

1 Simulations

1.1 4x4 Crossbar

Bestimmen und begründen Sie wie die Werte zustande kommen

$$Avg = \sum_{n=0}^N \frac{\binom{N}{n} (N - n + 1)! \cdot C_n}{\sum_{k=0}^N \binom{N}{k} (N - k + 1)!} \quad (1)$$

- avg Input Queue Length
 - Generated Packets
 - * size: $s = 512 \text{ b}$
 - * send interval: $t = 512 \text{ ns}$
 - Connection to XBar
 - * data rate: 1 Gbps

$$r_{\text{gen}} = \frac{512 \text{ b}}{512 \text{ ns}} = 1 \text{ Gbps} \quad (2)$$

$$C_n = 250 \text{ Mbps} \cdot n \quad (3)$$

$$r_{\text{q fill, eff}} = \sum_{n=1}^4 \frac{\binom{4}{n} (5 - n)! \cdot C_n}{141} \quad (4)$$

$$= 347.5 \text{ Mbps} \quad (5)$$

$$\Rightarrow l_{\text{q, avg}} = \frac{t_{\text{sim}} \cdot r_{\text{q fill, eff}}}{2 \cdot s} = 3394 \quad (6)$$

- avg End-to-End Latency
 - (no delays inside buffers, because the generated data rate, even for all 4 apps, is lower than a single data rate channels maximum throughput)
 - minimum
 - * App \rightarrow C \rightarrow Inport \rightarrow C \rightarrow Outport \rightarrow C \rightarrow App
 - * delay for packet: $t_{\text{delay}} = 512 \text{ ns}$ (per `DatarateChannel C`)
- $$\Rightarrow t_{e2e, \text{min}} = 1.536 \mu\text{s} \quad (7)$$
- maximum

at end of simulation, inport buffer full, all in to one out

$$l_{\text{inportq, max}} = 7031 \quad (8)$$

$$r_{\text{dequeue, min}} = 250 \text{ Mbps} \quad (9)$$

$$t_{\text{in queue, max}} = \frac{t_{q, \text{inport}} \cdot s}{r_{\text{dequeue}}} = 3.599 \text{ ms} \quad (10)$$

$$(11)$$

- on avg

$$t_{e2e, \text{ avg}} = t_{\text{in queue, max}}/2 = 1.8 \text{ ms} \quad (12)$$

- avg Arbiter Request Queue Length

- cases

$$C_n = n - 1 \quad 1 \leq n \leq 4 = N \quad (13)$$

- on avg (analogously to t_{e2e})

$$\Rightarrow l_{arbq, avg} = \sum_{n=1}^4 \frac{\binom{4}{n} (5-n)! \cdot C_n}{141} = 0.39 \quad (14)$$

- avg Arbiter Request Queue Time

- cases

$$C_n = 512 \text{ ns} * n \quad (15)$$

- on avg

$$\Rightarrow t_{arbq, avg} = \sum_{n=1}^4 \frac{\binom{4}{n} (5-n)! \cdot C_n}{141} \quad (16)$$

$$= \quad (17)$$

- avg Output Buffer Queue Length

- Generated Packets

- * size: $s = 512 \text{ b}$

- * send interval: $t = \text{uniform}(1 \mu\text{s}, 10 \mu\text{s}) = 5.5 \mu\text{s}$

- Connection to and in XBar

- * data rate: 1 Gbps

$$r_{\text{generated}} = \frac{512 \text{ b}}{5.5 \mu\text{s}} = 93 \text{ Mbps} \quad (18)$$

$$\Rightarrow l_{q, avg} = 0 \quad (19)$$

- avg Throughput

input and output buffers always empty

$$r_{\text{cross}} = 4 * r_{\text{gen}} \quad (20)$$

$$= 372 \text{ Mbps} \quad (21)$$

1.2 Throughput vs. Ports

- with our formula 0.6, 0.44, 0.347518, 0.286498, 0.243324, 0.211277, 0.186604, 0.167047, 0.151176, 0.138045,
 \downarrow 0.127005, 0.117594, 0.109478, 0.102407, 0.0961928, 0.0906881, 0.0857783, 0.0813722, 0.077396,
 \downarrow 0.0737899, 0.0705046, 0.067499, 0.0647391, 0.0621958, 0.0598446, 0.0576646, 0.0556378, 0.0537485,
 \downarrow 0.0519832, 0.0503302, 0.048779

1.3 Throughput vs. Injection Rate

- not in saturation until delay ≤ 832 ns
- saturation point $r_{sat} = \frac{512\text{b}}{832\text{ns}} = 615$ Mbps
- The main reason will be the arbiter. It can only process packets at about 600 Mbps on avg., therefore bottlenecking the rest of the system

1.4 Throughput vs. Bandwidth

- Network in saturation until Bandwidth > 1600 Mbps
- throughput at saturation point $r_{sat} = 1$ Gbps
- If the Arbiter can only work at 62% of the bandwidth, we reach this saturation point if 62% of the bandwidth is 1 Gbps, therefore the required bandwidth is 1613 Gbps

2 Optimizations

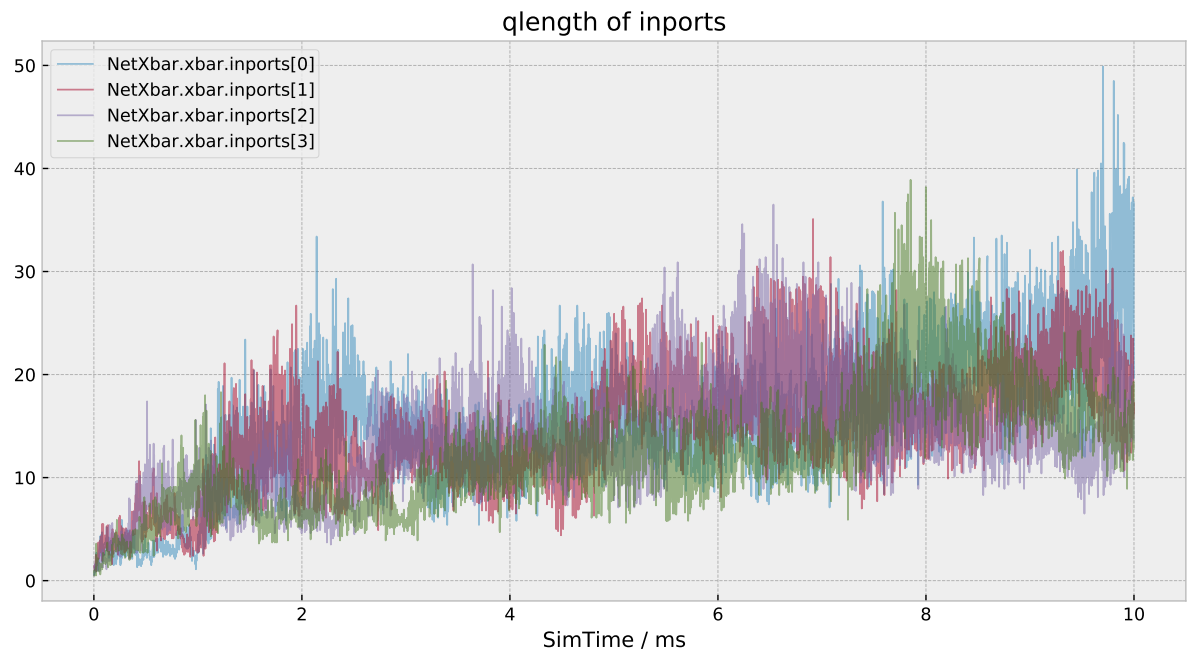


Abbildung 1: sim

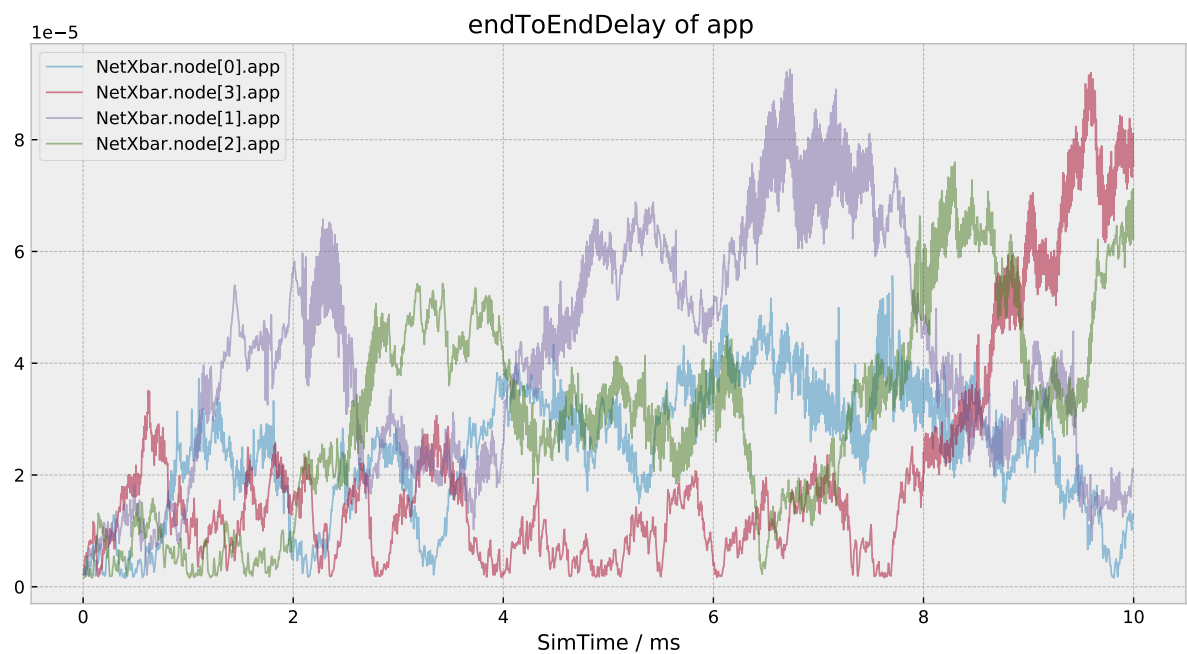


Abbildung 2

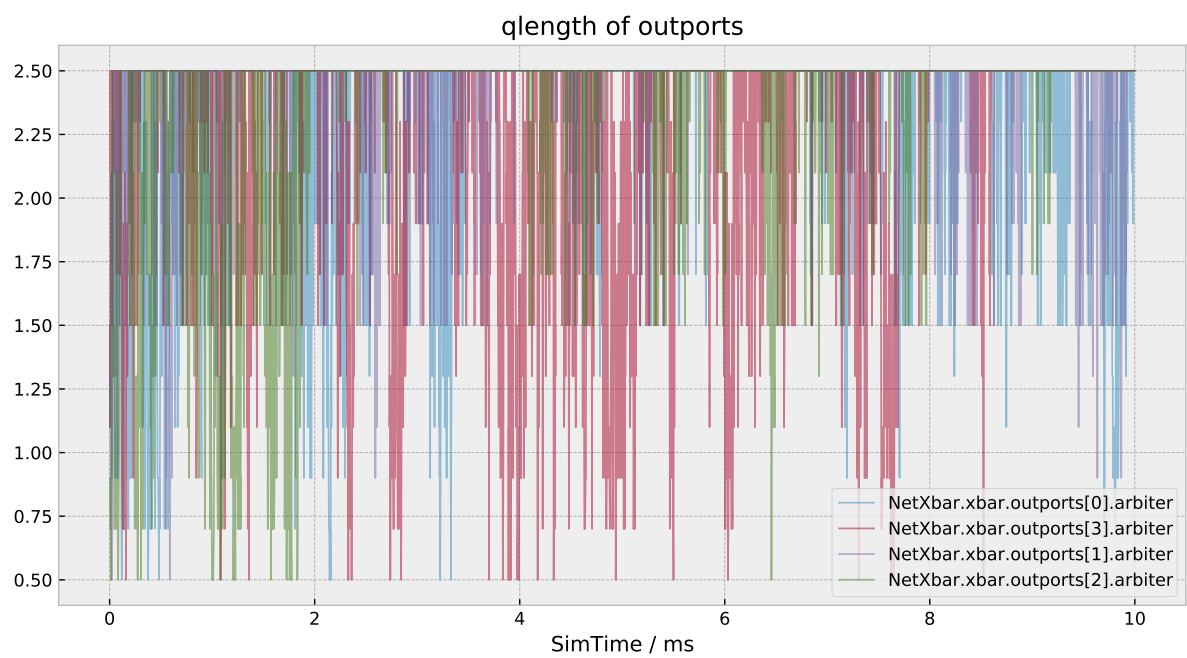


Abbildung 3