

DANIEL GONZALEZ CEDRE

DISCRETE MATHEMATICS

| 8TH OF JANUARY, 2026

These notes are intended for cs173 at the University of Illinois Urbana-Champaign.
Copyright © 2026 Daniel Gonzalez Cedre

Contents

<i>Logic</i>	3
<i>0 Language</i>	4
<i>0.1 A Brief History of...</i>	4
<i>0.2 Syntax and Semantics</i>	4
<i>0.3 A Recurring Theme</i>	4
<i>1 Zeroth-Order Logic</i>	5
<i>1.1 Truth Values</i>	5
<i>1.2 Logical Connectives</i>	5
<i>Negations</i>	5
<i>Conjunctions and Disjunctions</i>	5
<i>Conditional Statements</i>	6
<i>1.3 A Modest Proposal</i>	6
<i>Equivalence</i>	6
<i>Axioms</i>	6
<i>Rules of Inference</i>	6
<i>Classical Syllogisms</i>	6
<i>2 First-Order Logic</i>	7
<i>2.1 A More Expressive Language</i>	7
<i>Forming Formulæ Well</i>	7
<i>2.2 Rules of Inference</i>	7
<i>3 Rhetoric</i>	8
<i>3.1 Quantified Expressions</i>	8
<i>3.2 Conditional Expressions</i>	8
<i>3.3 Conjunctions</i>	8

3.4	<i>Disjunctions</i>	8
3.5	<i>Nonconstructive Arguments</i>	8
	<i>Mathematics</i>	9
4	<i>Foundations</i>	10
4.1	<i>Sets</i>	10
	<i>Existence</i>	10
	<i>Extensionality</i>	10
	<i>Pairing</i>	11
	<i>Separation</i>	11
	<i>Union</i>	11
	<i>Power</i>	11
	<i>Regularity</i>	11
	<i>Replacement (?)</i>	11
	<i>Choice (?)</i>	11
5	<i>Arithmetic</i>	12
6	<i>Ancient Number Theory</i>	13
7	<i>Combinatorics</i>	14
8	<i>Asymptotics</i>	15
9	<i>Infinity</i>	16

Notation

SYNTAX	PRONUNCIATION	TYPE	SEMANTICS
$x := y$	"Define x to be y ."	def.	Lets the label " x " refer to y .
$x = y$	" x equals y ."	pred.	Asserts x and y are identical.
$x \in y$	" x is an element of y ."	pred.	Asserts that x is inside of y .
$p \vdash q$	" p proves q ."	judg.	We can syntactically derive q from the assumption p .
$p \Rightarrow q$	" p implies q ."	judg.	
$p \equiv q$	" p is equivalent to q ."	judg.	The statements p and q are semantically interchangeable.
$p \Leftrightarrow q$	" p if and only if q ."	judg.	
\top	"True."	prop.	Asserts truth.
\perp	"False."	prop.	Asserts falsity.
$\neg p$	"Not p ."	prop.	Asserts the negation of p .
$p \wedge q$	" p and q ."	prop.	Asserts both p and q .
$p \vee q$	" p or q ."	prop.	Asserts at least one of p or q .
$p \rightarrow q$	"If p , then q ."	prop.	Asserts q whenever p is true.
$\forall x(\varphi(x))$	"For all x , x satisfies φ ."	prop.	Asserts φ about every object.
$\exists x(\varphi(x))$	"There exists an x such that x satisfies φ ."	prop.	Asserts φ about some object.
$(\forall x \in y)(\varphi(x))$	"For all x in y , x satisfies φ ."	prop.	Asserts φ about every x in y .
$(\exists x \in y)(\varphi(x))$	"There exists an x in y such that x satisfies φ ."	prop.	Asserts φ about some x in y .
$\{x \mid \varphi(x)\}$	"the collection of all x such that x satisfies φ "	obj.	The collection of all and only the objects with property φ .
$\{x \in y \mid \varphi(x)\}$	"the set of all x in y such that x satisfies φ "	obj.	The set of all and only the elements of y with property φ .
$f : x \rightarrow y$	" f is a function from x to y "	judg.	A function f whose domain is x and codomain is y .
$\mathbb{P}(x)$	"the power set of x "	obj.	The set of all subsets of x .
$\text{suc}(x)$	"the successor of x "	obj.	The "next thing after" x .
$\llbracket n \rrbracket$	"bracket n "	obj.	The set $\{m \in \mathbb{N} \mid m < n\}$ of the first n natural numbers.

Table 1. An overview of some important notation used throughout this text. Most of this notation is standard, but some things (e.g., \mathbb{P} , suc , $\llbracket n \rrbracket$) are not yet standardized.

Short for $\forall x(x \in y \rightarrow \varphi(x))$.

Short for $\exists x(x \in y \wedge \varphi(x))$.

We only refer to a collection as a "set" when we can prove it exists.

Depending on context, the expression $f : x \rightarrow y$ can refer to an *object* (i.e., a *function* named f whose domain is the set x and codomain is the set y), or it can be a *proposition* (i.e., the *assertion* that f is a function whose domain is the set x and codomain is the set y).

Writing Systems

GREEK	NAME	IPA	HEBREW	NAME	IPA
A α	<i>alpha</i>	[a]	א	<i>aleph</i>	[ʔ]
B β	<i>beta</i>	[v]	ב	<i>beth</i>	[v]
Γ γ	<i>gamma</i>	[ɣ]	ג	<i>gimel</i>	[ɣ]
Δ δ	<i>delta</i>	[ð]	ד	<i>daleth</i>	[ð]
E ε	<i>epsilon</i>	[e]	ה	<i>he</i>	[h]
Z ζ	<i>zeta</i>	[z]	ו	<i>waw</i>	[v]
H η	<i>eta</i>	[i]	ז	<i>zayin</i>	[z]
Θ θ	<i>theta</i>	[θ]	ח	<i>cheth</i>	[χ]
I ι	<i>iota</i>	[i]	ט	<i>teth</i>	[t]
K κ	<i>kappa</i>	[k]	י	<i>yod</i>	[j]
Λ λ	<i>lambda</i>	[l]	כ	<i>kaf</i>	[χ]
M μ	<i>mu</i>	[m]	ל	<i>lamed</i>	[l]
N ν	<i>nu</i>	[n]	מ	<i>mem</i>	[m]
Ξ ξ	<i>xi</i>	[ks]	נ	<i>nun</i>	[n]
O ο	<i>omicron</i>	[o]	ס	<i>samech</i>	[s]
Π π	<i>pi</i>	[p]	ע	<i>ayin</i>	[ʕ]
P ρ	<i>rho</i>	[r]	פ	<i>fe</i>	[f]
Σ σ	<i>sigma</i>	[s]	צ	<i>tsadi</i>	[ts]
T τ	<i>tau</i>	[t]	ק	<i>qof</i>	[k]
Υ υ	<i>upsilon</i>	[i]	ר	<i>resh</i>	[ʁ]
Φ φ	<i>phi</i>	[pʰ]	ש	<i>shin</i>	[ʃ]
X χ	<i>chi</i>	[kʰ]	ת	<i>thav</i>	[θ]
Ψ ψ	<i>psi</i>	[ps]			
Ω ω	<i>omega</i>	[ɔ:]			

Table 2. The Greek alphabet. Each glyph is shown in uppercase and lowercase along with its common English name and modern IPA pronunciation.

Table 3. The Hebrew abjad. Only non-final variants of each glyph are shown.

Logic

O

Language

*"No language is justly studied merely as an aid to other purposes.
It will in fact better serve other purposes, philological or
historical, when it is studied for love, for itself."*

—J. R. R. Tolkien

The purpose of these notes is to uh...

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequa doleamus.

0.1 A Brief History of...

It is worth reflecting on \forall and \exists and ω and \aleph .

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequa doleamus animo, cum corpore dolemus, fieri tamen permagna accessio potest, si aliquod aeternum et infinitum impendere malum nobis opinemur. Quod idem licet transferre in voluptatem, ut.

0.2 Syntax and Semantics

A language provides a way of encoding ideas as sequences of symbols.¹

0.3 A Recurring Theme



Figure 1. A fragment taken from the *Oxyrhynchus papyri*, dated to approximately 100 AD, depicting proposition 5 from book 2 of Euclid's *Elements*.

¹For our purposes, we will restrict our attention only to *written* languages.

1

Zeroth-Order Logic

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magnam aliquam quaerat voluptatem.

1.1 *Truth Values*

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magnam aliquam quaerat voluptatem. Ut enim aequo doleamus animo, cum corpore dolemus, fieri tamen permagna accessio potest, si aliquod aeternum et infinitum impendere malum nobis opinemur. Quod idem licet transferre in voluptatem, ut postea.

1.2 *Logical Connectives*

Negations

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magnam aliquam quaerat voluptatem.

Conjunctions and Disjunctions

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magnam aliquam quaerat voluptatem.

Conditional Statements

1.3 A Modest Proposal

Equivalence

Axioms

Rules of Inference

Classical Syllogisms

2

First-Order Logic

2.1 A More Expressive Language

Forming Formulae Well

2.2 Rules of Inference

3

Rhetoric

3.1 Quantified Expressions

3.2 Conditional Expressions

3.3 Conjunctions

3.4 Disjunctions

3.5 Nonconstructive Arguments

Mathematics

4

Foundations

“Finally I am becoming stupider no more.”

—Paul Erdős

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequo doleamus.

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequo doleamus.

4.1 Sets

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequo doleamus.

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore et dolore magna aliquam quaerat voluptatem. Ut enim aequo doleamus.

Existence

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Extensionality

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Pairing

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Separation

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Union

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Power

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Regularity

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Replacement (?)

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

Choice (?)

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

5

Arithmetic

6

Ancient Number Theory

7

Combinatorics

8

Asymptotics

9

Infinity

Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore. Lorem ipsum dolor sit amet, consectetur adipiscing elit, sed do eiusmod tempor incididunt ut labore.

$\rightarrow \emptyset$

$$\emptyset \sim \emptyset \sim \emptyset$$

$$(0 \leq 1) \wedge \left(2^{2^{2^2}} \geq \sqrt{\sqrt{2}} \right)$$

$$p \leftrightarrow q$$

$$\varphi \Rightarrow \psi$$

$$\varphi \Leftrightarrow \psi$$

$$\Gamma\equiv(\xi\wedge\chi)\vee\rho$$

$$M\models\mathbb{R}$$

$$p_0, \dots, p_{n-1} \vdash q_0, \dots, q_{m-1}$$

$$p_0, \dots, p_{n-1} \vdash q_0, \dots, q_{m-1}$$

$$(\forall n \in \mathbb{N})(s(n) > n)$$

$$(\forall n)(s(n) > n)$$