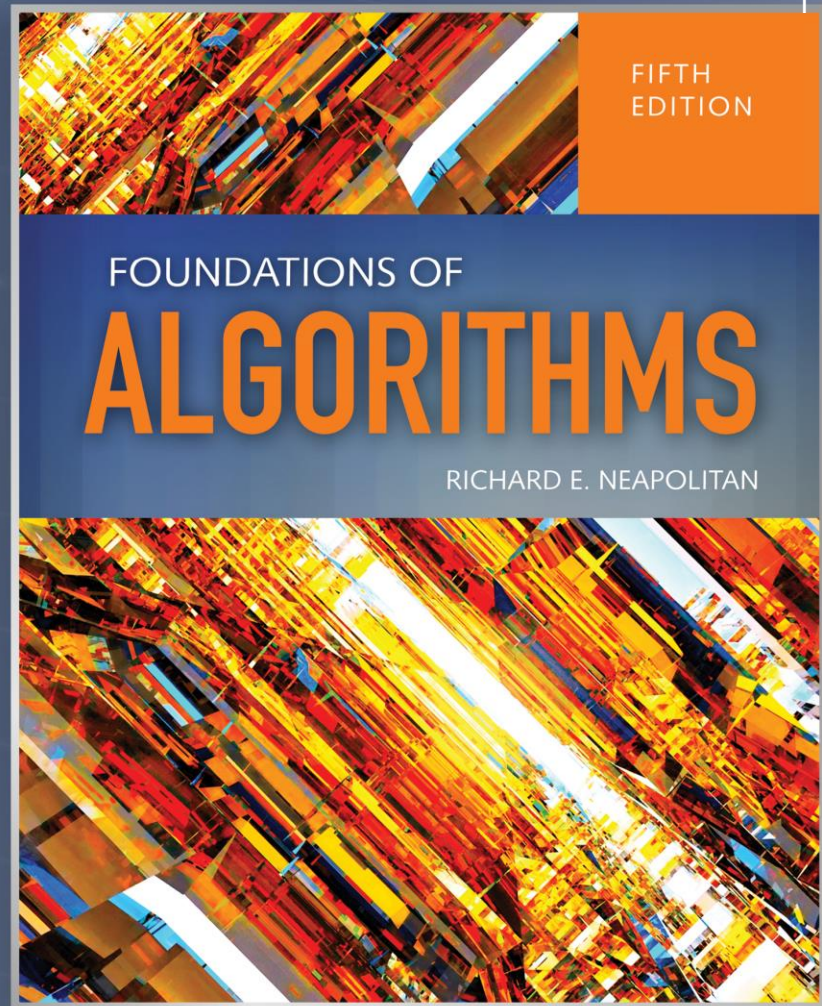


Algorithms:

Efficiency, Analysis,
and Order

Chapter 1



Objectives

- Analyze techniques for solving problems
- Define an algorithm
- Define growth rate of an algorithm as a function of input size
- Define Worst case, average case, and best case complexity analysis of algorithms
- Classify functions based on growth rate
- Define growth rates: Big O, Theta, and Omega

Methodology

- Approach to solving a problem
- Independent of Programming Language
- Independent of Style
- Sequential Search versus Binary Search
- Which technique results in the most efficient solution?

Problem?

- A question to which an answer is sought
- Parameters
 - Input to the problem
 - Instance: a specific assignment of values to the input parameters
- Algorithm:
 - Step-by-step procedure
 - Solves the Problem

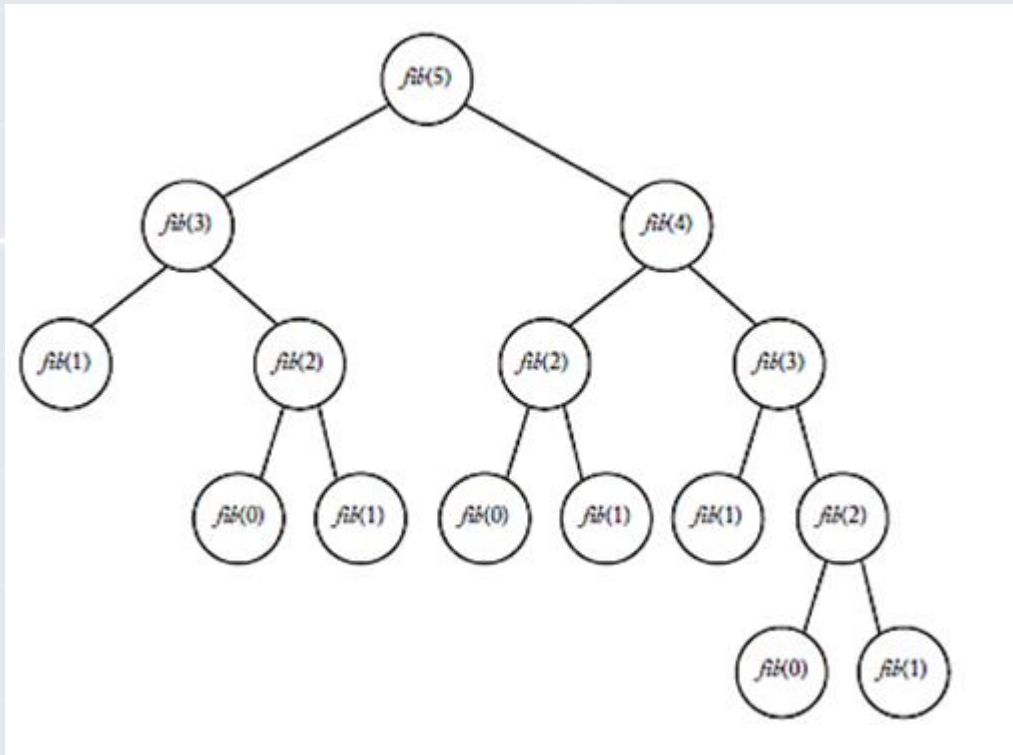
Sequential Search vs Binary Search – Worst Case

- Input Array S size n
- $X \notin S$
- Sequential Search: n operations
- Binary Search: $\lg n + 1$ operations

Fibonacci: Iterative vs Recursive

- $\text{Fib}_0 = 0$
- $\text{Fib}_1 = 1$
- $\text{Fib}_n = \text{Fib}_{n-1} + \text{Fib}_{n-2}$
- Calculate the nth Fibonacci Term:
 - Recursive calculates $2^{n/2}$ terms
 - Iterative calculates $n+1$ terms

Recursion Tree for the 5th Fibonacci Term



Comparison of Recursive and Iterative Solutions

n	$n + 1$	$2^{n/2}$	Execution Time Using Algorithm 1.7	Lower Bound on Execution Time Using Algorithm 1.6
40	41	1,048,576	41 ns*	1048 μ s†
60	61	1.1×10^9	61 ns	1 s
80	81	1.1×10^{12}	81 ns	18 min
100	101	1.1×10^{15}	101 ns	13 days
120	121	1.2×10^{18}	121 ns	36 years
160	161	1.2×10^{24}	161 ns	3.8×10^7 years
200	201	1.3×10^{30}	201 ns	4×10^{13} years

*1 ns = 10^{-9} second.

†1 μ s = 10^{-6} second.

Complexity Analysis

- Define Basic Operation
- Count the number of times the basic operation executes for each value of the input size
- Maybe dependent on input size (sequential search)
- Every-case time complexity analysis

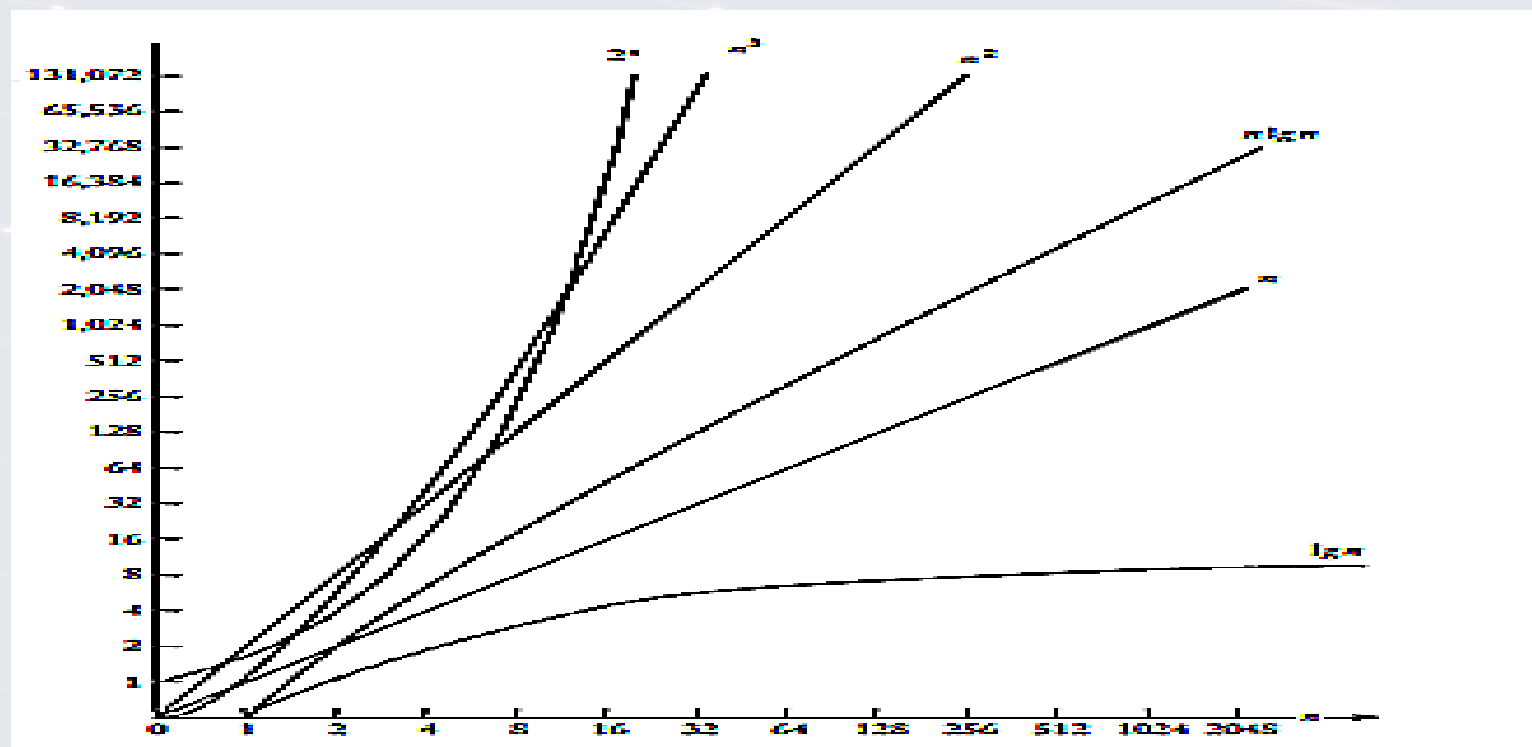
Complexity Analysis – Large n

- Worst Case
- Every Case
- Average Case
- Best Case

Order – Classes of Functions – Growth Rate

- $\Theta(f)$ – At the rate of f
- $O(f)$ – At most as fast as f
- $\Omega(f)$ - At least as fast as f
- n grows more slowly than $n^3 \Rightarrow n \in O(n^3)$
- n^3 grows faster than $n \Rightarrow n^3 \in \Omega(n)$
- By definition n and $2n$ grow at the same rate $\Rightarrow 2n \in \Theta(n)$

Growth Rates of Common Complexity Functions



Big O

- For a given complexity function $f(n)$, $O(f(n))$ is the set of complexity functions $g(n)$ for which there exists some positive real constant c and some nonnegative integer N such that for all $n \geq N$,
- $g(n) \leq c \times f(n)$
- $g(n) \in O(f(n))$

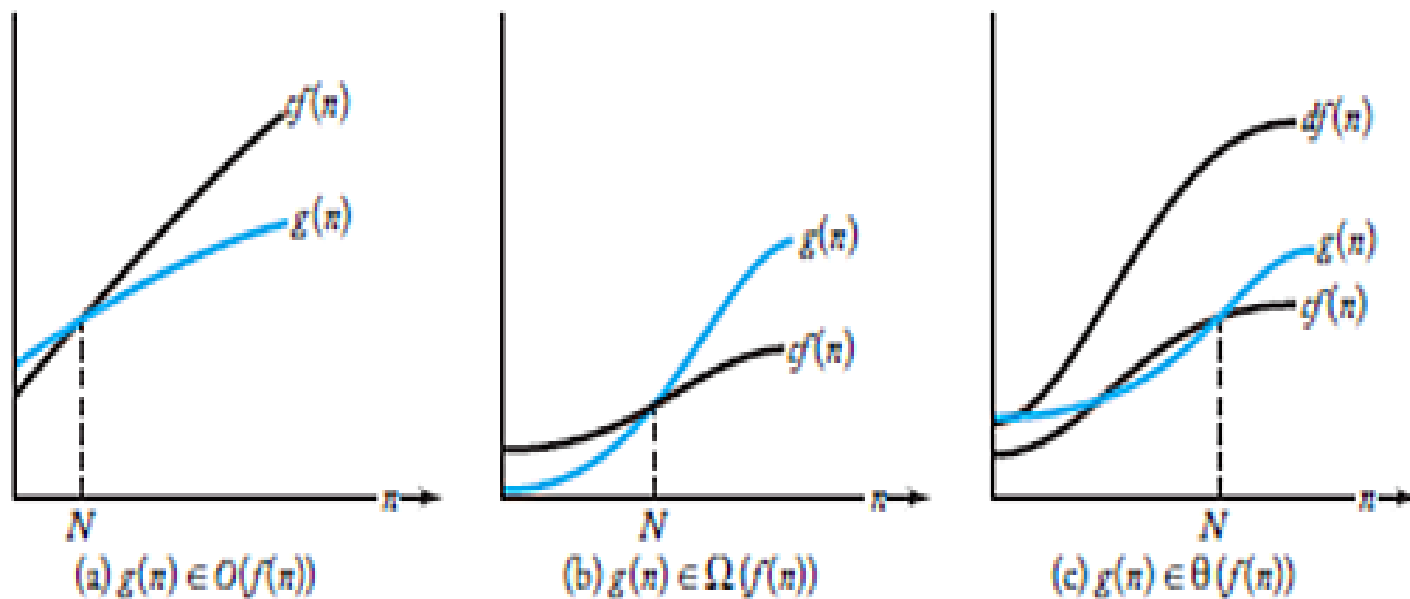
Omega

- For a given complexity function $f(n)$, $\Omega(f(n))$ is the set of complexity functions $g(n)$ for which there exists some positive real constant c and some nonnegative integer N such that, for all $n \geq N$,
- $g(n) \geq c \times f(n)$.
- $g(n) \in \Omega(f(n))$

Theta

- For a given complexity function $f(n)$,
- $\theta(f(n)) = O(f(n)) \cap \Omega(f(n))$
- This means that $\theta(f(n))$ is the set of complexity functions $g(n)$ for which there exists some positive real constants c and d and some nonnegative integer N such that, for all $n \geq N$,
- $c \times f(n) \leq g(n) \leq d \times f(n)$.
- $g(n) \in \theta(f(n))$

Big O, Omega, Theta



Limit Definitions for Big O, Theta, and Omega

- $\lim_{n \rightarrow \infty} g(n)/f(n) = c \Rightarrow g(n) \in \theta(f(n))$ for $c > 0$ and $c \neq \infty$
- $\lim_{n \rightarrow \infty} g(n)/f(n) = c \Rightarrow g(n) \in O(f(n))$ where $c \in \text{Set of all positive real numbers union } 0$
- $\lim_{n \rightarrow \infty} g(n)/f(n) = \infty$ OR $\lim_{n \rightarrow \infty} g(n)/f(n) = c > 0 \Rightarrow g(n) \in \Omega(f(n))$