

1 Modeling

Formal specification of a buffer with an infinite number of states.

1.1 Rigid data types

Example 1. *Specification of lists of arbitrary elements:*

```
spec!  LISTS
pr  BOOL
sorts Elt List .
op err : -> Elt .
op empty : -> List .
op _;_ : Elt List -> List .
op _in_ : Elt List -> Bool .
eq-1  $\forall L \cdot L ; \text{err} = L$ 
eq-2  $\forall E \cdot E \text{ in empty} = \text{false}$ 
eq-3  $\forall E, E' \cdot (E \text{ in } E' ; L) \text{ if } E = E'$ 
eq-4  $\forall E, E' \cdot E \text{ in } E' ; L = E \text{ in } L \text{ if } \neg(E = E')$ 
```

1.2 Nominals

Example 2. *Specification of nominals:*

```
spec!  NOMINAL
sort Nominal .
op init : -> Nominal .
op next : Nominal -> Nominal .
```

1.3 Flexible data types

Example 3. *Specification of the attributes read and del:*

```
spec BUFFER[LISTS, NOMINAL]
op read : List -> Elt .
op del : List -> List .
eq-5  $\forall Z \cdot @_Z \text{read}(\text{empty}) = \text{err}$ 
eq-6  $\forall Z \cdot @_Z \text{del}(\text{empty}) = \text{empty}$ 
eq-7  $\forall E, L \cdot @_{\text{init}} \text{read}(E ; L) = E$ 
eq-8  $\forall E, L \cdot @_{\text{init}} \text{del}(E ; L) = L$ 
eq-9  $\forall Z, E, L \cdot @_{\text{next}(Z)} \text{read}(E ; L) = @_Z \text{read}(L)$ 
eq-10  $\forall Z, E, L \cdot @_{\text{next}(Z)} \text{del}(E ; L) = E ; \text{del}(Z, L)$ 
```

2 Formal verification

The property we are interested in proving formally is $\Gamma \vdash \forall L, E \cdot \exists Z \cdot \text{read}(Z, L) = E \text{ if } (E \text{ in } L) = \text{true}$, where Γ consists of all equations in BUFFER. We proceed by induction on the structure of L .

- induction base: $\Gamma \vdash \forall E \cdot \exists Z \cdot \text{read}(Z, \text{empty}) = E \text{ if } (E \text{ in empty}) = \text{true}$
- induction step: