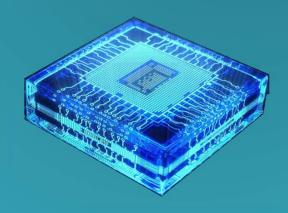




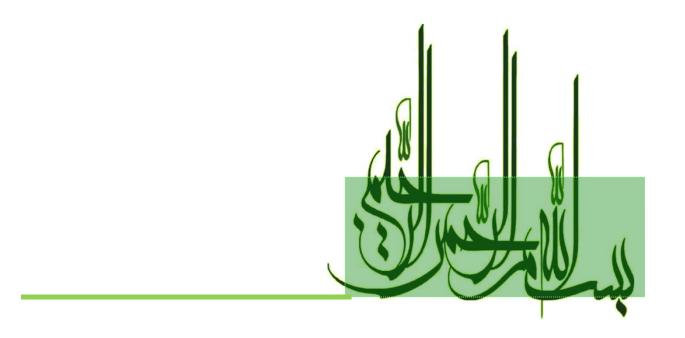
Microprocessors and Assembly language

Isfahan University of Technology (IUT)



Jump And Call

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Topics

- Introduction to jump and call
- Jump
- Call
 - Stack
 - Calling a function
- Time Delay



Jump and Call

- CPU executes instructions one after another.
 - For example in the following C program, CPU first executes the instruction of line 3 (adds b and c), then executes the instruction of line 4.

```
void main ()
   a = b + c;
   c = 2;
   d = a + c;
```



Jump and Call

- CPU executes instructions one after another.
 - For example in the following C program, CPU first executes the instruction of line 3 (adds b and c), then executes the instruction of line 4.

```
void main ()
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   c -= 2;
   d = a + c;
```



Jump and Call

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```
void main ()
   a = b + c;
   c -= 2;
   d = a + c;
```



- But sometimes we need the CPU to execute, an instruction other than the next instruction. For example:
 - When we use a conditional instruction (if)
 - When we make a loop
 - When we <u>call</u> a function



- Example 1: Not executing the next instruction, because of condition.
 - In the following example, the instruction of line 6 is not executed.

```
1  void main ()
2  {
3    int a = 2;
4    int c = 3;
5    if (a == 8)
6        c = 6;
7    else
8        c = 7;
9    c = a + 3;
}
```



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 - In the following example, the order of execution is as follows:
 - Line 4
 - Line 5
 - Again, line 4
 - Again line 5
 - Line 6

```
1  void main ()
2  {
3    int a, c = 0;
4    for(a = 2; a < 4; a++)
5        c += a;
6        a = c + 2;
7    }
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```



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```



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 - In the following example, the instruction of line 6 is not executed after line 5.

```
Code
     void func1 ();
     void main ()
         int a = 2, c = 3;
         func1 ();
         c = a + 3;
     void func1 () {
         int d = 5 / 2;
10
11
```



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11
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- In the assembly language, there are
 2 groups of instructions that make the CPU execute
 an instruction other than the next instruction.
- The instructions are:
 - Jump: used for making loop and condition
 - Call: used for making function calls



Jump

• Jump changes the Program Counter (PC) and causes the CPU to execute an instruction other than the next instruction.



Jump

There are 2 kinds of Jump

- Unconditional Jump: When CPU executes an unconditional jump, it jumps unconditionally (without checking any condition) to the target location.
 - Example: RJMP and JMP instructions
- Conditional Jump: When CPU executes a conditional jump, it checks a condition, if the condition is true then it jumps to the target location; otherwise, it executes the next instruction.



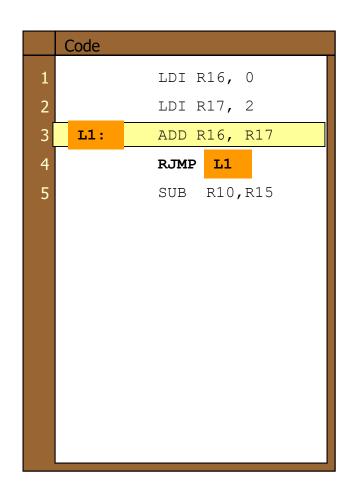
Unconditional Jump in AVR

 There are 3 unconditional jump instructions in AVR:

JMP, RJMP, and IJMP

Note

- We label the location where we want to jump, using a unique name, followed by ':'
- Then, in front of the jump instruction we mention the name of the label.
- This causes the CPU to jump to the location we have labeled, instead of executing the next instruction.





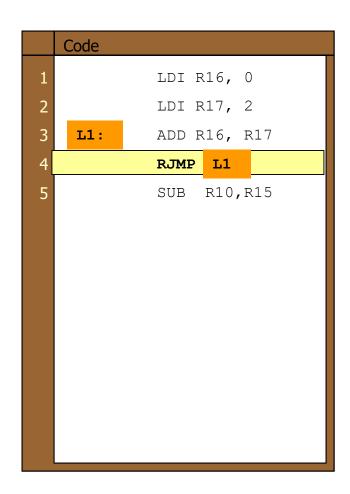
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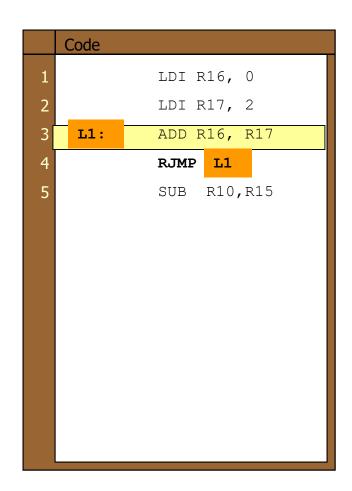
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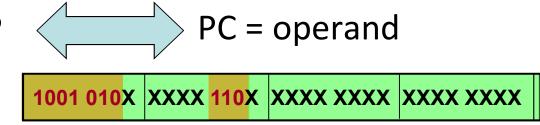


Ways of specifying the jump target

- There are 3 ways to provide the jump address:
 - PC = operand
 - PC = PC + operand
 - PC = Z register



• JMP



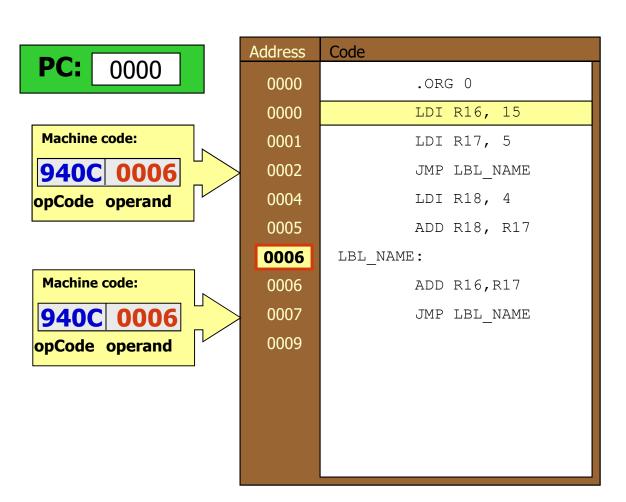
– Example:

1001 0100	0000 1100	0000 0000	0000 0110
-----------	-----------	-----------	-----------

- Operand = 00000000000000000110
- 4 byte(why?)

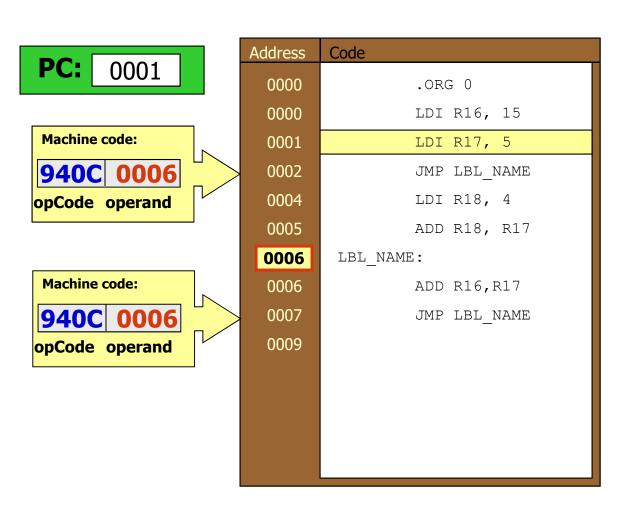


- In JMP, the operand, contains the address of the destination
- When an JMP is executed:
 - PC is loaded with the operand value



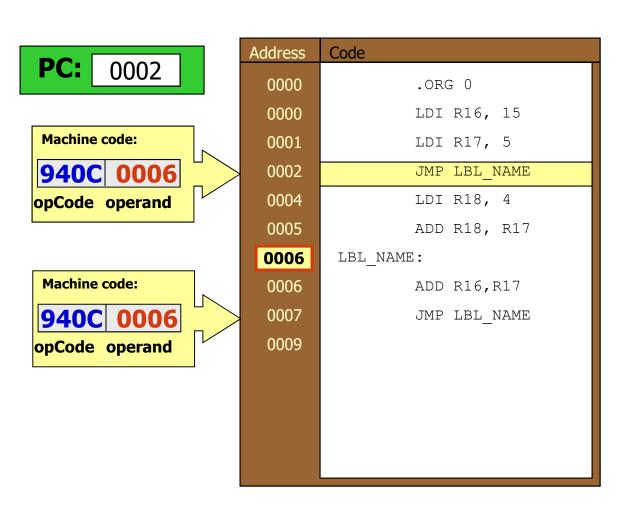


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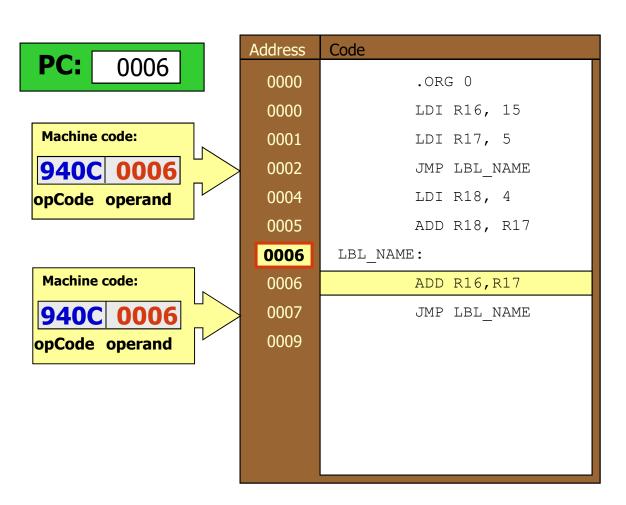
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JMP

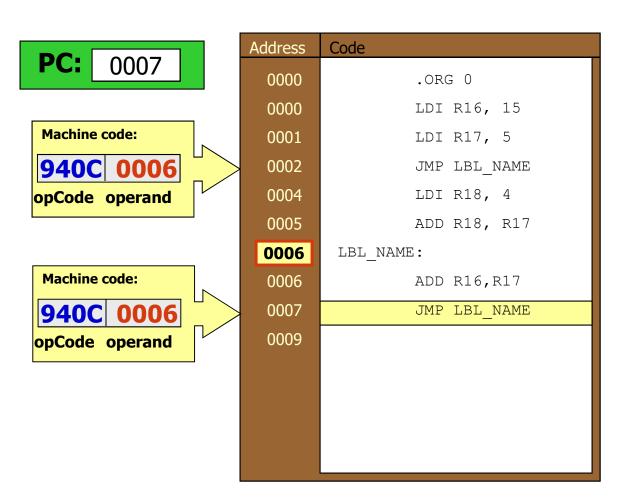
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JMP

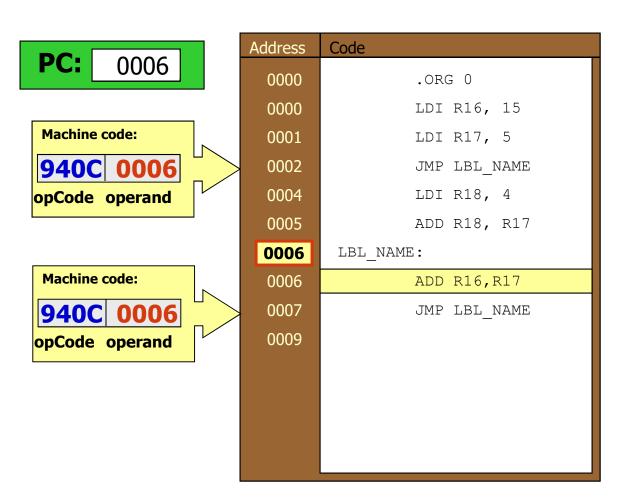
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JMP

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RJMP (Relative jump)

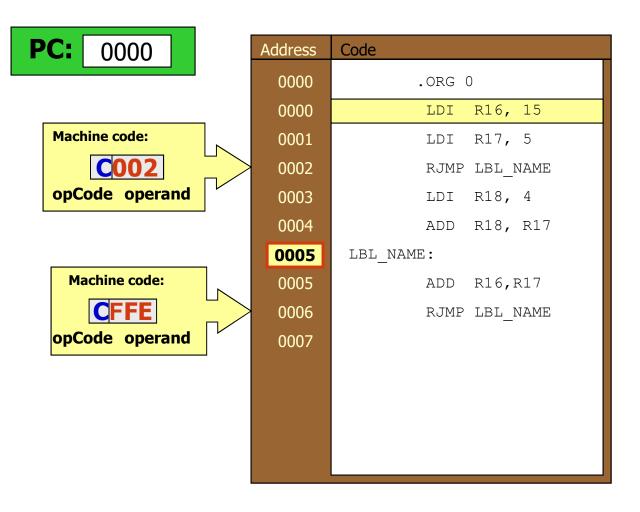
1100 XXXX XXXX XXXX

-Example: 1100 0000 0000 0110

- Operand = 00000000110
- PC = PC + 00000000110

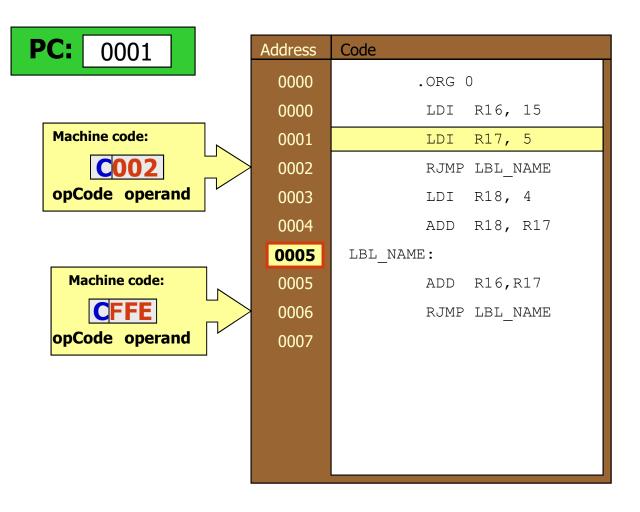


- When RJMP is executed:
 - The operand
 will be added
 to the current
 value of PC



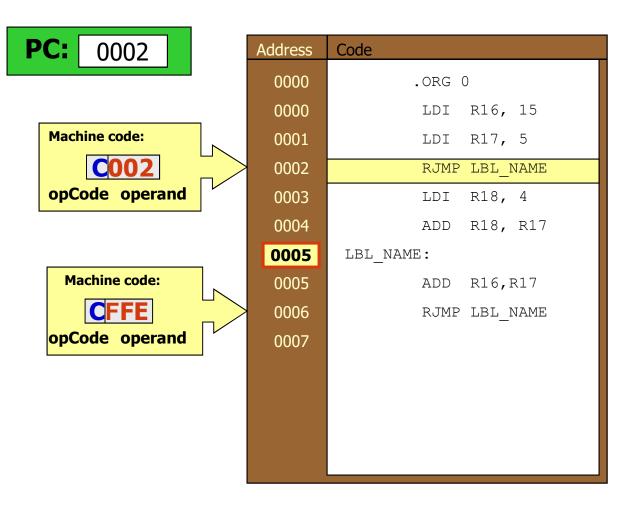


- When RJMP is executed:
 - The operand
 will be added
 to the current
 value of PC



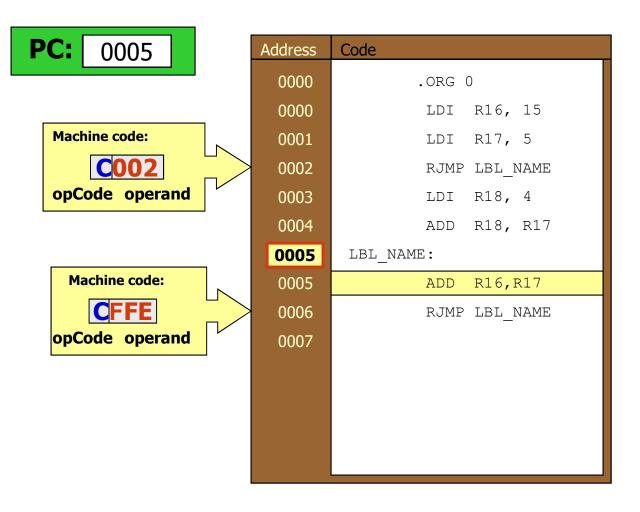


- When RJMP is executed:
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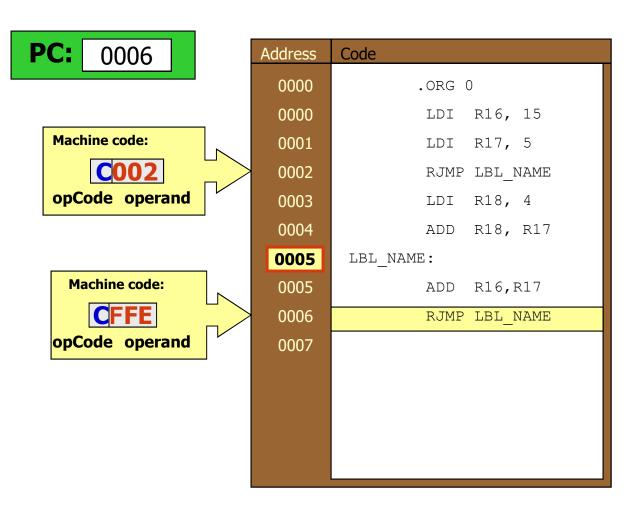


- When RJMP is executed:
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 value of PC



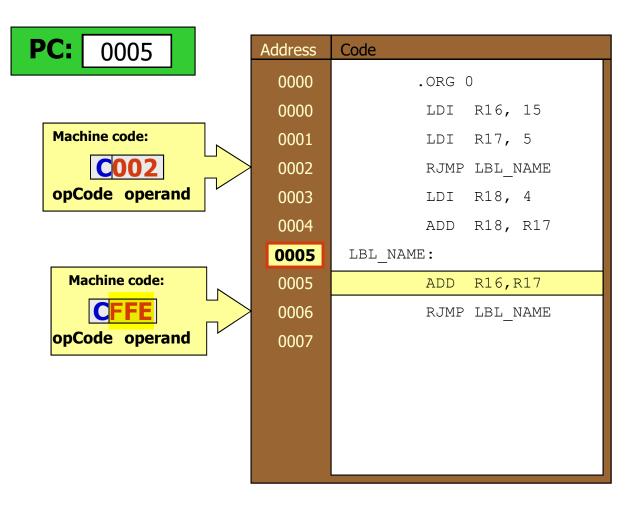


- When RJMP is executed:
 - The operand
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 value of PC





- When RJMP is executed:
 - The operand
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 value of PC





RJMP Backward

• RJMP PC = PC + operand

1100 XXXX XXXX XXXX

Negative Operand(Range Address?)



IJMP (Indirect jump)

• IJMP PC = Z register (R30:R31)

- The instruction has no operand.
- the Program counter is loaded with the contents of Z register.
 - For example, if Z points to location 100, by executing IJMP, the CPU jumps to location 100.



Conditional Jump in AVR

SREG: I T H S V N Z C

- Relative Jump
- 2 Byte jump
- Jump Address (-64: +64)

1111 01XX XXXX X000



Conditional Jump in AVR

SREG: I T H S V N Z C

• The conditional jump instructions in AVR are as follows:

Instruction	Abbreviation of	Comment
BREQ IbI	Branch if Equal	Jump to location lbl if $Z = 1$,
BRNE /b/	Branch if Not Equal	Jump if $Z = 0$, to location <i>lbl</i>
BRCS IbI	Branch if Carry Set	Jump to location <i>lbl</i> , if C = 1
BRLO <i>lbl</i>	Branch if Lower	
BRCC <i>lbl</i>	Branch if Carry Cleared	Jump to location <i>lbl</i> , if C = 0
BRSH <i>lbl</i>	Branch if Same or Higher	
BRMI <i>lbl</i>	Branch if Minus	Jump to location lbl, if $N = 1$
BRPL <i>lbl</i>	Branch if Plus	Jump if N = 0
BRGE <i>lbl</i>	Branch if Greater or Equal	Jump if S = 0
BRLT <i>lbl</i>	Branch if Less Than	Jump if S = 1
BRHS <i>lbl</i>	Branch if Half Carry Set	If H = 1 then jump to <i>lbl</i>
BRHC Ibl	Branch if Half Carry Cleared	if H = 0 then jump to IbI
BRTS	Branch if T flag Set	If T = 1 then jump to lbl
BRTC	Branch if T flag Cleared	If T = 0 then jump to lbl
BRIS	Branch if I flag set	If I = 1 then jump to lbl
BRIC	Branch if I flag cleared	If I = 0 then jump to lbl

Usages of Conditional jump

- Conditions
- Loop



Conditions

- When b is subtracted from a:
 - The result is zero, when a is equal to b

а

Carry will be set when a < b

- b

Note: Carry set when we borrow in high level concept

SREG: I T H S V N Z C

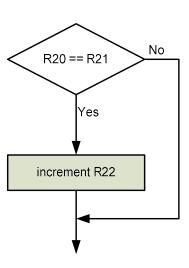


Example 1

• Write a program that if R20 is equal to R21 then R22 increases.

• Solution:

```
SUB R20,R21 ; Z will be set if R20 == R21
BRNE NEXT ; if Not Equal jump to next
INC R22
NEXT:
```



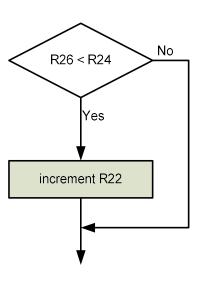


Example 2

Write a program that if R26 < R24 then R22 increases.

• Solution:

```
SUB R26,R24 ;C will be set when R26 < R24
BRCC L1 ;if Carry cleared jump to L1
INC R22
L1:
```





Example 3

Write a program that if R26 >= R24 then R22 increases.

```
• Solution:
```

L1:

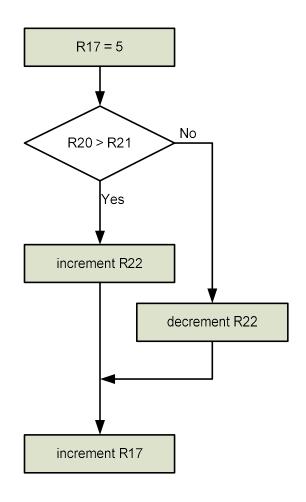
```
SUB R26,R24 ;C will be cleared when R26 >= R24
BRCS L1 ;if Carry set jump to L1
INC R22
```



R26 >= R24

Example 4: IF and ELSE

```
int main ( )
{
   R17 = 5;
   if (R20 > R21)
      R22++;
   else
      R22--;
   R17++;
}
```



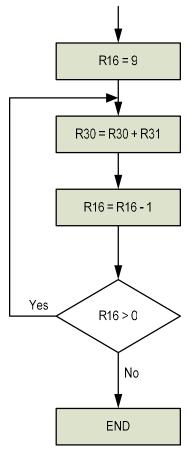


Example 4: IF and ELSE

```
int main ( )
                                                         R17 = 5
  R17 = 5;
   if (R20 > R21)
    R22++;
   else
                                                        R20 > R21
   R22--;
  R17++;
                                                           Yes
                                                       increment R22
   LDI R17,5
   SUB R21, R20 ;C is set when R20>R21
                                                                 decrement R22
   BRCC ELSE LABEL ; jump to else if cleared
   INC R22
   JMP NEXT
ELSE LABEL:
                                                       increment R17
   DEC R22
NEXT:
   INC R17
```



• Write a program that executes the instruction "ADD R30,R31" 9 times.

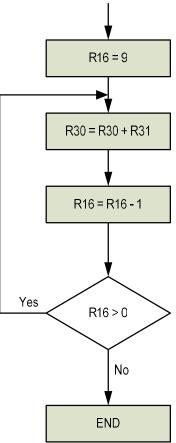




 Write a program that executes the instruction "ADD R30,R31" 9 times.

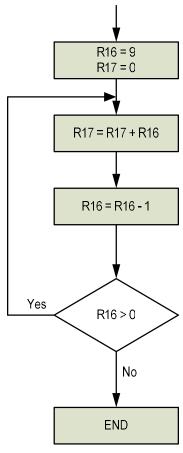
Solution:

```
.ORG 00
LDI R16,9 ;R16 = 9
L1: ADD R30,R31
DEC R16 ;R16 = R16 - 1
BRNE L1 ;if Z = 0
L2: RJMP L2 ;Wait here forever
```





• Write a program that calculates the result of 9+8+7+...+1





• Write a program that calculates the result of 9+8+7+...+1

• Solution:

```
.ORG 00

LDI R16, 9 ;R16 = 9

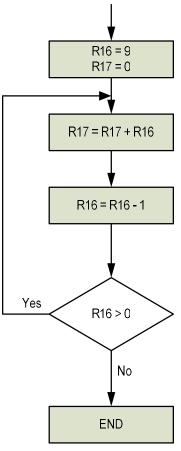
LDI R17, 0 ;R17 = 0

L1: ADD R17,R16 ;R17 = R17 + R16

DEC R16 ;R16 = R16 - 1

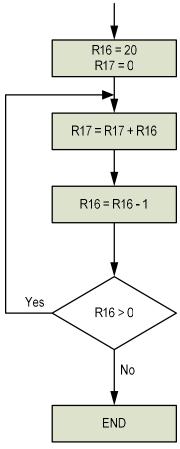
BRNE L1 ;if Z = 0

L2: RJMP L2 ;Wait here forever
```





• Write a program that calculates the result of 20+19+18+17+...+1





• Write a program that calculates the result of 20+19+18+17+...+1

Solution:

```
.ORG 00

LDI R16, 20 ;R16 = 20

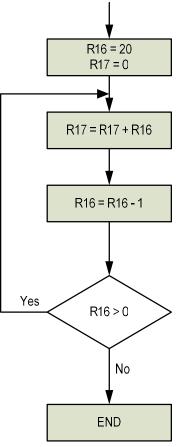
LDI R17, 0 ;R17 = 0

L1: ADD R17,R16 ;R17 = R17 + R16

DEC R16 ;R16 = R16 - 1

BRNE L1 ;if Z = 0

L2: RJMP L2 ;Wait here forever
```

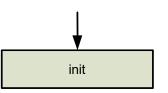




```
for (init; condition; calculation)
{
   do something
}
```

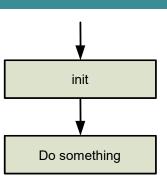


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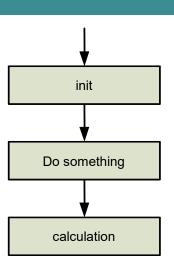


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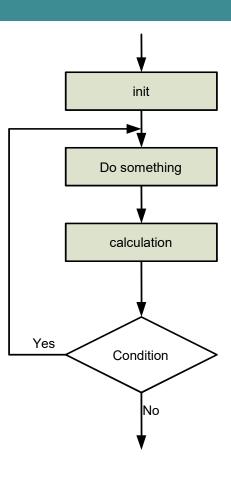


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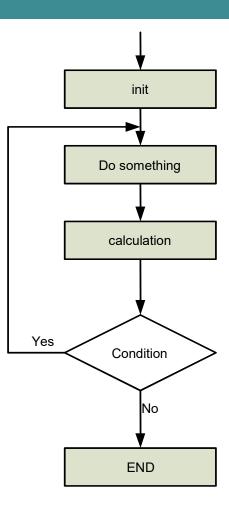


```
for (init; condition; calculation)
{
   do something
}
```





```
for (init; condition; calculation)
{
   do something
}
```





- Write a program that calculates 1+3+5+...+27
- Solution:

```
LDI R20,0

LDI R16,1

L1:ADD R20,R16

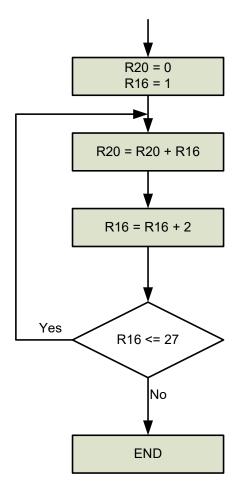
LDI R17,2

ADD R16,R17;R16 = R16 + 2

LDI R17,27;R17 = 27

SUB R17,R16

BRCC L1 ;if R16 <= 27 jump L1
```





More Points

- TST Rd
 - Update Statues Register based of Rd



Call Topics

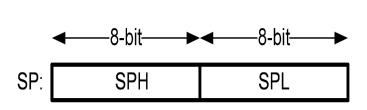
- Stack, Push and Pop
- Calling a function

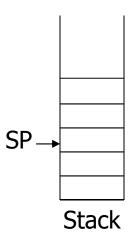


• POP
$$Rd$$

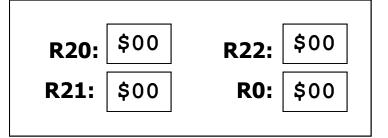
$$SP = SP + 1$$

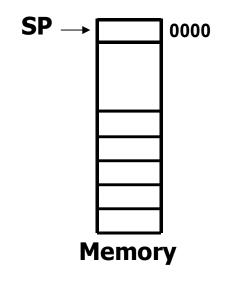
$$Rd = [SP]$$

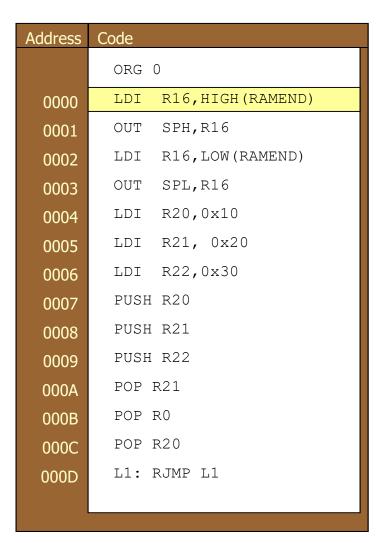




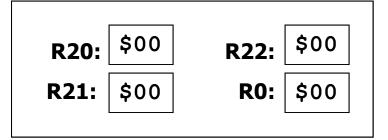


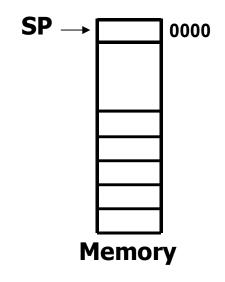


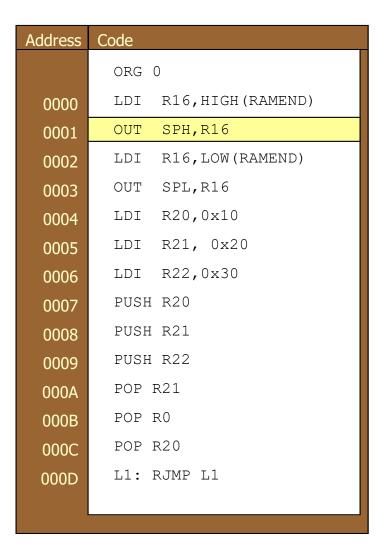




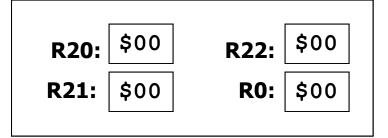


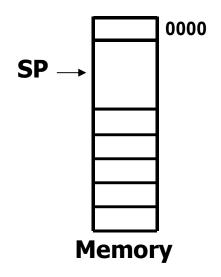








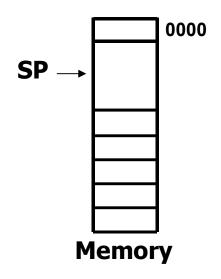




Address	Code
	ORG 0
0000	LDI R16,HIGH(RAMEND)
0001	OUT SPH,R16
0002	LDI R16, LOW (RAMEND)
0003	OUT SPL,R16
0004	LDI R20,0x10
0005	LDI R21, 0x20
0006	LDI R22,0x30
0007	PUSH R20
8000	PUSH R21
0009	PUSH R22
000A	POP R21
000B	POP RO
000C	POP R20
000D	L1: RJMP L1

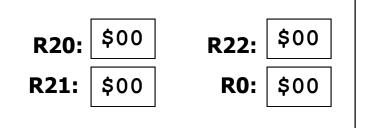


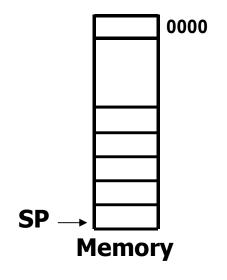


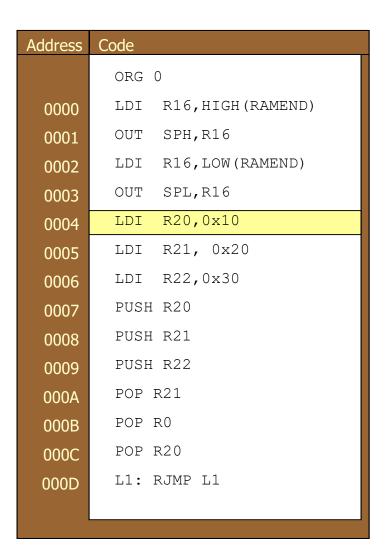


Address	Code		
	ORG 0		
0000	LDI R16,HIGH(RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,0x10		
0005	LDI R21, 0x20		
0006	LDI R22,0x30		
0007	PUSH R20		
0008	PUSH R21		
0009	PUSH R22		
000A	POP R21		
000B	POP RO		
000C	POP R20		
000D	L1: RJMP L1		

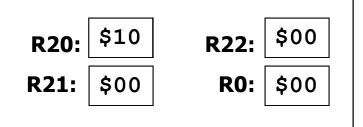


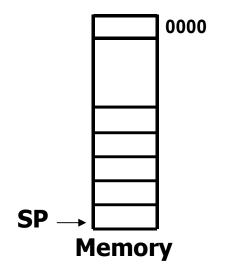


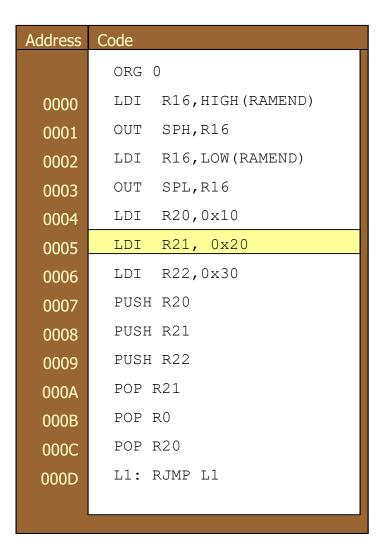




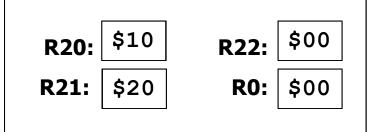


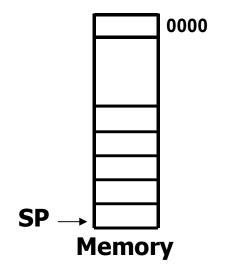


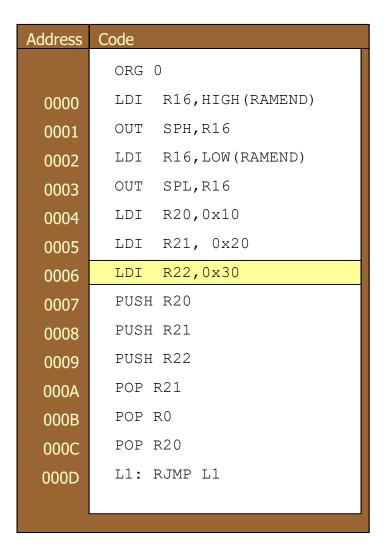




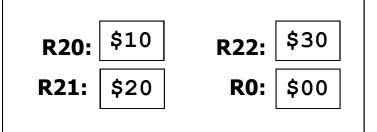


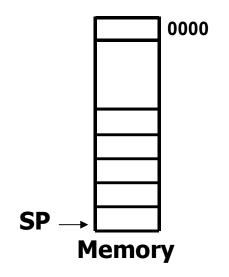


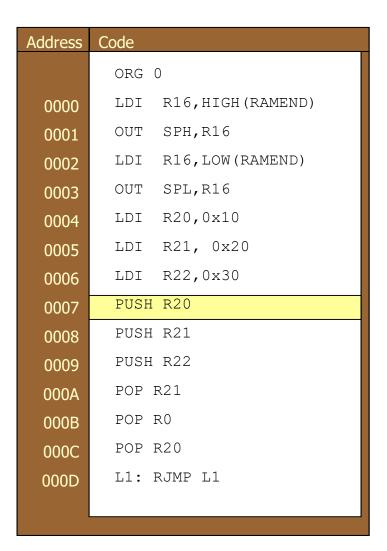




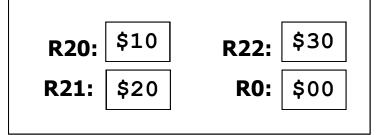


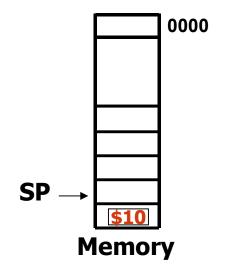


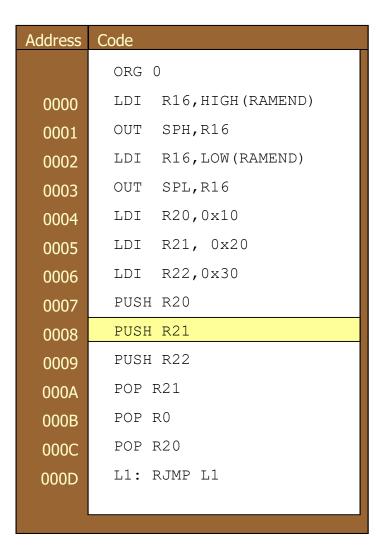






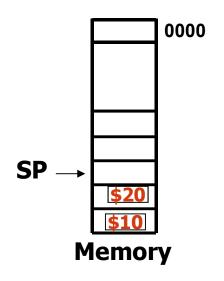


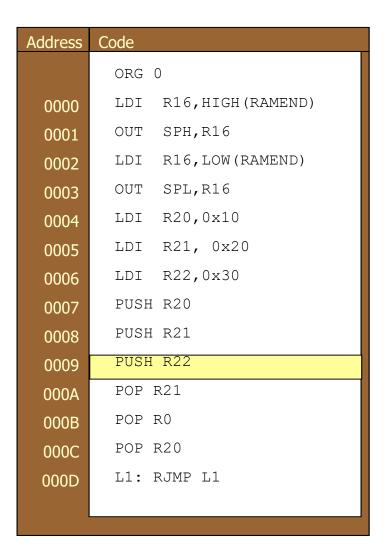




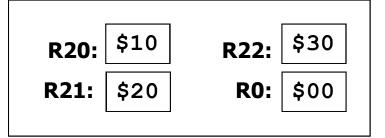


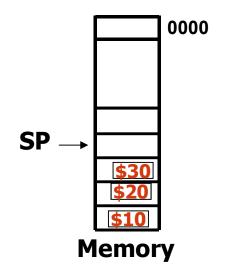








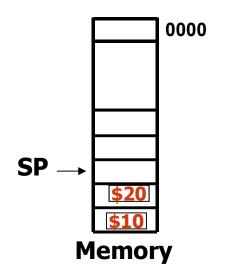


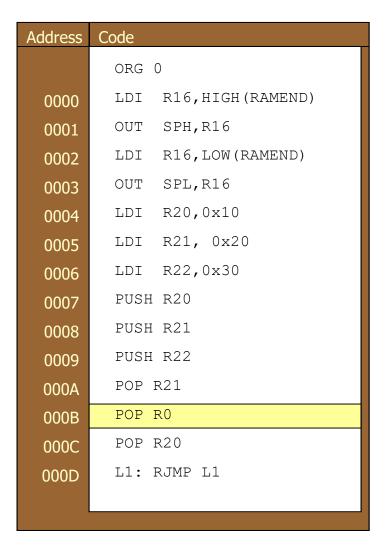


Address	Code
	ORG 0
0000	LDI R16,HIGH(RAMEND)
0001	OUT SPH,R16
0002	LDI R16,LOW(RAMEND)
0003	OUT SPL,R16
0004	LDI R20,0x10
0005	LDI R21, 0x20
0006	LDI R22,0x30
0007	PUSH R20
8000	PUSH R21
0009	PUSH R22
000A	POP R21
000B	POP RO
000C	POP R20
000D	L1: RJMP L1



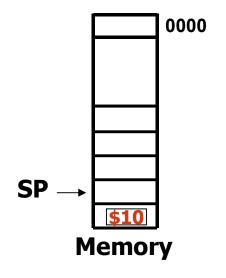


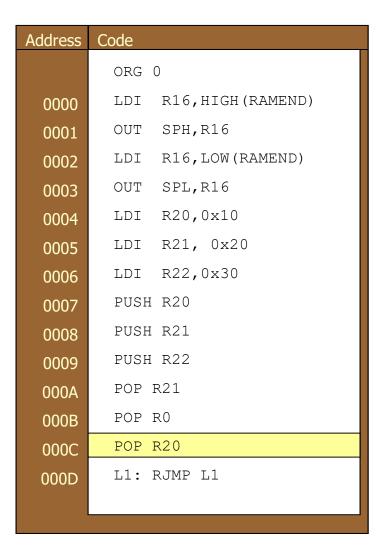






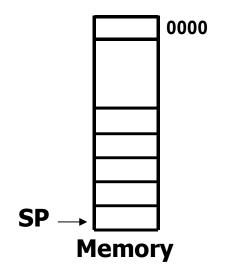


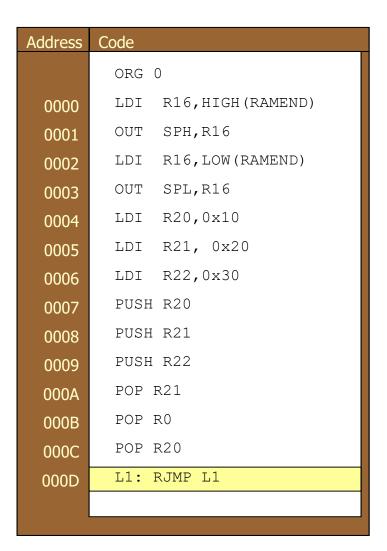






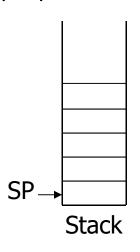








- To execute a call:
 - Address of the next instruction is saved
 - PC is loaded with the appropriate value

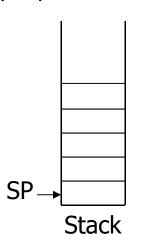




Address	Code		
0000	LDI R16,HIGH(RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,15		
0005	LDI R21,5		
0006	CALL FUNC_NAME		
8000	INC R20		
0009	L1: RJMP L1		
000A	FUNC_NAME:		
000A	ADD R20,R21		
000B	SUBI R20,3		
000C	RET		
000D			



- To execute a call:
 - Address of the next instruction is saved
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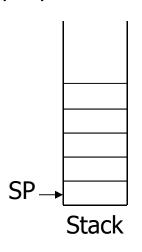




Address	Code		
0000	LDI R16, HIGH (RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,15		
0005	LDI R21,5		
0006	CALL FUNC_NAME		
8000	INC R20		
0009	L1: RJMP L1		
000A	FUNC_NAME:		
000A	ADD R20,R21		
000B	SUBI R20,3		
000C	RET		
000D			



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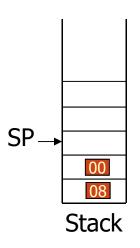


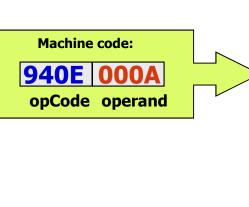


Address	Code		
0000	LDI R16,HIGH(RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,15		
0005	LDI R21,5		
0006	CALL FUNC_NAME		
0008	INC R20		
0009	L1: RJMP L1		
000A	FUNC_NAME:		
000A	ADD R20,R21		
000B	SUBI R20,3		
000C	RET		
000D			

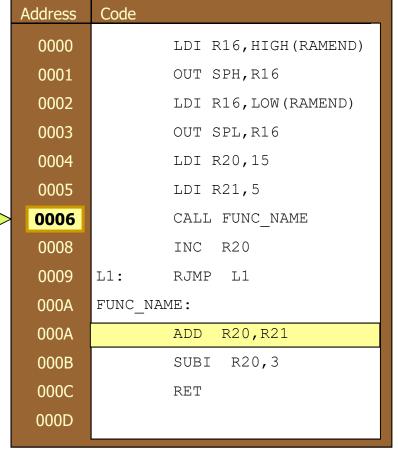


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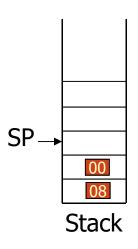


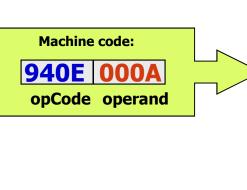
000A



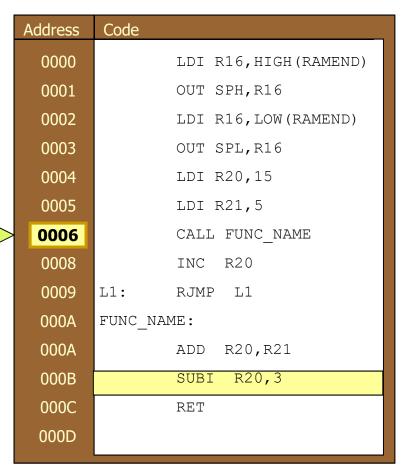


- To execute a call:
 - Address of the next instruction is saved
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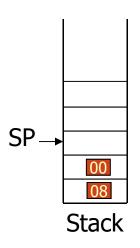


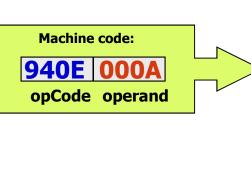
PC:	000B	
• •	0000	





- To execute a call:
 - Address of the next instruction is saved
 - PC is loaded with the appropriate value



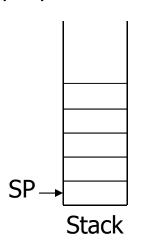


PC: 000C

Address	Code		
0000	LDI R16,HIGH(RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,15		
0005	LDI R21,5		
0006	CALL FUNC_NAME		
8000	INC R20		
0009	L1: RJMP L1		
000A	FUNC_NAME:		
000A	ADD R20,R21		
000B	SUBI R20,3		
000C	RET		
000D			



- To execute a call:
 - Address of the next instruction is saved
 - PC is loaded with the appropriate value

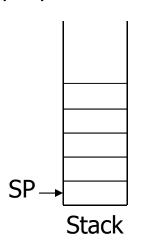




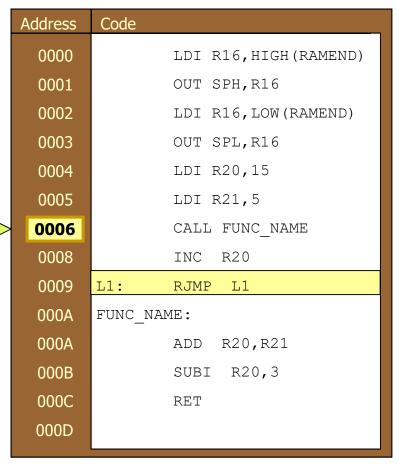
Address	Code		
0000	LDI R16,HIGH(RAMEND)		
0001	OUT SPH,R16		
0002	LDI R16,LOW(RAMEND)		
0003	OUT SPL,R16		
0004	LDI R20,15		
0005	LDI R21,5		
0006	CALL FUNC_NAME		
8000	INC R20		
0009	L1: RJMP L1		
000A	FUNC_NAME:		
000A	ADD R20,R21		
000B	SUBI R20,3		
000C	RET		
000D			



- To execute a call:
 - Address of the next instruction is saved
 - PC is loaded with the appropriate value





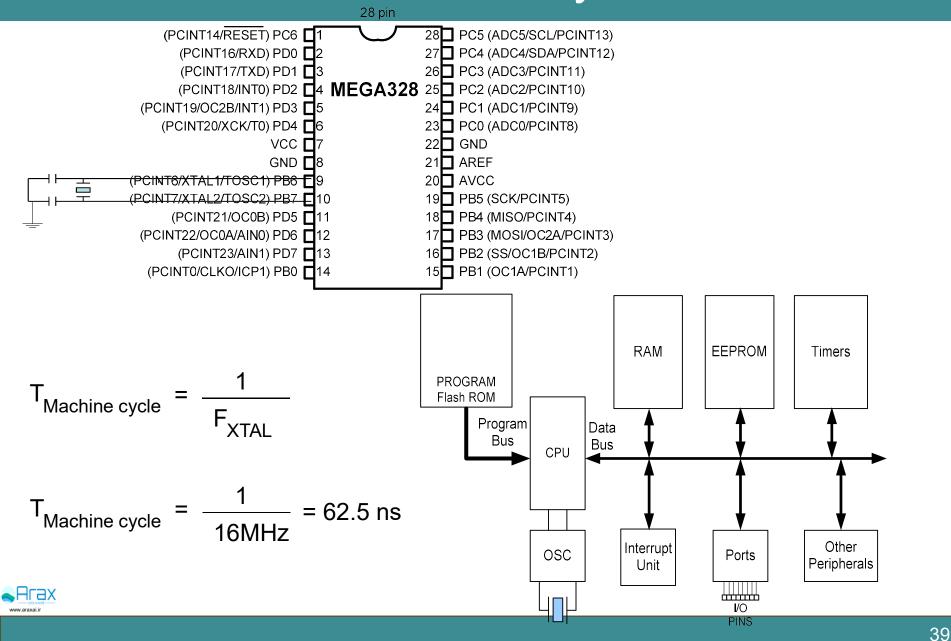




Some More Point

- RCALL (Relative CALL)
- ICALL (Indirect CALL)





```
LDI R16, 19
LDI R20, 95
LDI R21, 5
ADD R16, R20
ADD R16, R21
```

Delay =
$$5 \times T_{\text{machine cycle}} = 5 \times 62.5 \text{ ns} = 312.5 \text{ ns}$$



		machine cycle
LDI	R16, 19	1
LDI	R20, 95	1
LDI	R21, 5	1
ADD	R16, R20	1
ADD	R16, R21	1
		5

Delay =
$$5 \times T_{\text{machine cycle}} = 5 \times 62.5 \text{ ns} = 312.5 \text{ ns}$$



			machine cycle
	LDI	R16, 100	1
AGAIN:	ADD	R17,R16	1
	DEC	R16	1
	BRNE	AGAIN	_1/2_



LDI R16, 100

AGAIN: ADD R17,R16

DEC R16

BRNE AGAIN

Branch penalty



			machine cycle
	LDI	R16, 100	1
AGAIN:	ADD	R17,R16	1
	DEC	R16	1
	BRNE	AGAIN	1/2



			machine cycle
	LDI	R16, 100	
AGAIN:	ADD	R17,R16	1
	DEC	R16	1
	BRNE	AGAIN	_1/2_



			machine	cycle
	LDI	R16, 100	1	
AGAIN:	ADD	R17,R16	1	*100
	DEC	R16	1	*100
	BRNE	AGAIN	1/2	*100



			machine cycle
	LDI	R16, 50	1
AGAIN:	NOP		1
	NOP		1
	DEC	R16	1
	BRNE	AGAIN	1/2



			machine cycle
	LDI	R16, 50	1
AGAIN:	NOP		1
	NOP		1
	DEC	R16	1
	BRNE	AGAIN	1/2



			machine cycle	
	LDI	R16, 50	1	
AGAIN:	NOP		1 *5	50
	NOP		1 *5	50
	DEC	R16	1 *5	50
	BRNE	AGAIN	1/2 *5	50



	LDI	R17, 20	machine cycle
L1:	LDI	R16, 50	1
L2:	NOP		1 1
	NOP		1
	DEC	R16	1
	BRNE	L2	1/2
	DEC	R17	1
	BRNE	L1	1/2



	LDI	R17, 20	machine cycle
L1:	LDI	R16, 50	1
L2:	NOP		1
	NOP		1
	DEC	R16	1
	BRNE	L2	1/2
	DEC	R17	1
	BRNE	L1	1/2



	LDI	R17, 20
L1:	LDI	R16, 50
L2:	NOP	
	NOP	
	DEC	R16
	BRNE	L2
	DEC	R17
	BRNE	L1

machine cy	<u>cle</u>
1	
1	
1	
1	
1	
1/2	
1	
1/2	



	LDI	R17, 20
L1:	LDI	R16, 50
L2:	NOP	
	NOP	
	DEC	R16
	BRNE	L2
	DEC	R17
	BRNE	L1

machine cycle 1 1 *20 1 *20 * 50 1 *20 * 50 1 *20 * 50 1/2 *20 * 50 1 *20 * 50 1 *20 * 50

1/2

*20

