Discrete Mathematics

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Chapter 7

Number Theory

7.1 Ancient Greece

Definition 7.1 ***(Divisibility).***
Given two integers $a, b \in \mathbb{Z}$, we say $a \mid b :\Leftrightarrow (\exists k \in \mathbb{Z})(ak = b)$. We read $a \mid b$ as a *divides* b, meaning ${}^{b}/_{a} \in \mathbb{Z}$. ⌟

Lemma 7.1 ***(Initial object).***
If $x \in \mathbb{Z}$, then $1 \mid x$. ⌟

Proof. Let $x \in \mathbb{Z}$ and observe that $1 \cdot x = x$. Therefore, $1 \mid x$ by definition. Q.E.D.

Lemma 7.2 ***(Terminal object).***
If $x \in \mathbb{Z}$, then $x \mid 0$. ⌟

Proof. Let $x \in \mathbb{Z}$ and observe that $0 \cdot x = 0$. Therefore, $x \mid 0$ by definition. Q.E.D.

Lemma 7.3 ***(Divisibility is a Partial Order).***
The following statements hold for all $a, b, c \in \mathbb{Z}$:

I. $a \mid a$

Proof. Let $a \in \mathbb{Z}$ and observe that $1 \cdot a = a$. Therefore, $a \mid a$ by definition. Q.E.D.

II. $\big((a \mid b) \wedge (b \mid a)\big) \;\Rightarrow\; |a| = |b|$

Proof. Let $a, b \in Z$ and suppose $a \mid b$ and $b \mid a$. Then, there exist $k_1, k_2 \in \mathbb{Z}$ such that $ak_1 = b$ and $bk_2 = a$ by definition. But then $bk_2 = (ak_1)k_2 = a$, so $ak_1k_2 = a$, yielding $k_1k_2 = 1$. Since the only integers with multiplicative inverses are 1 and -1, we have $\{k_1, k_2\} \subseteq \{1, -1\}$, so $a = b$ or $a = -b$. Thus, $|a| = |b|$. Q.E.D.

III. $\big((a \mid b) \wedge (b \mid c)\big) \;\Rightarrow\; a \mid c$

Proof. Let $a, b, c \in Z$ and suppose $a \mid b$ and $b \mid c$. Then, there exist $k_1, k_2 \in \mathbb{Z}$ such that $ak_1 = b$ and $bk_2 = c$. This yields $ak_1k_2 = c$. Since $k_1, k_2 \in \mathbb{Z}$, we observe $k_1k_2 \in \mathbb{Z}$ and conclude $a \mid c$ by definition. Q.E.D.

⌟

Lemma 7.4 ***(Useful facts).***
The following statements hold for all $a, b, c \in \mathbb{Z}$:

I. $\big((a \mid b) \wedge (a \mid c)\big) \;\Rightarrow\; a \mid b + c$

II. $a \mid b \;\Rightarrow\; (\forall \ell \in \mathbb{Z})\big(a \mid b\ell\big)$

III. $\big((a \mid b) \wedge (b \neq 0)\big) \;\Rightarrow\; |a| \leqslant |b|$

The proofs of the above lemmata are left as exercises to the reader. ⌟

Corollary 7.1.
Given $a, b, c \in \mathbb{Z}$, if $a \mid b$ and $a \mid c$, then $(\forall \ell_1, \ell_2 \in \mathbb{Z})\big(a \mid \ell_1 b + \ell_2 c\big)$.

Definition 7.2 ***(Primality).***
We say that a natural number $p \in \mathbb{N}$ is *prime* $:\Leftrightarrow$ $(p > 1)$ and $(\forall n \in \mathbb{N})\big(n \mid p \;\Rightarrow\; n \in \{1, p\}\big)$.
We say $n \in \mathbb{N}$ is *composite* $:\Leftrightarrow$ n is not prime.

Lemma 7.5 ***(Fundamental Lemma of Arithmetic).***
If $n \in \mathbb{N}$ and $n > 1$, then $(\exists p \in \mathbb{N})\big(p \text{ is prime} \wedge p \mid n\big)$.

Proof. TODO Q.E.D.

Theorem 7.1 ***(Fundamental Theorem of Arithmetic).***
Every natural number greater than 1 has a *unique* prime factorization. Formally, for every natural number $n \in \mathbb{N}_{\geqslant 2}$ greater than 1, there exist *unique, distinct* primes $p_1, \ldots p_\ell \in \mathbb{N}_+$ with *unique* exponents $k_1, \ldots k_\ell \in \mathbb{N}_+$ such that

I. $\big(\forall i, j \in \{1, \ldots \ell\}\big)\big(i \neq j \;\Rightarrow\; p_i \neq p_j\big)$
II. $\big(\forall i \in \{1, \ldots \ell\}\big)(p_i \text{ is prime})$
III. $n = p_1^{k_1} p_2^{k_2} \ldots p_\ell^{k_\ell}$.

Theorem 7.2 ***(Euclid's Theorem).***
There are infinitely-many prime numbers.

Proof. TODO Q.E.D.

Definition 7.3 ***(Greatest Common Divisor).***
Given two integers $a, b \in \mathbb{Z}$, we say that $g \in \mathbb{Z}$ is the *greatest common divisor* (*a.k.a. greatest common factor*) of a and b $:\Leftrightarrow$

$$\big(g \mid a\big) \wedge \big(g \mid b\big) \wedge (\forall h \in \mathbb{Z})\Bigg(\Big(\big(h \mid a\big) \wedge \big(h \mid b\big)\Big) \;\Rightarrow\; h \mid g\Bigg).$$

Notice that, since $(\forall x)\big(1 \mid x\big)$, every pair of integers shares a common factor. Since common factors of a and b are bounded above by $\min\{a, b\}$, that means the set of all common factors of a and b is nonempty and bounded above, so it has a maximal element. Therefore, the greatest common divisor of any two integers always exists.

Definition 7.4 ***(Co-Primality).***
We say that two integers $a, b \in \mathbb{Z}$ are *co-prime* $:\Leftrightarrow$ their greatest common divisor is 1.

Theorem 7.3 ***(Euclid's Division Theorem).***
If $a, b \in \mathbb{Z}$, then there exist two *unique* integers $q, r \in \mathbb{Z}$ such that

$$a = bq + r \text{ and } 0 \leqslant r < b.$$

Here, q is called the *quotient* when a is divided by b, and r is the *remainder*, as illustrated by ${}^{a}/_{b} = q + {}^{r}/_{b}$.

Algorithm 7.1 ***(Euclid's Division Algorithm).***
We can find the greatest common divisor of two integers by recursively computing

$$\gcd(a, b) := \begin{cases} a & \text{if } b = 0 \\ \gcd(b, r) \begin{cases} \text{where} & a = bq + r \\ \text{and} & 0 \leqslant r < b \\ \text{and} & q, r \in \mathbb{Z}. \end{cases} & \text{if } b \neq 0 \end{cases}$$

This algorithm correctly computes the greatest common divisor of two arbitrary integers.

7.2 Modular Arithmetic

Definition 7.5 ***(Modular Congruence).***
Let $m \in \mathbb{N}_+$ and let $x, y \in \mathbb{Z}$. We say that $x \equiv y \pmod{m} \;:\Leftrightarrow\; m \mid x - y$. We read the sentence $x \equiv y \pmod{m}$ in English as "x is congruent to y modulo m." This expresses the idea that x and y have the *same remainder* after division by m, as we can see below.

$$\left.\begin{aligned} x &= q_x m + r \\ y &= q_y m + r \end{aligned}\right\} \begin{aligned} &\Leftrightarrow\; x - y = (q_x m + r) - \big(q_y m + r\big) \\ &\Leftrightarrow\; x - y = \big(q_x - q_y\big)m + (r - r) \\ &\Leftrightarrow\; x - y = \big(q_x - q_y\big)m \\ &\Leftrightarrow\; m \mid x - y \end{aligned}$$

⌟

Exercise 7.1.
Let $m \in \mathbb{N}_+$ and $w, x, y, z \in \mathbb{Z}$. The following are some useful facts about modular congruence.

I. $x \equiv y \pmod{m} \;\Rightarrow\; x + z \equiv y + z \pmod{m}$.

II. $\Big(\big(w \equiv z \pmod{m}\big) \wedge \big(x \equiv y \pmod{m}\big)\Big) \;\Rightarrow\; wx \equiv yz \pmod{m}$.

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Theorem 7.4 ***(Modular Congruence is an Equivalence Relation).***
Let $m \in \mathbb{N}_+$ and $x, y, z \in \mathbb{Z}$. The following are true.

I. $x \equiv x \pmod{m}$

II. $x \equiv y \pmod{m} \;\Rightarrow\; y \equiv x \pmod{m}$

III. $\Big(\big(x \equiv y \pmod{m}\big) \wedge \big(y \equiv z \pmod{m}\big)\Big) \;\Rightarrow\; x \equiv z \pmod{m}$

⌟

Definition 7.6 ***(Modular Residue Classes).***
Let $m \in \mathbb{N}_+$ and let $a \in \mathbb{Z}$. The set of solutions to the *linear congruence* $x \equiv a \pmod{m}$ is denoted by

$$[a]_m \coloneqq \big\{x \in \mathbb{Z} \;\big|\; x \equiv a \pmod{m}\big\}.$$

Each of these is known as an *equivalence class* of *residues* modulo m, indicating that all the integers in that class have remainder congruent to a after division by m.
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Definition 7.7 ***(Modular Rings).***
Let $m \in \mathbb{N}_+$. We define the *modular ring* of size m (*a.k.a.* the *cyclic group*) of size m by

$$\mathbb{Z}/m\mathbb{Z} \coloneqq \big\{[x]_m \;\big|\; x \in \mathbb{Z}\big\}$$

and we define *modular addition* and *modular multiplication* on its elements by

$$\begin{aligned} [x]_m + [y]_m &\coloneqq [x + y]_m \\ [x]_m \cdot [y]_m &\coloneqq [xy]_m. \end{aligned}$$

⌟

Theorem 7.5 ***(Bézout's Identity).***
Given $x, y \in \mathbb{Z}$, there exist $k_1, k_2 \in \mathbb{Z}$ such that

$$xk_1 + yk_2 = \gcd(x, y).$$

⌟

Algorithm 7.2 ***(Extended Euclidean Division Algorithm).***
We can find the greatest common divisor *and* the Bézout coefficients of two integers by recursively computing

$$\gcd(a,b) := \begin{cases} (a,1,0) & \text{if } b = 0 \\ (d,t,s-qt) \begin{cases} \text{where} & (d,s,t) = \operatorname{egcd}(b,r) \\ \text{and} & a = bq + r \\ \text{and} & 0 \leqslant r < b \\ \text{and} & q,r \in \mathbb{Z}. \end{cases} & \text{if } b \neq 0 \end{cases}$$

This algorithm correctly computes the greatest common divisor of two arbitrary integers.

Definition 7.8 ***(Euler's Totient Function).***
We define *Euler's totient function* $\varphi : \mathbb{N} \to \mathbb{N}$ by the number of integers $1 \leqslant z < n$ relatively prime with $n \in \mathbb{N}$

$$\varphi(n) := \left|\left\{ z \in \mathbb{Z} \;\middle|\; (1 \leqslant z < n) \wedge \big(\gcd(z,n) > 1\big) \right\}\right|$$

Lemma 7.6.
If $p \in \mathbb{N}_+$ is prime, then $\varphi(p) = p - 1$.

Theorem 7.6.
Let $x, y \in \mathbb{Z}$. If $\gcd(x,y) = 1$, then $\varphi(xy) = \varphi(x)\varphi(y)$.

Theorem 7.7 ***(Férmat's Little Theorem).***
Let $p \in \mathbb{N}_+$ be prime and $a \in \mathbb{Z}$. Then, $a^p \equiv p \pmod{p}$. Further, if $\gcd(a,p) = 1$, then $a^{p-1} \equiv 1 \pmod{p}$.

Theorem 7.8 ***(Euler's Theorem).***
Let $n \in \mathbb{N}_+$ and $a \in \mathbb{Z}$. If $\gcd(a,n) = 1$, then

$$a^{\varphi(n)} \equiv 1 \pmod{n}.$$

Algorithm 7.3 ***(RSA Encryption).***
The RSA* cryptosystem is an algorithm for performing *asymmetric* (*a.k.a. public key*) encryption. Its security is reliant on two key observations:

I. It takes roughly $e^{\left(\sqrt[3]{64/9}\right)(\ln n)^{1/3}(\ln \ln n)^{2/3}}$ time to factor $n \in \mathbb{N}$ into a product of primes.

II. There is no known way of finding the k^{th} root of x in $\mathbb{Z}/n\mathbb{Z}$ faster than by factoring n.

The algorithm has two stages, with **private** information that must be kept **secret** or **destroyed**, and **public** information that is shared through **insecure channels**.

Key Generation

1. Pick two (large) prime numbers $\boldsymbol{p}$ and $\boldsymbol{q}$. Compute $\boldsymbol{n} := pq$.
2. Compute $\boldsymbol{\varphi(n)} = \varphi(pq) = \varphi(p)\varphi(q) = (p-1)(q-1)$ using Theorem 7.6.
3. Pick a random number $\boldsymbol{e}$ in the range $1 < e < \varphi(n)$ relatively prime with $\varphi(n)$.
4. When checking that $\gcd\big(e, \varphi(n)\big) = 1$ in the previous step, use the *Extended Euclidean Algorithm* from Algorithm 7.2 to simultaneously obtain $\boldsymbol{d}$ satisfying $ed \equiv 1 \pmod{\varphi(n)}$.
5. Publish (e,n) publicly while keeping (d,n) secret.

The **public encryption key** is the pair $(\boldsymbol{e}, \boldsymbol{n})$, and the **private decryption key** is the pair $(\boldsymbol{d}, \boldsymbol{n})$. All other private information should be immediately *destroyed* for security.

*Named after Rivest, Shamir, and Adleman, the three coauthors of the original 1977 paper.

Message Passing: Encryption

1. Your friend takes a message m, which is a (binary) number, and—treating it like a string—chops it up into substrings $m = m_0 m_1 \ldots m_k$ so that each is in the range $0 < m_i < n$ and $\gcd(m_i, n) = 1$.
2. For each sub-message m_i, your friend computes the encrypted sub-message c_i by $c_i \equiv m_i^e \pmod{n}$, where $0 < c_i < n$.
3. He then sends c_i over an insecure channel to you.

Message Passing: Decryption

1. Receive c_i, which has possibly been intercepted by other parties.
2. Decrypt the encrypted message with your private key by computing $m_i \equiv c_i^d \pmod{n}$ such that $0 < m_i < n$.

Appendix A

The Algebra of Modular Arithmetic

Definition A.1 ***(Some Basic Algebra).***
Suppose we have a set G with a binary operation on $\boldsymbol{+} : G \times G \to G$ defined on it. We say this is a *monoid* if there exists an *identity element* e_0 such that

i. $(\forall g \in G)(e_0 \boldsymbol{+} g = g \boldsymbol{+} e_0 = g)$
ii. $(\forall g, h, k \in G)\big(g \boldsymbol{+} (h \boldsymbol{+} k) = (g \boldsymbol{+} h) \boldsymbol{+} k\big)$

We call G under $\boldsymbol{+}$ a *group* if we also have

iii. $(\forall g \in G)(\exists h \in G)(g \boldsymbol{+} h = e_0)$

We call G a *commutative* group* if we additionally have

iv. $(\forall g, h \in G)(g \boldsymbol{+} h = h \boldsymbol{+} g)$

If we then define another binary operation $\bullet : G \times G \to G$, then we call G with these two operations a *ring* if we can find another identity element $e_1 \in G$ such that

v. G is a monoid under $\bullet$ with identity e_1
vi. $(\forall g, h, k \in G)\big(g \bullet (h \boldsymbol{+} k) = (g \bullet h) \boldsymbol{+} (g \bullet k)\big)$
vii. $(\forall g, h, k \in G)\big((g \boldsymbol{+} h) \bullet k = (g \bullet k) \boldsymbol{+} (h \bullet k)\big)$

Finally, we say that G is a *field* if we also have

viii. $(\forall g \in G)\big(g \neq e_0 \;\Rightarrow\; (\exists h \in G)(g \bullet h = e_1)\big)$

⌟

Lemma A.1.
G with $\boldsymbol{+}$ and $\bullet$ is a field *iff* G with $\boldsymbol{+}$ is a group and $G \setminus \{e_0\}$ with $\bullet$ is a group. ⌟

Theorem A.1.
If $n \in \mathbb{N}_+$, then ${}^{\mathbb{Z}}\!/_{n\mathbb{Z}}$ forms a ring under modular arithmetic.
If $n \in \mathbb{N}_+$, then ${}^{\mathbb{Z}}\!/_{n\mathbb{Z}}$ forms a field under modular arithmetic *iff* n is prime. ⌟

Definition A.2 ***(Order).***
The *order* $|G|$ of a group G is its cardinality. The *order* $|g|$ of $g \in G$ is the smallest $n \in \mathbb{N}$ such that $g^{n+1} = e$. ⌟

Theorem A.2.
For any group G and any $g \in G$, we have that $|g|$ divides $|G|$. ⌟

*Usually referred to as an *Abelian group*, after Niels Henrik Abel.