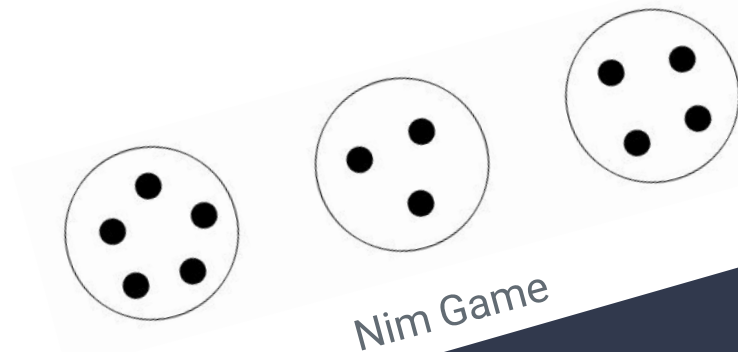


# IMPARTIAL GAMES

- Normal-Game
- Same moves
  - Identity of the player does not matter
  - There are no positions type L or R

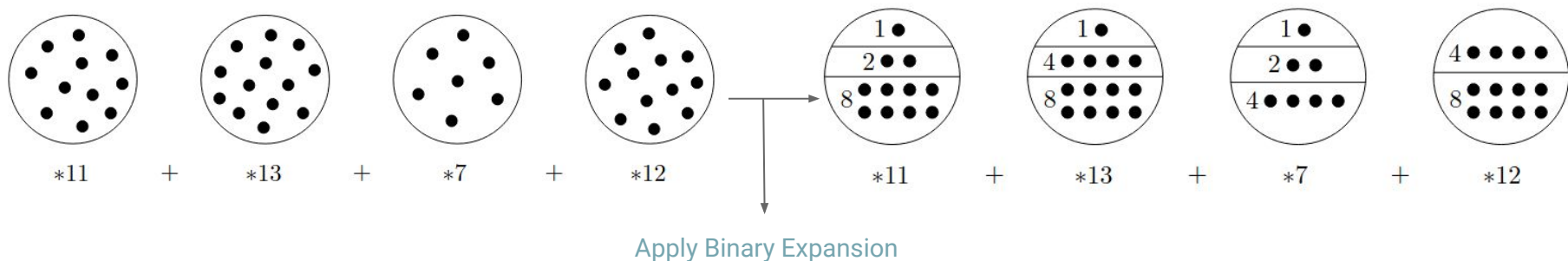


## Binary Expansion

- Representation of  $n$  as sum of distinct powers of two
  - $45 = 32 + 13 = 32 + 8 + 5 = 32 + 8 + 4 + 1$

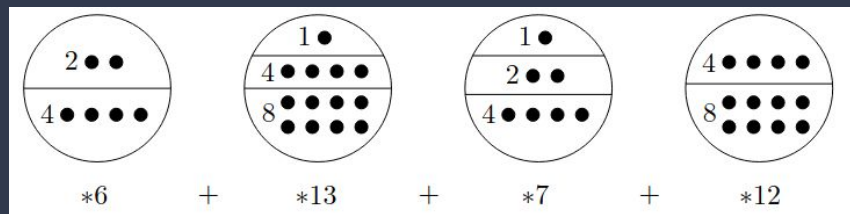
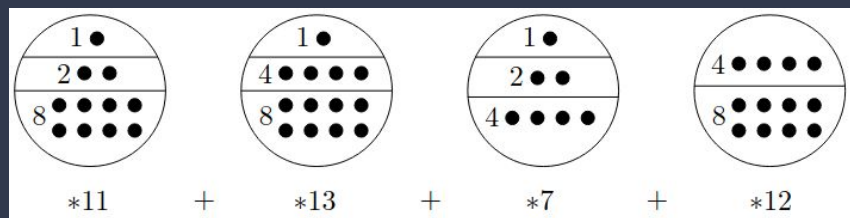
# Type of Positions

- Let  $a$  be a non-negative integer and let  $*a$  denote a Nim position of a single pile of  $a$  stones



- A position  $*a_1 + \dots + *a_k$  is **balanced** if, for every power of two, the total **number of subpiles** of that size is **even**

# Balancing a Position



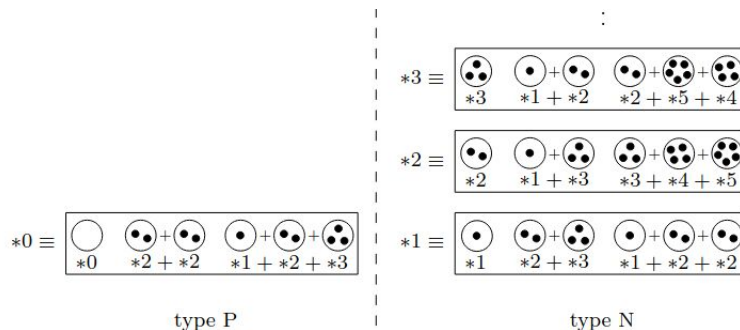
KEY → Take  $2^m$  the largest power of 2 which odd number of subpiles

- Balanced position → Type P
- Unbalanced position → Type N

- The Nim-Sum of nonnegative integers  $a_1 \oplus \dots \oplus a_k$  is a non-negative integer  $b$  s.t. if  $2^j$  appears in the binary expansion of  $b \leftrightarrow 2^j$  appears an odd number of time in the expansion  $a_1, \dots, a_k$ .

$$\circ \quad 13 \oplus 19 \oplus 10 = (8+4+1) \oplus (16+2+1) \oplus (8+2) = 4+16 = 20$$

- If  $b = a_1 \oplus \dots \oplus a_k$ , then  $*a_1 + \dots + *a_k = b$ , and  $b$  is called **Nimber**.



# The Sprague-Grundy Theorem

- For normal-play games, a position  $\alpha = \{\beta_1, \dots, \beta_k \mid \gamma_1, \dots, \gamma_m\}$  (Louise | Richard)
  - Redundant notation for impartial games.
- For  $S = \{a_1, \dots, a_n\}$  of non-negative integers, we define Minimal EXcluded values (MEX) of  $S$  as the smallest integer  $b$  which is not in  $S$  (i.e.  $\{0, 1, 2, 5, 8\}$  is 3)
- The MEX Principle → Let  $\alpha = \{\alpha_1, \dots, \alpha_k\}$  be a position in an impartial game and let  $\alpha_i \equiv^* a_i$ ,  $1 \leq i \leq k$ , then,  $\alpha \equiv b$ , where  $b$  is the MEX of the set  $\{\alpha_1, \dots, \alpha_k\}$
- The Sprague-Grundy Theorem → Every position in an impartial game is equivalent to a number.
- We define a winning move in an impartial game to be any move to a position of type P. Note that winning moves only exists from positions of type N
  - In Nim, balancing always provides a winning move.

# MEX Principle Examples

**Example 3.14** (Chop). Here we work out the nimber equivalents for some small positions in Chop.

$$\begin{aligned}
 \square &= \{ \} \equiv *0 \\
 \square\square &= \{ \square \} \equiv \{ *0 \} \equiv *1 \\
 \square\square\square &= \{ \square, \square\square \} \equiv \{ *0, *1 \} \equiv *2 \\
 \square\square\square\square &= \{ \square, \square\square, \square\square\square \} \equiv \{ *0, *1, *2 \} \equiv *3 \\
 \begin{array}{|c|} \hline \square \\ \hline \end{array} &= \{ \square \} \equiv \{ *0 \} \equiv *1 \\
 \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array} &= \{ \square\square, \begin{array}{|c|} \hline \square \\ \hline \end{array} \} \equiv \{ *1 \} \equiv *0 \\
 \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array} &= \{ \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array}, \begin{array}{|c|} \hline \square \\ \hline \end{array}, \square\square\square \} \equiv \{ *0, *1, *2 \} \equiv *3 \\
 \begin{array}{|c|c|c|c|} \hline \square & \square & \square & \square \\ \hline \end{array} &= \{ \begin{array}{|c|c|c|} \hline \square & \square & \square \\ \hline \end{array}, \begin{array}{|c|c|} \hline \square & \square \\ \hline \end{array}, \begin{array}{|c|} \hline \square \\ \hline \end{array}, \square\square\square\square \} \equiv \{ *0, *1, *3 \} \equiv *2
 \end{aligned}$$

**Example 3.15** (Chomp). Here we work out the nimber equivalents for some small positions in Chomp.

$$\begin{aligned}
 \boxtimes &= \{ \} \equiv *0 \\
 \boxtimes\boxtimes &= \{ \boxtimes \} \equiv \{ *0 \} \equiv *1 \\
 \begin{array}{|c|} \hline \boxtimes \\ \hline \end{array} &= \{ \boxtimes \} \equiv \{ *0 \} \equiv *1 \\
 \boxtimes\boxtimes\boxtimes &= \{ \boxtimes, \boxtimes\boxtimes \} \equiv \{ *0, *1 \} \equiv *2 \\
 \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array} &= \{ \boxtimes\boxtimes, \begin{array}{|c|} \hline \boxtimes \\ \hline \end{array} \} \equiv \{ *1 \} \equiv *0 \\
 \begin{array}{|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array} &= \{ \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array}, \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array}, \boxtimes\boxtimes\boxtimes \} \equiv \{ *0, *1, *2 \} \equiv *3 \\
 \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array} &= \{ \boxtimes\boxtimes, \begin{array}{|c|} \hline \boxtimes \\ \hline \end{array}, \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array} \} \equiv \{ *0, *1 \} \equiv *2 \\
 \begin{array}{|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array} &= \{ \begin{array}{|c|} \hline \boxtimes \\ \hline \end{array}, \boxtimes\boxtimes\boxtimes, \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array} \} \equiv \{ *1, *2, *3 \} \equiv *0 \\
 \begin{array}{|c|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array} &= \{ \begin{array}{|c|} \hline \boxtimes \\ \hline \end{array}, \boxtimes\boxtimes\boxtimes, \begin{array}{|c|c|} \hline \boxtimes & \boxtimes \\ \hline \end{array}, \begin{array}{|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array}, \begin{array}{|c|c|c|c|} \hline \boxtimes & \boxtimes & \boxtimes & \boxtimes \\ \hline \end{array} \} \equiv \{ *0, *1, *2, *3 \} \equiv *4
 \end{aligned}$$

# NIM Code in Python

[GitHub Code](#)

1. Ask for players name
2. Use random number generator to generate between 2 to 5 piles and to generate between 1 to 8 stones in each pile.
3. Display the board in the following fashion: (Note: O is the letter capital O)  
    Pile 1 : O O O  
    Pile 2 : O  
    ...
4. Ask player 1 for the move – the pile number and the number of stones to remove. You need to make sure that if a user enters something invalid, whether the pile or stone number, you ask for input again until valid input is entered. Also, the program is to give a suggestion to player 1 as to the number of stones to remove and from which pile in order to win the game.
5. Repeat step 4
6. Repeat step 5 for player 2.
7. When the game terminates, you need to proclaim the winner.
8. You should ask, if the players want to play the game again; if so, repeat steps 1-9