

Cheat-Proof Payout for Centralized and Distributed Online Games

By

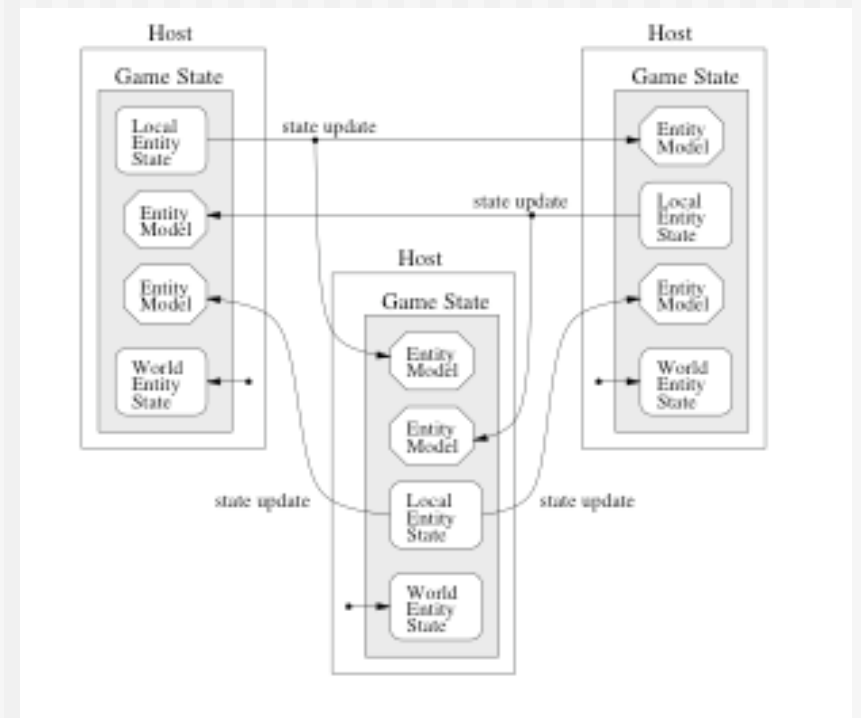
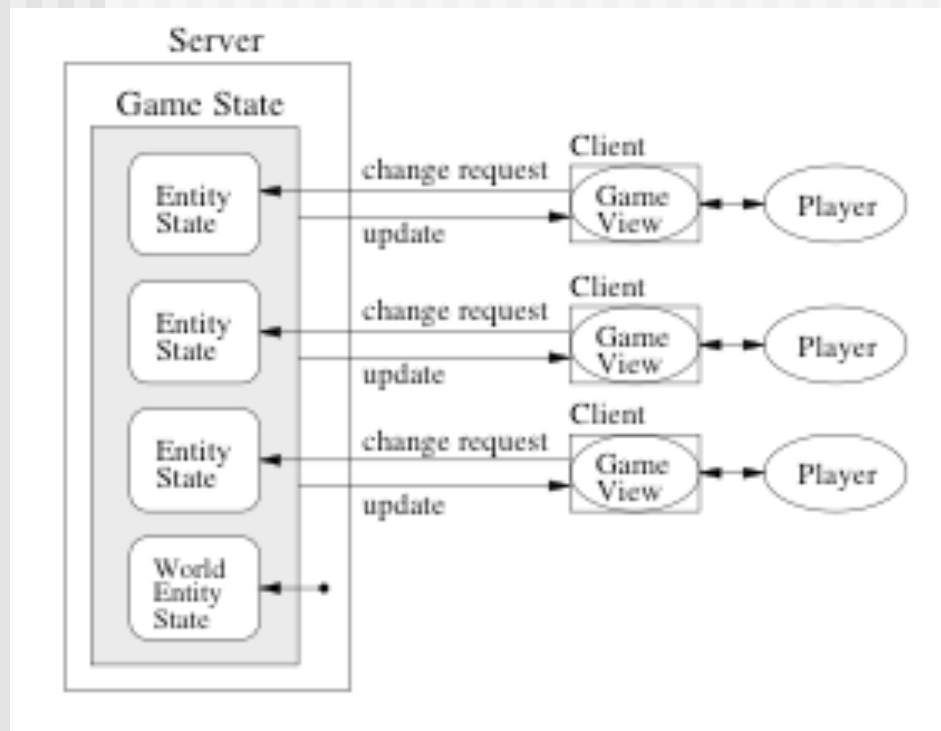
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Game System - Assumptions

- Grant Cheaters:
 - Read, insert, modify, and block network messages
 - Game software is readable by user - including **game state** and **cryptographic keys**
 - Which means security by obscurity won't work.

Game System - Centralized or Distributed



Common Synchronization Techniques Detrimental

- Example 1: Dead Reckoning
 - Potential for cheating - “suppress-correct cheat”
- Example 2: Stop-and-wait Simulations
 - Potential for cheating - “lookahead cheat”

Dead Reckoning- suppress-correct cheat

- Assume n steps in game simulation allowed to be dead-reckoned before connection is lost
- What if we drop $n-1$ packets?
- Then we have an up to date view of the world, but everyone else has an old view of us
- Anyone remember their proposed cheat?

Dead Reckoning- suppress-correct cheat

- Cheat Scenario:
- Player S chases player A
- Player S drops $n-1$ updates
- Meanwhile A uses dead reckoning to guess, but never REALLY knows where S is.
- S sends fabricated n th update that places S right behind A

Dead Reckoning- suppress-correct cheat

- As long as S sends a plausible update every nth step, A cannot differentiate cheating from sluggish link
- “We can conclude that fair play is indistinguishable from cheating when dead reckoning is used.” (p. 4)

Stop-and-wait Simulations - lookahead cheat

- Stop-and-wait simulations like AoE
- Eliminates possibility of “suppress-correct cheat” because there is no dead reckoning
- However stop-and-wait simulations allow for another kind of cheat
- Anyone remember this proposed cheat?

Stop-and-wait Simulations - lookahead cheat

- Wait for everyone to send their updates, then send your update after taking into account their information
- Scenario: B makes a lethal shot at A
- A has a cheating agent that sends the decision to raise shields in time (because it looked-ahead)

Solution Version 1.0: Lockstep Protocol

- Prevents suppress-correct and lookahead cheats
- Anyone remember what the basic idea was?

Solution Version 1.0: Lockstep Protocol

- Basic Idea:
- Turn into 2-step process:
 - First, send hash of your decision
 - Second, once everyone has sent the hash, send the decisions
 - If $\text{hash}(\text{decisions}) == \text{hash_sent}$ then good, else CHEATER!
- We've prevented the proposed cheats, but any Problems?

Solution Version 1.0: Lockstep Protocol

- Main problem is that we've just doubled the communication, and synchronization time since EVERYONE has to send the hash, before EVERYONE sends their decisions.
- This will slow us down - we're already battling time for reasonable user experience.
- Solution: Asynchronous version of Lockstep called "AS".

AS -

Asynchronous Synchronization

- Basic idea is to restrict interaction of each player to an “aura” or SOI - sphere of influence
- If our player can't be interacting with another player, then don't bother synchronizing game play with them
- If we can interact (SOIs intersect) then lockstep with them

AS -

Basic Algorithm Outline

- Compute local host's current state
- Send hash of it out
- Process hashes sent by remote hosts
- Foreach remote_host
 - Accept state of remote hosts (up to current time)
 - Compute potential influence (SOI's) intersections
 - IF no intersection of SOIs
 - THEN local host is not waiting for remote_host
 - ELSE IF hash(remote_host) accepted
 - THEN local host is not waiting for remote_host
 - ELSE local host is waiting for remote_host
- IF not waiting for any remote_host send local state

AS

- Prevents cheating same way lockstep does
- Prevents suppress-correct cheat because no dead reckoning
- Prevents lookahead cheat because we get hashes before we get real state (we don't send our state, until we've received hashes from everyone else)

AS -

Secret Possessions

- Can't really give knowledge of a “secret possession” to opposing players (assumption that they can see any information sent them)
- Instead we give them a “promise” - basically a hash or encrypted form of the real thing (if encrypted, the key is not known to opposing players)
- Player then signs the promise and gives it back - to prevent repudiation that the secret was promised.

AS -

Secret Possessions

- Now, a player can have the promise and competing players won't know he has the secret possession (I guess everyone has a "promise" field or something)
- The game can be played
- Afterward, the promise can be verified against the "secret possession" to detect cheating (prevents opponents from cheating, only checks for local player cheating - not preventing it)

AS - MMO's

- AS like algorithm can be used for MMO's
- One cheating problem with MMO's is sending position information, or other information, to everyone - "it might affect player strategy" - this would alert of ambushes or the like.
- We assume the world is divided into regions
- How do we tell if a player is in the same region without exposing my region (which may warn of an ambush)?
- Proposed algorithm can do this.

AS -

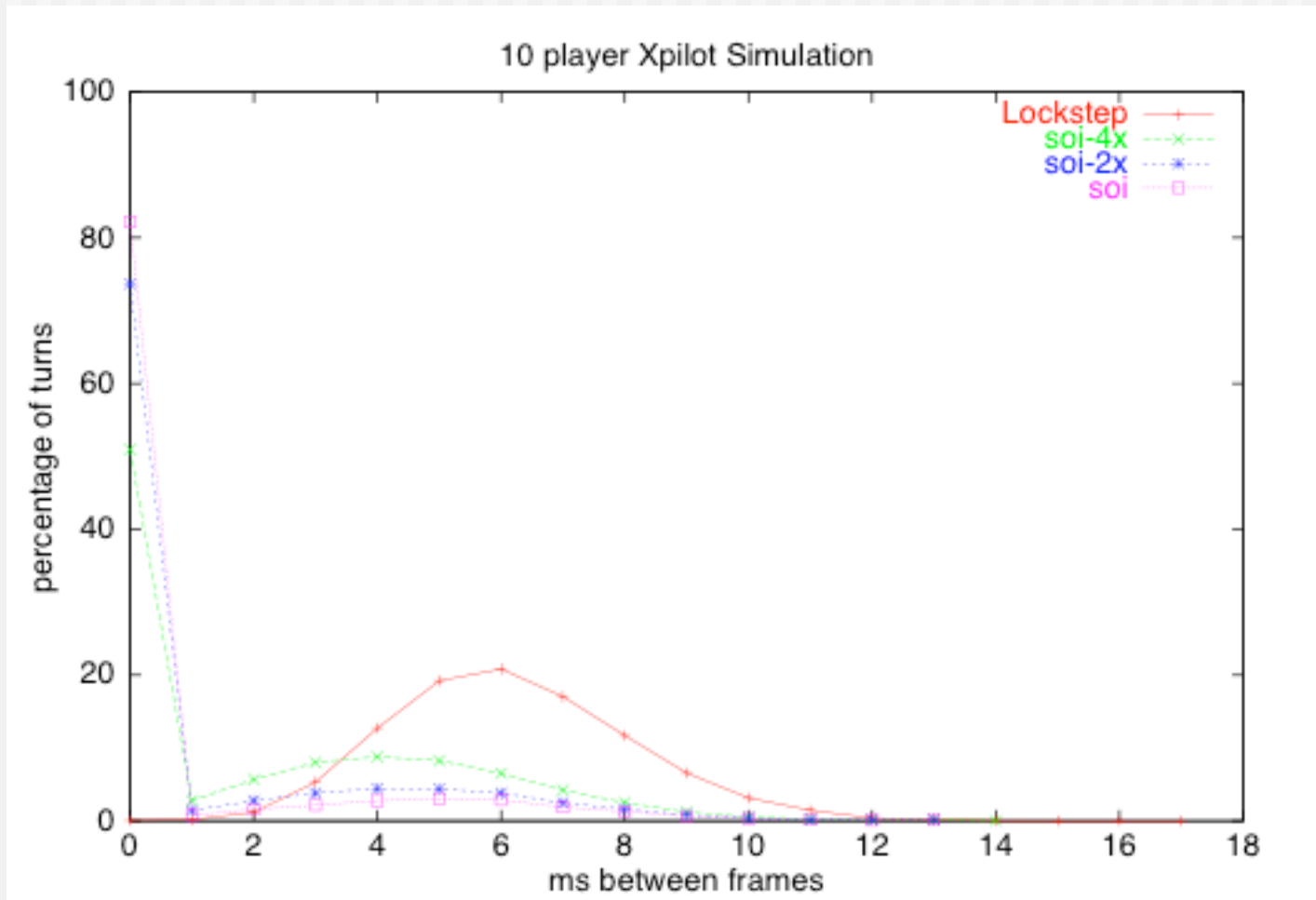
MMO's - Hidden Positions

- Player A generates rand R_a and sends B $h(R_a)$
- Player B generates rand R_b and sends A $h(R_b)$
- Both compute $R = R_a \text{ XOR } R_b$ (? which means they exchange R_a and R_b as well right ?)
- A computes $z = E_k(x+R)$, sends z to B
- B computes $w = E_{k'}(y+R)$, sends w to A
- A computes $w' = E_k(w)$, sends w' to B
- B computes $z' = E_{k'}(z)$, sends z' to A
- A & B have z' & w' , if $z' == w'$ then $x == y$ (same region)
- (because $z' = E_{k'}(E_k(x+R))$, $w' = E_k(E_{k'}(y+R))$) only difference can be from x and y)

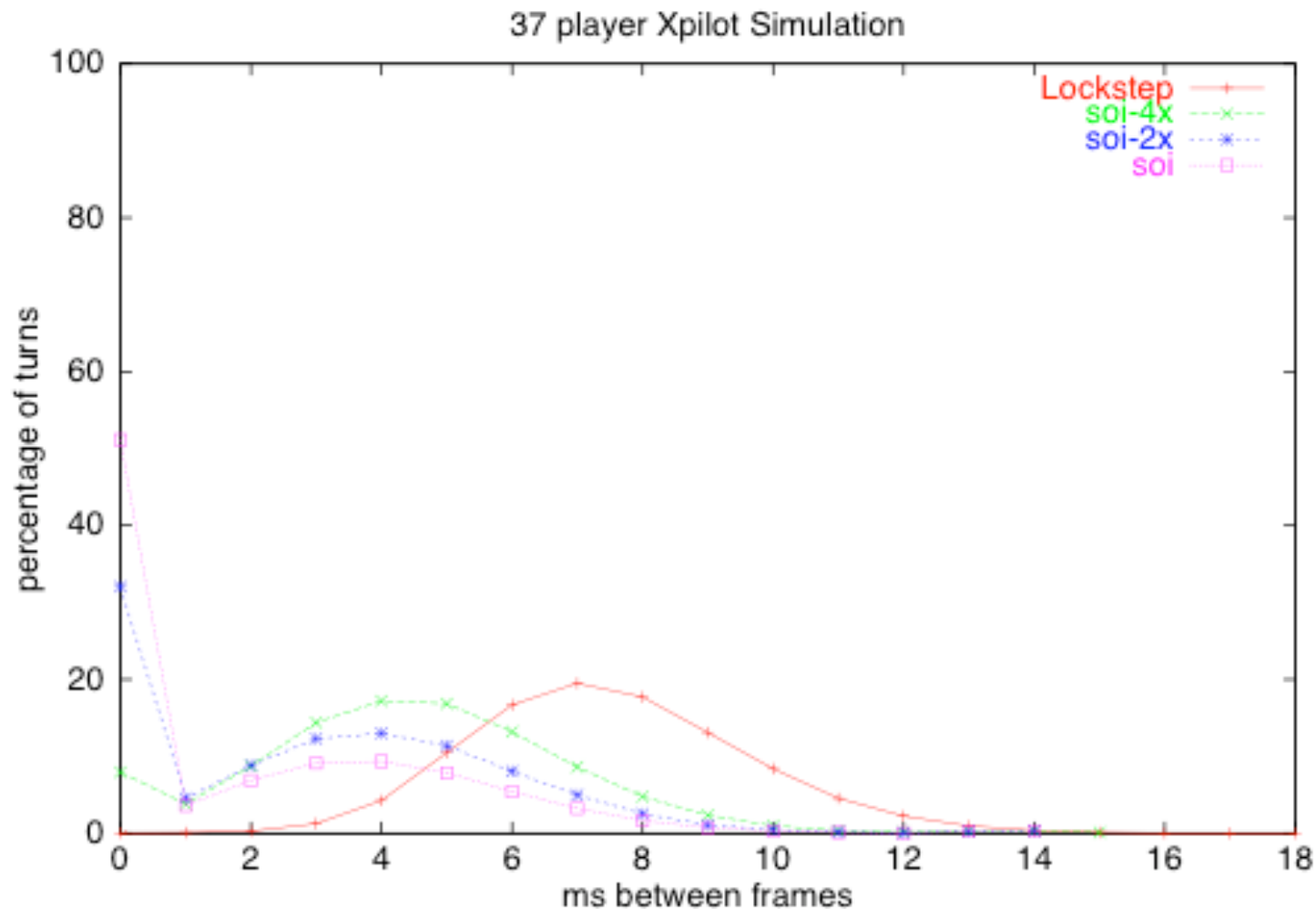
AS - Performance Comparison

- The reason we switched to AS was mostly due to performance
- Performance comparison used Xpilot game and 10,18,30, and 37 players
- Comparison was to see any improvement in frame rendering using AS over lockstep
- Results showed that AS always kept the ms between frames lower overall

AS - Performance Comparison



AS - Performance Comparison



Conclusions

- AS helps prevent two proposed cheats
- AS performs better than standard lock-step algorithm
- However, “it is clear dead reckoning will perform better...”(p.8)

Conclusions

- I appreciated their approach to make the game design resilient to hackers
- (traditionally just hack back at the hackers)
- But, they only really addressed a couple of possible cheats - it would have been nice to see more considered
- They claim it will scale for MMO's
- Comments?