Monotone Catenary Degree in Numerical Monoids

D. Gonzalez, C. Wright, J. Zomback

August 4, 2015



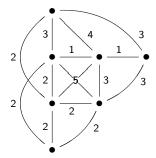
Definitions

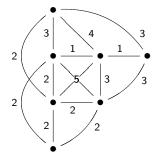
- Definitions
- Arithmetic Monoids

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- Generalized Arithmetic Monoids

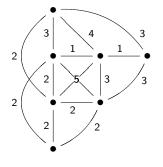
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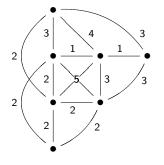




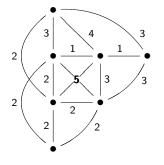
 Remove edges with the largest distances



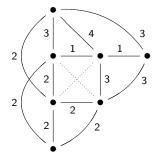
- Remove edges with the largest distances
- Continue until disconnected



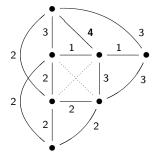
- Remove edges with the largest distances
- Continue until disconnected
- Final distance removed is catenary degree



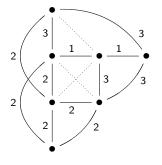
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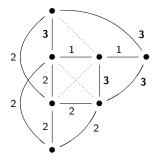
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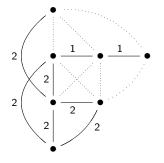
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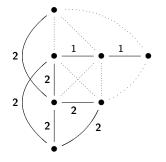
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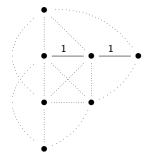
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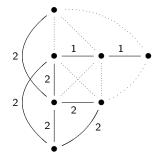
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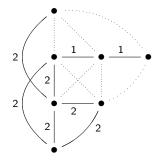
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$$c(m) = 2$$

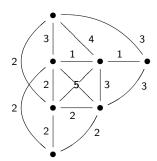
How does it differ?

 Only factorizations of the same length

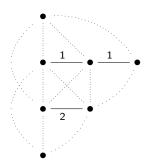
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- Minimum N-chain within a given length

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- Maximum over all such minima

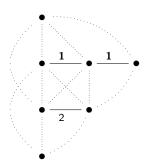
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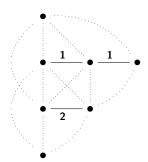
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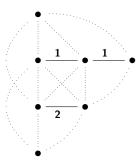
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$$c_{eq}(m) = \max\{1, 2\} = 2$$

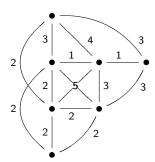
How does it differ?

 Only factorizations of adjacent lengths

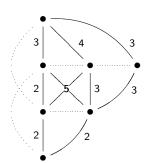
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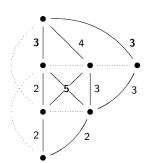
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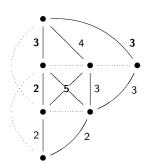
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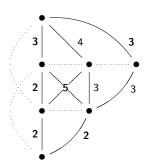


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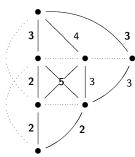
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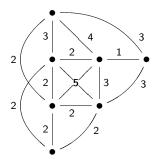
$$c_{adj}(m) = \max\{3, 2, 2\} = 3$$

How does it differ?

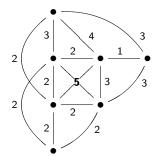
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- Only monotone chains between elements

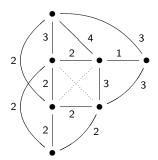
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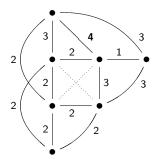
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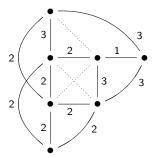
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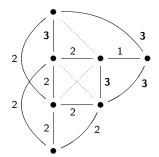
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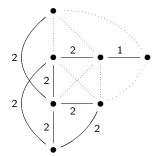
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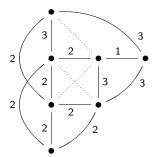
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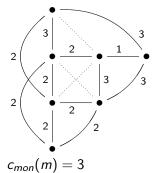
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We seek to investigate the inequality between monotone and regular catenary degree.

Consider a monoid M generated by a finite arithmetic sequence $a, a + d, \ldots, a + kd$, so $M = \langle a, a + d, \ldots, a + kd \rangle$. We will refer to a monoid of this form as an *arithmetic monoid*.

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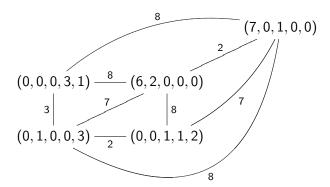
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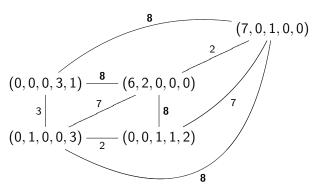
⟨5, 7, 9⟩

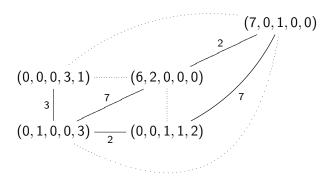
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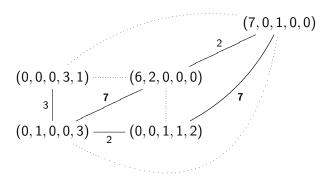
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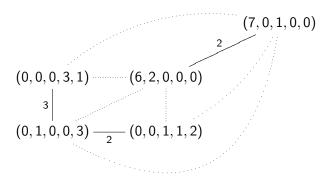
- ⟨5, 7, 9⟩
- $\langle 11, 15, 19, 23 \rangle$

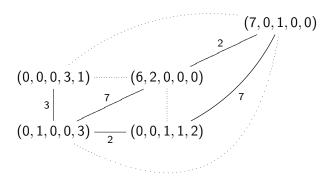


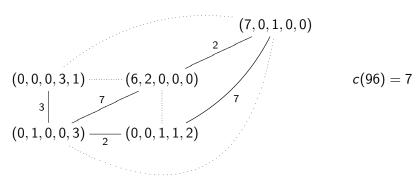












$$I = 8: (6, 2, 0, 0, 0) \xrightarrow{2} (7, 0, 1, 0, 0)$$

$$\begin{vmatrix} 7 & & | 7 \\ & & | 7 \end{vmatrix}$$

$$I = 4: (0, 0, 0, 3, 1) \xrightarrow{3} (0, 1, 0, 0, 3) \xrightarrow{2} (0, 0, 0, 1, 2)$$

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$$c_{mon}(96) = c(96) = 7.$$

Main Result for Arithmetic Sequences

Theorem

Given an arithmetic monoid $M = \langle a, a+d, \ldots, a+kd \rangle$, for all elements $m \in M$,

$$c_{mon}(m) = c(m).$$

As a result of this, $c_{mon}(M) = c(M)$.

Consider $M = \langle a, ah + d, \dots, ah + kd \rangle$. We will refer to a monoid of this form as a *generalized arithmetic monoid*. We will discuss generalized arithmetic monoids in embedding dimension three: $M = \langle a, ah + d, ah + 2d \rangle$.

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Examples:

• (3, 11, 16)

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$$\langle 3, 11, 16 \rangle$$

 $a = 3, h = 2, d = 5$

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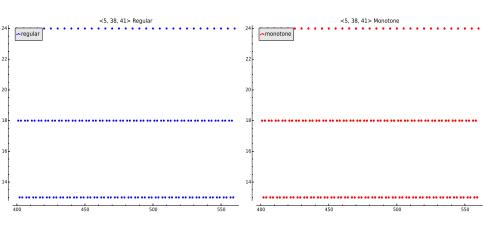
 $M = \langle a, ah + d, ah + 2d \rangle.$

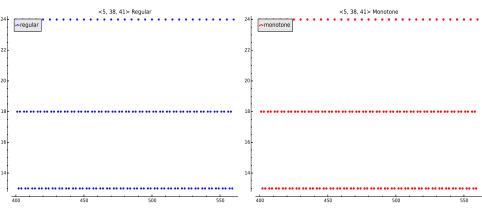
- $\langle 3, 11, 16 \rangle$ a = 3, h = 2, d = 5
- ⟨5, 38, 41⟩

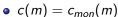
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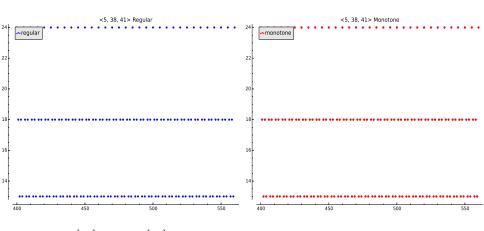
- $\langle 3, 11, 16 \rangle$ a = 3, h = 2, d = 5
- $\langle 5, 38, 41 \rangle$ a = 5, h = 7, d = 3





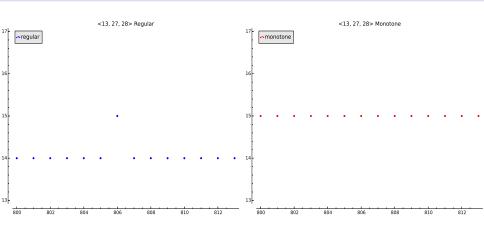




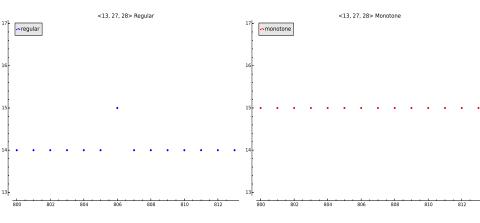


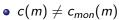
- $c(m) = c_{mon}(m)$
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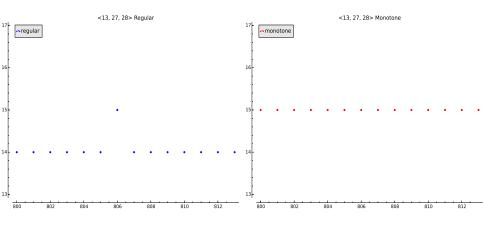






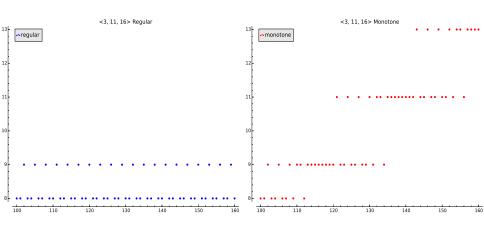


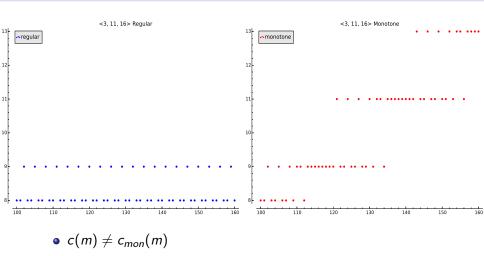




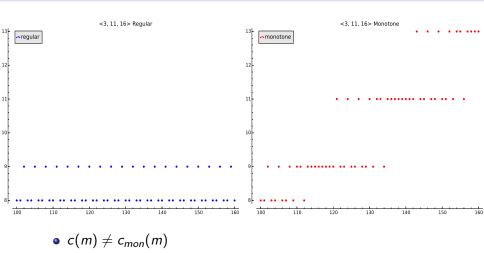
- $c(m) \neq c_{mon}(m)$
- $c(M) = c_{mon}(M)$







• $c(M) < c_{mon}(M)$





Conjecture

If
$$gcd(h-1, d) > 1$$
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Case 1: h < d $c(M) < c_{mon}(M)$

Case 2: $h \ge d$ and $c(M) < c_{eq}(M)$ $c(M) < c_{mon}(M)$

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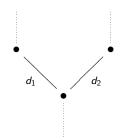
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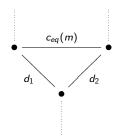
Case 3: $h \ge d$ and $c(M) = c_{eq}(M)$ $c(M) = c_{mon}(M)$

• Recall that $c_{mon}(m) = \max\{c_{eq}(m), c_{adj}(m)\}.$

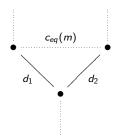




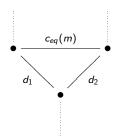
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- Recall that $c_{mon}(m) = \max\{c_{eq}(m), c_{adj}(m)\}.$
- $c_{eq}(M) = \frac{ah+2d-a}{\gcd(h-1,d)}$
- if gcd(h-1,d) = 1, then $c_{eq}(m) \ge c_{adj}(m)$
- if gcd(h-1,d) > 1, then $c_{eq}(m) < c_{adj}(m)$

When is $c_{mon}(M) > c(M)$?

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If we have

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When is $c_{mon}(M) > c(M)$?

If we have

- $c_{eq}(M) > c_{adj}(M)$
- Whenever $|z_1|=|z_2|=I$ we can find $|z_3|< I$ such that $d(z_1,z_3)< c_{eq}(M)$ and $d(z_2,z_3)< c_{eq}(M)$

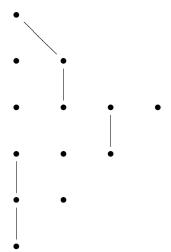
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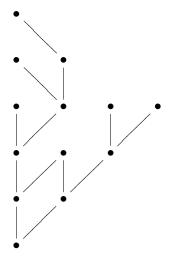
Then $c_{mon}(M) > c(M)$.

Is $c_{mon}(M) > c(M)$?



- $c_{mon}(M) = c_{eq}(M)$.
- Under our restrictions, $c_{eq}(M) > c_{adj}(M)$

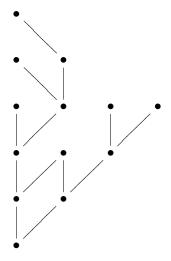
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$$c_{mon}(M) = c_{eq}(M)$$

- Under our restrictions, $c_{eq}(M) > c_{adj}(M)$
- We can move from two factorizations of length I to one of a lower length with $d < c_{eq}(M)$

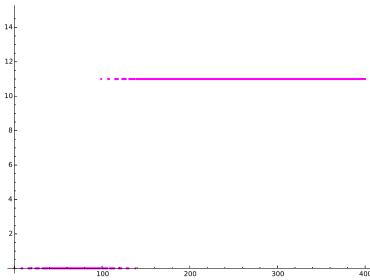
Is $c_{mon}(M) > c(M)$?



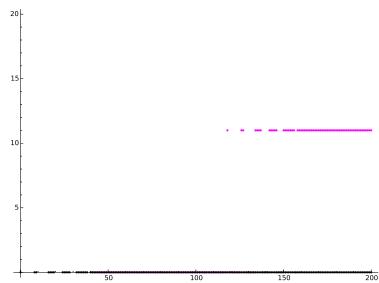
•
$$c_{mon}(M) = c_{ea}(M)$$

- Under our restrictions, $c_{eq}(M) > c_{adj}(M)$
- We can move from two factorizations of length I to one of a lower length with $d < c_{eq}(M)$
- $c_{mon}(M) > c(M)$

Equivalent Catenary Degree in $\langle 8, 9, 19 \rangle$



Equivalent Catenary Degree in (8, 9, 19)



Equivalent Catenary Degree in $\langle n_1, n_2, n_3 \rangle$

Theorem

Let M be a minimally generated numerical monoid $\langle n_1, n_2, n_3 \rangle$.

Then
$$c_{eq}(M) = \frac{n_3 - n_1}{\gcd(n_2 - n_1, n_3 - n_1)}$$

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- If $m = n_2 \left(\frac{n_3 n_1}{\gcd(n_2 n_1, n_3 n_1)} \right) + x$ for some $x \in M$, then $c_{eq}(m) = c_{eq}(M)$.
- If $m = n_2 \left(\frac{n_3 n_1}{\gcd(n_2 n_1, n_3 n_1)} \right) + x$ for some $x \notin M$, then $c_{eq}(m) = 0$.

Equivalent Catenary Degree in $\langle n_1, n_2, n_3 \rangle$

Theorem

For all $m \in M$ such that $m > n_2 \left(\frac{n_3 - n_1}{\gcd(n_2 - n_1, n_3 - n_1)} \right) + \mathcal{F}(M)$, $c_{eq}(m) = c_{eq}(M)$. Furthermore, $c_{eq}(m)$ can take on only two values: 0 and $c_{eq}(M)$.

Moving Between Factorizations

$$M = \langle n_1, n_2, n_3 \rangle$$

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 and let $|z_1|=|z_2|$. Then
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What if we want to move between factorization lengths?

$\mathsf{Theorem}$

Consider a factorization z = (b, c, d) of an element $s \in M$ with |z| = x. We can write any factorization z_i of s where $|z_i| = x - I$ as $z_i = \left(b + \frac{-\ln_1 + k(n_3 - n_1)}{n_2 - n_1} - k - l, c + \frac{\ln_1 - k(n_3 - n_1)}{n_2 - n_1}, d + k\right).$

$$M = \langle na, na + n, 2na + nx + 1 \rangle$$
 with $x \geq 2$

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•
$$c_{eq}(M) = na + nx + 1$$

For certain families of monoids, we can guarantee that our two conditions for $c_{mon}(M) > c(M)$ hold.

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Recall that $c_{mon}(M) = \max\{c_{eq}(M), c_{adj}(M)\}.$

•
$$c_{eq}(M) = na + nx + 1$$

•
$$c_{adi}(M) < na + nx + 1$$

$$\langle na, na + n, 2na + nx + 1 \rangle$$

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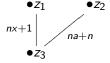
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 with $x \ge 2$

$$\begin{array}{c|c}
\bullet Z_1 & \bullet Z_2 \\
nx+1 & na+n \\
\bullet Z_3
\end{array}$$

$$\langle na, na + n, 2na + nx + 1 \rangle$$

$$M = \langle na, na + n, 2na + nx + 1 \rangle$$
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Theorem

Let $M = \langle na, na + n, 2na + nx + 1 \rangle$ for $n \in \mathbb{N}$ and $x \geq 2$. Then $c_{mon}(M) > c(M)$.

$$\langle a, a+1, \mathcal{F}\langle a, a+1 \rangle \rangle$$

We have shown that in many cases, $c_{mon}(M) > c(M)$.



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We have shown that in many cases, $c_{mon}(M) > c(M)$. How big can the difference be between $c_{mon}(M)$ and c(M)?

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$$\langle a, a+1, \mathcal{F}\langle a, a+1 \rangle \rangle$$

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$$c_{mon}(M) = a^2 - 2a - 1$$

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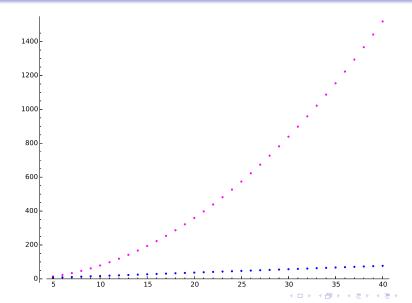
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•
$$c_{mon}(M) = a^2 - 2a - 1$$

•
$$c(M) = 2a - 3$$

•
$$c_{mon}(M) - c(M) = a^2 - 4a - 4$$

$$\langle a, a+1, \mathcal{F}\langle a, a+1 \rangle \rangle$$



$$c_{mon}(M) - c(M)$$

Theorem

The difference between the monotone and regular catenary degrees of a monoid can be arbitrarily large.

Takeaways

• In some monoids, namely those generated by arithmetic sequences, $c_{mon}(M) = c(M)$.

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- In generalized arithmetic monoids, we can have either $c_{mon}(M) > c(M)$ or $c_{mon}(M) = c(M)$.

Takeaways

- In some monoids, namely those generated by arithmetic sequences, $c_{mon}(M) = c(M)$.
- In generalized arithmetic monoids, we can have either $c_{mon}(M) > c(M)$ or $c_{mon}(M) = c(M)$.
- In general, we expect that $c_{mon}(M) > c(M)$. In fact, the difference between the two can grow arbitrarily large.

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