1 Simulations

1.1 4x4 Crossbar

Bestimmen und begründen Sie wie die Werte zustande kommen

$$Avg = \sum_{n=0}^{N} \frac{\binom{N}{n} (N-n+1)! \cdot C_n}{\sum_{k=0}^{N} \binom{N}{k} (N-k+1)!}$$
 (1)

- avg Input Queue Length
 - Generated Packets
 - * size: $s = 512 \,\mathrm{b}$
 - * send interval: $t = 512 \,\mathrm{ns}$
 - Connection to XBar
 - * data rate: 1 Gbps

$$r_{\rm gen} = \frac{512 \,\mathrm{b}}{512 \,\mathrm{ns}} = 1 \,\mathrm{Gbps}$$
 (2)

$$C_n = 250 \,\text{Mbps} \cdot n \tag{3}$$

$$r_{\text{q fill,eff}} = \sum_{n=1}^{4} \frac{\binom{4}{n} (5-n)! \cdot C_n}{141}$$
 (4)

$$= 347.5 \,\mathrm{Mbps} \tag{5}$$

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$$\Rightarrow l_{\text{q,avg}} = \frac{t_{\text{sim}} \cdot r_{\text{q fill, eff}}}{2 \cdot s} = 3394$$

$$(5)$$

- avg End-to-End Latency
 - (no delays inside buffers, because the generated data rate, even for all 4 apps, is lower than a single data rate channels maximum throughput)
 - minimum
 - * App \rightarrow C \rightarrow Inport \rightarrow C \rightarrow Outport \rightarrow C \rightarrow App
 - * delay for packet: $t_{delay} = 512ns$ (per DatarateChannel C)

$$\Rightarrow t_{e2e,min} = 1.536 \,\mu\text{s} \tag{7}$$

- maximum

at end of simulation, inport buffer full, all in to one out

$$l_{inporta,max} = 7031 \tag{8}$$

$$r_{dequeue,min} = 250 \,\text{Mbps}$$
 (9)

$$t_{\text{in queue, max}} = \frac{t_{q,inport} \cdot s}{r_{dequeue}} = 3.599 \,\text{ms}$$
 (10)

(11)

- on avg

$$t_{\text{e2e. avg}} = t_{\text{in queue. max}}/2 = 1.8 \,\text{ms}$$
 (12)

- avg Arbiter Request Queue Length
 - cases

$$C_n = n - 1 \qquad 1 \le n \le 4 = N \tag{13}$$

- on avg (analogously to t_{e2e})

$$\Rightarrow l_{arbq,avg} = \sum_{n=1}^{4} \frac{\binom{4}{n} (5-n)! \cdot C_n}{141} = 0.39$$
 (14)

- avg Arbiter Request Queue Time
 - cases

$$C_n = 512 \,\mathrm{ns} * n \tag{15}$$

- on avg

$$\Rightarrow t_{arbq,avg} = \sum_{n=1}^{4} \frac{\binom{4}{n} (5-n)! \cdot C_n}{141}$$
 (16)

$$= (17)$$

- avg Output Buffer Queue Length
 - Generated Packets
 - * size: $s = 512 \,\mathrm{b}$
 - * send interval: $t = \text{uniform}(1 \,\mu\text{s}, 10 \,\mu\text{s}) = 5.5 \,\mu\text{s}$
 - Connection to and in XBar
 - * data rate: 1 Gbps

$$r_{\text{generated}} = \frac{512 \,\text{b}}{5.5 \,\mu\text{s}} = 93 \,\text{Mbps} \tag{18}$$

$$\Rightarrow l_{\text{q,avg}} = 0 \tag{19}$$

• avg Throughput

input and output buffers always empty

$$r_{cross} = 4 * r_{gen} \tag{20}$$

$$= 372 \,\mathrm{Mbps} \tag{21}$$

1.2 Throughput vs. Ports

- with ourr formula 0.6, 0.44, 0.347518, 0.286498, 0.243324, 0.211277, 0.186604, 0.167047, 0.151176, 0.138045,
 - $\begin{array}{l} \hbox{\it i.} \ 0.127005, \ 0.117594, \ 0.109478, \ 0.102407, \ 0.0961928, \ 0.0906881, \ 0.0857783, \ 0.0813722, \\ 0.077396, \end{array}$
 - $\begin{smallmatrix} 0.0737899,\ 0.0705046,\ 0.067499,\ 0.0647391,\ 0.0621958,\ 0.0598446,\ 0.0576646,\ 0.0556378,\ 0.0537485, \end{smallmatrix}$
 - j. 0.0519832, 0.0503302, 0.048779

1.3 Throughput vs. Injection Rate

- not in saturation until delay $\leq 832 \,\mathrm{ns}$
- saturation point $r_{sat} = \frac{512 \text{ b}}{832 \text{ ns}} = 615 \text{ Mbps}$
- The main reason will be the arbiter. It can only process packets at about 600 Mbps on avg., therefore bottlenecking the rest of the system

1.4 Throughput vs. Bandwidth

- Network in aturation until Bandwith $> 1600 \,\mathrm{Mbps}$
- throughput at saturation point $r_{sat} = 1 \text{ Gbps}$
- If the Arbiter can only work at 62% of the bandwidth, we reach this saturation point if 62% of the bandwidth is 1 Gbps, therefore the required bandwith is 1613 Gbps

2 Optimizations