(i, j) : 1 \le j \le n\}, \tag{3.5}$$

is traversed in counter-clockwise order and when an opening intersection angle is reached, the corresponding disk is added to the active set; and when a closing one is seen, the corresponding

disk is removed from the active set. This way, finding an optimal solution can be achieved by keeping the weight of the active disks as well as the best one found so far.

This simple traversal, however, would not handle the case with $\Gamma_+(i,j) > \Gamma_-(i,j)$ which would cause the algorithm to remove a disk wihtout first adding it leading to a possible non-optimal solution. This can be fixed by traversing the following set instead

$$B_i = A_i \cup \{\Gamma_+(i,j) + 2\pi : 1 \le j \le n\} \cup \{\Gamma_-(i,j) + 2\pi : 1 \le j \le n\}, \tag{3.6}$$

which basically duplicates the list of intersections. This works because for any pair of disks, such that $\Gamma_+(i,j) > \Gamma_-(i,j)$, B_i will contain $\Gamma_-(i,j) < \Gamma_+(i,j) + 2\pi$. This is shown in Figure 7, where the intersections list is duplicated simulating the traversal in B_i (note the indication to where the traversal starts as well as the positive and negative signs representing when a intersection with another disk opens and closes, respectively). It can be seen that without the duplication, the algorithm would find 2 as the value of an optimal solution, which is wrong.

Figure 7 – The intersections list of a disk with three other disks.
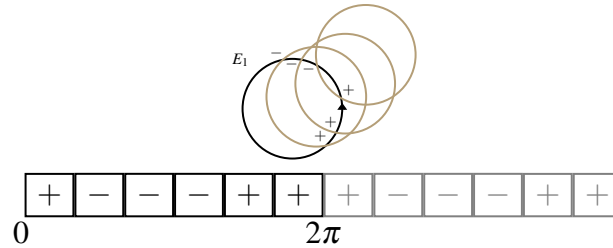
Source: Elaborated by the author.

Theorem 1. Algorithm 1 for solving the Maximum Clique Problem has a $\mathcal{O}((n+K)\log n)$ run time complexity, where K is the number of intersections of the n disks.

Proof. Finding every intersection can be done in $\mathcal{O}((n+K)\log n)$ by a plane sweep, the method is described in Bentley and Ottmann (1979). Because the traversal is made in counter-clockwise order, the intersection points have to be sorted by their intersection angles, so an additional $\mathcal{O}(K\log K)$ pre-processing is needed. All the other operations can be done in constant time. Therefore, the final algorithm complexity is $\mathcal{O}((n+K)\log n)$. $\qquad\square$

If a simpler implementation is desired, or the number of intersections is large, determining the set I_i (the set of disks that intersect with D_i, defined in Algorithm 1) can be simply done in $\mathcal{O}(n^2)$, making the algorithm have a worst-case complexity of $\mathcal{O}(n^2 \log n)$.

Procedure 1 – Algorithm for $MCD(\mathcal{P}, 1)$ with unit weights

---

Input: A set of points $\mathcal{P} = \{p_1, \ldots, p_n\}$.
Output: The maximum number of points that can be covered by a unit disk.

1: **procedure** $MCD_1(\mathcal{P})$
2:      $Q_{best} \leftarrow \{\}$
3:      $ans \leftarrow$ center of D_1
4:      **for all** $p_i \in \mathcal{P}$ **do**
5:          Let D_i be the disk with center at p_i
6:          Let I_i be the set of disks that intersect with D_i
7:          $A \leftarrow \{\}$ ▷ The multiset of intersection angles with D_i.
8:          **for all** $j \in I_i$ **do**
9:              $A \leftarrow A \cup \{\Gamma_+(i,j) \cup \Gamma_-(i,j)\}$
10:          **end for**
11:          $Q \leftarrow \{D_i\}$ ▷ The set of active disks.
12:          **for** $cnt = 1..2$ **do**   ▷ Do it twice.
13:              **for** $a \in A$ **do**   ▷ Assuming A is sorted.
14:                  Let D_a be the disk that intersects D_i at angle a.
15:                  **if** a is a starting angle **then**
16:                      $Q \leftarrow Q \cup \{D_a\}$
17:                  **else**
18:                      $Q \leftarrow Q \setminus \{D_a\}$
19:                  **end if**
20:                  **if** $|Q_{best}| < |Q|$ **then**
21:                      $Q_{best} \leftarrow Q$
22:                      $ans \leftarrow$ point corresponding to the intersection angle a
23:                  **end if**
24:              **end for**
25:          **end for**
26:      **end for**
27:      **return** ans
28: **end procedure**

3.4   Multiple disks, $MCD(\mathcal{P},m)$

In (BERG; CABELLO; HAR-PELED, 2006), a $\mathcal{O}(n^{2m-1} \log n)$ algorithm for $MCD(\mathcal{P}, m)$ is developed as a sub-routine for its $(1 - \varepsilon)$-approximation algorithm. Firstly, they solve the sub-problem $MCD(\mathcal{P}, 2)$ in $\mathcal{O}(n^3 \log n)$. Then for the rest of the points that are not in that solution, it uses the algorithm developed in (CHAZELLE; LEE, 1986) for the one-disk case, checking every possible solution for every one of the disks left.

Also, in (HE *et al.*, 2015) an heuristic method for large-scale *MCD* is proposed. It uses a probabilistic algorithm called mean-shift which is a gradient ascent method proven to converge to a local density maxima of any probability distribution. The mean-shift is utilized to find good candidates of centers for the unit disks, then the method backtracks to find the best assignment. The results showed that the greedy algorithm achieves an optimal coverage in some instances,

but for some other ones it has a 15 percent worse coverage ratio.

It can be seen that in any solution of $MCD(\mathcal{P}, m)$, a disk placed at a point q that covers at least one point $p \in \mathcal{P}$ has a correspondence to the Maximum Weight Clique Problem: the point q is inside an intersection area of at least one disk and that area is bounded by some disk, which means it will be checked by Algorithm 1 as a candidate to be an optimal solution. The number of points Algorithm 1 goes through is $\mathcal{O}(n^2)$, then checking every possible center for every ellipse yields a $\mathcal{O}(n^{2m})$ run-time complexity. This algorithm is described in Chapter 4 for the ellipses case.

The choice of developing a different method for the problem, instead of using the one from (BERG; CABELLO; HAR-PELED, 2006), is taken for the sake of simplicity, considering both algorithms achieve similar bounds.

CHAPTER

4

MAXIMAL COVERING BY ELLIPSES

In this chapter, the version of the Planar Maximal Covering by Ellipses Problem where every ellipse has to be parallel to the coordinate axis will be introduced and a method for it will be developed.

4.1 Axis-Parallel

The maximal planar covering using axis-parallel ellipses was first introduced in (CAN-BOLAT; MASSOW, 2009) which proposed a mixed integer non-linear programming method for the problem. This first approach showed to be not that efficient as it could not find an optimal solution for some instances within a timeout defined by them. To obtain solutions, not necessarily optimal ones, for the instances which the exact method showed inefficiency, a heuristic technique called Simulated Annealing was used to develop another method. Comparisons were made, which showed that the second approach was able to obtain good solutions, compared to the optimal ones found for some of the instances, within a good run-time.

The second work found in the literature was (ANDRETTA; BIRGIN, 2013), which developed a method that breaks the problem into smaller ones fixing the set of points an ellipse is going to cover. For each set of points fixed as the points an ellipse is going to cover, a small optimization problem is solved to find out if there is a location where the ellipse can be placed, so to cover the set of fixed points. To enumerate the possible solutions and then find an optimal one, the method defined a data structure that stores every set of points an ellipse can cover. This method showed better results and was able to find optimal solutions for the instances that the first method could not get as well as for new created instances.

4.1.1   Notation and definition of the problem

Axis-parallel ellipses are defined as the set of points that satisfy Equation 2.7. Therefore, all it takes to describe one is a pair of positive real numbers $(a,b) \in \mathbb{R}^2_{>0}$, also called the shape parameters, and a center point $q \in \mathbb{R}^2$.

Firstly, the case with only one ellipse is considered, an instance of the problem is denoted as $MCE(\mathcal{P},a,b)$ where $\mathcal{P}$ is a set of points and $(a,b) \in \mathbb{R}^2_{>0}$ is a pair of real numbers called the shape parameters of an ellipse. In the general case, every point has weights, but firstly this detail will be ignored and the weights are assumed to be unitary. The notation used here is similar to the one introduced on Chapter 3, the ellipse with an undefined center is referred to as E. To denote the ellipse with center set to be at point q, $E(q)$ is used. Also, the set of points covered by $E(q)$ is denoted by $E(q) \cap \mathcal{P}$, which indirectly defines $E(q)$ to be the set of points that satisfy Equation 2.8 (in other words $E(q)$ is the coverage region defined by the ellipse with shape parameters (a,b), located at center q). Also, to denote the points on the border of $E(q)$ we use the notation $\tilde{E}(q)$, therefore $p \in \tilde{E}(q)$ translates to p being on the border of $E(q)$. Hence, the problem can be defined as follows.

Definition 3. Let $MCE(\mathcal{P},a,b)$ be an instance of the maximal covering by one ellipse, with E being an ellipse with shape parameters $(a,b) \in \mathbb{R}^2_{>0}$, an optimal solution of $MCE(\mathcal{P},a,b)$ is given by

$$\max_q |\mathcal{P} \cap E(q)|. \tag{4.1}$$

4.1.2   One Disk algorithm adaptation

An adaptation of Algorithm 1 is obtained by just replacing the function that finds the intersection points between two disks by a function that finds the intersection points between two ellipses E_i and E_j. It can be seen in Figure 8 that the intersection points and their correspondents $\Gamma_-(i,j)$ and $\Gamma_+(i,j)$ functions behave the same way as in the disks case.

The intersection of two ellipses as well as determining $\Gamma_-(i,j)$ and $\Gamma_+(i,j)$ is described thoroughly in Appendix A.

Figure 8 – Three ellipses and their intersection points

Source: Elaborated by the author.

4.1.3 Multiple ellipses

The multiple ellipses case is handled using the same idea of the multiple disks case. The only difference is that an instance of the multiple ellipses may contain ellipses of different shapes, which does not happen for the disks case as every disk has the same radius. For this reason, a different pre-processing has to be done for every one of them.

An instance of the multiple ellipses case is denoted as $MCE(\mathcal{P}, \mathcal{E})$, with $\mathcal{P}$ being a set of n points and $\mathcal{E} = \{E_1, \ldots, E_m\}$ being a set of m ellipses, each one with shape parameters $(a_i, b_i) \in \mathbb{R}^2_{>0}, i = 1 \ldots m$. Also, without loss of generality, the weight of every point is assumed to be unitary.

Definition 4. Let $MCE(\mathcal{P}, \mathcal{E})$ be an instance of the maximal covering by ellipses, an optimal solution is given by

$$\max_{q_1, \ldots, q_m} \left| \bigcup_{i=1}^{m} \mathcal{P} \cap E_i(q_i) \right|. \tag{4.2}$$

Algorithm 3 describes the adapted version of the maximal disk covering algorithm for the ellipses case. In Algorithm 2, the MCE_1 procedure returns every possible set of points that an ellipse with shape parameters (a, b) can cover. With that, procedure MCE does the backtracking process, assigning every possible cover to every ellipse.

As stated in Lemma 1, MCE_1 runs in $\mathcal{O}(n^2 \log n)$, where n is the number of points. The number of sets of points an ellipse can cover, however, is $\mathcal{O}(n^2)$. Note that the $\log n$ being part of the complexity is due to sorting the set A. If MCE_1 is called only in a pre-process phase storing its return for every ellipse, a $\mathcal{O}(n^{2m})$ run-time complexity can be achieved.

Also, it can be seen that the unitary weights assumption can be easily removed through replacing the way the answer is updated: the weights of the covered points should be added to the answer instead of the number of covered points, this could be done by keeping an extra variable along with every possible set of points an ellipse can cover that is returned by MCE_1.

It is worth noting that some easy improvements, which do not change the algorithm's overall complexity, can be made in the implementation. For example, if an ellipse covers two sets of points X and Y, with $X \subset Y$, then set X can be ignored by the algorithm because of the positive weights constraint. Also, if two ellipses have their centers with Euclidean distance greater than their semi-major parameter, they for sure do not intersect. Depending on the input, this observation can make the algorithm not go through the whole list of ellipses every time it needs to determine the ellipses pairwise intersections.

Procedure 2 – Algorithm for $MCE(\mathcal{P}, a, b)$ with unit weights

---

Input: A set of points $\mathcal{P} = \{p_1, \ldots, p_n\}$, and the shape parameters (a, b) of an ellipse.
Output: A collection of subsets of $\mathcal{P}$ that the ellipse is able to cover.

  1: **procedure** $MCE_1(\mathcal{P}, a, b)$
  2:      $Z \leftarrow \{\}$ ▷ A collection of subsets of $\mathcal{P}$, each being a possible coverage.
  3:      **for all** $p_i \in \mathcal{P}$ **do**
  4:          Let E_i be the ellipse with center at p_i and parameters (a, b)
  5:          Let I_i be the set of ellipses that intersect with E_i
  6:          $A \leftarrow \bigcup_{j \in I_i} \{\Gamma_+(i, j) \cup \Gamma_-(i, j)\}$ ▷ The multiset of intersection angles with E_i.
  7:          $Z \leftarrow Z \cup \{\{p_i\}\}$
  8:          $Cov \leftarrow \{p_i\}$ ▷ The set of active disks.
  9:          **for** $cnt = 1..2$ **do**   ▷ Do it twice.
10:              **for** $a \in A$ **do**   ▷ Assuming A is sorted.
11:                 Let p_a be the point represented by the ellipse that intersects E_i at angle a.
12:                 **if** a is a starting angle **then**
13:                     $Cov \leftarrow Cov \cup \{p_a\}$
14:                 **else**
15:                     $Cov \leftarrow Cov \setminus \{p_a\}$
16:                 **end if**
17:                 $Z \leftarrow Z \cup \{Cov\}$
18:              **end for**
19:          **end for**
20:      **end for**
21:      **return** Z
22: **end procedure**

Procedure 3 – Algorithm for $MCE(\mathcal{P}, \mathcal{E})$ with unit weights

---

Input: A set of points $\mathcal{P} = \{p_1, \ldots, p_n\}$, and a set of ellipses $\mathcal{E} = \{E_1, \ldots, E_m\}$.
Output: The maximum number of points that can be covered by the set of ellipses $\mathcal{E}$.

  1: **procedure** $MCE(\mathcal{P}, \mathcal{E}, j = 1)$
  2:      **if** $j = m + 1$ **then**
  3:          **return** 0
  4:      **end if**
  5:      $ans \leftarrow 0$
  6:      **for** $E \in \mathcal{E}$ **do**
  7:          Let (a, b) be the shape parameters of E
  8:          $Q \leftarrow MCE_1(\mathcal{P}, a, b)$
  9:          **for** $Cov \in Q$ **do**
10:              $ans \leftarrow \max\{ans, |Cov| + MCE(\mathcal{P} \setminus Cov, \mathcal{E}, j + 1)\}$ ▷ Calls the procedure for the
next ellipse.
11:          **end for**
12:      **end for**
13:      **return** ans
14: **end procedure**

MAXIMAL COVERING BY ELLIPSES WITH ROTATION

This section introduces the covering problem that does not have any axis-parallel constraint for the ellipses, we denote this version of the problem as Maximal Covering by Ellipses with Rotation (MCER). The removal of this constraint introduces a new variable which is responsible for determining the rotation angle of every ellipse.

An instance of the non-axis-parallel is defined exactly like the axis-parallel one on Chapter 4. It is given by a set of demand points $\mathcal{P} = \{p_1, \ldots, p_n\}$, with every point having a unitary weight, and a set of ellipses $\mathcal{E} = \{E_1, \ldots, E_m\}$, with fixed shape parameters $(a_i, b_i) \in \mathbb{R}^2_{>0}$, $i \in \{1, \ldots, n\}$ with the additional condition that $a_i > b_i$. Given an instance of *MCER*, we define $Q := (q_1, \ldots, q_m) \in \mathbb{R}^{2m}$ to be the centers of each ellipse, $\Theta := (\theta_1, \ldots, \theta_m) \in [0, \pi)^m$ to be the angle of rotation of each ellipse and $E_i(q_i, \theta_i)$ to be the coverage region of ellipse E_i with its center at point q_i rotated by angle θ_i, which is given by Equation 2.11. Also, $\mathcal{P} \cap E_i(q_i, \theta_i)$ is denoted as the set of points covered by E_i on this configuration. Therefore MCER is defined as the problem of determining Q and Θ (placing and rotating each ellipse) to maximize the number of points covered by the m ellipses, which is given by Equation 5.1.

$$\max_{Q, \Theta} \left| \bigcup_{i=1}^{m} \mathcal{P} \cap E_i(q_i, \theta_i) \right|. \tag{5.1}$$

An additional notation is used on this chapter, $\tilde{E}_i(q_i, \theta_i)$ is defined to be the set of points on the border of $E_i(q_i, \theta_i)$, specially, the operation $\mathcal{P} \cap \tilde{E}_i(q_i, \theta_i)$ is used to refer to the set of points from $\mathcal{P}$ that lie on the border of $E_i(q_i, \theta_i)$.

Proposition 1. Let $(\mathcal{P}, \mathcal{E})$ be an instance of MCER. In an optimal solution of MCER, for any $E_j \in \mathcal{E}$, such that $|\mathcal{P} \cap E_j(q_j, \theta_j)| \geq 2$, there is q'_j such that $\mathcal{P} \cap E_j(q'_j, \theta_j) = \mathcal{P} \cap E_j(q_j, \theta_j)$ and $\mathcal{P} \cap \tilde{E}_j(q'_j, \theta_j) \geq 2$.

Demonstration. First, the angle of rotation can be ignored as it does not change.

Let $A = \mathcal{P} \cap E_j(q_j, \theta_j)$ be the set of points covered by E_j and $X = \cap_{p \in A} E_j(p, \theta_j)$ be the region of intersection of ellipses centered at each point from A.

As it was shown on Chapter 4, X is a region that is limited by arcs of ellipses. As this region is the non-empty intersection of more than one ellipse, there are at least two of these arcs that encounter at one point creating a vertex. Selecting any of these vertices as q'_j will make $|\mathcal{P} \cap \tilde{E}_j(q'_j, \theta_j)| \geq 2$.

<div align="right">□</div>

What Proposition 1 is saying is that any optimal solution for MCER can be transformed into another optimal solution where every ellipse covers the same set of points and those which cover more than one point has two points on their border. Also, this equivalent optimal solution can be always achieved by just translating the ellipses. An example is shown on Figure 9 for one ellipse.

A lot of the ideas developed in this chapter are based on fixing two points on the border of an ellipse, which is why Definition 5 introduces a new notation for angles that an ellipse rotated by it can be translated into a center, such that it contains two fixed points.

<div align="center">Figure 9 – An optimal solution before and after applying Proposition 1.</div>
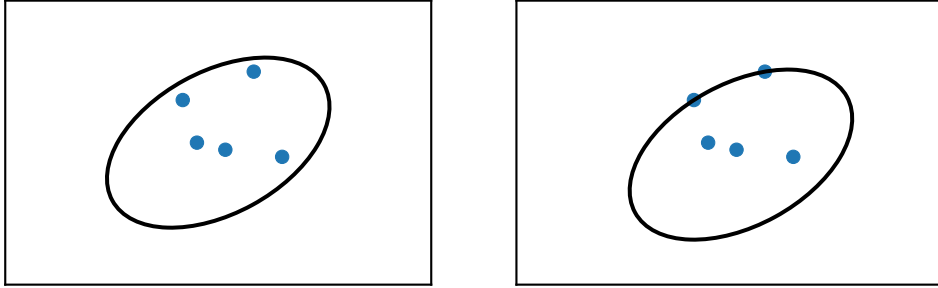
<div align="center">Source: Elaborated by the author.</div>

Definition 5. Let E be an ellipse and $u, v \in \mathbb{R}^2$. An angle $\theta \in [0, \pi)$ is said to be (E, u, v)-feasible if there is $q \in \mathbb{R}^2$ such that $\{u, v\} \subset \tilde{E}(q, \theta)$.

On Figure 10 two examples for Definition 5 are shown. On one of them, an ellipse is rotated by $\pi/4$ and it is located such that it contains the two fixed points on its border. That means $\pi/4$ is a (E, u, v)-feasible angle. On the other example, the ellipse is rotated by $\pi/2$ and there is not a center where the ellipse can be placed so it contains the two fixed points–they are too far apart. This makes $\pi/2$ a not (E, u, v)-feasible angle.

Lemma 1. Let $(\mathcal{P}, \mathcal{E})$ be an instance of MCER, in an optimal solution, for any $E_j \in \mathcal{E}$, such that $|\mathcal{P} \cap E_j(q_j, \theta_j)| > 2$, at least one of the two cases is true:

Figure 10 – A (E,u,v)-feasible angle and a not (E,u,v)-feasible angle.

Source: Elaborated by the author.

1. There is q', θ', and $\{u,v,w\} \subset \mathcal{P} \cap E_j(q_j,\theta_j)$, such that $\{u,v,w\} \subset \tilde{E}(q',\theta')$.

2. Let $A = \mathcal{P} \cap E_j(q_j,\theta_j)$, and $u,v \in A$ such that there exists $\hat{q}_j$ such that $\{u,v\} \subset \tilde{E}_j(\hat{q}_j,\theta_j)$ and $A \subset E_j(\hat{q}_j,\theta_j)$. Then for any (E_j,u,v)-feasible angle $\theta \in [0,2\pi]$, there exists $\bar{q}_j$ such that $\{u,v\} \subset \tilde{E}_j(\bar{q}_j,\theta)$ and $A \subset E_j(\bar{q}_j,\theta)$.

The first case of Lemma 2 is saying that there is another optimal solution which has E_j covering the same set of points, but with three points on its border.

The second case of Lemma 2 says that after fixing a pair of points on the border of E_j maintaining the covered set, for any angle that allows the two points to stay on the border of E_j, there is a center that maintains the covered set the same.

On Figure 11 both cases of Lemma 2 are shown. There, it can be seen that for the second case, it does not matter which feasible angle by which the ellipse is rotated, the third point will always be inside the coverage area. Also, an example of the first case is shown where there are three points lie exactly on the border of the ellipse.

The idea to prove Lemma 2 is that after fixing u,v on the border of E_j, which is possible by Proposition 1, the movement of rotation and translation while keeping u,v on the border is continuous. Because of that, the negation of case two implies case one and vice versa.

If we define an equivalence relation between optimal solutions as: S_1 is equivalent to S_2 if they both cover the same set of points, we can use Lemma 2 and Proposition 1 to identify the equivalence classes. Let S be any optimal solution, in $[S]$ (its equivalence class) there is another solution where any ellipse $E_j \in \mathcal{E}$ falls in at least one of the cases below:

- E_j covers only one point.

Figure 11 – An example of Lemma 2.
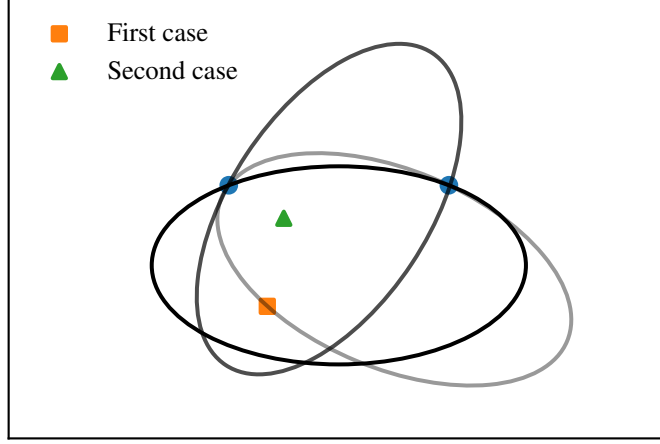
Source: Elaborated by the author.

- E_j covers more than one point with u and v being on the border of E_j. Note that there could be infinitely many of solutions like that, however, Lemma 2 guarantees that any (E_j, u, v)-feasible angle yields an equivalent optimal solution.

- E_j covers more than two points with three of them on its border.

The following two sections will treat the second and third cases (the first case is trivial). Going through every possibility of an ellipse falling in any of the three cases guarantees that an optimal solution is found.

5.1   Ellipse by two points

Let E be an ellipse with shape parameters $(a, b) \in \mathbb{R}^2_{>0}$ and $u, v \in \mathbb{R}^2$, one wants to find a (E, u, v)-feasible angle $\theta \in [0, \pi]$ and every center $q \in \mathbb{R}^2$ such that $\{u, v\} \subset \tilde{E}(q, \theta)$.

For a fixed angle, finding every center such that two points are on the border of the ellipse is done on Appendix A, from there we know that there could be at most 2 of such centers. The only thing left to be done is finding a feasible angle. It turns out that the angle that makes the major-axis of the ellipse to be aligned with the line that passes through u and v will be a feasible angle if, and only if the set of feasible angles is not empty. This can be seen geometrically as other angles achieve a lesser maximum distance between the two points on the border.

5.2   Ellipse by three points

CHAPTER

6

ELLIPSE BY THREE POINTS

The problem of finding every center and angle of rotation of a fixed shape ellipse which makes it have three points on its border is presented in this section. Even though its simple statement–it is short and uses only basic mathematical concepts–we were not able to find any work on it, or even on related problems. As a result, starting from scratch, we ended up trying a handful of approaches with most of them failing on the way. We try to give a review of some of those, going through the issues with the failing ones. At the end we propose an algorithm that takes around 6^3 operations to compute every solution of the problem. Also, we run some experiments to analyze its efficiency in terms of numerical precision and CPU time.

6.1 Definition

We call the problem introduced in this section Ellipse by Three points (E3PNT). An instance of it is given by three points $u, v, w \in \mathbb{R}^2$ and the shape parameters of an ellipse $(a,b) \in \mathbb{R}^2_{>0}$, with $a > b$.

Let E be the coverage area of an ellipse with the given shape parameters (a,b), then a solution of E3PNT is a pair $(q, \theta) \in \mathbb{R}^2 \times [0, \pi)$, such that $\{u, v, w\} \subset \partial E(q, \theta)$. As a last remark, because of its application on Chapter 5, it is only useful a method would only be useful to our case if it can encounter every solution of E3PNT. This requirement, of course, makes the problem more challenging.

6.2 Transforming the problem

Initially, E3PNT is a problem with 3 unknown variables: q_x, q_y, and θ. In this section, we transform E3PNT into the problem of finding the roots of an univariate function using the known problem of determining the circumcenter of a triangle. We also shown that the number of solutions of E3PNT is at most 6.

Figure 12 – Transforming an ellipse into a circle. T1, T2, and T3 represent the steps of the transformation.

Source: Elaborated by the author.

To make it simpler, let us translate the system, so the point u is at $(0,0)$. Then, we assume that the ellipse is actually axis-parallel and the points are the ones rotating. When an angle is found such that the three points lie on the border of the axis-parallel ellipse, a linear transformation can be applied to compress the x-axis by $\frac{b}{a}$, transforming the ellipse into a circle of radius b. This transformation can be seen on Figure 12 where a solution of the E3PNT is transformed into a solution of the problem of finding a circumscribed circle of a triangle. This process can be parametrized by the angle of rotation of the points, as described by

$$\varphi(p,\theta) = \begin{bmatrix} \frac{b}{a} & 0 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \cos\theta & \sin\theta \\ -\sin\theta & \cos\theta \end{bmatrix} \begin{bmatrix} p_x \\ p_y \end{bmatrix}, \tag{6.1}$$

because of the invertibility of linear transformations, solutions for E3PNT can be obtained by reversing these transformations.

After that transformation, the problem to be solved is finding a circumscribed circle of the triangle formed by the points $\{(0,0),\varphi(v,\theta),\varphi(w,\theta)\}$, such that the circle's radius is b. Therefore, the only unknown of this problem is the angle of rotation. This idea reduces the number of variables from 3 to only 1, however, this would only be useful if checking the existence of a circumscribed circle with radius b is a convenient problem.

It turns out that for three non-colinear fixed points, there is always an unique circumscribed circle. Also, finding this unique circle can be done analytically, and its radius R can be computed through the expression

$$R = \frac{\|\varphi(v,\theta)\|_2 \|\varphi(w,\theta)\|_2 \|\varphi(v,\theta) - \varphi(w,\theta)\|_2}{4A(\theta)}, \tag{6.2}$$

with $A(\theta)$ being the area of the triangle formed by the points $\{(0,0),\varphi(v,\theta),\varphi(w,\theta)\}$–note that

the transformation does not preserve distance or area. For more details about Equation 6.2 see Johnson and Young (1960, p. 189).

Imposing the radius to be equal b and squaring to eliminate the square roots present in the Euclidean distance, a function $\xi : [0, 2\pi) \mapsto \mathbb{R}_{>0}$ is defined as

$$\xi(\theta) = 16b^2 A(\theta)^2 - \|\varphi(v, \theta)\|_2^2 \|\varphi(w, \theta)\|_2^2 \|\varphi(v, \theta) - \varphi(w, \theta)\|_2^2, \tag{6.3}$$

in such a way that its zeros determine solutions for E3PNT.

Lastly, two questions about $\xi(\theta)$ that arise are: is its set of roots finite? And, if its set of roots is finite, can they be found analytically?

6.2.1 The number of solutions is limited

One of the steps of the method developed on Chapter 5 is to iterate over every solution of E3PNT for every triplet of points. Of course doing that is only possible for a finite number of solutions. Also, given that it is limited, discovering a bound for it is essential for determining the efficiency of the algorithm.

Back on Chapter 2 real trigonometric polynomials were introduced. It was stated that a n-degree polynomial can have up to $2n$ distinct roots. It turns out that ξ is a real trigonometric polynomial of degree 6 and it can be written in the format given by Equation 2.19 which implies that it can have up to 12 distinct roots. To show that, just note that it is possible to write $\|\varphi(v, \theta)\|_2^2$ and $A(\theta)^2$ in that form

$$\|\varphi(v, \theta)\|_2^2 = (v_x \frac{b}{a} \cos\theta + v_y \frac{b}{a} \sin\theta)^2 + (v_y \cos\theta - v_x \sin\theta)^2 \tag{6.4}$$

$$A(\theta)^2 = \frac{1}{4} \det \begin{pmatrix} v_x \frac{b}{a} \cos\theta + v_y \frac{b}{a} \sin\theta & v_y \cos\theta - v_x \sin\theta \\ w_x \frac{b}{a} \cos\theta + w_y \frac{b}{a} \sin\theta & w_y \cos\theta - w_x \sin\theta \end{pmatrix}^2. \tag{6.5}$$

It is also possible to see that the term of higher the degree of ξ is the multiplication of the three squared lengths of the triangle formed by the points $\{(0, 0), \varphi(v, \theta), \varphi(w, \theta)\}$. This multiplication has the same degree of $(\|\varphi(v, \theta)\|_2^2)^3$, and because $\|\varphi(v, \theta)\|_2^2$ has degree 2, the degree of ξ is 6.

Because ellipses are symmetrical with respect to their major-axis, and any rotation in the interval $[0, \pi)$ is identical to a rotation in $[\pi, 2\pi)$, the number of different solutions is cut in half. Therefore, the number of angles of rotation and centers that an ellipse of fixed shape can be placed, so it has three fixed points on its border is limited to 6.

6.3 An attempt using the conic general equation

The idea of this approach was to use the six-parameter conic equation to represent an ellipse. This equation is given by

$$Ax^2 + Bxy + Cy^2 + Dy + Ex + F = 0. \tag{6.6}$$

This equation actually represents any conic, for it to be an ellipse the condition $B^2 - 4AC < 0$ must be satisfied.

Setting the first point to be the origin, we get $F = 0$, using the other two points, it is possible to write D and E in terms of A, B, C. As any multiple of Equation 6.6 represents the same conic, we can set B to be equal 1. Then, we end up with two variables, A and C, and still need to impose that the final equation represents an ellipse and its major-axis and minor-axis have the predefined value. Let $\Delta = 4AC - B^2 = 4AC - 1$, and assume $F = 0$, then the expressions for both major-axis and minor-axis respectively are

$$a^2 = \frac{2\dfrac{AE^2 - BDE + CD^2}{\Delta}}{A + C - \sqrt{1 + (A - C)^2}} \tag{6.7}$$

$$b^2 = \frac{2\dfrac{AE^2 - BDE + CD^2}{\Delta}}{A + C + \sqrt{1 + (A - C)^2}}. \tag{6.8}$$

These two equations define two curves in $\mathbb{R}^2$ with A and C being the chosen variables. The solutions lie in the set of intersection of these curves. Finding this set was judged to be non-trivial and probably could be approximated numerically, however, we decided not to further pursue this approach.

Another idea which has been explored was working with the ratio $\frac{a^2}{b^2}$ which becomes an expression that allows A to be written as a function of C. This function appeared, at first we thought, to be monotonic, we tried to develop a method based on that, however, cases where the function does not behave as nicely were found. It is likely that developing a method to approximate solutions working with this function is possible, but we decided not to continue on this track.

6.4   An approximation method

One of the most useful techniques when dealing with complicated functions is approximation. They appear in various methods whenever a derivative or integral needs to be calculated or for example, in our case, when the roots of a function need to be determined. In general, one has a function f that is part of a family of functions $\mathscr{A}$ and wants to select a simpler function f^* from a set of functions $\mathscr{A}^*$, such that f^* is close enough to f (POWELL, 1981, p. 3). For this problem, the approximation of $\xi(\theta)$ on the interval $[0, \pi)$ is considered. The approximation set of functions is going to be the set of n-degree Chebyshev polynomials which the roots can be found through determining the eigenvalues of a $n \times n$ matrix.

6.4.1   Chebyshev polynomial

Chebyshev polynomials are widely used in Numerical Analysis in areas like numerical integration, polynomial approximation, and ordinary and partial differential equations. They are also very useful in practice and are present in extension libraries in Python, MATLAB and C.

Because of the scope of this work, only a brief introduction of Chebyshev polynomials of the first kind and its usage in polynomial interpolation is given. For a more thorough work on the subject, please check the book by Mason and Handscomb (2003).

We refer to $T_n : [-1,1] \mapsto [-1,1]$ as the n-degree Chebyshev polynomial of the first kind, and it is defined as

$$T_n(x) = \cos(n \arccos x). \tag{6.9}$$

It is important to mention that this definition can be extended to the whole real line. Using some trigonometric identities, T_n can also be expressed as a recurrence relation

$$T_n(x) = 2x T_{n-1}(x) - T_{n-2}(x). \tag{6.10}$$

An important property worth bringing up is that Chebyshev polynomials are orthogonal and form a basis for the polynomial space. This implies that any p_n of degree up to n can be expressed as a truncated Chebyshev series

$$p_n(x) = \sum_{j=0}^{n} a_j T_j(x). \tag{6.11}$$

One of the greatest qualities of Chebyshev polynomials is its numerical stability. Gautschi (1979) showed that the matrix that maps polynomials onto its coefficients written in the power form[1] has a condition number that grows exponentially with n. On the other hand, the matrix that converts polynomials to the Chebyshev basis as Equation 6.11, has a linear condition number bounded by $\sqrt{2}n$.

6.4.2   Chebyshev interpolation

Polynomial interpolation is a form of approximating a function by a polynomial of degree n that passes through $n+1$ chosen points. In fact, this polynomial is unique and it is determined by Lagrange's formula

$$f_n(x) = \sum_{j=0}^{n} f(x_j) \frac{\prod_{k \neq j}^{n+1} (x - x_k)}{\prod_{k \neq j}^{n+1} (x_j - x_k)}, \tag{6.12}$$

with f being the function to be approximated, and f_n the unique n-degree polynomial that passes through $\{(x_j, f(x_j)) : j = 0, 1, \ldots n\}$. Because of the uniqueness of interpolant polynomials, there is a direct link between the quality of an approximation and the points chosen to interpolate.

---

[1]   A polynomial is in the power form or the monomial form if it can be written as $\sum_{j=0}^{n} a_j x^j$.

As a matter of fact, depending on the points one chooses, even increasing the degree of the interpolation makes the approximation worsen. This is known as Runge's phenomenon and an example can be seen in Powell (1981, p. 37) where uniformly spaced points are chosen to interpolate the function $f(x) = (1+x^2)^{-1}$ on the interval $[-5,5]$.

That is where Chebyshev interpolation comes in. Instead of choosing $n+1$ arbitrary points, the $n+1$ roots of T_{n+1}, which are also known as Chebyshev Nodes, are chosen as the interpolation points. The $n+1$ Chebyshev nodes are given by

$$x_j = \cos\left(\frac{\pi(j - \frac{1}{2})}{n+1}\right), \tag{6.13}$$

for $j = 1, \ldots, n+1$. This particular choice defeats Runge's phenomenon and provides a convergent approximation. Note that, if the domain of the function to be interpolated is defined on a range other than $[-1,1]$, let us say $[a,b]$, then the transformation

$$\hat{x}_j = \frac{a+b}{2} + \frac{b-a}{2}x_j \tag{6.14}$$

can be done to map it to the Chebyshev Nodes' domain $[-1,1]$.

Then, the Chebyshev interpolation of a function $f : [a,b] \mapsto \mathbb{R}$ can be determined using Lagrange's formula and the points $\hat{x}_1, \ldots, \hat{x}_n$. As it was mentioned in Chapter 2, finding the roots of a polynomial written in the monomial form can be done by determining the eigenvalues of a so-called Frobenius companion matrix. For small n this works fine, however, converting the polynomial obtained by Equation 6.12 to the power form, as n grows, becomes a very ill-conditioned problem. An alternative method can be found in Boyd (2013) where the Chebyshev interpolation is calculated directly as a truncated Chebyshev series, as in Equation 6.11, in $\mathcal{O}(n^2)$. Also, given a polynomial written in the Chebyshev basis, a $n \times n$ matrix can be constructed, such that its eigenvalues are the roots of that polynomial. Boyd (2013) refers to this matrix as the Chebyshev-Frobenius companion matrix.

Therefore, the whole process of interpolating and finding the roots can be done using only Chebyshev polynomials, which have great numerical stability. Also, Chebyshev-Frobenius matrices have the same property as companion matrices which allows their eigenvalues to be found by a QR algorithm. Summing the two steps, a $\mathcal{O}(n^3)$ algorithm can be achieved, with n being the degree of the interpolation.

The last question that needs to be addressed is how close the roots of the Chebyshev interpolant f_n are to the roots of ξ?

Even though ξ is complicated enough, in a sense that finding its roots directly is no trivial task, it is a very well-behaved function: it is analytic and has infinitely many continuous and integrable derivatives. This satisfy all the requirements of the result in Gottlieb and Orszag (1977, p. 28) which says that if a function has m continuous and integrable derivatives on a closed interval, then its absolute difference between the Chebyshev truncate series is $\mathcal{O}(n^{-m})$.

Also, in Battles and Trefethen (2004) a theorem is presented stating that if a function is analytic on a neighborhood of $[-1, 1]$, then the convergence is $\mathcal{O}(C^n)$, for some $C < 1$.

To choose the degree of the interpolation we use the last coefficient rule-of-thumb introduced by Boyd (2001, p. 50). There is no guarantee that this method will choose n, such that f_n is close enough to ξ everywhere on $[0, \pi)$, nonetheless, in practice, it is taken to be a good estimate for the error

$$r_n = \max_{0 \le \theta < \pi} |f_n(\theta) - \xi(\theta)|, \tag{6.15}$$

which measures how far the interpolation is at the point it worst approximates.

6.5 Converting ξ into a polynomial

In Chapter 2, a brief introduction is given on how to get the roots of a polynomial. For that reason, we discuss two ways of converting ξ into a polynomial in this section. Symbolic computing was used to compute the polynomials, the Python external library called SymPy was utilized (see Meurer *et al.* (2017) for more details).

6.5.1 Real polynomial

The first attempt was using the identity $x = \tan \frac{\theta}{2}$ from which it is possible to construct a 12-degree polynomial. At first, the root-finding algorithm described on Chapter 2 seemed to work fine and return every solution of E3PNT, however, we later found out that for some instances, priorly known roots were not being found. The cause was not for sure identified, but a good guess would be that for angles which are greater than $\frac{\pi}{4}$, x starts growing too rapidly which could lead to numerical instability.

6.5.2 Complex polynomial

The second approach is based on a idea published on Boyd (2006) which uses the identities on Equation 6.16 to convert real trigonometric polynomials, which is the case of ξ, into univariate complex polynomials in order to obtain its roots using the companion matrix scheme. This approach is preferable as it preserves the numerical stability of the original real trigonometric polynomial–more details about this can be found in Weidner (1988) where it is said that computing the roots of a real trigonometric polynomial through this transformation does not yield loss of accuracy.

$$\cos \theta = \frac{e^{i\theta} + e^{-i\theta}}{2} \tag{6.16}$$

$$\sin \theta = \frac{e^{i\theta} - e^{-i\theta}}{2i}. \tag{6.17}$$

It is possible to show that with that substitution and changing the variable to $z = e^{i\theta}$, we obtain the following function $g : \mathbb{S} \mapsto \mathbb{C}$, with $\mathbb{S}$ being the unit complex circle ($\mathbb{S} = \{z \in \mathbb{C} : |z| = 1\}$):

$$g(z) = \sum_{k=0}^{12} c_k z^{k-6}, \tag{6.18}$$

for some $c_0, \dots, c_{12} \in \mathbb{C}$. As the equalities on Equation 6.16 are valid for any $\theta \in \mathbb{R}$, the function $g'(\theta) = g(e^{i\theta})$ and the function ξ are equivalent since $g(e^{i\theta}) = \xi(\theta)$ for any $\theta \in [0, 2\pi)$.

Now there is only two things preventing g from being a polynomial: firstly, z shows up with negative exponents (up to -6); secondly its domain is the unit circle $\mathbb{S}$, not the whole complex set $\mathbb{C}$. The first issue can be fixed by multiplying g by z^6, note that this does not create further problems as $0 \notin \mathbb{S}$. The second issue is removed by simply extending the domain from $\mathbb{S}$ to $\mathbb{C}$, as $\mathbb{S} \subset \mathbb{C}$, roots outside the unit circle could appear which requires a further check for every root that is found. Finally, from g, the polynomial $h : \mathbb{C} \mapsto \mathbb{C}$ is defined as follows:

$$h(z) = z^6 g(z) = \sum_{k=0}^{12} c_k z^k. \tag{6.19}$$

By its definition it is possible to see that every root of g is also a root of h, and conversely, every root $\hat{z}$ of h which is in $\mathbb{S}$, is also a root of g. Lastly, every root of g will correspond to a root of ξ though their angles on the unit circle.

6.5.2.1 Further improvements

It is possible to make another reduction to obtain a degree-6 polynomial, this can lower the size of the companion matrix utilized on obtaining the roots of h from $12^2 = 144$ to $6^2 = 36$. As it has been mentioned before, an ellipse is symmetric with respect to its own major or minor axis. In other words, rotating an ellipse by $\theta \in [0, \pi)$ is the same as rotating the same ellipse by $\pi + \theta$. On the other hand, on Chapter 2, it was stated that the angle of the complex numbers z and $-z$ has the same symmetry with each other as the aforementioned symmetry of ellipses:

$$angle(-z) = \pi + angle(z).$$

From that, as $g(e^{i\theta}) = \xi(\theta)$ for every $\theta \in [0, 2\pi)$, we conclude that $g(-z) = g(z)$. This implies that h is, in fact, an even polynomial, or that $h(-z) = h(z)$ is true for every $z \in \mathbb{C}$:

$$h(-z) = (-z)^6 g(-z) = z^6 g(z). \tag{6.20}$$

Therefore, all the odd degree coefficients of h are 0 and we can define the 6-degree polynomial $f : \mathbb{C} \mapsto \mathbb{C}$ with the substitution $y = z^2$:

$$f(y) = \sum_{k=0}^{6} c_{2k} y^k. \tag{6.21}$$

Then from every root $\hat{y}$ of f, two roots of h can be obtained: $\sqrt{\hat{y}}$ and $-\sqrt{\hat{y}}$. As the angle of one of the roots will not be between $[0, \pi)$ we can ignore on of them. Note that the the square root of $\hat{y}$ does not need to be calculated, as only the angles are needed and they can be obtained by:

$$angle(\sqrt{z}) = angle(z)/2. \tag{6.22}$$

Finally, using the QR algorithm mentioned on Chapter 2 a $\mathcal{O}(n^3)$ algorithm can be constructed for E3P. Because n will always be 6, we can say that we have a $\mathcal{O}(1)$ algorithm.

It is also worth mentioning that a pattern on the coefficients of f was identified, and maybe for future work it can be used for further improvements. Analyzing the polynomials produced for several instances, the following seems to be true:

$$c_k = \overline{c_{6-k}}, \tag{6.23}$$

for $k = 0, \ldots, 6$. For now, we neither have any ideas on how Equation 6.23 could be proved nor how it could be used for further improving our algorithm.

6.6 Numerical Experiments

In this section we run some experiments to verify how big the interpolation degree must be for a good precision to be achieved. Let f_n be the n-degree Chebyshev interpolation of ξ, we define the interpolation error r_n as:

Then, we want to determine the smallest n, such that $r_n \leq \varepsilon$, for a predefined ε. Unfortunately, calculating r_n involves taking samples from both functions on the whole interval, which is not viable computationally. In Boyd (2001, p. 50) a rule-thumb for estimating the interpolation error is given. It is worth mentioning that this rule

It uses a theorem that states that r_n is limited by the sum of the coefficients of the Chebyshev series that were removed by the truncation. The rule-of-thumb for estimating r_n is given by:

$$r_n \approx |a_n|. \tag{6.24}$$

Note that

Obtaining an expression for it is not trivial, and sampling the whole interval is not viable computationally.

Therefore, an estimation of r_n is used to measure how good is the approximation.

Also we adopt two suggestions from (BOYD, 2013). The first is to divide the interval into K subintervals to achieve a precision without having to increase the degree of the interpolant too much. The second is to use a couple of Newton's iterations to refine the roots found.

Figure 13 – The interpolation error measured on roots that were found.
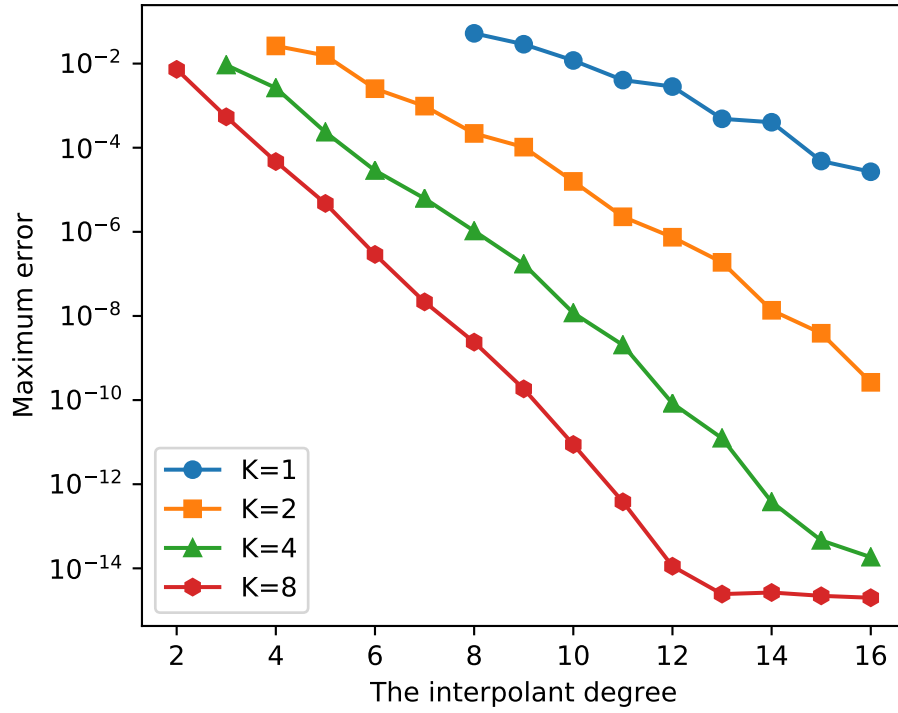
Let θ^* be a root of f_n, then we measure the error associated with that root as $|\xi(\theta^*)|$. For the numerical experiments, we considered every triplet of points from an instance with 25 points. Then for some values of K and n we define the error as the maximum error for every root that were found.

Figure 14 – The interpolation error measured on roots that were found after three Newton's iterations.

Source: Elaborated by the author.

CHAPTER

7
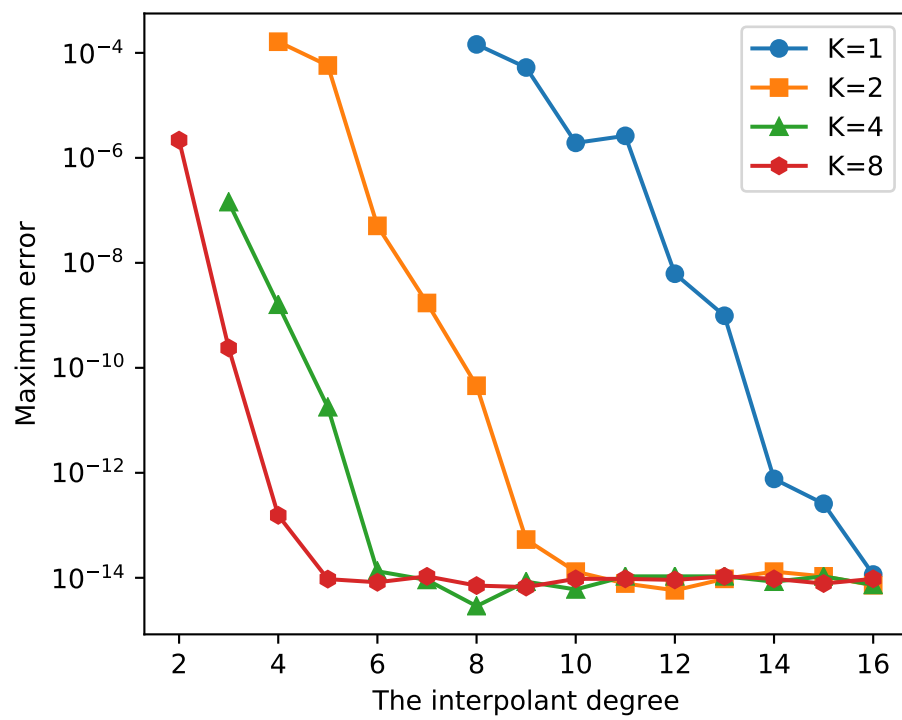
FUTURE WORK

To take advantage of the great amount of work found in the literature, we decided to first introduce the planar maximal covering by disks problem, develop a method for it, and just then adapt it for the ellipses case. It turned out that because of the similarities between the two problems, adapting was possible and actually very simple. This made the method developed by us have a very different approach than the ones in (ANDRETTA; BIRGIN, 2013) and (CANBOLAT; MASSOW, 2009). The next step is to implement it and compare the results that (ANDRETTA; BIRGIN, 2013) obtained.

For the next step of our master's research we set the following objectives as primary:

- Study the $(1-\varepsilon)$-approximation method for the planar covering with disks in (BERG; CABELLO; HAR-PELED, 2006) and develop an adapted version of the algorithm for ellipses with the same time complexity of $\mathcal{O}(n\log n)$.

- Develop an exact method for the version of the problem introduced in (ANDRETTA; BIRGIN, 2013) where the ellipses can be freely rotated.

The following goals are set as secondary:

- Develop a probabilistic approximation algorithm based on (ARONOV; HAR-PELED, 2008) which proposed a Monte Carlo approximation for the problem of finding the deepest point in a arrangement of regions. The method runs in $\mathcal{O}(n\varepsilon^2 \log n)$ and can be applied to solve the case with one ellipse. The case with more than one ellipse is left as a challenge for us for the next steps of our research.

- In (HE *et al.*, 2015), the task of finding every center candidate, after eliminating all the non-essential ones, is done in $\mathcal{O}(n^5)$ run-time complexity. We want to generalize this for the elliptical distance function and achieve a better run-time complexity. We also intend to use the mean-shift algorithm to try to develop a greedy version for the ellipses version.

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APPENDIX

A

INTERSECTIONS OF TWO ELLIPSES

In this appendix the intersection of two ellipses with the same shape parameters $(a,b) \in \mathbb{R}^2_{>0}$ is described with more detail, as well as determining the functions $\Gamma_+(i,j)$ and $\Gamma_-(i,j)$ for two ellipses that intersect.

A.1 Intersection

Let E_1 and E_2 be two ellipses that the intersection will be determined here. Without loss of generality, let us assume that E_1 is at the origin and E_2 is located at the center $(h,k) \in \mathbb{R}^2$. Their equations are given by

$$\frac{x^2}{a^2} + \frac{y^2}{b^2} = 1 \qquad (E_1)$$

$$\frac{(x-h)^2}{a^2} + \frac{(y-k)^2}{b^2} = 1 \qquad (E_2)$$

As they are both equal to one we can get the following

$$b^2 x^2 + a^2 y^2 = b^2(x-h)^2 + a^2(y-k)^2$$
$$b^2(-2xh + h^2) + a^2(-2yk + k^2) = 0$$
$$x(2hb^2) = b^2 h^2 + a^2(-2yk + k^2)$$
$$x = y\frac{-2ka^2}{2hb^2} + \frac{b^2 h^2 + a^2 k^2}{2hb^2}$$

Which can be rewritten as

$$x = y\alpha + \beta$$

with the constants α and β being

$$\alpha = \frac{-2ka^2}{2hb^2}$$
$$\beta = \frac{b^2h^2 + a^2k^2}{2hb^2}$$

Then replacing it back to the equation of E_1 we get

$$\frac{(y\alpha + \beta)^2}{a^2} + \frac{y^2}{b^2} = 1$$
$$b^2(y\alpha + \beta)^2 + y^2a^2 - a^2b^2 = 0$$
$$y^2(b^2\alpha^2 + a^2) + y(2\beta\alpha b^2) + b^2\beta^2 - a^2b^2 = 0$$

Which is a second degree polynomial, therefore, E_1 and E_2 intersect if, and only if the roots of the polynomial are real. The intersection points itself can be obtained by solving the polynomial for y and applying its value onto the $x = y\alpha + \beta$ equation.

A.1.1 Determining $\Gamma_+(i,j)$ and $\Gamma_-(i,j)$

Let us assume that E_1 and E_2, each one with shape parameters $(a,b) \in \mathbb{R}^2_{>0}$, intersect at p_1 and p_2. Then, to determine $\Gamma_+(1,2)$ and $\Gamma_-(1,2)$, we need to first determine the angles of intersection of p_1 and p_2 on E_1. For that, we will use the curve defined in Equation 2.9 because it is easier to work with angles here.

Given a point (x,y), to find the angle it makes with the major axis, from Equation 2.9, we can get that

$$\frac{y - q_y}{x - q_x} = \frac{b}{a}\tan t$$
$$t = \arctan\left(\frac{a}{b}\frac{y - q_y}{x - q_x}\right)$$

As the image of arctan is $[-\frac{\pi}{2}, \frac{\pi}{2}]$, we need to check the sign of $x - q_x$ to determine the angle in $[0, 2\pi]$. After that, we can get the two angles that represent the intersection points p_1 and p_2 on E_1.

To find out which one of the angles are $\Gamma_+(1,2)$, we need to go further and determine the derivative of $\gamma(t)$ which is going to be used to determine the vectors tangent to the ellipses at the intersection points.

$$\gamma'(t) = \begin{cases} x'(t) = -a\sin t \\ y'(t) = b\cos t \end{cases} \tag{A.1}$$

Let γ_1 and γ_2 be the curves describing E_1 and E_2 respectively. Also, let s_1 be the angle, such that $\gamma_1(s_1) = p_1$, and t_1 be the angle, such that $\gamma_2(t_1) = p_1$. Then, the tangent vectors to the E_1 and E_2 at p_1 are $\gamma_1'(s_1)$ and $\gamma_2'(t_1)$ respectively.

Figure 15 – Determining $\Gamma_+(1,2)$
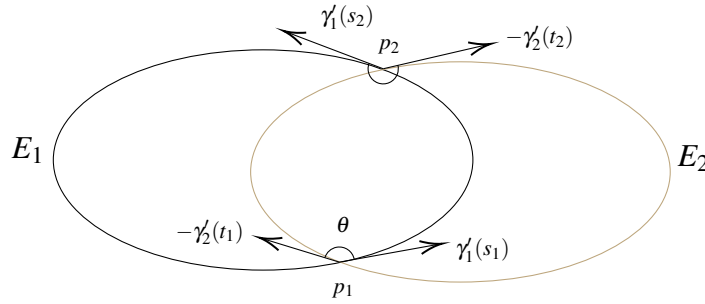
Source: Elaborated by the author.

The following lemma states a relation between s_1 and $\Gamma_+(1,2)$

Lemma 2. Let θ be the angle between $\gamma_1'(s_1)$ and $-\gamma_2'(t_1)$. Then, $\theta \leq \pi$ if, and only if $\Gamma_+(1,2) = s_1$.

Instead of a formal proof of Lemma 3, a graphical explanation using Figure 15 is provided.

First, let us state some facts that can also be seen in Figure 15

- $E_1 \cap E_2$ is convex and bounded by two arcs, one from each ellipse.

- Starting at any of the intersection points, one of the $E_1 \cap E_2$ arcs will be clockwise-oriented and the other, counter-clockwise-oriented. In Figure 15, for example, it is clear that only the E_1 arc starting at p_1, ending at p_2, is counter-clockwise-oriented.

- The counter-clockwise-oriented arc starting at $\Gamma_+(1,2)$ is from the ellipse E_1.

Let us assume that p_1 is the intersection point which is the opening angle $\Gamma_+(1,2)$. Then, the vectors $\gamma_1'(s_1)$ and $-\gamma_2'(t_1)$ are tangent to the $E_1 \cap E_2$ area at point p_1. Because of the convexity of $E_1 \cap E_2$, the angle between $\gamma_1'(s_1)$ and $-\gamma_2'(t_1)$ has to be less than or equal to π (see Figure 15), which is what Lemma 3 says. It is easy to prove the converse by proving the contra-positive assuming that p_1 is the point which determines the angle $\Gamma_-(1,2)$.

Lastly, in Figure 15, it can be seen that if one the intersection points is classified as $\Gamma_+(1,2)$ the other will necessarily be classified as $\Gamma_-(1,2)$. This gives us all we need to implement Algorithm 3.