Discrete Optimization

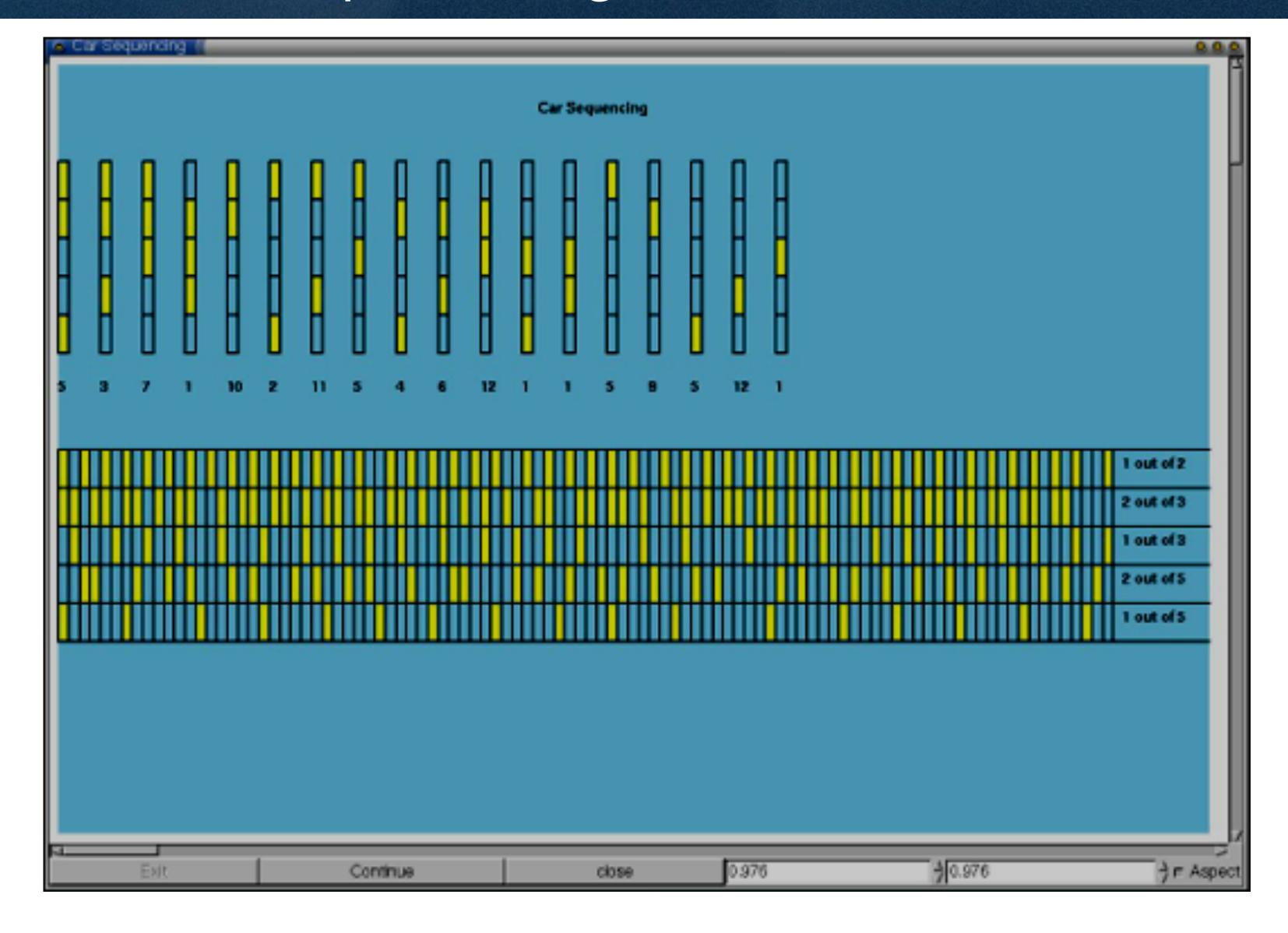
Local Search: Part II

Goals of the Lecture

- Local search
 - -swaps

- Cars on an assembly line
- Cars require specific options
 - -leather seats, moonroof
- Capacity constraints on the production units
 - at most 2 out of 5 successive cars can require a moonroof
- Sequence all the cars such that the capacity constraints are satisfied



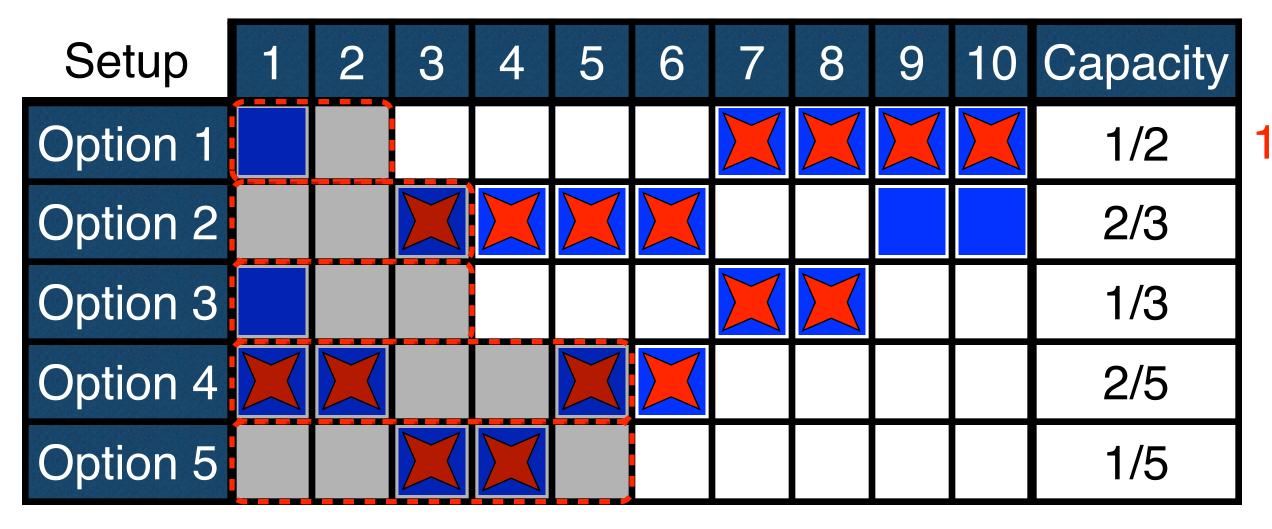


Swaps

- Neighborhood
 - -swap two configurations on the assembly line
- Search strategy
 - find a configuration that appears in violations
 - swap that configuration with another configuration to minimize the number of violations

Slots	1	2	3	4	5	6	7	8	9	10	Demand
Class 1											1
Class 2											1
Class 3											2
Class 4											2
Class 5											2
Class 6											2

			_	_		
Options	1	2	3	4	5	Demand
Class 1	yes		yes	yes		1
Class 2				yes		1
Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	



Slots	1	2	3	4	5	6	7	8	9	10	Demand
Class 1											1
Class 2											1
Class 3											2
Class 4											2
Class 5											2
Class 6											2
								¥			
Setup	1	2	3	4	5	6	7	8	9	10	Capacity
Option 1											1/2
Option 2											2/3

Options	1	2	3	4	5	Demand
Class 1	yes		yes	yes		1
Class 2				yes		1
Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	

3

2

2

1/3

2/5

1/5

2

3

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Option 3

Option 4

Option 5

1	2	3	4	5	6	7	8	9	10	Demand
										1
										1
										2
										2
										2
										2
		1 2	1 2 3	1 2 3 4	1 2 3 4 5	1 2 3 4 5 6 0	1 2 3 4 5 6 7 1	1 2 3 4 5 6 7 8 <t< td=""><td>1 2 3 4 5 6 7 8 9 Image: Control of the contro</td><td>1 2 3 4 5 6 7 8 9 10 1</td></t<>	1 2 3 4 5 6 7 8 9 Image: Control of the contro	1 2 3 4 5 6 7 8 9 10 1

Options	1	2	3	4	5	Demand
Class 1	yes		yes	yes		1
Class 2				yes		1
Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	

Setup	1	2	3	4	5	6	7	8	9	10	Capacity
Option 1											1/2
Option 2											2/3
Option 3											1/3
Option 4											2/5
Option 5											1/5

1

1

2

 C

Slots	1	2	3	4	5	6	7	8	9	10	Demand
Class 1											1
Class 2											1
Class 3											2
Class 4											2
Class 5											2
Class 6											2
					Y						

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Class 1	yes		yes	yes		1
Class 2				yes		1
Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	

Setup	1	2	3	4	5	6	7	8	9	10	Capacity
Option 1											1/2
Option 2											2/3
Option 3											1/3
Option 4											2/5
Option 5											1/5

1

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Class 1											1
Class 2											1
Class 3											2
Class 4											2
Class 5											2
Class 6											2

Options	1	2	3	4	5	Demand
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Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	

Setup	1	2	3	4	5	6	7	8	9	10	Capacity
Option 1											1/2
Option 2											2/3
Option 3											1/3
Option 4											2/5
Option 5											1/5

Slots	1	2	3	4	5	6	7	8	9	10	Demand
Class 1											1
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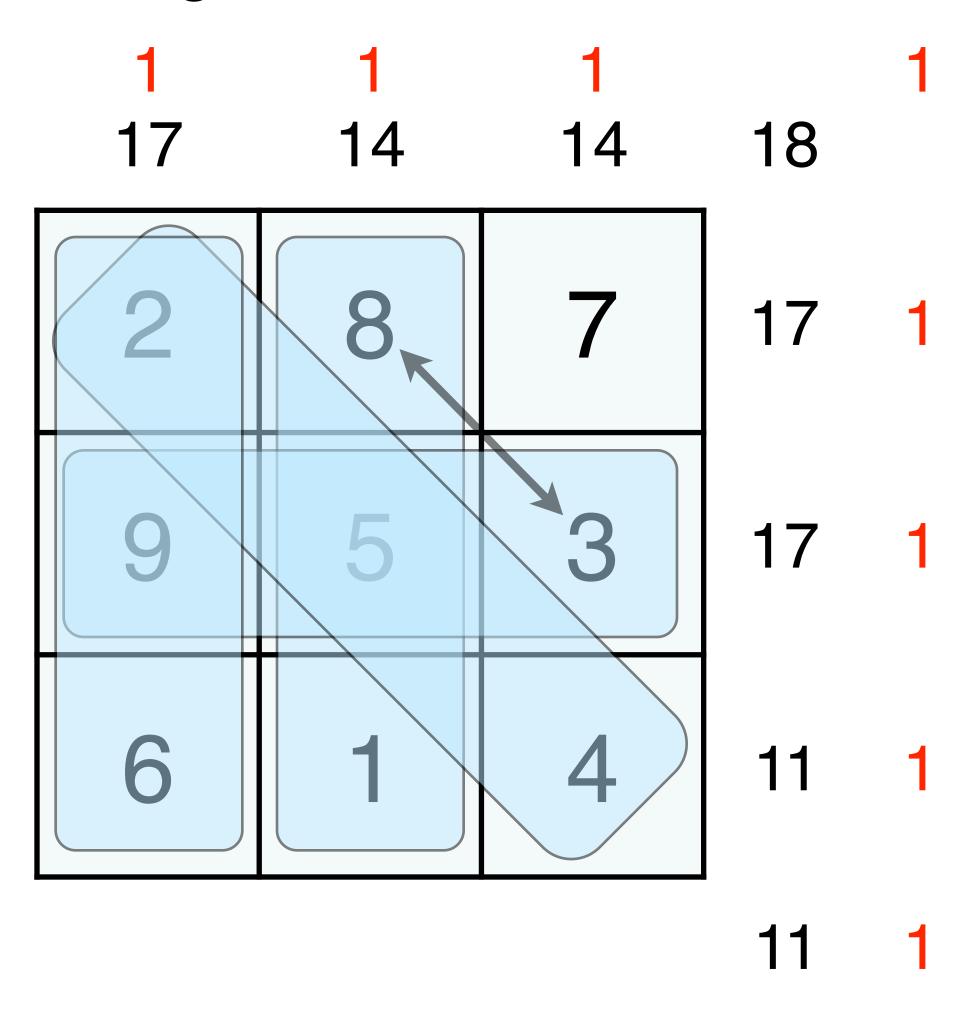
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Class 1	yes		yes	yes		1
Class 2				yes		1
Class 3		yes			yes	2
Class 4		yes		yes		2
Class 5	yes		yes			2
Class 6	yes	yes				2
Capacity	1/2	2/3	1/3	2/5	1/5	

Setup	1	2	3	4	5	6	7	8	9	10	Capacity
Option 1											1/2
Option 2											2/3
Option 3											1/3
Option 4											2/5
Option 5											1/5

Swaps

- Neighborhood
 - -swap two configurations on the assembly line
- Why swaps, not assignments
 - automatically maintain the demand constraint, that is the correct number of configurations is indeed produced
 - hard constraints
 - always feasible during the search
 - -soft constraints
 - may be violated during the search

```
range R = 1..n;
range D = 1..n^2;
int T = n*(n^2+1)/2;
var{int} s[R,R] in D;
solve {
   forall(i in R) {
      sum(j in R) s[i,j] = T;
      sum(j in R) s[j,i] = T;
   sum(i in R) s[i,i] = T;
   sum(i in R) s[i,n-i+1] = T;
   alldifferent(all(i in R, j in R) s[i,j]);
```



1 17	19	1 19	18	1
2	3	7	12	1
9	5	8	22	1
6	1	4	11	1
			11	1

Neighborhood for the Magic Square Problem

- Swapping the value of two cells
 - the alldifferent constraint as a hard constraint
 - swaps are used in many permutation problems
 - the inequalities are the soft constraints
- What the violations?
 - -0/1 violations would be pretty useless
 - many moves would not change the violations
 - -purely random walk
- ► For an equation I = r
 - -use abs(I r) as a measure of violations
 - -drives the search much more effectively

2 17	1 14	1 14	18	3	= 21
2	8	7	17	2	
9	5	3	17	3	
6	1	4	11	4	
			11	4	

1 16	0 15	1 14	18	3	= 14
2	9	7	18	3	
8	5	3	16	1	
6	1	4	11	4	
			11	4	

1 16	0 15	1 14	15	0	= 5
2	9	4	15	0	
8	5	3	16	1	
6	1	7	14	1	
			14	1	

0 15	0 15	0 15	15	0	= 0
2	9	4	15	0	
7	5	3	15	0	
6	1	8	15	0	
			15	0	

Until Next Time

Citations

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