# Discrete Optimization

Linear Programming: Part III

# Goals of the Lecture

- Linear programming
  - -the Simplex algorithm

# BFS and the Naive Algorithm

- 1. An optimal solution is located at a vertex.
- 2. A vertex is a Basic Feasible Solution (BFS).
- Naive algorithm
  - generate all basic feasible solutions
    - select m basic basic variables and perform Gaussian elimination
    - test whether it is feasible
  - -select the BFS with the best cost
- How many basic solutions?

$$\frac{n!}{m!(n-m)!}$$

# Outline of the Simplex Algorithm

Goal: You want to solve a linear program

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# The Simplex Algorithm

- Local search algorithm
  - -move from BFS to BFS
  - -guaranteed to find the global optimum
    - because of convexity

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- Local search algorithm
  - -move from BFS to BFS
  - -guaranteed to find the global optimum
    - because of convexity
- ► Key idea: how to move fro BFS to BFS.

- How to move to another BFS
  - select a non-basic variable with a negative coefficient: entering variable
  - introduce this variable in the basis by removing a basic variable: *leaving variable*
  - perform Gaussian elimination
- Local move: swap a basic and a non-basic variables

Not a BFS: I cannot select the leaving variable arbitrarily!

- ► How to choose the leaving variable?
- we must maintain feasibility

$$l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}}$$

### The Local Move

- Moving from BFS to BFS
  - -select the entering variable xe
    - non-basic variable with negative coefficients in the right-hand side
  - select the leaving variable x<sub>1</sub> to maintain feasibility

$$l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}}$$

- -apply Gaussian elimination
  - eliminate x<sub>e</sub> from the right-hand side
- This operation is called pivoting in linear programming
  - -pivot(e,l)

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# The Simplex Algorithm

A BFS is optimal if its objective function, after having eliminated all the basic variables, is of the form

$$c_0 + c_1 x_1 + \ldots + c_n x_n$$

with

$$c_i \ge 0 \ (1 \le i \le n).$$

min 
$$x_1 + x_2 + x_3 + x_4 + x_5$$
 subject to 
$$3x_1 + 2x_2 + x_3 = 1$$
$$5x_1 + x_2 + x_3 + x_4 = 3$$
$$2x_1 + 5x_2 + x_3 + x_5 = 4$$

First, find a BFS and eliminate the basic variable from the objective function

 $x_5 = 3 + x_1 - 3x_2$ 

```
min 6 - 3x_1 - 3x_2 subject to x_3 = 1 - 3x_1 - 2x_2 x_4 = 2 - 2x_1 + x_2 x_5 = 3 + x_1 - 3x_2
```

min 
$$6 - 3x_1 - 3x_2$$
 subject to  $x_3 = 1 - 3x_1 - 2x_2$   $x_4 = 2 - 2x_1 + x_2$   $x_5 = 3 + x_1 - 3x_2$ 

min  $6 - 3x_1 - 3x_2$  subject to  $x_3 = 1 - 3x_1 - 2x_2$  $x_4 = 2 - 2x_1 + x_2$  $x_5 = 3 + x_1 - 3x_2$ 

min 
$$6 - 3x_1 - 3x_2$$
 subject to 
$$x_3 = 1 - 3x_1 - 2x_2$$

$$x_4 = 2 - 2x_1 + x_2$$

$$x_5 = 3 + x_1 - 3x_2$$
min subject to 
$$x_2 = \frac{1}{2} - \frac{3}{2}x_1 + \frac{3}{2}x_3$$

$$x_4 = \frac{5}{2} - \frac{7}{2}x_1 - \frac{1}{2}x_3$$

$$x_5 = \frac{3}{2} + \frac{11}{2}x_1 + \frac{3}{2}x_3$$

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# Look Ma, My Local Move is Improving

- Assume that, in the BFS,
  - $-b_1>0, ..., b_m>0$
  - there exists an entering variable e with ce < 0
  - -there exists a leaving variable I
- then the move pivot(e,l) is improving

# The Simplex Algorithm

while  $\exists 1 \leq i \leq n : c_i < 0$  do choose e such that  $c_e < 0$ ;  $l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}};$  pivot(e,l);

# The Simplex Algorithm

```
while \exists 1 \leq i \leq n : c_i < 0 do choose e such that c_e < 0; l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}}; pivot(e,l);
```

- Assume that during the execution
  - -b<sub>1</sub>, ..., b<sub>m</sub> are always strictly positive
  - the objective function is bounded by below
- then the simplex algorithm terminates with an optimal solution

Selecting the leaving variable

$$l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}}$$

min 
$$6 - 3x_1 - 3x_2$$
 subject to  $x_3 = 1 + 3x_1 - 2x_2$   $x_4 = 2 + 2x_1 + x_2$   $x_5 = 3 + x_1 - 3x_2$ 

► There are no leaving variables

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$$l = \underset{i:a_{ie} < 0}{\operatorname{arg-min}} \frac{b_i}{-a_{ie}}$$

► There are no leaving variables

What is the basic solution?

$$\{x_1 = 0; x_2 = 0; x_3 = 1; x_4 = 2; x_5 = 3\}$$

- ► What happens if I increase the value of x<sub>1</sub>
  - the solution remains feasible
  - the value of the objective can decrease arbitrarily

► What if some b<sub>i</sub> becomes zero?

min 
$$5 - 3x_1 - 3x_2$$
 subject to  $x_3 = 0 - 3x_1 - 2x_2$   $x_4 = 2 - 2x_1 + x_2$   $x_5 = 3 + x_1 - 3x_2$ 

What is the leaving variable?

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We need to find a new way to guarantee termination

# Termination of the Simplex Algorithm

### Many approaches

- -Bland rule
  - select always the first entering variable lexicographically
- Lexicographic pivoting rule
  - break ties when selecting the leaving variable by using a lexicographic rule

$$l = \underset{i:a_{ie} < 0}{\operatorname{arg-lex-min}} \frac{b_i}{-a_{ie}}$$

- Perturbation methods

min 
$$x_1 + x_2 + x_3 + x_4 + x_5$$
 subject to 
$$3x_1 + 2x_2 + x_3 = 1$$
$$5x_1 + x_2 + x_3 + x_4 = 3$$
$$2x_1 + 5x_2 + x_3 + x_5 = 4$$

► How do I find my first BFS?

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min 
$$c_1x_1 + \ldots + c_nx_n$$
 subject to 
$$a_{11}x_1 + \ldots + a_{1n}x_n = b_1$$
 
$$\ldots$$
 
$$a_{m1}x_1 + \ldots + a_{mn}x_n = b_m$$

► Introduce artificial variables

Introduce artificial variables

- ► I have an easy BFS
- But to another problem

► Introduce artificial variables

min 
$$c_1x_1 + \ldots + c_nx_n$$
 subject to 
$$a_{11}x_1 + \ldots + a_{1n}x_n + y_1 = b_1$$
 
$$\ldots$$
 
$$a_{i1}x_1 + \ldots + a_{in}x_n + y_i = b_i$$
 
$$\ldots$$
 
$$a_{m1}x_1 + \ldots + a_{mn}x_n + y_m = b_m$$

► I have an easy BFS



But to another problem

► Introduce artificial variables

► I have an easy BFS



But to another problem



### Two-Phase Method

- ► First find a BFS
  - if there is one
- ► Then find an optimal BFS

### Two-Phase Method

### ► First find a BFS

- ► Feasible if the objective value reaches 0
  - all the yi are thus zeros
    - I have a BFS without the yi
    - really? always?

# Until Next Time