MIPS Architecture

Registers

The MIPS processor has 32 general-purpose registers, plus one for the program counter (called PC) and two for the results of the multiplication and division operations, called HI and LO, for the high 32 bits and the low 32 bits of the answer. The following chart summarizes the registers' usage.

Number	Name	Use
\$0		always holds the value 0
\$1	\$at	reserved by the assembler
\$2 \$3	\$v0 \$v1	expression evaluation and function results
\$4 \$7	\$a0 \$a3	*first 4 function parameters
\$8 \$15	\$t0 \$t7	*temporaries
\$16 \$23	\$s0 \$s7	saved values
\$24 \$25	\$t8 \$t9	*temporaries
\$26 \$27	\$k0 \$k1	reserved for use by operating system
\$28	\$gp	global pointer
\$29	\$sp	stack pointer
\$30	\$s8	saved value
\$31	\$ra	return address

A * in the use column means the values in those registers are **not** preserved across procedure calls.

Opcode Formats

The MIPS processor uses 3 different types of instructions.

I-Type (Immediate) Instructions

bits 31 ... 26 opcode

bits 25 ... 21 source register

bits 20 ... 16 target (destination) register

bits 15 ... 0 immediate operand

J-Type (Jump) Instructions

bits 31 ... 26 opcode

bits 25 ... 0 target (destination) offset

R-Type (Register) Instructions

bits 31 ... 26 opcode

bits 25 ... 21 source register

bits 20 ... 16 target (source) register

bits 15 ... 11 destination register

bits 10 ... 6 shift amount

bits 5... 0 function information

Assembler Directives

.data [addr]: Indicates beginning of data section. If addr is provided, then the location counter is set to addr.

.text [addr]: Indicates beginning of code section. If addr is provided, then the location counter is set to addr.

.asciiz <string>: Allocates memory for string of chars. Terminated with '\0', and padded with '\0' to word boundary.

.space <number> : Allocates <number> of bytes of memory. <number> is increased to word boundary

.word [value]: Allocates a word. If value is provided, then the word is set to that value.

.end: Indicates end of program. (This is ignored).

Selection from MIPS-32 Instruction Set

	Load/Store Instructions	Multiply/Divide Instructions			
LB	Load Byte	MULT	Multiply		
LBU	Load Byte Unsigned	MULTU	Multiply Unsigned		
LH	Load Halfword	DIV	Divide		
LHU	Load Halfword Unsigned	DIVU	Divide Unsigned		
LW	Load Word	MFHI	Move From HI		
LWL	Load Word Left	MTHI	Move To HI		
LWR	Load Word Right	MFLO	Move From LO		
SB	Store Byte	MTLO	Move to LO		
SH	Store Halfword				
SW	Store Word	Ju	mp & Branch Instructions		
SWL	Store Word Left	J	Jump		
SWR	Store Word Right	JAL	Jump and Link		
	Arithmetic Instructions	JR	Jump to Register		
	(ALU Immediate)	JALR	Jump and Link Register		
ADDI	Add Immediate	BEQ	Branch on Equal		
ADDIU	Add Immediate Unsigned	BNE	Branch on Not Equal		
SLTI	Set on Less Than Immediate	BLEZ	Branch on Less than or Equal to Zero		
SLTIU	Set on Less than Immediate Unsigned	BGTZ	Branch on Greater than Zero		
ANDI	AND Immediate	BLTZ	Branch on Less than Zero		
ORI	OR Immediate	BGEZ	Branch on Zero		
XORI	Exclusive OR Immediate	BLTZAL	Branch on Less than Zero and Link		
Hold	Arithmetic Instructions	BGEZAL	Branch on Zero and Link		
	(3-operand, Register Type)	Shift and Special Instructions			
ADD	Add	SLL	Shift Left Logical		
ADDU	Add Unsigned	SLLV	Shift Left Logical, Variable		
SUB	Subtract	SRA	Shift Right Arithmetic		
SUBU	Subtract Unsigned	SRAV	Shift Right Arithmetic, Variable		
SLT	Set on Less Than	SRL	Shift Right Logical		
SLTU	Set on Less Than Unsigned	SRLV	Shift Right Logical, Variable		
AND	Bitwise And	BREAK	Break		
OR	Bitwise OR	SYSCALL	System Call		
XOR	Bitwise exclusive OR				
NOR	NOR				
		1			

Rules on Delays and Interlocks

- There is one delay slot after any branch or jump instruction, i.e., the following instruction is executed even if the branch is taken. That following instruction must not be itself a jump or branch.
- There is one delay slot after a "load" no matter what size is being loaded. That is, the instruction after a "load" must *not* use the register being loaded.
- Multiplication will place its results in the LO and HI registers after an undefined number of following instructions have executed. There's a hardware interlock to stall further multiplications, divisions, or move from LO or HI to execute until the operation is finished.
- Division is like multiplication but most likely slower.

MIPS Opcodes and Formats

These are synopses of many of the core MIPS instructions. Not all instructions are listed; in particular, those involving traps, floats, or memory management are omitted.

ADD rd, rs, rt add Opcode: 000000 Func: 100000

Adds rs and rt, puts result into rd. Exception on overflow.

ADDI rt, rs, immediate add, immediate Opcode: 001000

Sign-extends the 16-bit immediate o 32 bits, adds it to rs, puts resultinto rt. Exception on overflow.

ADDIU rt, rs, immediate add, unsigned immediate Opcode: 001001

Sign-extends the 16-bit immediate o 32 bits, adds it to rs, puts resultinto rt. Never causes an overflow.

ADDU rd, rs, rt add, unsigned Opcode: 000000 Func: 100001

Adds rs and rt, puts result into rd. Never causes an overflow.

AND rd, rs, rt and Opcode: 000000 Func: 100100

Bitwise and's rs and rt, puts result into rd.

ANDI rt, rs, immediate and, immediate Opcode: 001100

Sign-extends the 16-bit immediateto 32 bits, bitwise ands it with rs, puts result into rt.

BEO rs, rt, offset branch equal Opcode: 000100

If rs == rt, branches to offset[after executing the following instruction]. For most assemblers, offset is a label.

BGEZ rs, offset branch greater-equal-zero Opcode: 000001 rt: 00001

If $rs \ge 0$, branches to offset[after executing the following instruction]. For most assemblers, offset is a label.

BGEZAL rs, offset branch greater-equal-zero, and link Opcode: 000001 rt: 10001

If $rs \ge 0$, branches to offset[after executing the following instruction]. For most assemblers, offset is a label. Always places address of following instruction into r31. Note that rs may not itself be r31. (This is a subroutine call instruction)

BGTZ rs, offset branch greater-than-zero Opcode: 000111

If rs > 0, branches to offset[after executing the following instruction]. For most assemblers, offset is a label.

BLEZ rs, offset branch less-equal-zero Opcode: 000110

If $rs \le 0$, branches to offset[after executing the following instruction]. For most assemblers, offset is a label.

BLTZ rs, offset branch less-than-zero Opcode: 000001 rt: 00000

If rs < 0, branches to offset[after executing the following instruction]. For most assemblers, offset is a label.

BLTZAL, rs offset branch less-than-zero, and link Opcode: 00001 rt: 10000

If rs < 0, branches to offset [after executing the following instruction].

For most assemblers, offsetis a label. Always places address of following instruction into r31.

Note that rs may not itself be r31. (This is a subroutine call function.)

BNE rs, rt, label branch not-equal Opcode: 000101

If $rs \neq rt$ branches to offset [after executing the following instruction].

For most assemblers, offset is a label.

BREAK break Opcode: 000000 func: 001101

Causes a Breakpoint exception that transfers control to the exception handler.

DIV rs, rt divide Opcode: 000000 func: 011010

Divides rsby rt, treating both as (signed) 2's complement numbers. Quotient goes into special register LO and remainder into special register HI. Get them via the MFHI and MFLO instructions.

No overflow exception occurs, and the result is undefined if rtcontains 0.

Note that divides take an undefined amount of time; other instructions will execute in parallel.

Opcode: 000000

func: 011011

MFHI and MFLO will interlock until the division is complete.

DIVU rs, rt divide, unsigned

Divides rsby rt, treating both as unsigned numbers. Quotient goes into special register LO and remainder into special register HI. Get them via the MFHI and MFLO instructions. No overflow

exception occurs, and the result is undefined if rt contains 0.

Note that divides take an undefined amount of time; other instructions will execute in parallel. MFHI and MFLO will interlock until the division is complete. This instruction never causes an

exception.

J label jump Opcode: 000010

Jump to label [after executing the following instruction].

JAL label jump and link Opcode: 000011

Jump to label[after executing the following instruction]. Places the address of the following instruction

into r31. (This is a subroutine-call instruction.)

JALR rd, rs jump and link, register Opcode: 000000 func: 001001

Jump to address contained in rs [after executing the following instruction]. Places address of following instruction into rd. Note that rs and rd may not be the same register. If rd is omitted in the

assembly language, it is register 31. (This is a subroutine-call instruction.)

JR rs jump, register Opcode: 000000 func: 001000

Jump to address contained in rs [after executing the following instruction].

LA rt, addr load address into register Pseudo instruction

This is a a pseudo instruction that is translated into:

lui \$rt, addr(16..31) followed by ori \$rt, \$rt, addr(0..15)

LB rt, offset(rs) load byte Opcode: 100000

Sign-extend the 16-bit offsetto 32 bits, and add it to rsto get an effective address. Load the byte

from this address into rt and sign-extend it to fill the entire register.

LBU rt, offset(rs) load byte, unsigned

Sign-extend the 16-bit offsetto 32 bits, and add it to rsto get an effective address. Load the byte from this address into rt and zero-extend it to fill the entire register.

LH rt, offset(rs) load halfword

Sign-extend the 16-bit offsetto 32 bits, and add it to rsto get an effective address. Load the halfword (16 bits) from this address into rt and sign-extend it to fill the entire register. Exception if odd address.

LHU rt, offset(rs) load halfword, unsigned

Sign-extend the 16-bit offsetto 32 bits, and add it to rsto get an effective address. Load the halfword (16 bits) from this address into rt and zero-extend it to fill the entire register. Exception if odd address.

LI rt, immediate load a 32-bit immediate into a register

Pseudo instruction

Opcode: 001111

Opcode: 100011

Opcode: 000000

Opcode: 000000

func: 010000

func: 010010

func: 010011

Opcode: 100100

Opcode: 100001

Opcode: 100101

This is a a pseudo instruction that is translated into:

lui \$rt, immediate(16..31) followed by ori \$rt, \$rt, immediate(0..15)

LUI rt, immediate load upper immediate

Put 16-bit immediate in the top half of rtand fill the bottom half withzeros.

LW rt, offset(rs) load word

Sign-extend the 16-bit offsetto 32 bits, and add it to rsto get an effective address. Load the word (32 bits) from this address into rt. Exception if address is not word-aligned.

move from HI MFHI rd

Move contents of special register Hlinto rd. Neither of the two instructions following this may modify the HIregister! Note that multiplication and division put results into HI. MFHI stalls until that operation is complete.

MFLO rd move from LO

Move contents of special register LOinto rd. Neither of the two instructions following this may modify the LO register! Note that multiplication and division put results into LO. MFLO stalls until that operation is complete.

MTHI rs move to HI

Opcode: 000000 func: 010001 Move contents of rsinto special register HI. May cause contents of LO to become undefined; no need to get specific here; just be sure to do MTLO too.

MTLO rs move to LO

Opcode: 000000 Move contents of rsinto special register LO. May cause contents of HI tobecome undefined; no need to get specific here; just be sure to do MTHI too.

MULT rs, rt multiply

Opcode: 000000 func: 011000 Multiplies rsby rt, treating both as (signed) 2's complement numbers. Low word of result goes into special register LO and high word into special register HI. Get them via the MFHI and MFLO instructions. No over- flow exception occurs.

Note that multiplies take an undefined amount of time; other instructions will execute in parallel. MFHI and MFLO will interlock until the multiplication is complete.

MULTU rs, rt multiply, unsigned

Opcode: 000000 func: 011001 Multiplies rsby rt, treating both as unsigned numbers. Low word of result goes into special register LO and high word into special register HI. Get them via the MFHI and MFLO instructions. No overflow exception occurs. Note that multiplies take an undefined amount of time; other instructions will execute in parallel. MFHI and MFLO will interlock until the multiplication is complete. This instruction never causes an exception.

NOP no-op Pseudo instruction

Do nothing for one cycle; good for filling a delay slot. Assemblers often use sll \$0, \$0, 0.

NOR rd. rs. rt Opcode: 000000 func: 100111 nor

Performs bitwise logical nor of rsand rt, putting result into rd.

Opcode: 000000 func: 100101 OR rd, rs, rt or

Performs bitwise logical or of rsand rt, putting result into rd.

ORI rt, rs, immediate or, immediate Opcode: 001101

Zero-extends 16-bit immediate to 32 bits, and bitwise ors it with rt, putting result into rd.

SB rt, offset(rs) store byte Opcode: 101000

Sign-extend the 16-bit offset to 32 bits, and add it to rs to get an effective address. Store least significant byte from rtinto this address.

SH rt, offset(rs) store halfword Opcode: 101001

Sign-extend the 16-bit offsetto 32 bits, and add it to rs to get an effective address. Store least significant byte from rt into this address. Exception if odd address.

func: 000000 SLL rd, rt, sa shift left logical Opcode: 000000

Shift contents of rt left by the amount indicated in sa, insertion zeroes into the emptied low order bits. Put the result into rd.

SLLV rd, rs, rt Opcode: 000000 func: 0001000 *shift left logical, variable*

Shift contents of rt left by the amount indicated in the bottom five bitsof rs (0..4), inserting zeros into the low order bits. Put result into rd.

SLT rd, rs, rt Opcode: 000000 func: 101010 set on less-than

If rs < rt with a signed comparison, put 1 into rd. Otherwise put 0 into rd.

SLTI rt, rs, immediate set on less-than, immediate Opcode: 001010

Sign-extend the 16-bit immediate to a 32-bit value. If rs is less than this value with a signed comparison, put 1 into rt. Otherwise put 0 into rt.

SLTIU rt, rs immediate set on less-than, immediate unsigned Opcode: 001011

Sign-extend the 16-bit immediateto a 32-bit value. If rs is less than this value with an unsigned comparison, put 1 into rt. Otherwise put 0 into rt.

SLTU rd, rs, rt set on less-than, unsigned Opcode: 000000 func: 101011

If rt < rs with an unsigned comparison, put 1 into rd. Otherwise put 0 into rd.

shift right arithmetic Opcode: 000000 func: 000011 SRA rd, rt, sa

Shift contents of rtright by the amount indicated by sa, sign-extending the high order bits.

Put result into rd.

Opcode: 000000 func: 000111 SRAV rd, rs, rt shift right arithmetic, variable

Shift contents of rt right by the amount indicated in the bottom five bitsof rs (0..4), signextending the high order bits. Put result into rd.

SRL rd, rt, sa shift right logical Opcode: 000000 func: 000010

Shift contents of rtright by the amount indicated in sa, zero-filling the high order bits. Put

result into rd.

SRLV rd, rs, rt shift right logical, variable Opcode: 000000 func: 000110 Shift contents of rt right by the amount indicated in the bottom five bits of rs(0..4), zero-filling the high order bits. Put result into rd.

SUB rd, rs, rt subtract Opcode: 000000 func: 100010

Put rs - rt into rd. Exception if overflow.

SUBU rd, rs, rt subtract unsigned Opcode: 000000 func: 100011

Put rs – rt into rd. Never causes exception.

SW rt, offset(rs) store word Opcode: 101011

Sign-extend the 16-bit offsetto 32 bits, and add it to rs to get an effective address. Store rt into this address. Exception if address is not word-aligned.

SYSCALL system call Opcode: 000000 func: 001100

Causes a System Call exception. For ECS 50 the response is based on the value in $v0. 1 = print_int$, $v0. 1 = print_int$, $v0. 1 = print_int$.

XOR rd, rs, rt exclusive or Opcode: 000000 func: 100110

Performs bitwise exclusive xor of rs and rt, putting result into rd.

XORI rt, rs, immediate xor immediate Opcode: 001110

Zero-extends 16-bit immediate to 32 bits, and bitwise exclusive xors it with rs, putting result

into rt.

MIPS Example: Initializing an Array

This shows an assembly language program which initializes the integer array *arr* such that each of its ten elements is equal to the index of that element. It also prints the value after it is inserted in the array.

```
# Written by: Matt Bishop and adapted by Sean Davis
# Registers used:
           $0 -- to get a 0 (standard usage)
#
           $a0 -- to choose system service
           $t0 -- index of arr
#
           $t1 -- offset of current element from base of arr
#
           $t2 -- temporary (usually holds result of comparison)
           .data 0x40
                                 # set start address of data section
arr:
           .space 40
                                 # allocate 10 words
                                 # set start address of instructions
           .text 0
           addu $t0, $0, $0
                                 # initialize index of array
init:
           sll $t1, $t0, 2
                                 # go to offset of next element
loop:
           sw $t0, arr($t1)
                                 # store integer into element
                                 # copy from $t0 to $a0
           add $a0, $t0, $0
                                 # set $v0 to print_integer code for syscall
           addi $v0, $0, 1
           syscall
                                 # print the integer in $a0
                                 # add one to current array index
           addiu $t0, $t0, 1
           slti $t2, $t0, 10
                                 # see if the index is 10 yet
           bne $t2, $0, loop
                                 # nope -- go back for another
                                 # for the delay slot
           addiu $v0, $0, 10
                                 # set $v0 to exit code for syscall
           syscall
                                 # exit
           .end
```

MIPS32 Encoding of the Opcode Field

ор	opcode bits 2826								
		0	1	2	3	4	5	6	7
bits	3129	000	001	010	011	100	101	110	111
0	000	SPECIAL δ	REGIMM δ	J	JAL	BEQ	BNE	BLEZ	BGTZ
1	001	ADDI	ADDIU	SLTI	SLTIU	ANDI	ORI	XORI	LUI
2	010	СОР0 δ	COP1 δ	COP2 θδ	COP1X ¹ δ	BEQL φ	BNEL φ	BLEZL φ	BGTZL φ
3	011	β	β	β	β	SPECIAL2 δ	JALX ε	8	SPECIAL3 ² δ⊕
4	100	LB	LH	LWL	LW	LBU	LHU	LWR	β
5	101	SB	SH	SWL	SW	β	β	SWR	CACHE
6	110	LL	LWC1	LWC2 θ	PREF	β	LDC1	LDC2 θ	β
7	111	SC	SWC1	SWC2 θ	*	β	SDC1	SDC2 θ	β

MIPS32 SPECIAL Opcode Encoding of Function Field

fur	nction	bits 20							
		0	1	2	3	4	5	6	7
bit	s 53	000	001	010	011	100	101	110	111
0	000	SLL ¹	ΜΟΥСΙ δ	SRL δ	SRA	SLLV	*	SRLV δ	SRAV
1	001	JR ²	JALR ²	MOVZ	MOVN	SYSCALL	BREAK	*	SYNC
2	010	MFHI	MTHI	MFLO	MTLO	β	*	β	β
3	011	MULT	MULTU	DIV	DIVU	β	β	β	β
4	100	ADD	ADDU	SUB	SUBU	AND	OR	XOR	NOR
5	101	*	*	SLT	SLTU	β	β	β	β
6	110	TGE	TGEU	TLT	TLTU	TEQ	*	TNE	*
7	111	β	*	β	β	β	*	β	β

MIPS32 REGIMM Encoding of rt Field

	rt bits 1816									
		0	1	2	3	4	5	6	7	
bits	2019	000	001	010	011	100	101	110	111	
0	00	BLTZ	BGEZ	BLTZL φ	BGEZL φ	*	*	*	3	
1	01	TGEI	TGEIU	TLTI	TLTIU	TEQI	*	TNEI	*	
2	10	BLTZAL	BGEZAL	BLTZALL φ	BGEZALL φ	*	*	*	*	
3	11	*	*	*	*	3	3	*	SYNCI ⊕	