



ICP Group Seminar

Robustness analysis of quantum algorithms against coherent control errors

January 30th, 2023

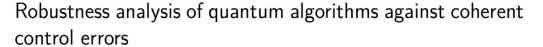
Daniel Fink



Background



- Collaboration with Julian Berberich
- Jointly working in SimTech PN8
- Paper under current development



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Noise poses a major obstacle in current quantum devices. Several recent studies have shown that coherent control errors, for which an ideal gate e^{-iH}

is perturbed by an additional error gate

1 Introduction

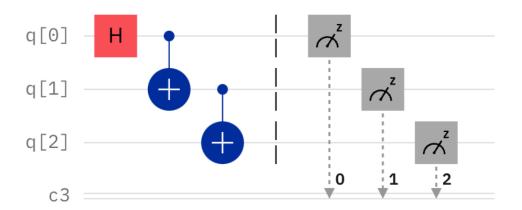
Quantum computing has emerged as a powerful tool to overcome limitations of classical computing and solve problems that were previously







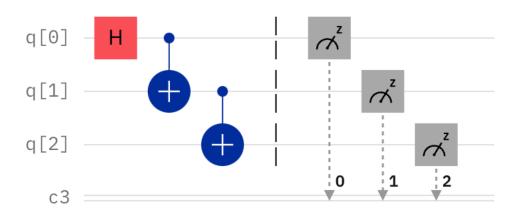








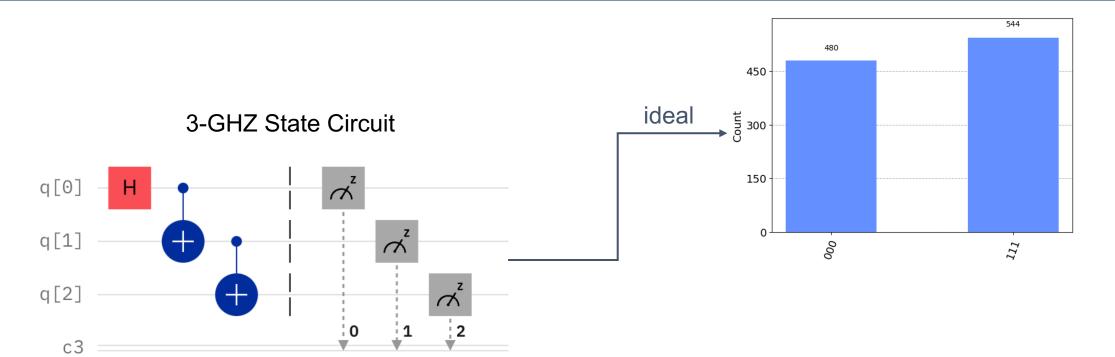
3-GHZ State Circuit



$$|\psi\rangle = \frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$$



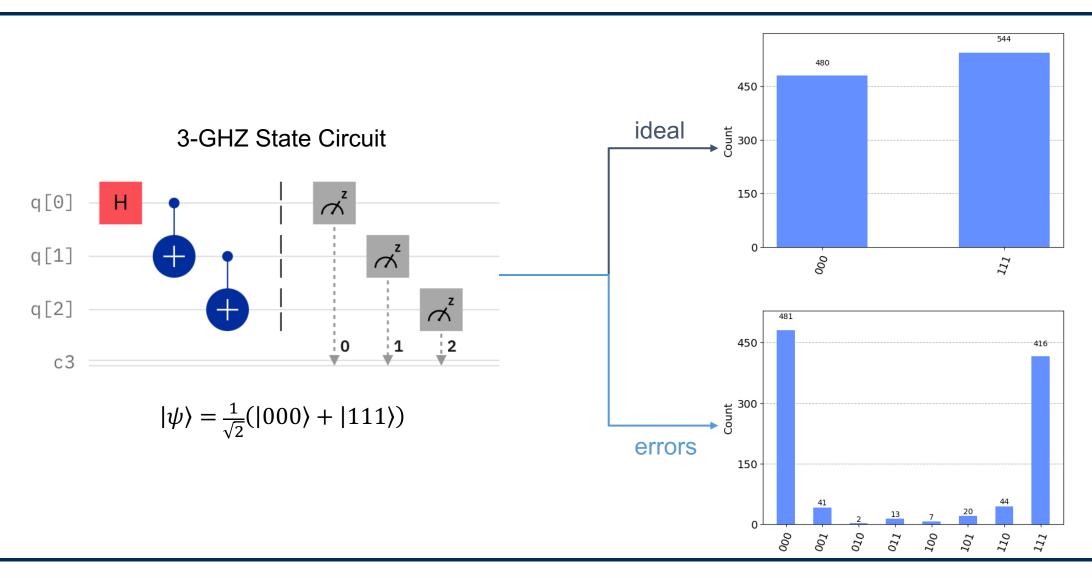




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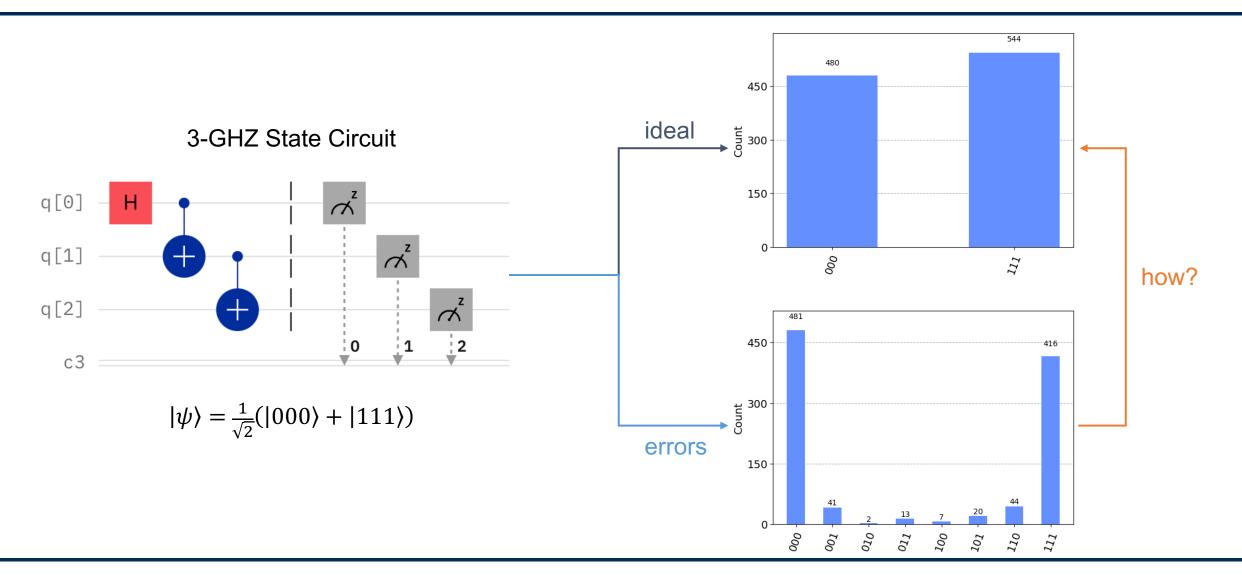














Agenda



- Quantum Computing
- Handling Errors: QEC & QEM
- Our approach: Lipschitz bounds
- Experiments
- Conclusion & Outlook







Qubits



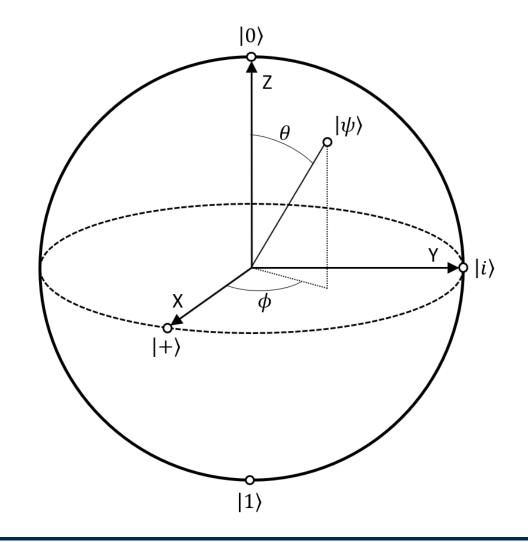
- General single qubit state: $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$, $|\alpha|^2 + |\beta|^2 = 1$
- $\rightarrow |\psi\rangle = \cos\frac{\theta}{2}|0\rangle + e^{i\phi}\sin\frac{\theta}{2}|1\rangle$, with $\theta, \phi \in \mathbb{R}$ (up to a global phase)



Bloch Sphere



$$|\psi\rangle = \cos\frac{\theta}{2}|0\rangle + e^{i\phi}\sin\frac{\theta}{2}|1\rangle$$





Quantum Gates



- Unitary operator $U: H \to H$, $U|\psi_0\rangle = |\psi\rangle$
- Examples: q[0] q[0] Y q[0] Z Pauli Gates

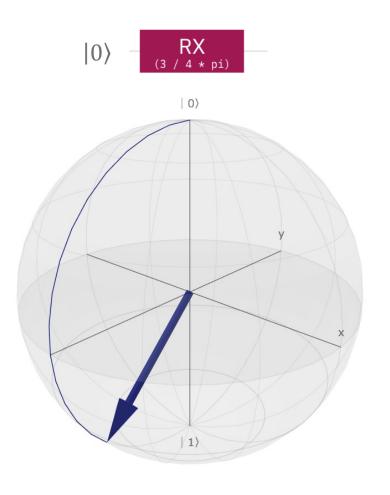






Quantum Gates







Quantum Circuits



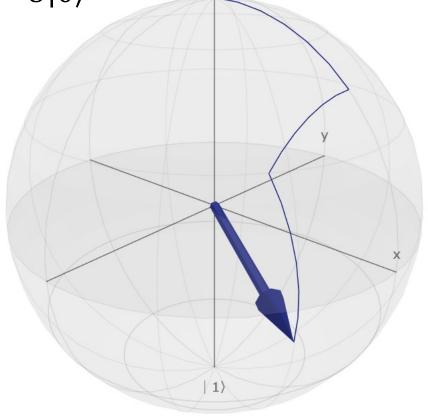
• Gates together form a circuit: $|\psi\rangle = U_N ... U_1 |0\rangle = U|0\rangle$

Example:

|0> -

RX (pi / 4)





0)



Quantum Circuits



• Gates together form a circuit: $|\psi\rangle = U_N ... U_1 |0\rangle = U|0\rangle$

Example:

 $|0\rangle$

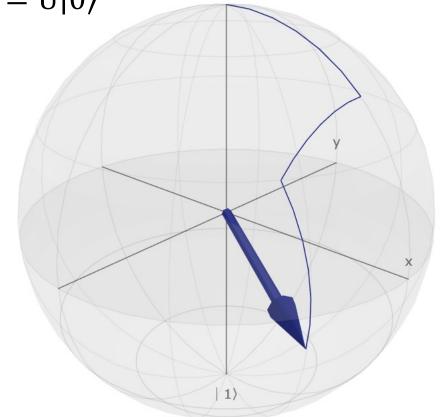








→ Transpilation is necessary

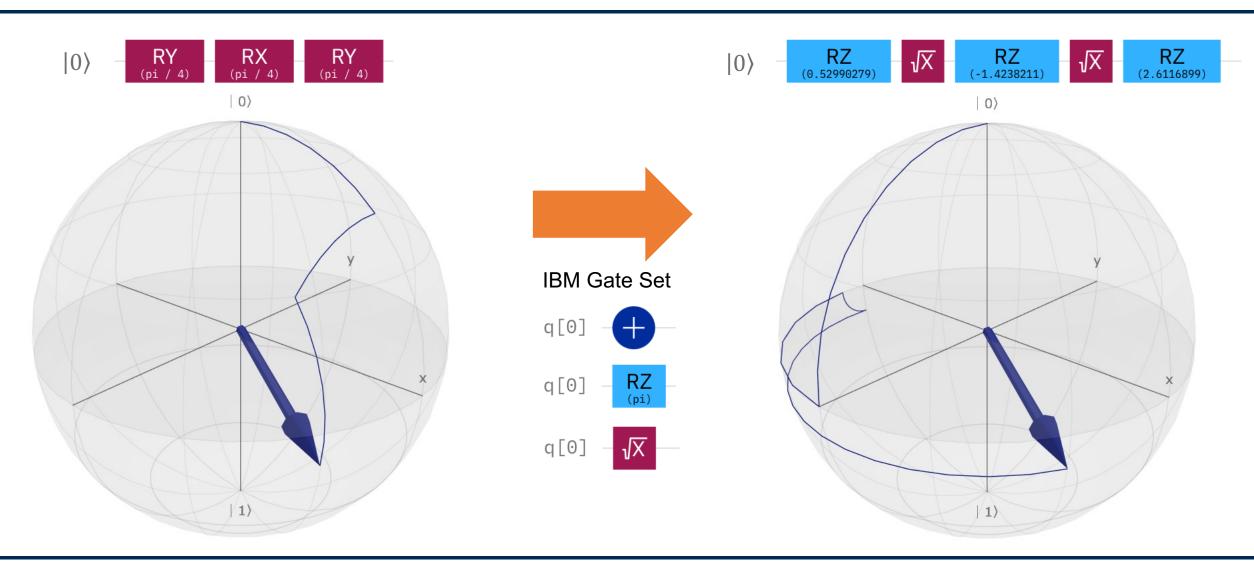


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Transpilation





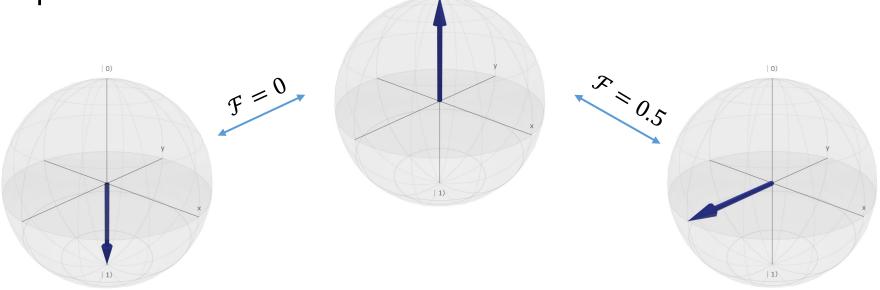


Comparing Quantum States



- Goal: compare two states $|\psi\rangle \leftrightarrow |\phi\rangle$
- Define the Fidelity $\mathcal{F}(|\psi\rangle, |\phi\rangle) = |\langle\psi|\phi\rangle| \in [0,1]$

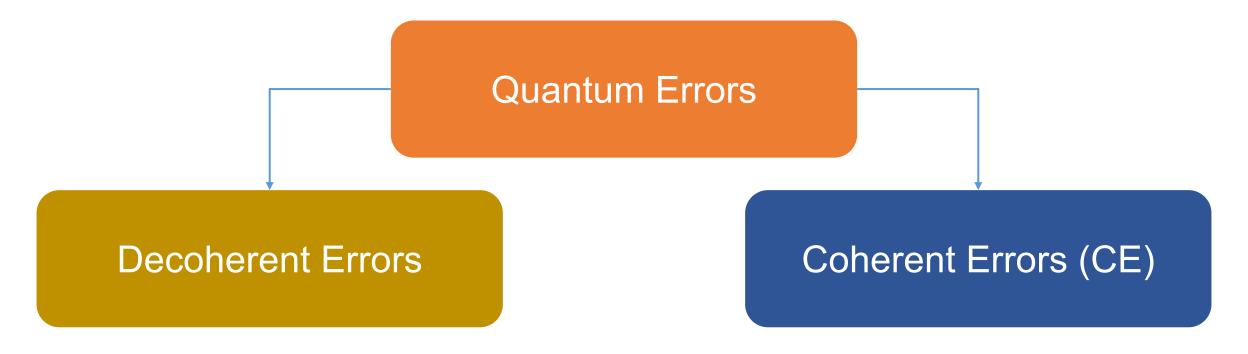






Errors in Quantum Computing





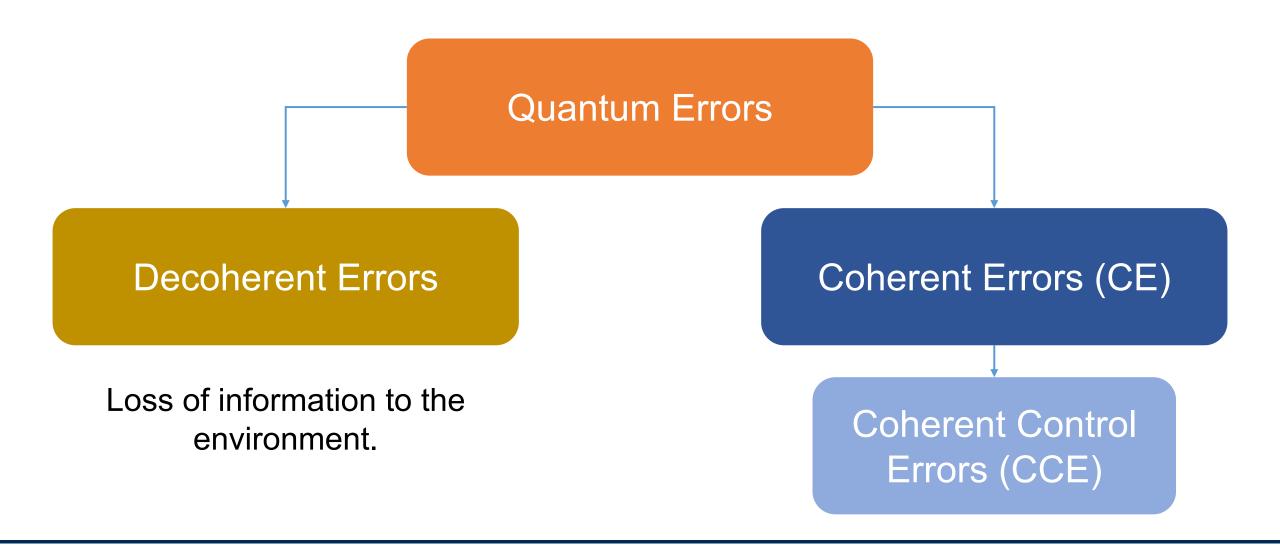
Loss of information to the environment.

Information is preserved, but unintentional operations have been applied.



Errors in Quantum Computing







Coherent Errors



- Every Unitary can be written as $U = e^{-iG}$
- $G = G^{\dagger}$ is a Hermitian operator (generator)
- E.g.: $RZ(\theta) = e^{-i\frac{\theta}{2}Z}$



Coherent Errors



- Every Unitary can be written as $U = e^{-iG}$
- $G = G^{\dagger}$ is a Hermitian operator (generator)
- E.g.: $RZ(\theta) = e^{-i\frac{\theta}{2}Z}$
- Coherent error: $e^{-i\frac{\theta}{2}Z} \rightarrow e^{-i\frac{\theta}{2}Z}e^{-iG}$

any noise generator G



Coherent Errors



- Every Unitary can be written as $U = e^{-iG}$
- $G = G^{\dagger}$ is a Hermitian operator (generator)
- E.g.: $RZ(\theta) = e^{-i\frac{\theta}{2}Z}$
- Coherent error: $e^{-i\frac{\theta}{2}Z} \rightarrow e^{-i\frac{\theta}{2}Z}e^{-iG}$
- Coherent control error: $e^{-i\frac{\theta}{2}Z} \rightarrow e^{-i(1+x)\frac{\theta}{2}Z}$

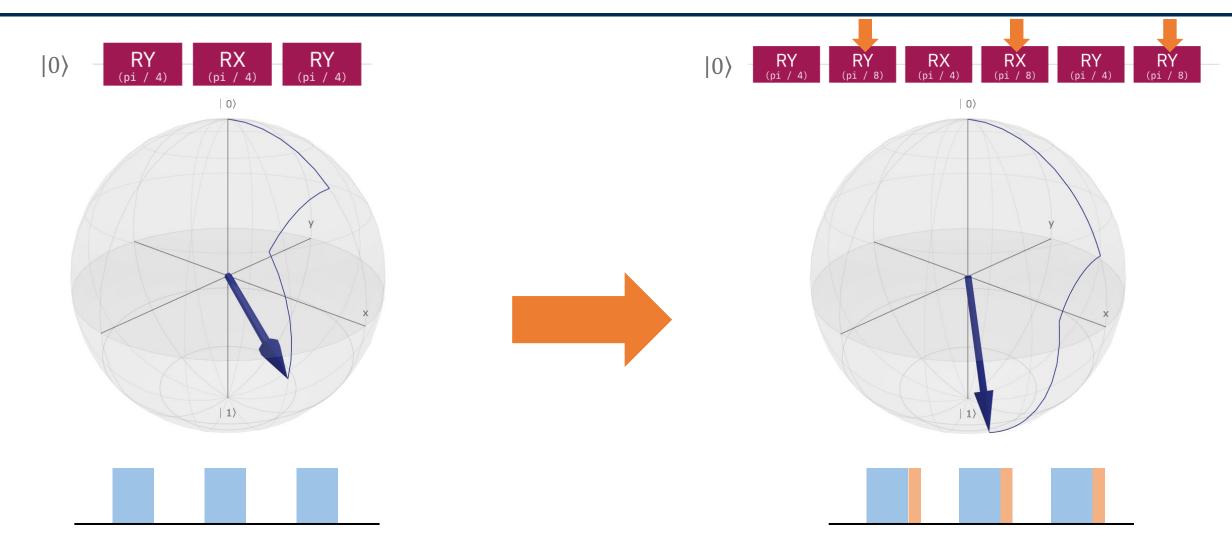
any noise generator G

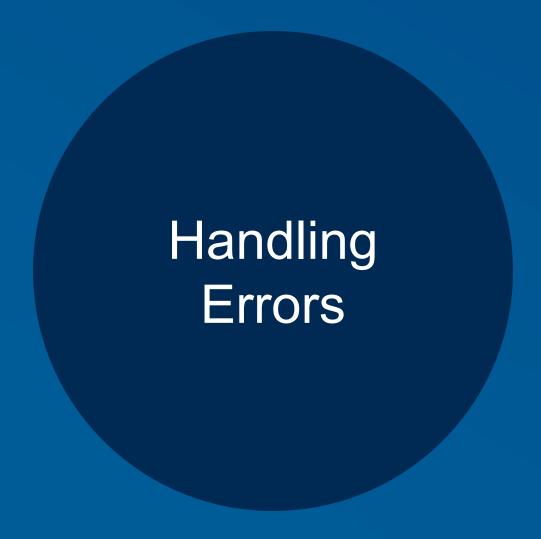
same generator



Coherent Control Errors



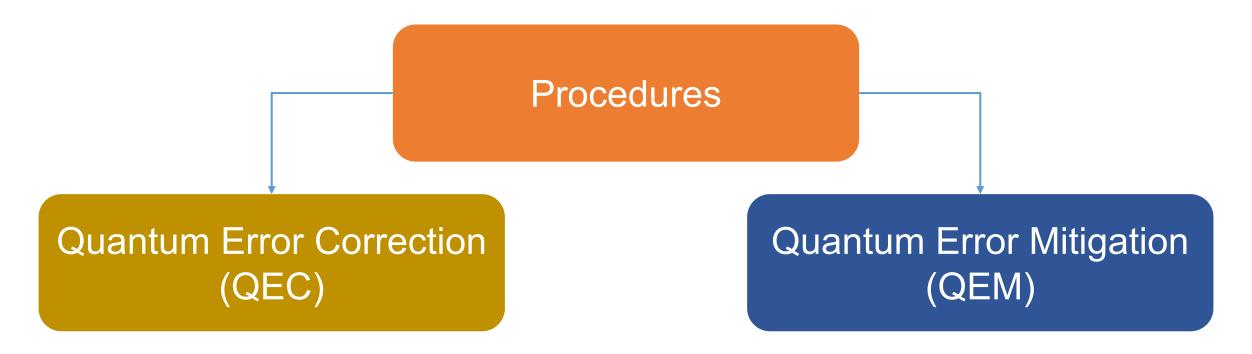






Handling Errors





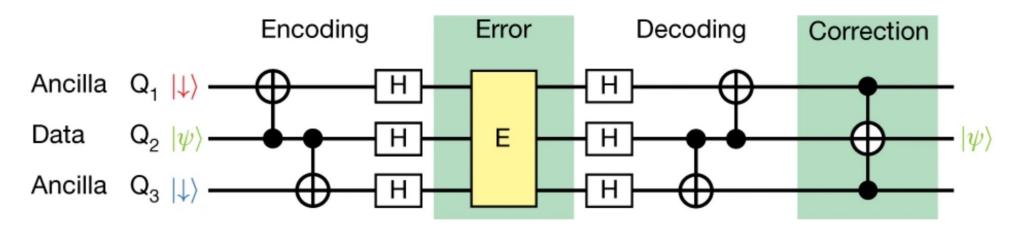
Add additional gates to detect and compensate for errors.

Reduce noise effects via classical post-processing.



Quantum Error Correction



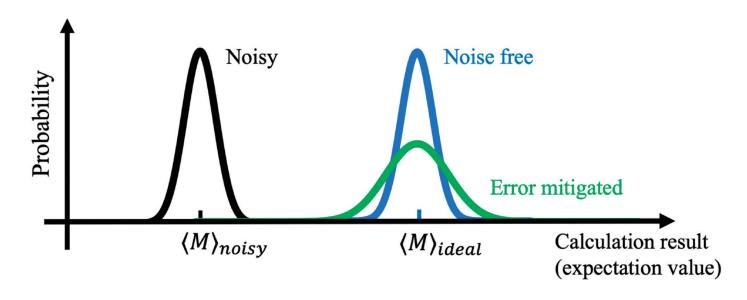


Takeda et al., "Quantum error correction with silicon spin qubits", Nature 608, 682–686 (2022)

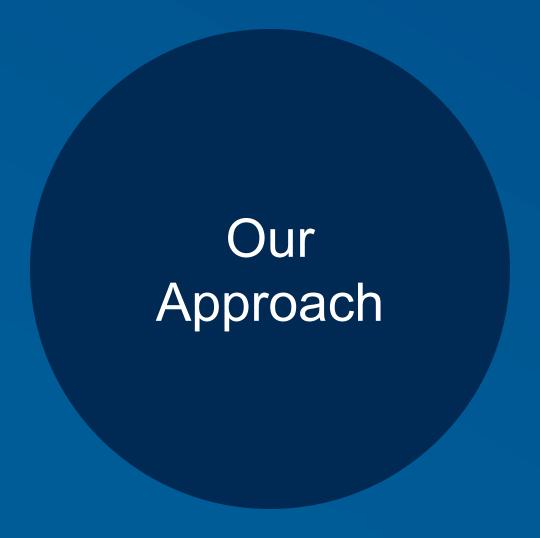


Quantum Error Mitigation





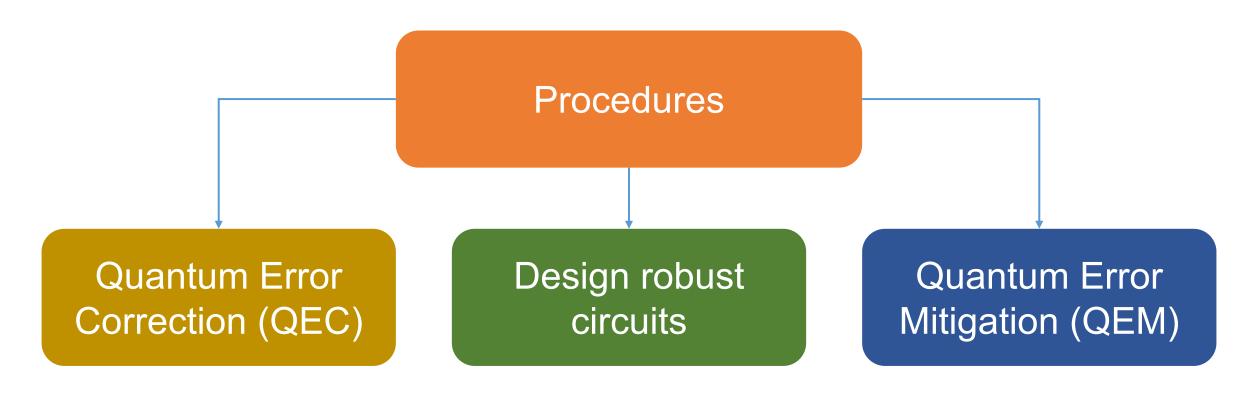
Endo et al., "Hybrid Quantum-Classical Algorithms and Quantum Error Mitigation", Journal of the Physical Society of Japan, 90, 032001 (2021)





Our Approach





Design robust circuits in the first place. Can be used together with QEC/QEM.





Noise-free circuit

$$|\hat{\psi}\rangle = \widehat{\mathbf{U}}_N \widehat{\mathbf{U}}_{N-1} \dots \widehat{\mathbf{U}}_1 |\psi_0\rangle$$

Noisy circuit

$$|\psi(x)\rangle = \widehat{\mathrm{U}}_N(x_N)\widehat{\mathrm{U}}_{N-1}(x_{N-1}) \dots \widehat{\mathrm{U}}_1(x_1)|\psi_0\rangle$$

with noise $x \in \mathbb{R}^N$

$$|\psi(0)\rangle = |\hat{\psi}\rangle$$

Noise level

$$\epsilon \in \mathbb{R}^+$$
 such that $||x||_2 < \epsilon$





Definition

A scalar L > 0 is a Lipschitz bound of $x \mapsto |\psi(x)\rangle$ if

$$\||\psi(x)\rangle - |\psi(x')\rangle\|_2 \le L\|x - x'\|_2$$

for all $x, x' \in \mathbb{R}^N$.





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for all $x, x' \in \mathbb{R}^N$.

- The minimal *L* is called the Lipschitz constant.
- L bounds the worst-case amplification of a perturbation x:

$$\||\psi(x)\rangle - |\hat{\psi}\rangle\|_2 \le L\|x\|_2$$





Theorem

For any $x \in \mathbb{R}^N$ with $||x||_2 < \epsilon$, and any initial state $|\psi_0\rangle$, we have

$$\left|\left\langle \psi(x)\right|\hat{\psi}\right\rangle\right| \geq 1 - \frac{L^2 \epsilon^2}{2},$$

with L > 0 being a Lipschitz bound of $x \mapsto |\psi(x)\rangle$.





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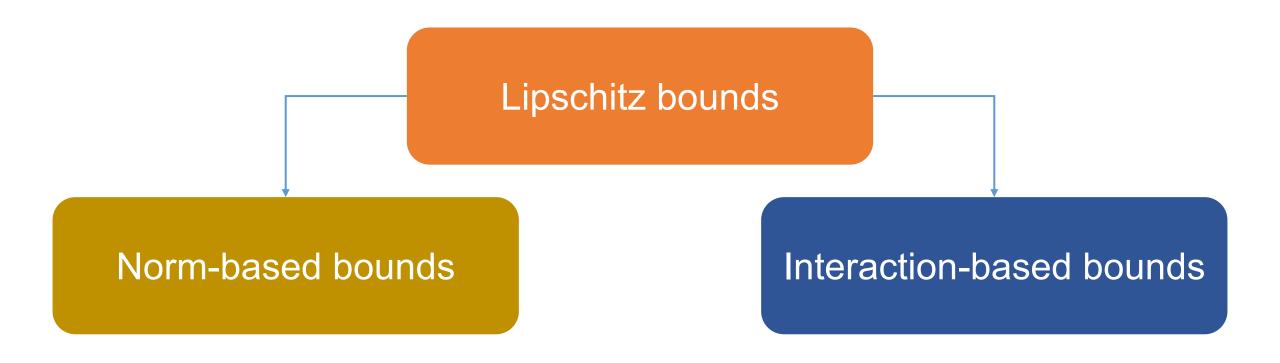
$$\left|\left\langle \psi(x)\right|\hat{\psi}\right\rangle\right| \geq 1 - \frac{L^2\epsilon^2}{2},$$

with L > 0 being a Lipschitz bound of $x \mapsto |\psi(x)\rangle$.

- The Lipschitz constant can be hard to compute.
- We show a way how to calculate Lipschitz bounds.









Lipschitz Bounds



Theorem: Norm-based bounds

The following is a Lipschitz bound of $x \mapsto |\psi(x)\rangle$:

$$L = \sum_{i=1}^{N} \|\mathbf{G}_i\|_2$$



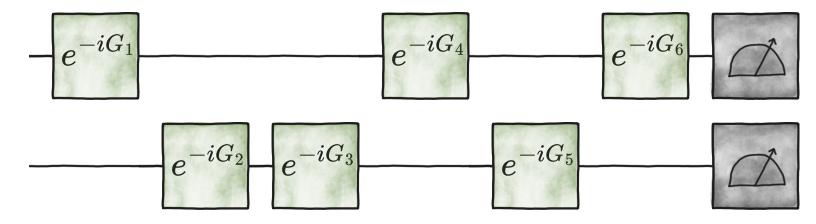
Lipschitz Bounds



Theorem: Norm-based bounds

The following is a Lipschitz bound of $x \mapsto |\psi(x)\rangle$:

$$L = \sum_{i=1}^{N} \|\mathbf{G}_i\|_2$$



$$L = ||G_1||_2 + ||G_2||_2 + ||G_3||_2 + ||G_4||_2 + ||G_5||_2 + ||G_6||_2$$



Lipschitz Bounds



Theorem: Interaction-based bounds

The following are Lipschitz bounds of $x \mapsto |\psi(x)\rangle$:

$$\sum_{i=1}^{\frac{N}{2}} \| [G_{2i-1} \quad G_{2i}] \|_{2}$$

If N is odd:
$$||G_N||_2 + \sum_{i=1}^{\frac{N-1}{2}} ||[G_{2i-1} \quad G_{2i}]||_2$$

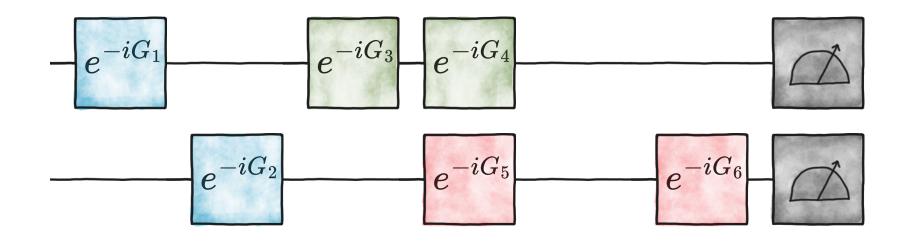
■ [A B] is a block matrix:





Lipschitz Bounds





$$L = ||G_1||_2 + ||G_2||_2 + ||[G_3 \quad G_4]||_2 + ||[G_5 \quad G_6]||$$





Which bounds are better?

$$||[G_1 \quad G_2]||_2 \le ||G_1||_2 + ||G_2||_2$$

$$||[G_1 \quad G_2]||_2 < ||G_1||_2 + ||G_2||_2$$

How big is the gap?

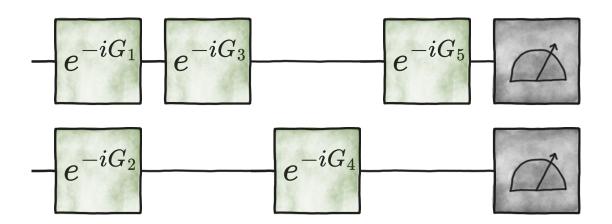
→ Current investigation

Design Guidelines



Design Guidelines



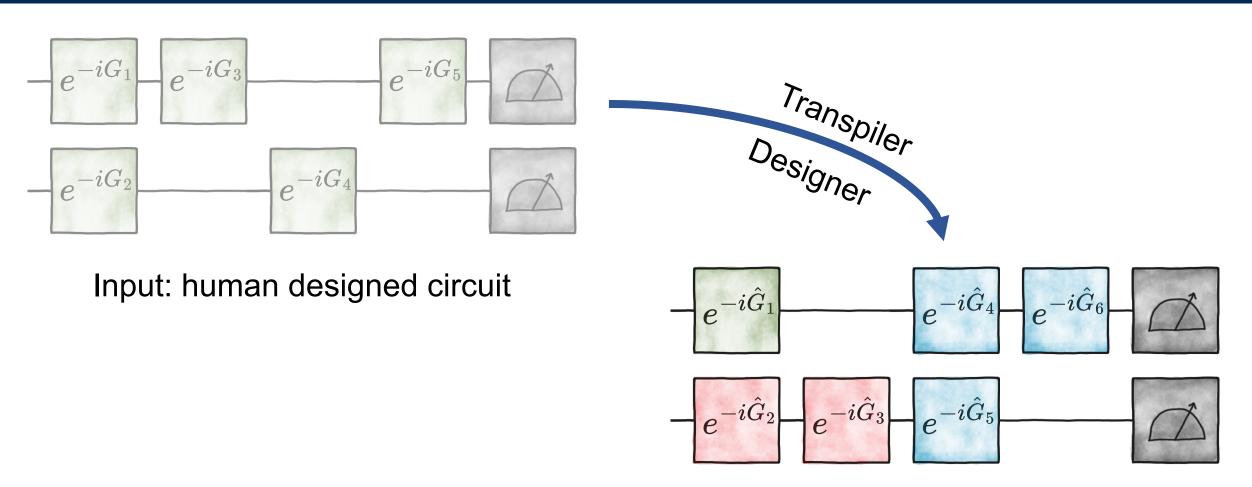


Input: human designed circuit



Design Guidelines





Output: transpiled resilient circuit





Experiments



- Target: compare circuits with different Lipschitz bounds
- Goal: show lower Lipschitz bounds imply robustness



Experiments

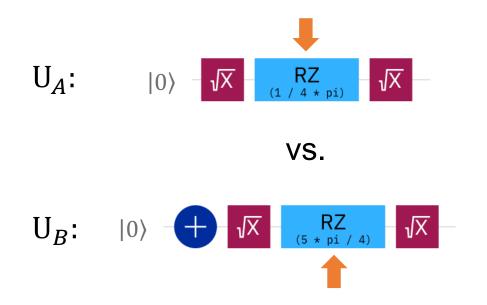


- Target: compare circuits with different Lipschitz bounds
- Goal: show lower Lipschitz bounds imply robustness
- Procedure:
 - 1. Choose a set of noise levels $\{\epsilon\} \subseteq \mathbb{R}^+$
 - 2. For each ϵ , draw several $x \in B_{\epsilon}(0) \subseteq \mathbb{R}^N$ uniformly
 - 3. Insert CCE gates: $e^{-i\frac{\theta}{2}Z} \rightarrow e^{-i(1+x_i)\frac{\theta}{2}Z}$
 - 4. Calculate $\mathcal{F} = \left| \left\langle \psi_{\epsilon}(x) | \hat{\psi} \right\rangle \right|$



Experiment I

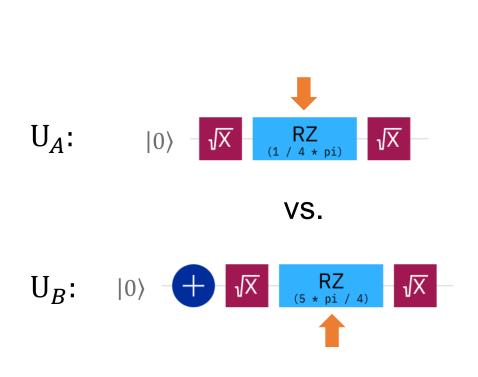


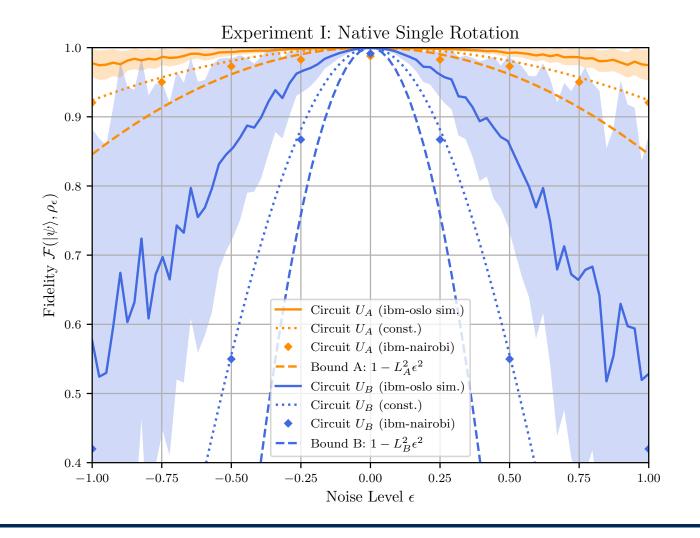




Experiment I



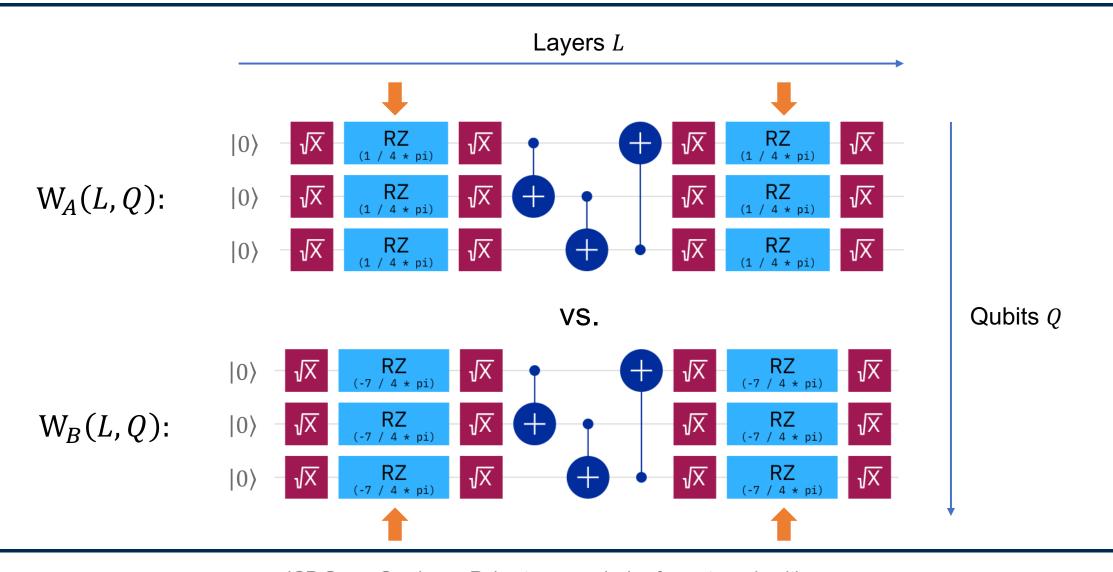






Experiment II

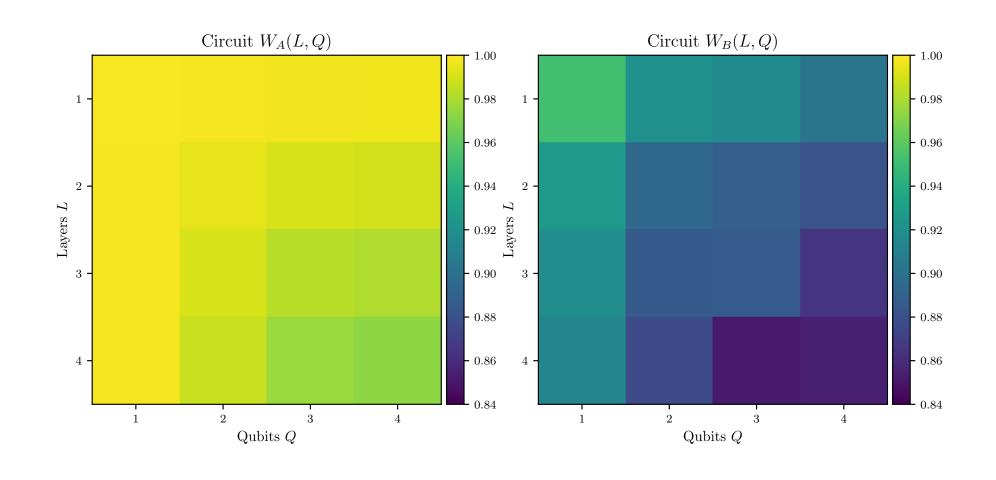






Experiment II





Conclusion & Outlook



Conclusion



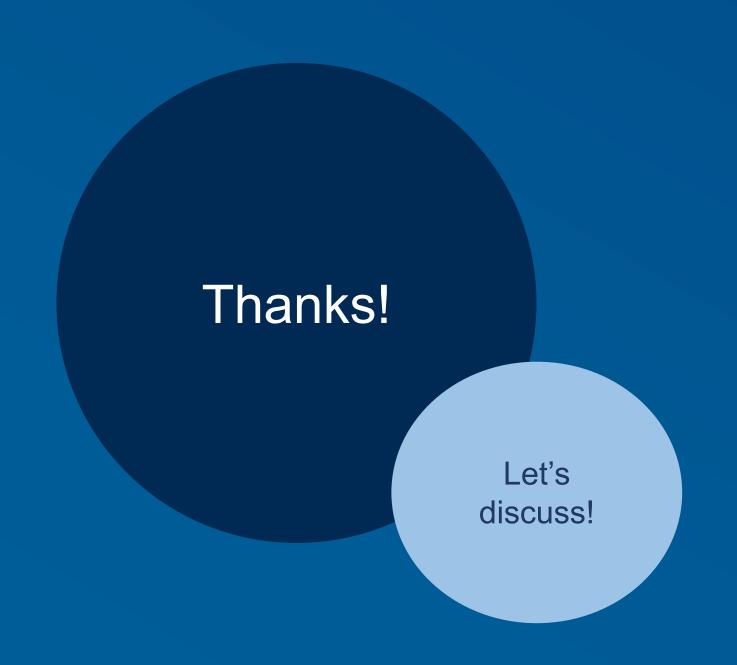
- Framework for robustness analysis for C(C)E
- Derived worst-case error bounds
- Defined guidelines for quantum algorithm design
- Performed numerical validation



Outlook



- Extend the framework to account for decoherent errors $|\psi\rangle \rightarrow \rho$
- Connect worst-case Lipschitz bounds with QEC/QEM
- Integrate the framework into QML
 - → Lipschitz bounds give rise to the use of regularization in QML









$$\begin{aligned} \| [H_1 \quad H_2] \|_2 &= \sqrt{\lambda_{\max}(H_1^{\dagger} H_1 + H_2^{\dagger} H_2)} \\ &\leq \sqrt{\lambda_{\max}(H_1^{\dagger} H_1) + \lambda_{\max}(H_2^{\dagger} H_2)} \\ &\leq \sqrt{\lambda_{\max}(H_1^{\dagger} H_1) + \sqrt{\lambda_{\max}(H_2^{\dagger} H_2)}} \\ &= \| H_1 \|_2 + \| H_2 \|_2. \end{aligned}$$





$$\lambda_{\max}(H_1^{\dagger}H_1 + H_2^{\dagger}H_2)$$

$$\leq \lambda_{\max}(H_1^{\dagger}H_1) + \lambda_{\max}(H_2^{\dagger}H_2)$$

This inequality is strict if and only if the eigenvectors corresponding to the maximum eigenvalues of $H_1^{\dagger}H_1$ and $H_2^{\dagger}H_2$ do not align, i.e.,

$$\underset{\|v\|_2=1}{\operatorname{argmax}} v^{\dagger} H_1^{\dagger} H_1 v \cap \underset{\|v\|_2=1}{\operatorname{argmax}} v^{\dagger} H_2^{\dagger} H_2 v$$
$$= \emptyset.$$





Theorem Reversed

If we want to guarantee a worst-case fidelity of no less than \mathcal{F} for any $x \in \mathbb{R}^n$ with $||x||_2 < \epsilon$ and any initial state $|\psi_0\rangle$, then the noise level must be bounded by

$$\epsilon \leq \frac{\sqrt{2}}{L} \sqrt{1 - \mathcal{F}}$$
.



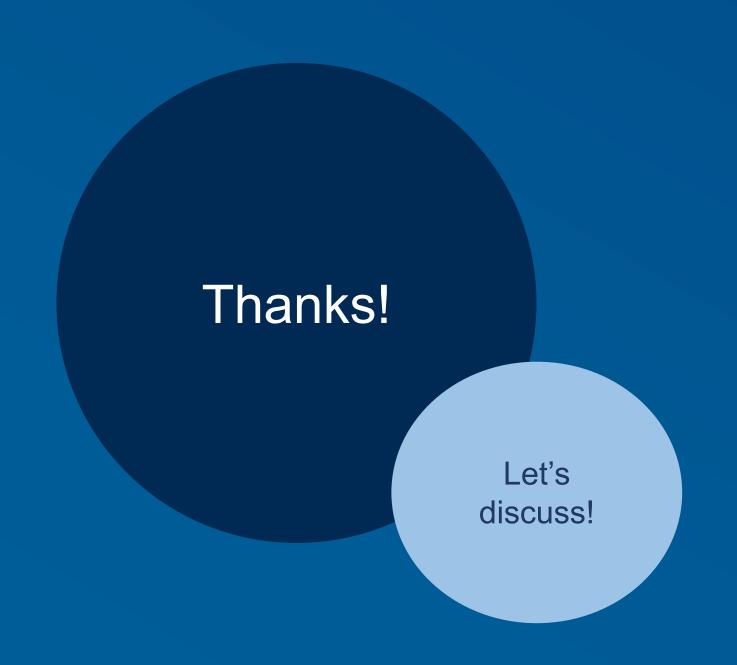


Theorem Reversed

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.

- ullet may be determined via calibration results.
- Can be used to determine if QEC is possible.









Norm-based bounds

Considering $||G||_2$ of each gate e^{-iG}

Interaction-based bounds

Considering subsequent gates $e^{-iG}e^{-iH}$





Problem

Find L such that, for any $x \in \mathbb{R}^N$ with $||x||_2 < \epsilon$, and any initial state $|\psi_0\rangle$, it holds that

$$\left|\left\langle \psi(x)\right|\hat{\psi}\right\rangle\right| \geq 1 - g(\epsilon, L),$$

where L > 0 depends only on the circuit components.





Problem

Find L such that, for any $x \in \mathbb{R}^N$ with $||x||_2 < \epsilon$, and any initial state $|\psi_0\rangle$, it holds that

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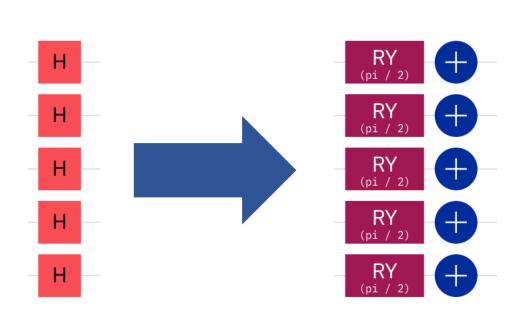
- This is a worst-case bound w.r.t. CCEs and the initial states.
- We show: L can be a Lipschitz bound of $x \mapsto |\psi(x)\rangle$.

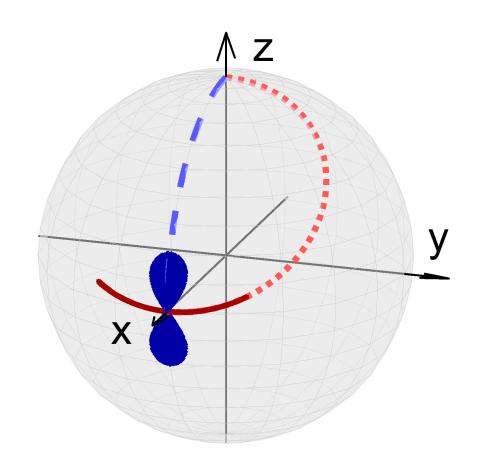


Design Guidelines





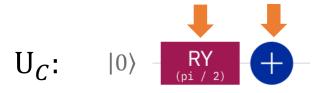






Experiment II





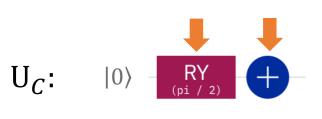
VS.

$$U_D$$
: $|0\rangle$



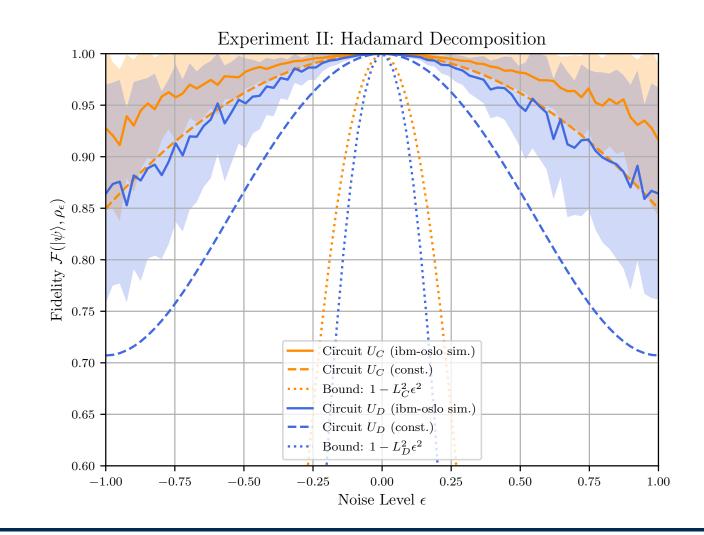
Experiment II





VS.

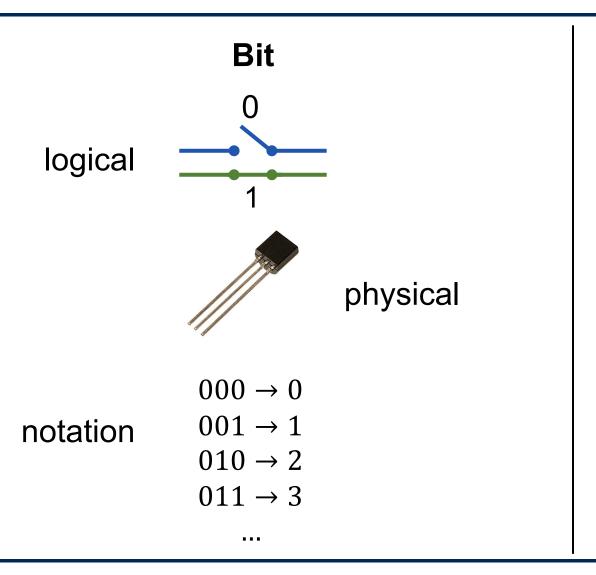
 U_D : $|0\rangle$





Qubits

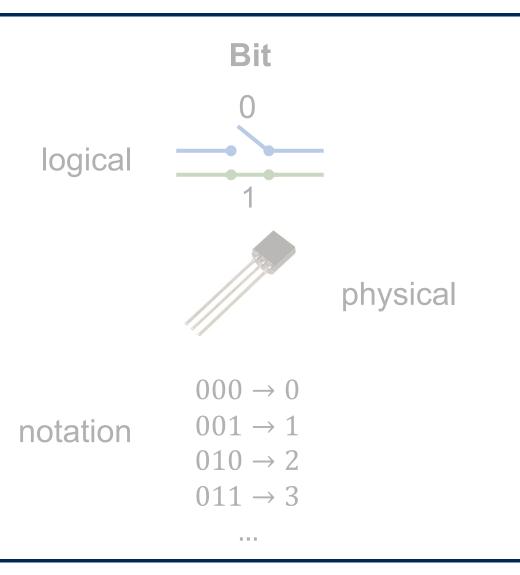


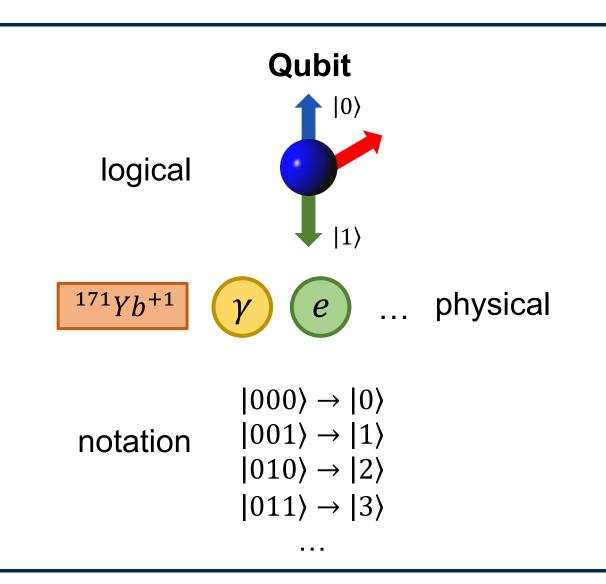




Qubits









Qubits



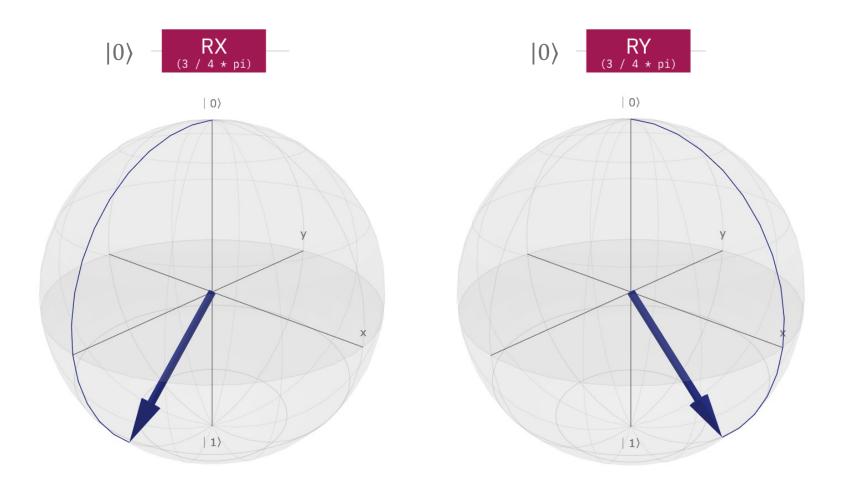
- Computational basis: $|0 \dots 0\rangle, \dots, |1 \dots 1\rangle \in H$
- General single qubit state: $|\psi\rangle = \alpha|0\rangle + \beta|1\rangle$, $|\alpha|^2 + |\beta|^2 = 1$

 $\blacksquare \text{ Measurement: } |\psi\rangle \to \begin{cases} 0 \text{ with } P(0) = |\alpha|^2 \\ \\ 1 \text{ with } P(1) = |\beta|^2 \end{cases}$



Quantum Gates







Quantum Gates



