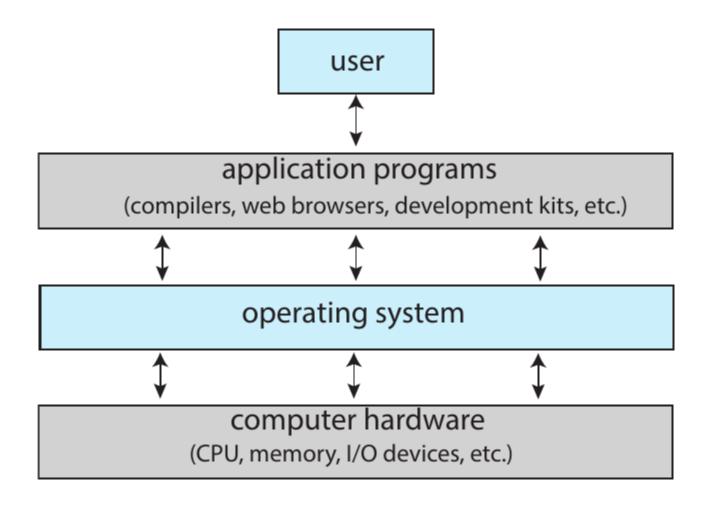
# **Operating Systems**

Isfahan University of Technology Electrical and Computer Engineering Department

Zeinab Zali

Session 2: Introduction to Operating System & Computer Environments

#### Four Components of a Computer System



## **Computer System Structure**

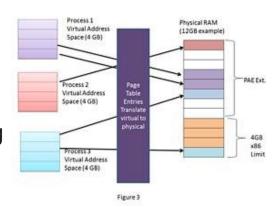
- Computer system can be divided into four components
  - Hardware provides basic computing resources
    - ▶ CPU, memory, I/O devices
  - Operating system
    - Controls and coordinates use of hardware among various applications and users
  - Application programs define the ways in which the system resources are used to solve the computing problems of the users
    - Word processors, compilers, web browsers, database systems, video games
  - Users
    - People, machines, other computers

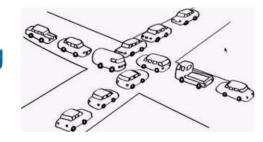
## Whats is an OS?

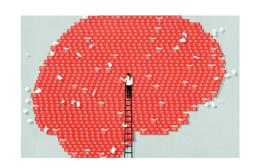
- There is a body of software, in fact, that is responsible for
  - making it easy to run programs (even allowing you to seemingly run many at the same time)
  - allowing programs to share memory and CPU
  - enabling programs to interact with devices
  - and other fun stuff
- That body of software is called the operating system (OS)
- OS is in charge of making sure the system operates correctly and efficiently in an easy-to-use manner

# OS three pieces

- Virtualization
  - A primary way the OS does its roles
    - •Specially its key role as a resource manager for managing two main resource **CPU** and **Memory**
- Concurrency
  - A conceptual term to refer to a host of problems that arise, and must be addressed, when working on many things at once in the same program
    - OS itself
    - Multi-thread or multi-process programs
- Persistence
  - \*storing any files the user creates in a reliable and efficient manner on the disks of the system







# Virtualization

- Process: OS abstraction of the processor, main memory and I/O devices for a running program
  - Multiple processes can concurrently run, each thinking itself as the exclusive user of the hardware.
- Virtual Memory: OS abstraction of Memory
  - Each process perceives the same picture of memory used only by itself (its address space).
- File: OS abstraction of I/O devices
  - All input and output in the system is performed by reading and writing files

# **CPU Virtualization**

 virtualizing the CPU: Turning a single CPU (or small set of them) into a seemingly infinite number of CPUs and thus allowing many programs to seemingly run at once

```
#include <stdio.h>
#include <stdlib.h>
#include "common.h"

int main(int argc, char *argv[])
{
    if (argc != 2) {
        fprintf(stderr, "usage: cpu <string>\n");
        exit(1);
    }
    char *str = argv[1];

    while (1) {
        printf("%s\n", str);
        Spin(1);
    }
    return 0;
}
```

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        fprintf(stderr, "usage: cpu <string>\n");
        exit(1);
    }
    char *str = argv[1];

    while (1) {
        printf("%s\n", str);
        Spin(1);
    }
    return 0;
}
```

```
zeinabzali:./cpu A & ./cpu B & ./cpu C
[1] 27878
[2] 27879
A
B
C
A
B
C
A
B
C
A
B
C
A
B
C
A
B
C
A
B
C
A
```

- Memory is just an array of bytes
  - to read memory, one must specify an address to be able to access the data stored there
  - to write (or update) memory, one must also specify the data to be written to the given address
- A program keeps all of its data structures in memory, and accesses them through various instructions
  - loads and stores
- Each instruction of the program is in memory too
  - thus memory is accessed on each instruction fetch

 virtualizing the Memory: Each process accesses its own private virtual address space which the OS somehow maps onto the physical memory of the machine

```
#include <unistd.h>
#include <stdio.h>
#include <stdlib.h>
#include "common.h"
int main(int argc, char *argv[]) {
    if (argc != 2) {
        fprintf(stderr, "usage: mem <value>\n");
        exit(1):
    int *p:
    p = malloc(sizeof(int));
    assert(p != NULL);
    printf("(%d) addr pointed to by p: %p\n", (int) getpid(), p);
    *p = atoi(argv[1]); // assign value to addr stored in p
    //p = atoi(argv[1]);
    while (1) {
        Spin(1);
        *p = *p + 1;
        printf("(%d) value of p: %d\n", getpid(), *p);
    return 0;
```

virtualizing the Memory: Each process accesses
its own private virtual address space which the
OS somehow maps onto the physical memory of
the machine setarch \$(uname --machine) --addr-no-randomize /bin/bash

```
#include <unistd.h>
#include <stdio.h>
                                              zeinabzali:./mem 10
#include <stdlib.h>
#include "common.h"
                                               (28692) addr pointed to by p: 0x555555756260
                                               (28692) value of p: 11
int main(int argc, char *argv[]) {
   if (argc != 2) {
                                               (28692) value of p: 12
       fprintf(stderr, "usage: mem <value>\n");
       exit(1);
                                               (28692) value of p: 13
                                                28692) value of p: 14
   int *p:
   p = malloc(sizeof(int));
   assert(p != NULL);
   printf("(%d) addr pointed to by p: %p\n", (int) getpid(), p);
   *p = atoi(argv[1]); // assign value to addr stored in p
   //p = atoi(argv[1]);
   while (1) {
       Spin(1);
       *p = *p + 1;
```

printf("(%d) value of p: %d\n", getpid(), \*p);

return 0;

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#include <stdio.h>
#include <stdlib.h>
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    *p = atoi(argv[1]); // assign value to addr stored in p
    //p = atoi(argv[1]);
    while (1) {
        Spin(1);
        *p = *p + 1;
        printf("(%d) value of p: %d\n", getpid(), *p);
    return 0;
```

```
zeinabzali:./mem 10 & ./mem 100 & ./mem 1000
[1] 28699
[2] 28700
(28699) addr pointed to by p: 0x555555756260
(28700) addr pointed to by p: 0x555555756260
(28701) addr pointed to by p: 0x555555756260
(28699) value of p: 11
(28700) value of p: 101
(28701) value of p: 1001
(28699) value of p: 12
(28700) value of p: 102
(28701) value of p: 1002
(28699) value of p: 13
(28700) value of p: 103
(28701) value of p: 1003
(28699) value of p: 14
(28700) value of p: 104
(28701) value of p: 1004
(28699) value of p: 15
(28700) value of p: 105
(28701) value of p: 1005
```

# Concurrency

 problems that arise, and must be addressed, when working on many things at once (i.e., concurrently)

```
#include <stdio.h>
#include <stdlib.h>
#include "common.h"
#include "common threads.h"
volatile int counter = 0;
int loops;
void *worker(void *arg) {
    int i:
    for (i = 0; i < loops; i++) {
        counter++;
    return NULL;
int main(int argc, char *argv[]) {
    if (argc != 2) {
        fprintf(stderr, "usage: threads <loops>\n");
        exit(1);
    loops = atoi(argv[1]);
    pthread t p1, p2;
    printf("Initial value : %d\n", counter);
    Pthread_create(&p1, NULL, worker, NULL);
    Pthread create(&p2, NULL, worker, NULL);
    Pthread join(p1, NULL);
    Pthread join(p2, NULL);
    printf("Final value : %d\n", counter);
    return 0:
```

# Concurrency

 problems that arise, and must be addressed, when working on many things at once (i.e., concurrently)

```
#include <stdio.h>
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    pthread t p1, p2;
    printf("Initial value : %d\n", counter);
    Pthread_create(&p1, NULL, worker, NULL);
    Pthread create(&p2, NULL, worker, NULL);
    Pthread join(p1, NULL);
    Pthread join(p2, NULL);
    printf("Final value : %d\n", counter);
    return 0:
```

```
zeinabzali:./threads 10
Initial value : 0
Final value : 20
zeinabzali:./threads 100
Initial value : 0
Final value : 200
zeinabzali:./threads 1000
Initial value : 0
Final value : 2000
zeinabzali:10000
10000: command not found
zeinabzali:./threads 10000
Initial value : 0
Final value : 0
Final value : 14991
zeinabzali:
```

# OS design goals

- build up some abstractions in order to make the system convenient and easy to use.
  - Abstractions are fundamental to everything we do in computer science
- high performance; minimizing the overheads of the OS
- Protection between applications, as well as between the OS and applications
  - making sure that the malicious or accidental bad behavior of one program does not harm others;
  - isolating processes
- Reliability; OS must run non-stop; when it fails, all applications running on the system fail as well
- Security; an extension of protection, really against malicious applications especially in these highly-networked times
- Others: mobility, energy efficiency, ...

## Sum up: What is an Operating System?

- ✓ A program that acts as an intermediary between a user of a computer and the computer hardware.
- Operating system operations
  - Process Management
  - Memory Management
  - Storage Management
  - I/O handling
  - Protection and Security
    - user-ID, Group-ID, permission
    - Viruses, attacks, intrusion

# **Computer Environments**

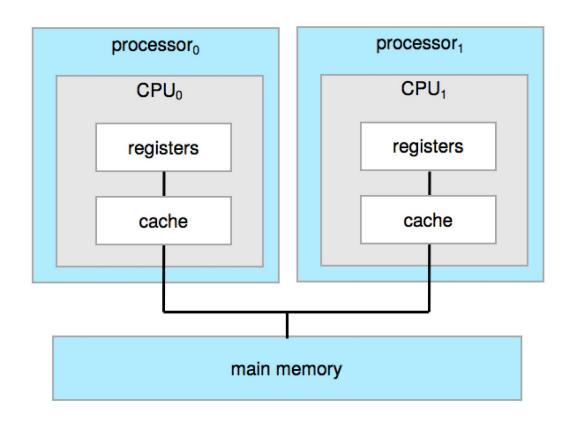


### **Multiprocessors**

- Multiprocessors systems growing in use and importance
  - Also known as parallel systems, tightly-coupled systems
  - Advantages include:
    - 1. Increased throughput
    - 2. Economy of scale
    - 3. **Increased reliability** graceful degradation or fault tolerance
  - Two types:
    - Asymmetric Multiprocessing each processor is assigned a specie task.
    - 2. Symmetric Multiprocessing each processor performs all tasks





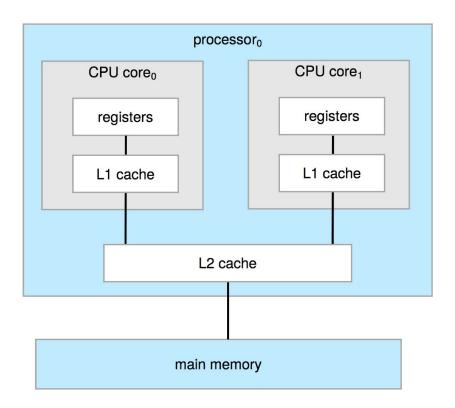






## **Dual-Core Design**

- Multi-chip and multicore
- Systems containing all chips
  - Chassis containing multiple separate systems



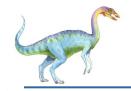




## **Distributed Systems**

- Collection of separate, possibly heterogeneous, systems networked together
  - Network is a communications path, TCP/IP most common
    - Local Area Network (LAN)
    - Wide Area Network (WAN)
    - Metropolitan Area Network (MAN)
    - Personal Area Network (PAN)
- Network Operating System provides features between systems across network
  - Communication scheme allows systems to exchange messages
  - Illusion of a single system





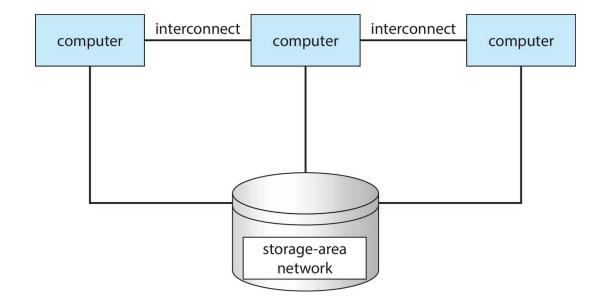
# **Clustered** Systems

- Like multiprocessor systems, but multiple systems working together
  - Usually sharing storage via a storage-area network (SAN)
  - Provides a high-availability service which survives failures
    - Asymmetric clustering has one machine in hot-standby mode
    - Symmetric clustering has multiple nodes running applications, monitoring each other
  - Some clusters are for high-performance computing (HPC)
    - Applications must be written to use parallelization
  - Some have distributed lock manager (DLM) to avoid conflicting operations

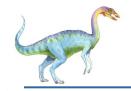




## **Clustered Systems**

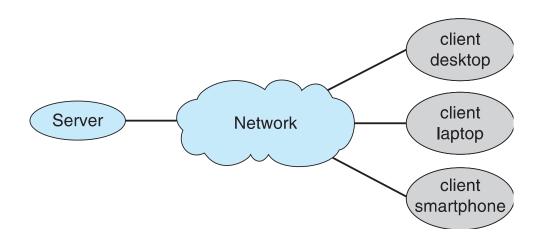




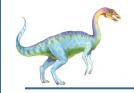


#### **Client Server**

- Client-Server Computing
  - Dumb terminals supplanted by smart PCs
  - Many systems now servers, responding to requests generated by clients
    - Compute-server system provides an interface to client to request services (i.e., database)
    - File-server system provides interface for clients to store and retrieve files

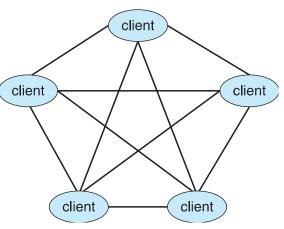




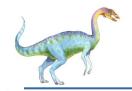


#### Peer-to-Peer

- Another model of distributed system
- P2P does not distinguish clients and servers
  - Instead all nodes are considered peers
  - May each act as client, server or both
  - Node must join P2P network
    - Registers its service with central lookup service on network, or
    - Broadcast request for service and respond to requests for service via discovery protocol
  - Examples include Napster and Gnutella,
     Voice over IP (VoIP) such as Skype







#### **Cloud Computing**

- Delivers computing, storage, even apps as a service across a network
- Logical extension of virtualization because it uses virtualization as the base for its functionality.
  - Amazon EC2 has thousands of servers, millions of virtual machines, petabytes of storage available across the Internet, pay based on usage



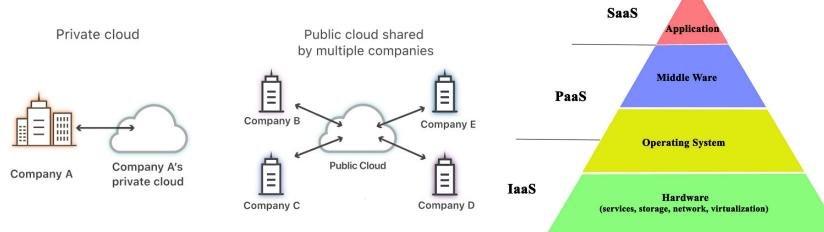




#### **Cloud Computing types**

- Public cloud available via Internet to anyone willing to pay
- Private cloud run by a company for the company's own use
- Hybrid cloud includes both public and private cloud components
- Software as a Service (SaaS) one or more applications available via the Internet (i.e., word processor)
- Platform as a Service (PaaS) software stack ready for application use via the Internet (i.e., a database server)

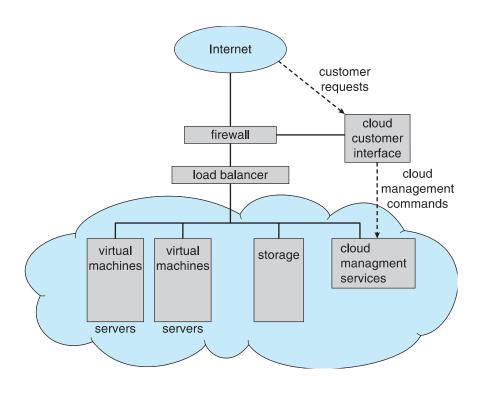
Infrastructure as a Service (laas) – servers or storage available over





#### **Cloud Computing Environments**

- Cloud computing environments composed of traditional OSes, plus VMMs, plus cloud management tools
  - Internet connectivity requires security like firewalls
  - Load balancers spread traffic across multiple applications







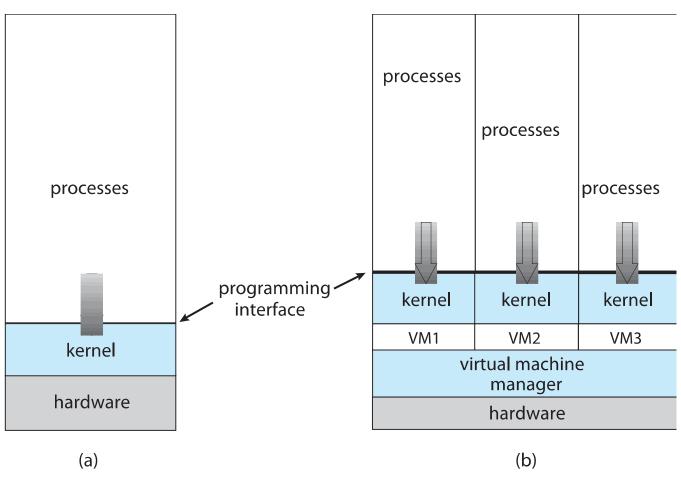
#### **Virtualization**

- Virualization is a technology that allows us to abstract the hardware of a single computer (the CPU, memory, disk drives, network interface cards, and so forth) into several different execution environments, thereby creating the illusion that each separate environment is running on its own private computer.
- A user of a virtual machine can switch among the various operating systems in the same way a user can switch among the various processes running concurrently in a single operating system.
- Emulation is simulating computer hardware in software,
- Broadly speaking, virtualization software is one member of a class that also includes emulation.
  - Emulation is typically used when the source CPU type is different from the target CPU type.





## **Computing Environments - Virtualization**





### Container-based Virtualization

