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#### 1 Data Structures

### 1.1 Disjoint Set Union

```
struct DSU
    int par[V], sz[V];
   DSU(){init(V);}
    void init(int n)
        for (int i = 0; i < n; i++)
        par[i] = i, sz[i] = 1;
   int find(int x)
        return x == par[x] ? x : (par[x] = find(par[x]));
   int getSize(int k){return sz[find(k)];}
   void unite(int x, int y)
        int u=find(x), v=find(y);
        if(u==v) return;
        if(sz[u]>sz[v]) swap(u, v);
        sz[v]+=sz[u];
        sz[u] = 0;
        par[u] = par[v];
};
```

# 2 Dynamic Programming

#### 2.1 Li Chao Tree

```
/*
    Objective
    (1) Line insert query (ax + b)
    (2) Max / Min on x = t query
    Current Implementation : Max query
*/
struct LiChao
{
    struct Line // Linear function ax + b
    {
        int a, b;
        int eval(int x)
```

```
{
        return a*x + b;
};
struct Node // [start, end] has line f
    int left, right;
    int start, end;
    Line f;
};
Node new_node(int a, int b)
    return \{-1,-1,a,b,\{0,-INF\}\};
    // for min, change -INF to INF
}
vector <Node> nodes;
void init(int min_x, int max_x)
    nodes.push_back(new_node(min_x, max_x));
void insert(int n, Line new_line)
    int xl = nodes[n].start, xr = nodes[n].end;
    int xm = (xl + xr)/2;
    Line llo, lhi;
    llo = nodes[n].f, lhi = new_line;
    if (llo.eval(xl) >= lhi.eval(xl))
        swap(llo, lhi);
    if (llo.eval(xr) <= lhi.eval(xr))</pre>
        nodes[n].f = lhi;
        // for min, lhi -> llo
        return;
    else if (llo.eval(xm) > lhi.eval(xm))
        nodes[n].f = 11o;
        // for min, llo -> lhi
        if (nodes[n].left == -1)
```

```
{
                nodes[n].left = nodes.size();
                nodes.push_back(new_node(x1,xm));
            insert(nodes[n].left, lhi);
            // for min, lhi -> llo
        }
        else
            nodes[n].f = lhi;
            // for min, lhi -> llo
            if (nodes[n].right == -1)
                nodes[n].right = nodes.size();
                nodes.push_back(new_node(xm+1,xr));
            insert(nodes[n].right,llo);
            // for min, llo -> lhi
        }
    void insert(Line f)
        insert(0, f);
    int get(int n, int q)
        // for min, max -> min, -INF -> INF
        if (n == -1) return -INF;
        int xl = nodes[n].start, xr = nodes[n].end;
        int xm = (xl + xr)/2;
        if (q > xm)
            return max(nodes[n].f.eval(q), get(nodes[n].right, q));
        else return max(nodes[n].f.eval(q), get(nodes[n].left, q));
    int get(int pt)
        return get(0, pt);
};
```

## 3 Mathematics

#### 3.1 Fast Fourier Transform

```
Multiply two polynomials in O(n \log n). For
#define _USE_MATH_DEFINES
#define sz(v) ((int)(v).size())
#define all(v) (v).begin(),(v).end()
typedef vector<int> vi;
typedef complex<double> base;
void fft(vector <base> &a, bool invert)
    int n = sz(a);
    for (int i=1, j=0; i < n; i++){
        int bit = n \gg 1;
        for (; j>=bit; bit>>=1) j -= bit;
        j += bit;
        if (i < j) swap(a[i],a[j]);</pre>
    for (int len=2;len<=n;len<<=1){
        double ang = 2*M_PI/len*(invert?-1:1);
        base wlen(cos(ang),sin(ang));
        for (int i=0; i < n; i+=len){
            base w(1);
            for (int j=0; j<len/2; j++){}
                base u = a[i+j], v = a[i+j+len/2]*w;
                 a[i+j] = u+v;
                 a[i+j+len/2] = u-v;
                 w *= wlen;
        }
    if (invert){
        for (int i=0;i<n;i++) a[i] /= n;
    }
}
void multiply(const vi &a,const vi &b,vi &res)
    vector <base> fa(all(a)), fb(all(b));
    int n = 1;
```

```
while (n < max(sz(a), sz(b))) n <<= 1;
    fa.resize(n); fb.resize(n);
    fft(fa,false); fft(fb,false);
    for (int i=0;i<n;i++) fa[i] *= fb[i];</pre>
    fft(fa,true);
    res.resize(n);
    for (int i=0; i < n; i++)
        res[i] = int(fa[i].real()+(fa[i].real()>0?0.5:-0.5));
//multiply(a, b, res)
//for n = b.size
//sum(a[i+j], b[n-1-j]) \rightarrow res[n-1+i]
3.2 Berlekamp Massey Algorithm
Find minimum linear recurrence from 3n first terms.
struct Berlekamp_Massey
{
    const int mod = 1000000007;
    using lint = long long;
    lint ipow(lint x, lint p){
        lint ret = 1, piv = x;
        while(p){
            if(p & 1) ret = ret * piv % mod;
            piv = piv * piv % mod;
            p >>= 1;
        return ret;
    vector<int> berlekamp_massey(vector<int> x){
        vector<int> ls, cur;
        int lf, ld;
        for(int i=0; i<x.size(); i++){</pre>
            lint t = 0:
            for(int j=0; j<cur.size(); j++){</pre>
                t = (t + 111 * x[i-j-1] * cur[j]) \% mod;
            if((t - x[i]) \% mod == 0) continue;
            if(cur.empty()){
                cur.resize(i+1);
                lf = i;
                1d = (t - x[i]) \% mod;
```

```
continue;
        }
        lint k = -(x[i] - t) * ipow(ld, mod - 2) % mod;
        vector<int> c(i-lf-1);
        c.push_back(k);
        for(auto &j : ls) c.push_back(-j * k % mod);
        if(c.size() < cur.size()) c.resize(cur.size());</pre>
        for(int j=0; j<cur.size(); j++){</pre>
            c[i] = (c[i] + cur[i]) \% mod;
        if(i-lf+(int)ls.size()>=(int)cur.size()){
            tie(ls, lf, ld) = make_tuple(cur, i, (t - x[i]) % mod);
       }
        cur = c;
    for(auto &i : cur) i = (i % mod + mod) % mod;
    return cur;
int get_nth(vector<int> rec, vector<int> dp, lint n){
    int m = rec.size();
    vector<int> s(m), t(m);
    s[0] = 1;
    if(m != 1) t[1] = 1;
    else t[0] = rec[0];
    auto mul = [&rec](vector<int> v, vector<int> w){
        int m = v.size();
        vector<int> t(2 * m);
        for(int j=0; j<m; j++){
            for(int k=0; k < m; k++){
                t[j+k] += 111 * v[j] * w[k] % mod;
                if(t[j+k] >= mod) t[j+k] -= mod;
            }
        for(int j=2*m-1; j>=m; j--){
            for(int k=1; k<=m; k++){</pre>
                t[j-k] += 111 * t[j] * rec[k-1] % mod;
                if(t[i-k] >= mod) t[i-k] -= mod;
            }
        t.resize(m);
        return t;
    };
```

```
while(n){
        if(n & 1) s = mul(s, t);
        t = mul(t, t);
        n >>= 1;
    lint ret = 0;
    for(int i=0; i<m; i++) ret += 111 * s[i] * dp[i] % mod;</pre>
    return ret % mod;
int guess_nth_term(vector<int> x, lint n){
    if(n < x.size()) return x[n];</pre>
    vector<int> v = berlekamp_massey(x);
    if(v.empty()) return 0;
    return get_nth(v, x, n);
struct elem{int x, y, v;};
vector<int> get_min_poly(int n, vector<elem> M)
    vector<int> rnd1, rnd2;
    mt19937 rng(0x14004);
    auto randint = [&rng](int lb, int ub){
        return uniform_int_distribution<int>(lb, ub)(rng);
    };
    for(int i=0; i<n; i++){
        rnd1.push_back(randint(1, mod - 1));
        rnd2.push_back(randint(1, mod - 1));
    }
    vector<int> gobs;
    for(int i=0; i<2*n+2; i++){
        int tmp = 0;
        for(int j=0; j<n; j++){
            tmp += 1ll * rnd2[j] * rnd1[j] % mod;
            if(tmp >= mod) tmp -= mod;
        gobs.push_back(tmp);
        vector<int> nxt(n);
        for(auto &i : M){
            nxt[i.x] += 1ll * i.v * rnd1[i.y] % mod;
            if(nxt[i.x] >= mod) nxt[i.x] -= mod;
        }
        rnd1 = nxt;
```

```
auto sol = berlekamp_massey(gobs);
        reverse(sol.begin(), sol.end());
        return sol:
    // Usage : guess_nth_term(first_values, n);
};
    Geometry
4.1 Point header
Use as typedef Point<double> P;
template \langle class T \rangle int sgn(T x) \{ return (x > 0) - (x < 0); \}
template <class T>
struct Point {
    typedef Point P_;
    T x, v;
    explicit Point(T x=0, T y=0) : x(x), y(y) {}
    bool operator<(P_ p) const { return tie(x,y) < tie(p.x,p.y); }</pre>
    bool operator==(P_ p) const { return tie(x,y)==tie(p.x,p.y); }
    P_{-} operator+(P_{-} p) const { return P_{-}(x+p.x, y+p.y); }
    P_ operator-(P_ p) const { return P_(x-p.x, y-p.y); }
    P_ operator*(T d) const { return P_(x*d, y*d); }
    P_{\text{operator}}/(T d) \text{ const } \{ \text{ return } P_{\text{operator}}/(x/d, y/d); \}
    T dot(P_ p) const { return x*p.x + y*p.y; }
    T cross(P_ p) const { return x*p.y - y*p.x; }
    T cross(P_ a, P_ b) const { return (a-*this).cross(b-*this); }
    T dist2() const { return x*x + y*y; }
    double dist() const { return sqrt((double)dist2()); }
    // angle to x-axis in interval [-pi, pi]
    double angle() const { return atan2(y, x); }
    P_ unit() const { return *this/dist(); } // makes dist()=1
    P_ perp() const { return P_(-y, x); } // rotates +90 degrees
    P_ normal() const { return perp().unit(); }
    // returns point rotated 'a' radians ccw around the origin
    P_ rotate(double a) const {
        return P_(x*cos(a)-y*sin(a),x*sin(a)+y*cos(a)); }
   friend ostream& operator<<(ostream& os, P_ p) {</pre>
        return os << "(" << p.x << "," << p.y << ")"; }
};
```

#### 4.2 Line / Line segment

```
lineDist: Returns the signed distance between point p and the line containing points
a and b. Positive value on left side and negative on right as seen from a towards b.
a==b gives nan.
lineInter, segInter: Unique intersection: (1, point), No intersection: (0, (0, 0)),
Infinitely many: (-1, (0, 0)).
segDist: Return shortest distance between p and segment [s, e].
double lineDist(const P& a, const P& b, const P& p)
    return (double)(b-a).cross(p-a)/(b-a).dist();
pair<int, P> lineInter(P s1, P e1, P s2, P e2)
    auto d = (e1 - s1).cross(e2 - s2);
    if (d == 0) // if parallel
        return \{-(s1.cross(e1, s2) == 0), P(0, 0)\};
    auto p = s2.cross(e1, e2), q = s2.cross(e2, s1);
    return \{1, (s1 * p + e1 * q) / d\};
}
double segDist(P& s, P& e, P& p)
    if (s==e) return (p-s).dist();
    auto d = (e-s).dist2(), t = min(d,max(.0,(p-s).dot(e-s)));
    return ((p-s)*d-(e-s)*t).dist()/d;
}
vector<P> segInter(P a, P b, P c, P d)
   auto oa = c.cross(d, a), ob = c.cross(d, b),
        oc = a.cross(b, c), od = a.cross(b, d);
   // Checks if intersection is single non-endpoint point.
   if (sgn(oa) * sgn(ob) < 0 && sgn(oc) * sgn(od) < 0)
       return {(a * ob - b * oa) / (ob - oa)};
   set<P> s:
   if (onSegment(c, d, a)) s.insert(a);
   if (onSegment(c, d, b)) s.insert(b);
   if (onSegment(a, b, c)) s.insert(c);
   if (onSegment(a, b, d)) s.insert(d);
   return {all(s)};
}
```

```
bool onSegment(P s, P e, P p) {
    return p.cross(s, e) == 0 \&\& (s - p).dot(e - p) <= 0;
}
4.3 Polygons
vector<P> convexHull(vector<P> pts)
    if (sz(pts) <= 1) return pts;
    sort(all(pts));
    vector<P> h(sz(pts)+1);
    int s = 0, t = 0;
    for (int it = 2; it--; s = --t, reverse(all(pts)))
        for (P p : pts) {
            while (t \ge s + 2 \&\& h[t-2].cross(h[t-1], p) \le 0) t--;
            h[t++] = p:
    return \{h.begin(), h.begin() + t - (t == 2 && h[0] == h[1])\};
}
//Returns two points with max distance on a convex hull
array<P, 2> hullDiameter(vector<P> S) {
    int n = sz(S), j = n < 2 ? 0 : 1;
    pair<11, array<P, 2>> res({0, {S[0], S[0]}});
    rep(i,0,j)
        for (;; j = (j + 1) \% n) {
            res = \max(\text{res}, \{(S[i] - S[j]).dist2(), \{S[i], S[j]\}\});
            if ((S[(i + 1) \% n] - S[i]).cross(S[i + 1] - S[i]) >= 0)
                 break:
    return res.second;
}
   • (-1, -1) if no collision,
   • (i, -1) if touching the corner i,
   • (i, i) if along side (i, i + 1),
   • (i, j) if crossing sides (i, i + 1) and (j, j + 1).
bool inPolygon(vector<P> &p, P a, bool strict = true) {
    int cnt = 0, n = sz(p);
    rep(i,0,n) {
```

```
P q = p[(i + 1) \% n];
        if (onSegment(p[i], q, a)) return !strict;
        //or: if (segDist(p[i], q, a) <= eps) return !strict;</pre>
        cnt \hat{} = ((a.y<p[i].y) - (a.y<q.y)) * a.cross(p[i], q) > 0;
    return cnt;
// Twice the signed area of polygon
template < class T>
T polygonArea2(vector<Point<T>>& v) {
   T = v.back().cross(v[0]);
    rep(i,0,sz(v)-1) = v[i].cross(v[i+1]);
    return a;
}
#define cmp(i,j) sgn(dir.perp().cross(poly[(i)%n]-poly[(j)%n]))
#define extr(i) cmp(i + 1, i) >= 0 && cmp(i, i - 1 + n) < 0
template <class P> int extrVertex(vector<P>& poly, P dir) {
    int n = sz(poly), lo = 0, hi = n;
   if (extr(0)) return 0;
    while (lo + 1 < hi) \{
        int m = (lo + hi) / 2;
        if (extr(m)) return m;
        int ls = cmp(lo + 1, lo), ms = cmp(m + 1, m);
        (ls < ms | | (ls == ms \&\& ls == cmp(lo, m)) ? hi : lo) = m;
   }
    return lo;
}
#define cmpL(i) sgn(a.cross(poly[i], b))
template <class P>
array<int, 2> lineHull(P a, P b, vector<P> poly) {
    int endA = extrVertex(poly, (a - b).perp());
    int endB = extrVertex(poly, (b - a).perp());
   if (cmpL(endA) < 0 \mid | cmpL(endB) > 0)
        return {-1, -1}:
    array<int, 2> res;
   rep(i,0,2) {
        int lo = endB, hi = endA, n = sz(poly);
        while ((lo + 1) % n != hi) {
```

```
int m = ((lo + hi + (lo < hi ? 0 : n)) / 2) % n;
            (cmpL(m) == cmpL(endB) ? lo : hi) = m;
        res[i] = (lo + !cmpL(hi)) % n;
        swap(endA, endB);
    if (res[0] == res[1]) return {res[0], -1};
    if (!cmpL(res[0]) && !cmpL(res[1]))
        switch ((res[0] - res[1] + sz(poly) + 1) \% sz(poly)) {
            case 0: return {res[0], res[0]};
            case 2: return {res[1], res[1]};
        }
    return res;
}
    Graph
5.1 Network Flow
5.1.1 Dinic's Maxflow
Fast max flow algorithm. Use maxflow and AddEdge.
struct Edge
{
    int u, v;
   ll cap, flow;
    Edge() {}
    Edge(int u, int v, ll cap): u(u), v(v), cap(cap), flow(0) {}
};
struct Dinic
    int N;
    vector<Edge> E;
    vector<vector<int>> g;
    vector<int> d, pt;
    Dinic(int N): N(N), E(O), g(N), d(N), pt(N) {}
    void AddEdge(int u, int v, ll cap)
        if (u != v)
            E.push_back(Edge(u, v, cap));
```

```
g[u].push_back(E.size() - 1);
        E.push_back(Edge(v, u, 0));
        g[v].push_back(E.size() - 1);
    }
}
bool BFS(int S, int T)
    queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    d[S] = 0;
    while(!q.empty())
        int u = q.front();
        q.pop();
        if (u == T) break;
        for (int k: g[u])
            Edge &e = E[k];
            if (e.flow < e.cap && d[e.v] > d[e.u] + 1)
                d[e.v] = d[e.u] + 1;
                q.push(e.v);
        }
    return d[T] != N + 1;
ll DFS(int u, int T, ll flow = -1)
    if (u == T || flow == 0) return flow;
    for (int &i = pt[u]; i < g[u].size(); i++)</pre>
        Edge &e = E[g[u][i]];
        Edge &oe = E[g[u][i]^1];
        if (d[e.v] == d[e.u] + 1)
            11 amt = e.cap - e.flow;
            if (flow != -1 \&\& amt > flow) amt = flow;
            if (ll pushed = DFS(e.v, T, amt))
```

```
e.flow += pushed;
                     oe.flow -= pushed;
                    return pushed;
                }
            }
        }
        return 0;
    }
    11 MaxFlow(int S, int T)
        11 total = 0;
        while (BFS(S, T))
            fill(pt.begin(), pt.end(), 0);
            while (ll\ flow = DFS(S,\ T))
                total += flow;
        return total;
};
5.1.2 Min cost Max Flow
Fast min cost max flow algorithm. Use solveMCMF and AddEdge.
solveMCMF returns pair of (max flow, cost of such flow).
const int MAXN = 1010;
struct MCMF
{
    struct edg { int pos, cap, rev, cost; };
    vector<edg> gph[MAXN];
    void clear(){
        for(int i=0; i<MAXN; i++) gph[i].clear();</pre>
    void AddEdge(int s, int e, int cap, int cst)
    // add edge (s to e, capacity = x, cost = c)
        gph[s].push_back({e, cap, (int)gph[e].size(), cst});
        gph[e].push_back({s, 0, (int)gph[s].size()-1, -cst});
    int dist[MAXN], pa[MAXN], pe[MAXN];
    bool inque[MAXN];
```

```
bool spfa(int src, int sink)
    memset(dist, 0x3f, sizeof(dist));
    memset(inque, 0, sizeof(inque));
    queue<int> que;
    dist[src] = 0;
    inque[src] = 1;
    que.push(src);
    bool ok = 0;
    while(!que.empty()){
        int x = que.front();
        que.pop();
        if(x == sink) ok = 1;
        inque[x] = 0;
        for(int i=0; i<gph[x].size(); i++)</pre>
            edg e = gph[x][i];
            if(e.cap > 0 && dist[e.pos] > dist[x] + e.cost)
                dist[e.pos] = dist[x] + e.cost;
                pa[e.pos] = x;
                pe[e.pos] = i;
                if(!inque[e.pos]){
                    inque[e.pos] = 1;
                    que.push(e.pos);
                }
            }
        }
    }
    return ok;
}
pii solveMCMF(int src, int sink){
    int MCMF_COST = 0;
    int MCMF_FLOW = 0;
    while(spfa(src, sink)){
        int cap = 1e9;
        for(int pos = sink; pos != src; pos = pa[pos]){
            cap = min(cap, gph[pa[pos]][pe[pos]].cap);
        }
        MCMF_COST += dist[sink] * cap;
        for(int pos = sink; pos != src; pos = pa[pos]){
```