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DISCRETE MATHEMATICS

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Notation

SYNTAX	PRONUNCIATION	TYPE	SEMANTICS
$x := y$	"Define x to be y ."	<i>def.</i>	Lets the label " x " refer to y .
$x = y$	" x equals y ."	<i>pred.</i>	Asserts x and y are identical.
$x \in y$	" x is an element of y ."	<i>pred.</i>	Asserts that x is inside of y .
$p \vdash q$ $p \Rightarrow q$	" p proves q ." " p implies q ."	<i>judg.</i>	We can syntactically derive q from the assumption p .
$p \equiv q$ $p \Leftrightarrow q$	" p is equivalent to q ." " p if and only if q ."	<i>judg.</i>	The statements p and q are semantically interchangeable.
\top	"True."	<i>prop.</i>	Asserts truth.
\perp	"False."	<i>prop.</i>	Asserts falsity.
$\neg p$	"Not p ."	<i>prop.</i>	Asserts the negation of p .
$p \wedge q$	" p and q ."	<i>prop.</i>	Asserts both p and q .
$p \vee q$	" p or q ."	<i>prop.</i>	Asserts at least one of p or q .
$p \rightarrow q$	"If p , then q ."	<i>prop.</i>	Asserts q whenever p is true.
$\forall x(\varphi(x))$	"For all x , x satisfies φ ."	<i>prop.</i>	Asserts φ about every object.
$\exists x(\varphi(x))$	"There exists an x such that x satisfies φ ."	<i>prop.</i>	Asserts φ about some object.
$(\forall x \in y)(\varphi(x))$	"For all x in y , x satisfies φ ."	<i>prop.</i>	Asserts φ about every x in y .
$(\exists x \in y)(\varphi(x))$	"There exists an x in y such that x satisfies φ ."	<i>prop.</i>	Asserts φ about some x in y .
$\{x \mid \varphi(x)\}$	"the collection of all x such that x satisfies φ "	<i>obj.</i>	The collection of all and only the objects with property φ .
$\{x \in y \mid \varphi(x)\}$	"the set of all x in y such that x satisfies φ "	<i>obj.</i>	The set of all and only the elements of y with property φ .
$f : x \rightarrow y$	" f is a function from x to y "	<i>judg.</i>	A function f whose domain is x and codomain is y .
$\mathbb{P}(x)$	"the power set of x "	<i>obj.</i>	The set of all subsets of x .
$\text{succ}(x)$	"the successor of x "	<i>obj.</i>	The "next thing after" x .
$\llbracket n \rrbracket$	"bracket n "	<i>obj.</i>	The set $\{m \in \mathbb{N} \mid m < n\}$ of the first n natural numbers.

Table 1. An overview of some important notation used throughout this text. Most of this notation is standard, but some things (e.g., \mathbb{P} , succ , $\llbracket n \rrbracket$) are not yet standardized.

Short for $\forall x(x \in y \rightarrow \varphi(x))$.

Short for $\exists x(x \in y \wedge \varphi(x))$.

We only refer to a collection as a "set" when we can prove it exists.

Depending on context, the expression $f : x \rightarrow y$ can refer to an *object* (i.e., a *function* named f whose domain is the set x and codomain is the set y), or it can be a *proposition* (i.e., the *assertion* that f is a function whose domain is the set x and codomain is the set y).

Writing Systems

GREEK	NAME	IPA
Α α	<i>alpha</i>	[a]
Β β	<i>beta</i>	[v]
Γ γ	<i>gamma</i>	[ɣ]
Δ δ	<i>delta</i>	[ð]
Ε ε	<i>epsilon</i>	[e]
Ζ ζ	<i>zeta</i>	[z]
Η η	<i>eta</i>	[i]
Θ θ	<i>theta</i>	[θ]
Ι ι	<i>iota</i>	[i]
Κ κ	<i>kappa</i>	[k]
Λ λ	<i>lambda</i>	[l]
Μ μ	<i>mu</i>	[m]
Ν ν	<i>nu</i>	[n]
Ξ ξ	<i>xi</i>	[ks]
Ο ο	<i>omicron</i>	[o]
Π π	<i>pi</i>	[p]
Ρ ρ	<i>rho</i>	[r]
Σ σ	<i>sigma</i>	[s]
Τ τ	<i>tau</i>	[t]
Υ υ	<i>upsilon</i>	[i]
Φ φ	<i>phi</i>	[p ^h]
Χ χ	<i>chi</i>	[k ^h]
Ψ ψ	<i>psi</i>	[ps]
Ω ω	<i>omega</i>	[ɔ:]

HEBREW	NAME	IPA
א	<i>aleph</i>	[ʔ]
ב	<i>beth</i>	[v]
ג	<i>gimel</i>	[ɣ]
ד	<i>daleth</i>	[ð]
ה	<i>he</i>	[h]
ו	<i>waw</i>	[v]
ז	<i>zayin</i>	[z]
ח	<i>cheth</i>	[χ]
ט	<i>teth</i>	[t]
י	<i>yod</i>	[j]
כ	<i>kaf</i>	[χ]
ל	<i>lamed</i>	[l]
מ	<i>mem</i>	[m]
נ	<i>nun</i>	[n]
ס	<i>samech</i>	[s]
ע	<i>ayin</i>	[ʕ]
פ	<i>fe</i>	[f]
צ	<i>tsadi</i>	[ts]
ק	<i>qof</i>	[k]
ר	<i>resh</i>	[ʁ]
ש	<i>shin</i>	[ʃ]
ת	<i>thav</i>	[θ]

Table 2. The Greek alphabet. Each glyph is shown in uppercase and lowercase along with its common English name and modern IPA pronunciation.

Table 3. The Hebrew abjad. Only non-final variants of each glyph are shown.

Logic

0

Language

*"No language is justly studied merely as an aid to other purposes.
It will in fact better serve other purposes, philological or
historical, when it is studied for love, for itself."*

—J. R. R. Tolkien

The purpose of these notes is to uh...

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0.1 A Brief History of...

It is worth reflecting on \forall and \exists and ω and \aleph .

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0.2 Syntax and Semantics

A language provides a way of encoding ideas as sequences of symbols.¹

0.3 A Recurring Theme

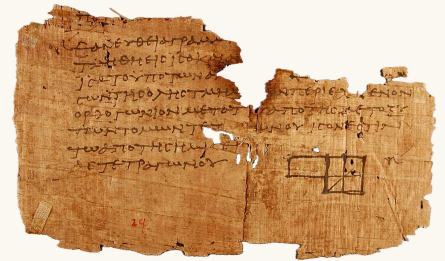


Figure 1. A fragment taken from the *Oxyrhynchus papyri*, dated to approximately 100 AD, depicting proposition 5 from book 2 of Euclid's *Elements*.

¹For our purposes, we will restrict our attention only to *written* languages.

1

Zeroth-Order Logic

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1.1 Truth Values

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1.2 Logical Connectives

Negations

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Conjunctions and Disjunctions

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*Conditional Statements***1.3** *A Modest Proposal**Equivalence**Axioms**Rules of Inference**Classical Syllogisms*

2

First-Order Logic

2.1 A More Expressive Language

Forming Formulæ Well

2.2 Rules of Inference

3

Rhetoric

3.1 *Quantified Expressions*

3.2 *Conditional Expressions*

3.3 *Conjunctions*

3.4 *Disjunctions*

3.5 *Nonconstructive Arguments*

Mathematics

4

Foundations

“Finally I am becoming stupider no more.”

—Paul Erdős

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4.1 Sets

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Existence

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Extensionality

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Pairing

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Separation

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Union

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Power

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Regularity

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Replacement (?)

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Choice (?)

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5

Arithmetic

6

Ancient Number Theory

7

Combinatorics

8

Asymptotics

9

Infinity

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$\rightarrow \emptyset$

$$\emptyset \sim \emptyset \sim \emptyset$$

$$(0 \leqslant 1) \wedge \left(2^{2^{2^2}} \geqslant \sqrt{\sqrt{2}} \right)$$

$$p \leftrightarrow q$$

$$\varphi \Rightarrow \psi$$

$$\varphi \Leftrightarrow \psi$$

$$\Gamma \equiv (\xi \wedge \chi) \vee \rho$$

$$M \models \mathbb{R}$$

$$p_0, \dots, p_{n-1} \vdash q_0, \dots, q_{m-1}$$

$$p_0, \dots, p_{n-1} \vdash q_0, \dots, q_{m-1}$$

$$(\forall n \in \mathbb{N})(s(n) > n)$$

$$(\forall n)(s(n) > n)$$