# Discrete Optimization

Linear Programming: Part VI

## Goals of the Lecture

- Linear programming
  - -where is this dual coming from?
  - -what does it mean?
  - what is useful for?

can we find an upper bound?

$$10x_1 + 2x_2 + 6x_3 + 16x_4 \le 110$$

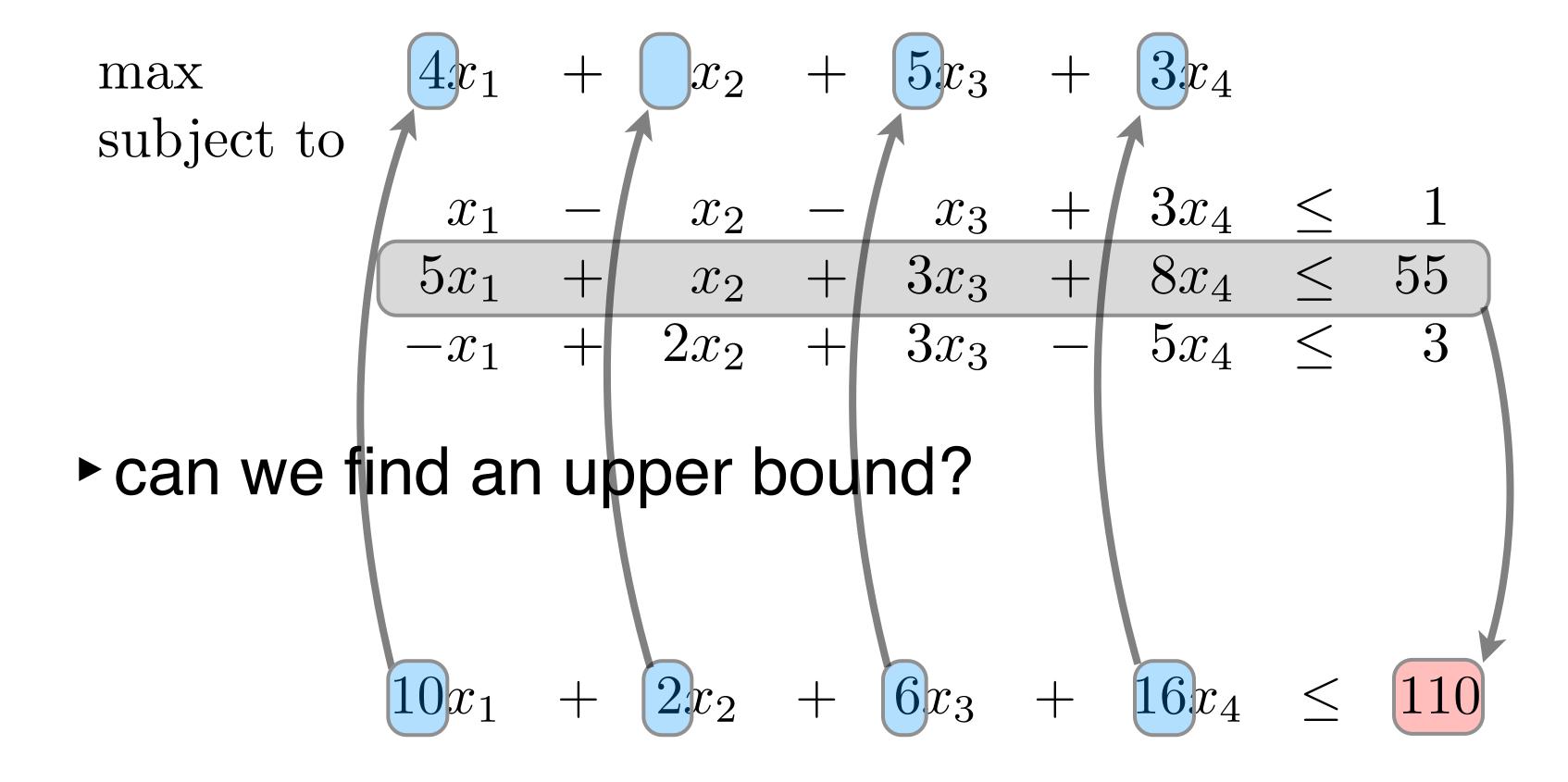
max 
$$4x_1 + x_2 + 5x_3 + 3x_4$$
 subject to 
$$x_1 - x_2 - x_3 + 3x_4 \le 1$$
 
$$5x_1 + x_2 + 3x_3 + 8x_4 \le 55$$
 
$$-x_1 + 2x_2 + 3x_3 - 5x_4 \le 3$$

► can we find an upper bound?

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$$y_1$$
 (  $x_1$  -  $x_2$  -  $x_3$  +  $3x_4$  ) +  $y_2$  (  $5x_1$  +  $x_2$  +  $3x_3$  +  $8x_4$  ) +  $y_3$  (  $-x_1$  +  $2x_2$  +  $3x_3$  -  $5x_4$  )  $\leq$   $y_1$  +  $55y_2$  +  $3y_3$ 

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minimize 
$$y_1 + 55y_2 + 3y_3$$

## Complementarity Slackness

#### What do these dual variables mean?

Let  $x^*$  and  $y^*$  be the optimal solutions to the primal and dual. The following conditions are necessary and sufficient for the optimality of  $x^*$  and  $y^*$ .

$$\sum_{j=1}^{n} a_{ij} x_j^* = b_i \quad \forall y_i^* = 0 \quad (1 \le i \le m)$$

and

$$\sum_{i=1}^{n} a_{ij} y_i^* = c_j \ \lor x_j^* = 0 \ (1 \le j \le n)$$

## Economic Interpretation

max 
$$\sum_{j=1}^{n} c_{j}x_{j}$$
 subject to 
$$\sum_{j=1}^{n} a_{ij}x_{j} \leq b_{i}+t_{i} \quad (1 \leq i \leq m)$$

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► for some small t<sub>i</sub>, this linear program has an optimal solution

$$z^* + \sum_{i=1}^m y_i^* t_i$$

optimal primal objective

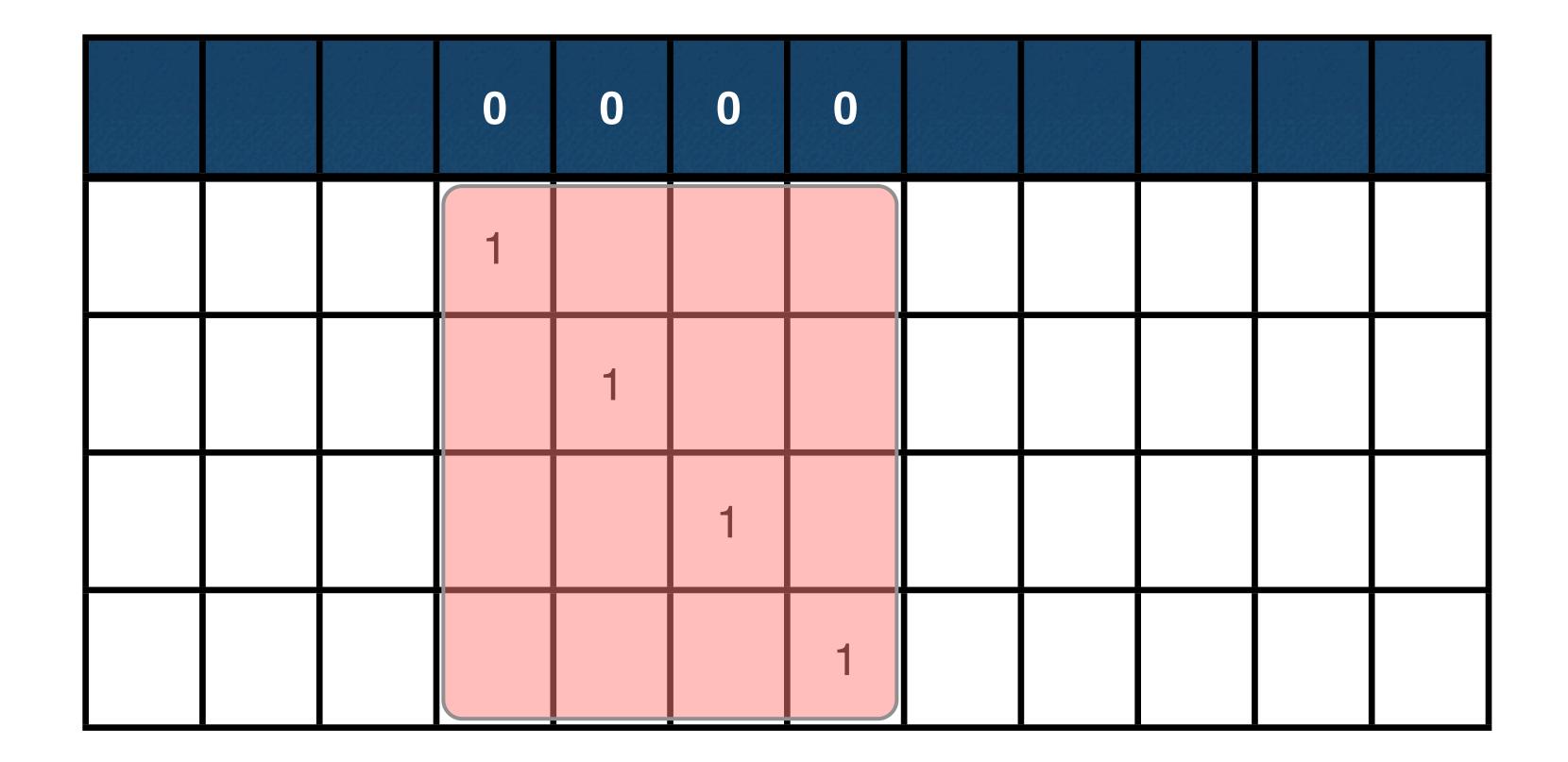
dual solution

# Duality in the Tableau

	0	0	0	0			
	1						
		1					
			1				
				1			

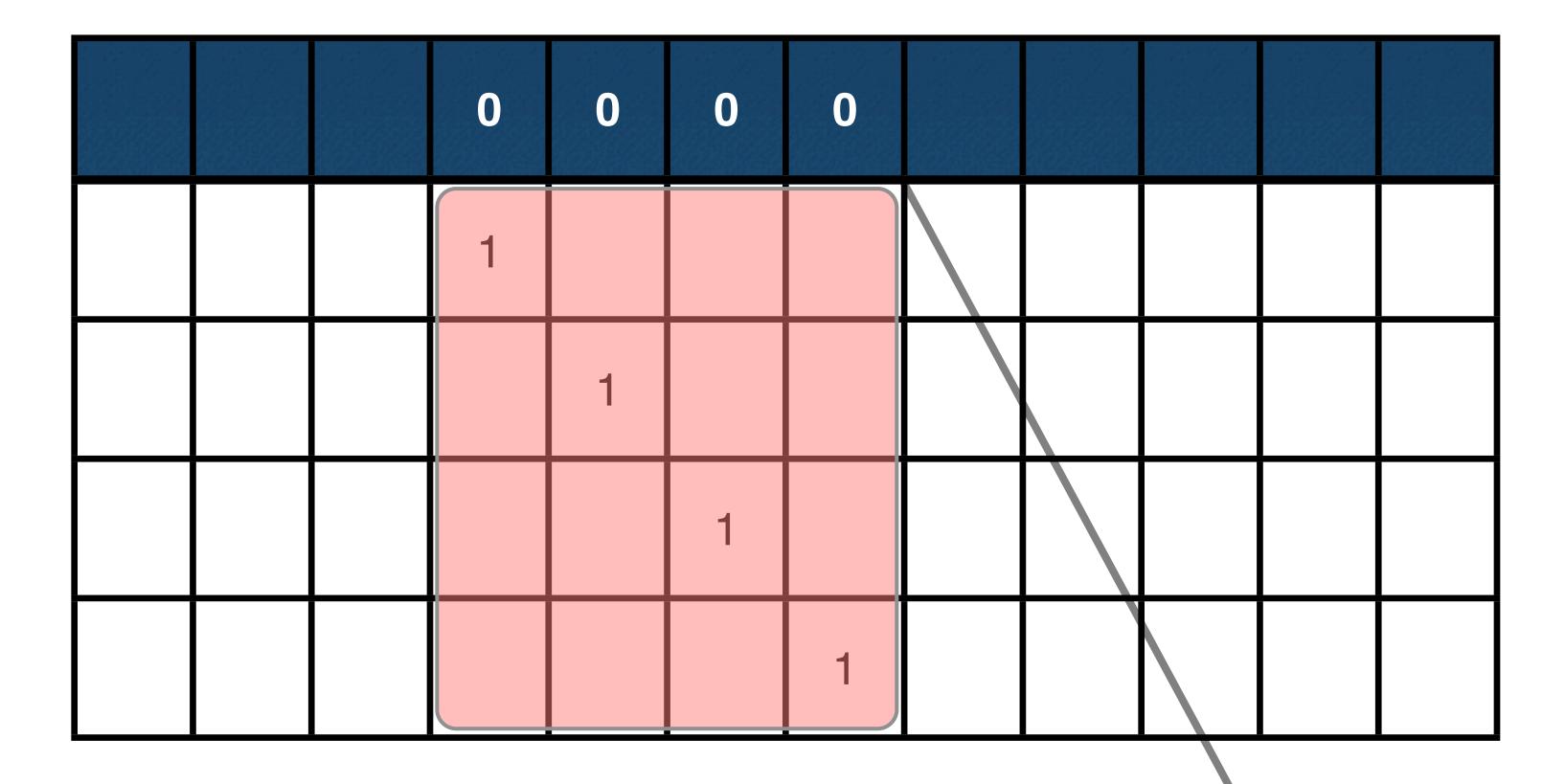
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# Duality in the Tableau



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## Duality in the Tableau



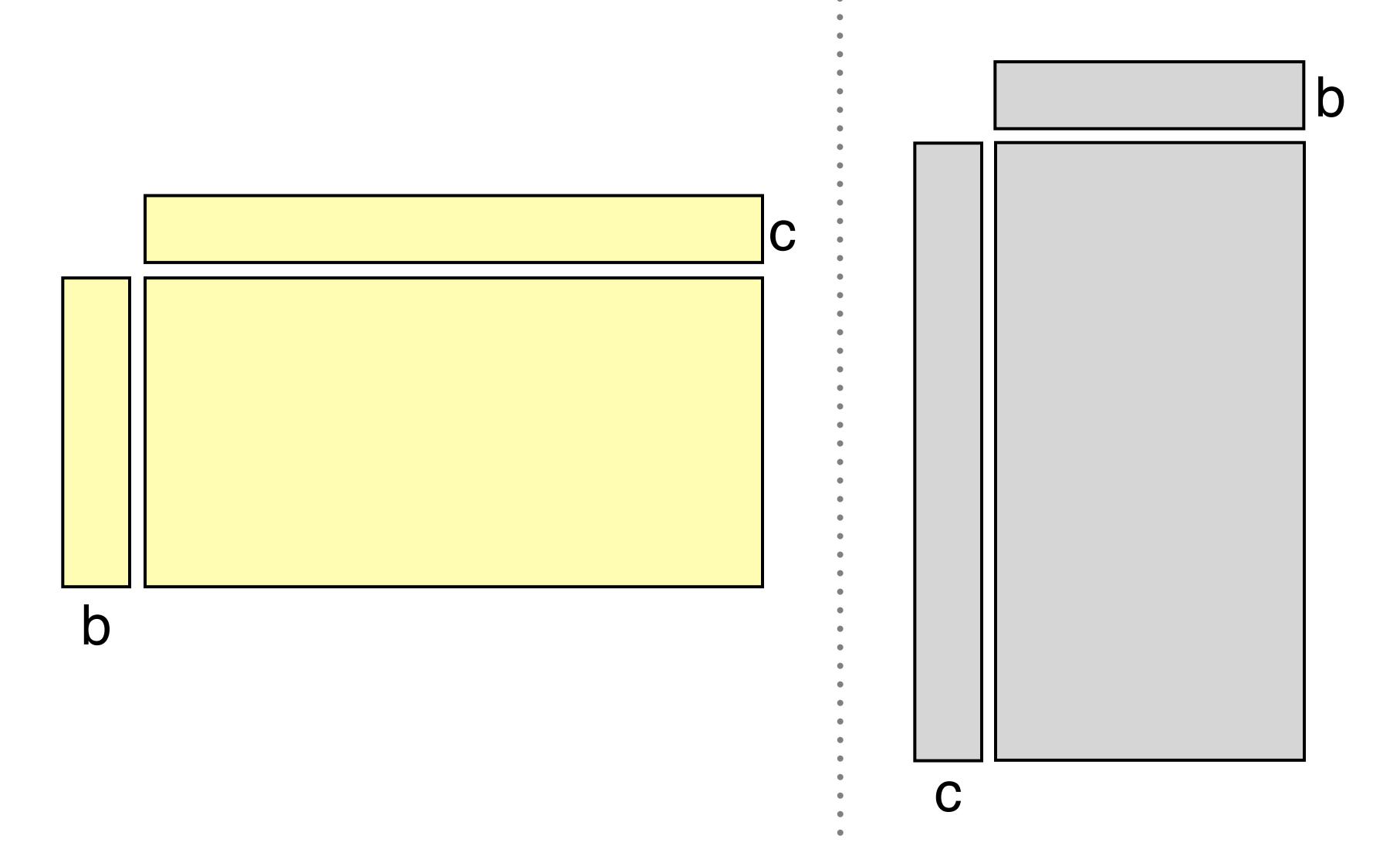
ullet at optimality, the cost become  $\,c_j - c_B A_B^{-1} A_j\,$ 

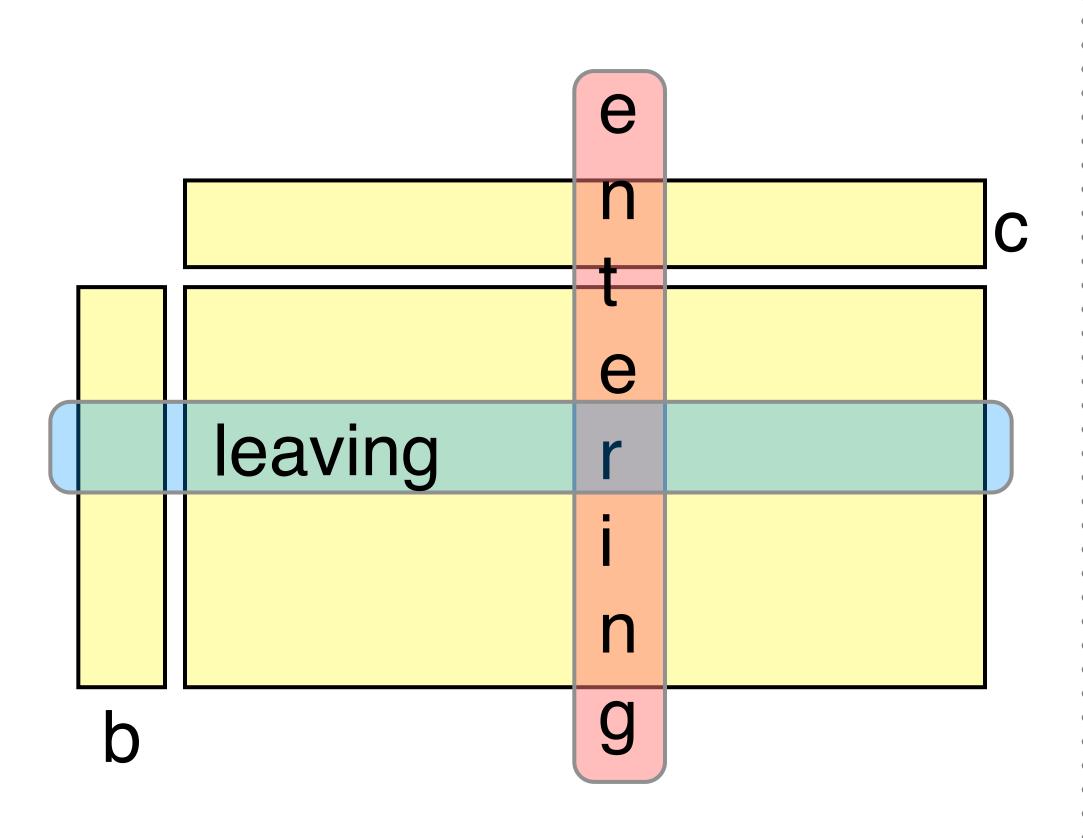
- ► Primal simplex
  - -primal always feasible; dual not

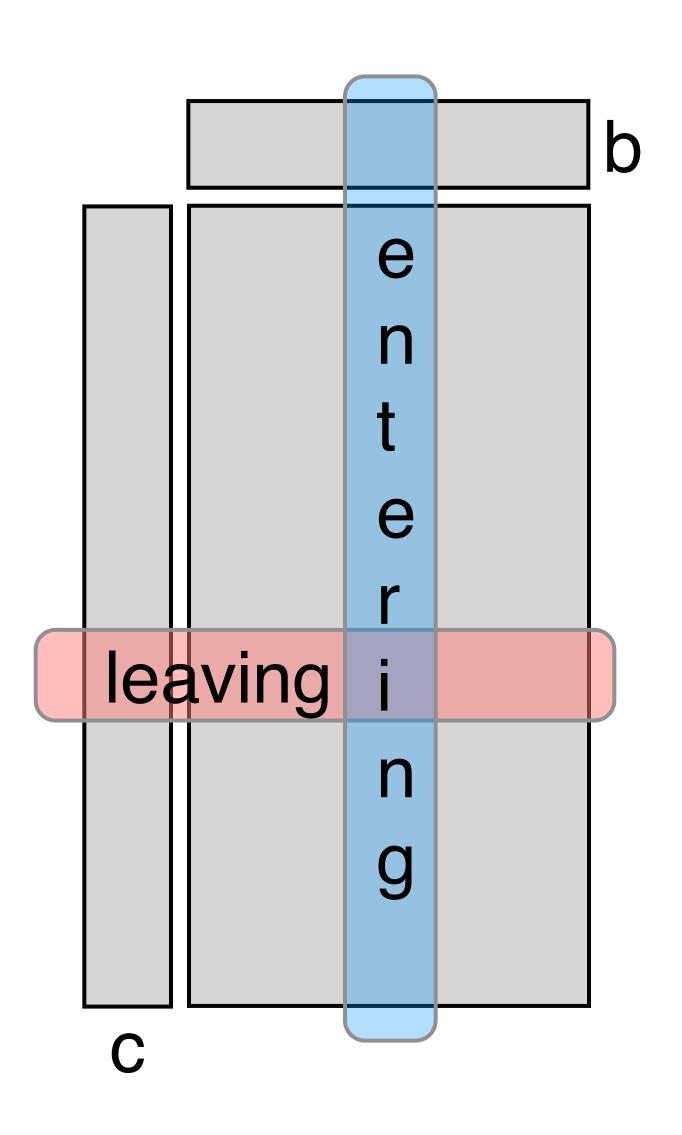
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- ► Use the same tableau







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## Until Next Time