

Problem Set 2

Discrete Mathematics

Due on the 4th of February, 2024

We say that a propositional formula is a *tautology* if it is logically equivalent to \top under any assignment of truth values to its variables.

- (5 pts) 1. Consider the following proof of $p \rightarrow (q \rightarrow r) \equiv (p \rightarrow q) \rightarrow r$.

Proof. Observe the following chain of reasoning.

$$\begin{aligned}
 p \rightarrow (q \rightarrow r) &\equiv p \vee \neg(q \rightarrow r) && \text{by conditional disintegration} \\
 &\equiv p \vee \neg(q \vee \neg r) && \text{by conditional disintegration} \\
 &\equiv p \vee \neg q \vee \neg r && \text{by associativity} \\
 &\equiv (p \vee \neg q) \vee \neg r && \text{by associativity} \\
 &\equiv (p \rightarrow q) \vee \neg r && \text{by conditional disintegration} \\
 &\equiv (p \rightarrow q) \rightarrow r && \text{by conditional disintegration}
 \end{aligned}$$

Therefore, $p \rightarrow (q \rightarrow r) \equiv (p \rightarrow q) \rightarrow r$.

Q.E.D.

Find all of the mistakes, if any, in this proof, and *explain why*.

- (40 pts) 2. Prove the claims below *without truth tables* for all propositions p, q, r .

- $p \rightarrow q \equiv \neg q \rightarrow \neg p$.
- $(p \wedge (p \rightarrow q)) \rightarrow q$ is a tautology.
- $(\neg q \wedge (p \rightarrow q)) \rightarrow \neg p$ is a tautology.
- $(p \rightarrow q) \rightarrow ((p \rightarrow \neg q) \rightarrow \neg p)$ is a tautology.

- (40 pts) 3. In this problem, we will progressively establish that the alternative axioms Hilbert proposed are all tautologies *without truth tables*. Here, the variables p, q , and r all represent arbitrary propositions.

- Show $p \rightarrow p$ is a tautology.
- Show $(p \rightarrow q) \rightarrow (\neg q \rightarrow \neg p)$ is a tautology.
- Show $p \rightarrow (q \rightarrow p)$ is a tautology.
- Show $(p \rightarrow (q \rightarrow r)) \rightarrow ((p \rightarrow q) \rightarrow (p \rightarrow r))$ is a tautology.

- (10 pts) 4. Show that \neg and \wedge are sufficient to express *any* proposition.

- (5 pts) 5. Is there a *single connective* capable of expressing *any* proposition?¹ Justify your answer with a proof.

You may rely on the following theorems throughout this problem set in addition to the axioms of classical logic.

- Uniqueness of Negations
- $\neg \top \equiv \perp$ and $\neg \perp \equiv \top$
- Double Negation
- Idempotency
- Domination
- De Morgan's Laws

¹ This does not necessarily have to be one of $\neg, \wedge, \vee, \rightarrow$, nor \leftrightarrow . You can define new logical connectives using truth tables.