1 Modeling

We present the specification of a mutual exclusion protocol.

1.1 Rigid data types

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Example 1 (Labels).
spec! LABEL
sort Label .
ops re wt cs : -> Label [ctor].
op _~_ : Label Label -> Bool [comm].
var L : Label .
eq (re ~ wt) = false .
eq (re \sim cs) = false .
eq (wt \sim cs) = false .
ceq true = false if re = wt .
ceq true = false if re = cs .
ceq true = false if wt = cs.
Example 2 (Process identifiers).
spec* PID
inc BOOL .
sort Pid .
op _~_ : Pid Pid -> Bool [comm].
vars I J : Pid .
eq I \tilde{} I = true .
ceq I = J if I ~ J [nonexec].
Example 3 (lists of process identifiers).
spec! SEQUENCE{X :: PID}
sort Sequence .
subsorts X$Pid < Sequence .
--- constructors
op empty : -> Sequence [ctor] .
op _,_ : Sequence Sequence -> Sequence [ctor id: empty assoc].
vars Q Q' : Sequence . var I : X$Pid .
op top : Sequence -> X$Pid .
eq top(empty) = empty .
eq top(I,Q) = I .
op get : Sequence -> Sequence .
eq get(empty) = empty .
eq get(I,Q) = Q.
ceq true = false if Q,I,Q' := empty .
ceq [lemma-top]: top(Q,I) = top(Q) if top(Q) :: X$Pid.
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1.2 Nominals

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Example 4 (Agents).
spec* AGENT
sort Agent
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Example 5 (Nominals).
   spec! NOMINAL{Y :: AGENT}
   sorts Sys.
   --- actions
   op init : -> Sys [ctor].
   ops want try exit : Sys Y$Agent -> Sys [ctor].
          Flexible data types
  1.3
  Example 6 (Mutual exclusion protocol).
    spec* QLOCK{X :: PID, Y :: AGENT}
    pr SEQUENCE{X} . pr NOMINAL{Y} . pr LABEL .
    --- observers
    op pid:\rightarrow X\$Pid --- extract pid from agents
    op \operatorname{sq} : \to \operatorname{Sequence} \operatorname{\operatorname{\mathsf{---}}} \operatorname{\mathsf{gives}} the waiting queue for each state
    op \operatorname{pc} : \to \operatorname{Label} --- indicates the label of each agent at a given state
    --- variables
    vars S S_1 S_2: Sys
    vars I J K: X$Pid
    vars A B C:Y$Agent
    var Q: Sequence
      --- restrictions ---
 (1) \forall A, S_1, S_2 \cdot @_{S_1} @_A \text{pid} = @_{S_2} @_A \text{pid} --- pid depends only of the agent
 (2) \forall A, B, S \cdot @_S @_A sq = @_S @_B sq --- sq depends only of the current state
      --- init ---
 (3) \forall A \cdot @_{init} @_{A} pc = re
 (4) @_{init} sq = empty
      --- want ---
 (5) \forall S, A, B \cdot @_{want(S,A)} @_B pc = wt if @_S @_A pc = re   A = B
 (6) \forall S, A, B \cdot @_{want(S,A)} @_B pc = @_S @_B pc \text{ if } \neg(A = B)
 (7) \forall S, A, B \cdot @_{want(S,A)} @_B pc = @_S @_B pc \text{ if } \neg (@_A @_S pc = re)
 (8) \forall S, A \cdot @_{want(S,A)} sq = (@_S sq), (@_A pid) if @_S @_A pc = re
 (9) \forall S, A \cdot @_{want(S,A)} sq = @_S sq if \neg (@_S @_A pc = re)
      --- try ---
(10) \forall S, A, B \cdot @_{try(S,A)} @_B pc = cs \text{ if } @_S @_A pc = wt \land (@_A pid), Q := @_S sq \land A = B
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(11) $\forall S, A, B \cdot @_{trv(S,A)} @_B pc = @_S @_B pc \text{ if } \neg (A = B)$

- $(12) \ \forall S,A,B \cdot @_{\texttt{try}(S,A)} \, @_B \, pc = @_S \, @_B \, pc \ \text{if} \ \neg (@_S \, @_A \, pc = \texttt{wt})$
- $(13) \ \forall S,A,B \cdot @_B @_{\texttt{try}(S,A)} \, pc = @_B \, @_S \, pc \ \text{if} \ \neg (\texttt{top}(@_S \, sq) = @_A \, pid)$
- (14) $\forall S, A \cdot @_{try(S,A)} sq = @_S sq$
- --- exit ---
- (15) $\forall S, A, B \cdot @_{\texttt{exit}(S,A)} @_B pc = \texttt{re if } @_S @_A pc = \texttt{cs } \bigwedge A = B$
- (16) $\forall S, A, B \cdot @_{\texttt{exit}(S,A)} @_B pc = @_A @_B pc \text{ if } \neg(A = B)$
- (17) $\forall S, A, B \cdot @_{\texttt{exit}(S,A)} @_B pc = @_A @_B pc \text{ if } \neg (@_S @_A pc = cs)$
- (18) $\forall S, A \cdot @_{\texttt{exit}(S,A)} sq = get(@_S sq) \text{ if } @_S @_A pc = cs$
- (19) $\forall S, A \cdot @_{\texttt{exit}(S,A)} sq = @_S sq \text{ if } \neg (@_S @_A pc = cs)$