

1. There are 6 primary classes in total:
 - **BinaryNode**: Represents a node stored in an unbalanced binary search tree. It contains the references to a city record, the left and the right subtrees.
 - **City**: Represents a city record in the database. A city record contains the information about it, which are its name, population and payload. Every city record is referenced by one of the node in the binary search trees.
 - **DatabaseBST**: Represent a generic Binary Search Tree that references to the database in an efficient way for later searching.
 - **NameComp**: Represent a comparator for the Name Binary Search Tree. It will be used for making decisions in directions (left or right) while going from the top to the bottom of the tree.
 - **PopulationComp**: Represent a comparator for the Name Binary Search Tree. It will be used to make decisions in directions (left of right) while going from the top to the bottom of the tree.
 - **bst**: The main program that reads the inputs and return the outputs accordingly.

2. In order to support the findKth method:
 - I created one more field in the BinaryNode class: `int leftSize`, which counts the number of nodes in the left subtree of the current node. The size of the left subtree is increased by 1 if a node is inserted onto the left of the current node. Similarly, the size of the left subtree is decreased by 1 if a node is deleted from the left of the current node. Some small modifications in the `insert()` and `delete()` methods helps me to achieve that without affecting the complexity of those 2 methods. Finally, in the `findKth()` method, I recursively go down the tree until `leftSize == K` (since it's 0 indexed) and return the result. Therefore, the overall complexity for the `findKth()` method is $O(\log n)$ with n is the number of nodes on the trees.

3. Average case running times of all operations (suppose that n is the number of nodes on the trees and the trees are kept being balanced):
 - **insert**: it keeps going down the tree until it finds an appropriate place to add a node in. The height of the tree is about $\log n$, so the complexity is $O(\log n)$.
 - **find**: it keeps going down the tree until it finds all of the desired duplicates. The height of the tree is about $\log n$, so the complexity is $O(\log n)$.
 - **findKth**: as explained above in 2., so the complexity is $O(\log n)$
 - **findRange**: finds all the city records whose values are within the range. It changes direction right away when there's some value that is out of the range. Worst case is $O(n)$ when the whole tree is within the given range.
 - **delete**: first, calls the find operation to obtain a list of all desired duplicates. Then, use the duplicates' addresses to remove all of them out of the two trees. The height of the tree is about $\log n$, so its complexity is $O(\log n)$.
 - **sort**: do an in-order traversal. The complexity is $O(n)$

- tree: also do an in-order traversal. The complexity is $O(n)$
- makenull: empty the tree. The complexity is $O(1)$