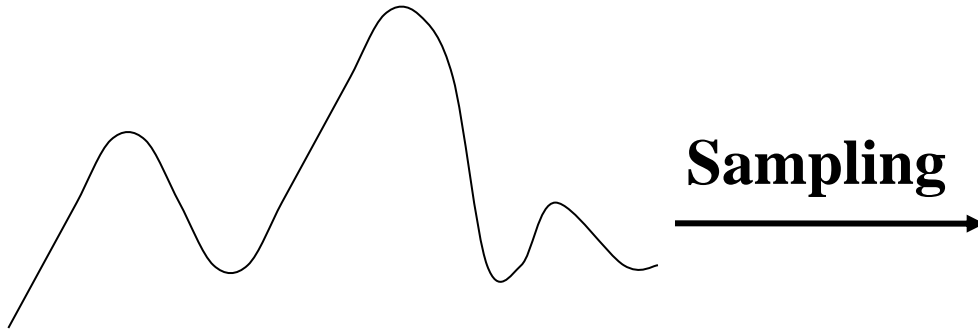


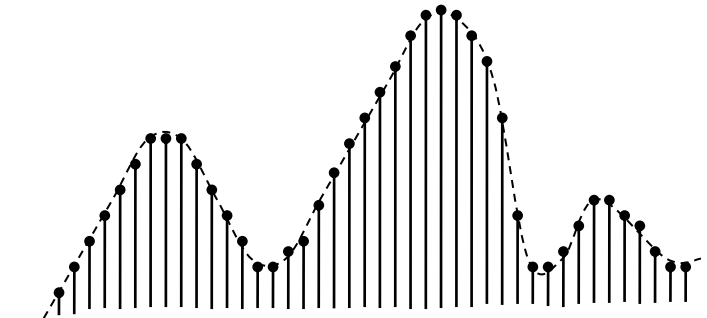
Signal compression: Quantization

1.1

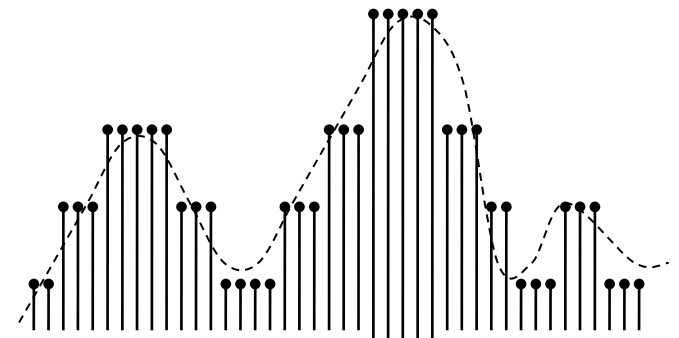
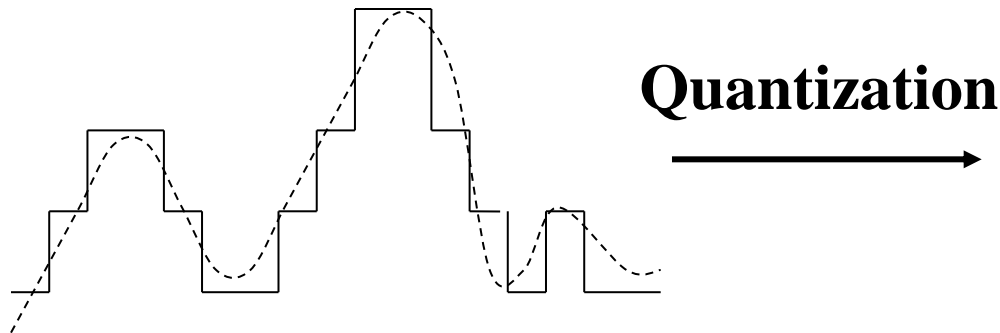
To store, transmit or manipulate signals, they first must be digitized.



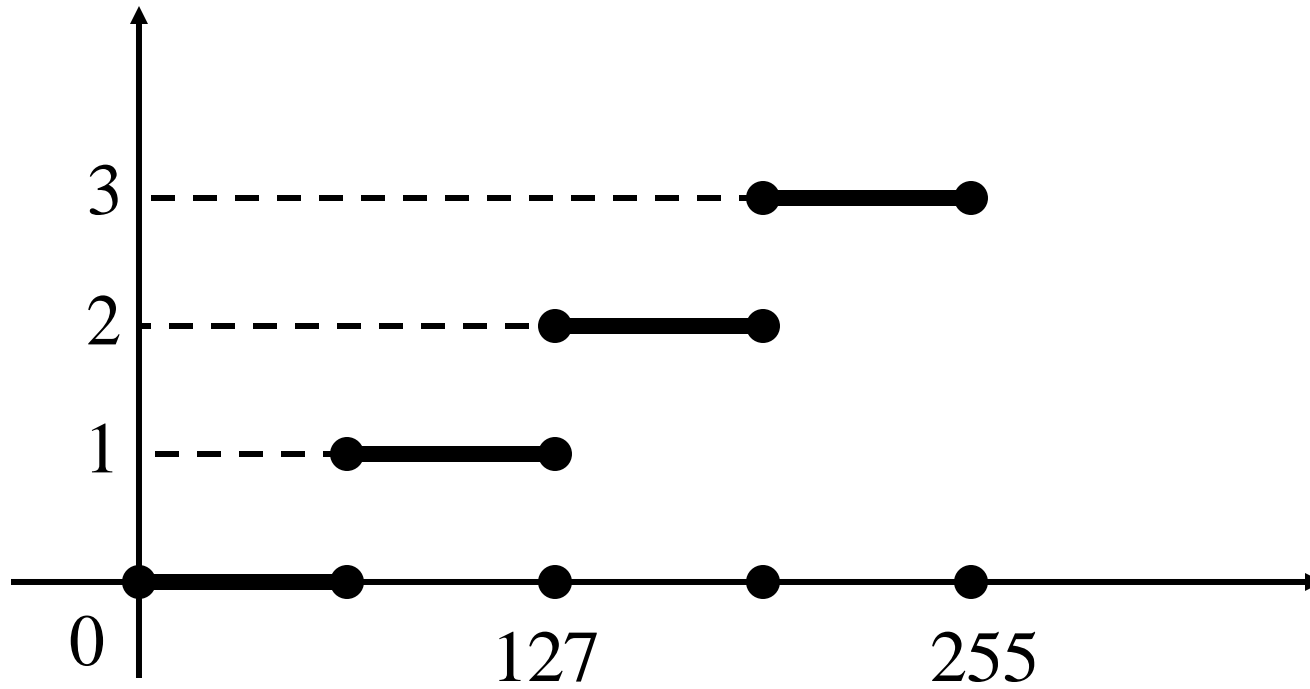
Amplitude-continuous
Time-continuous



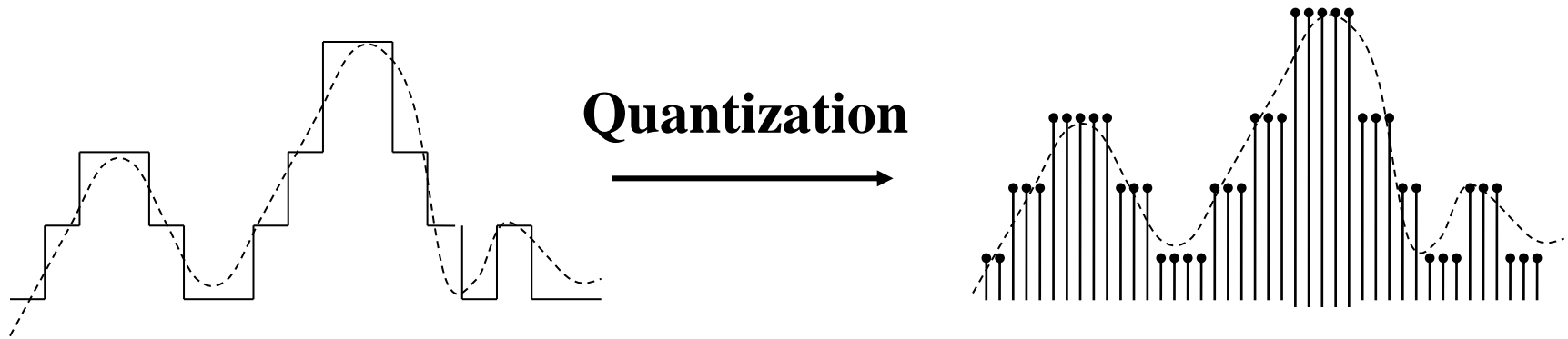
Amplitude-continuous,
Time-discrete



Amplitude-discrete,
Time-discrete



Matlab demo



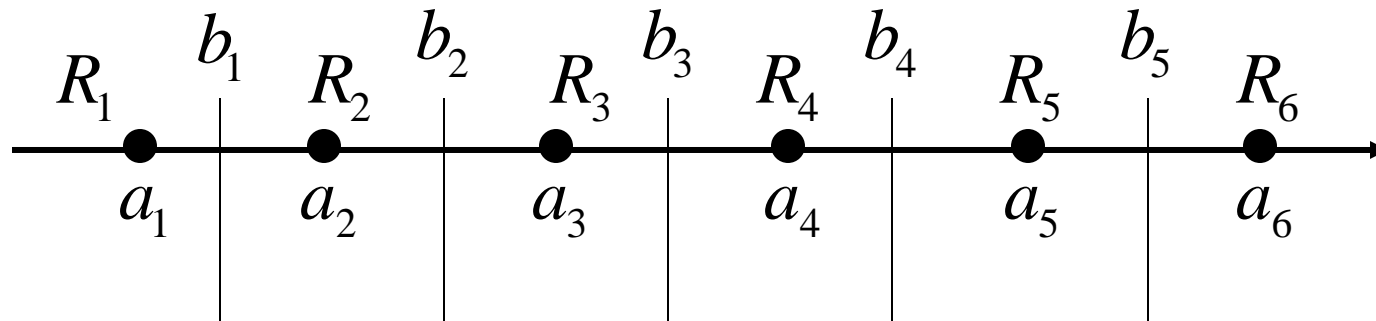
In coming sequence U_1, U_2, \dots is mapped into a sequence of discrete random variables V_1, V_2, \dots . The objective is that, for each m , V_m should represent U_m with as little distortion as possible.

When an analog random variable U is quantized into a discrete random variable $V=V(U)$, the mean-squared distortion (MSE, mean-squared error) is defined by

$$E\left[U - V^2 \right]$$

A scalar quantizer

A scalar quantizer partitions the set of real numbers into M quantizer regions (intervals): R_1, R_2, \dots, R_M containing representation points a_j and separated by boundaries b_1, \dots, b_{M-1}



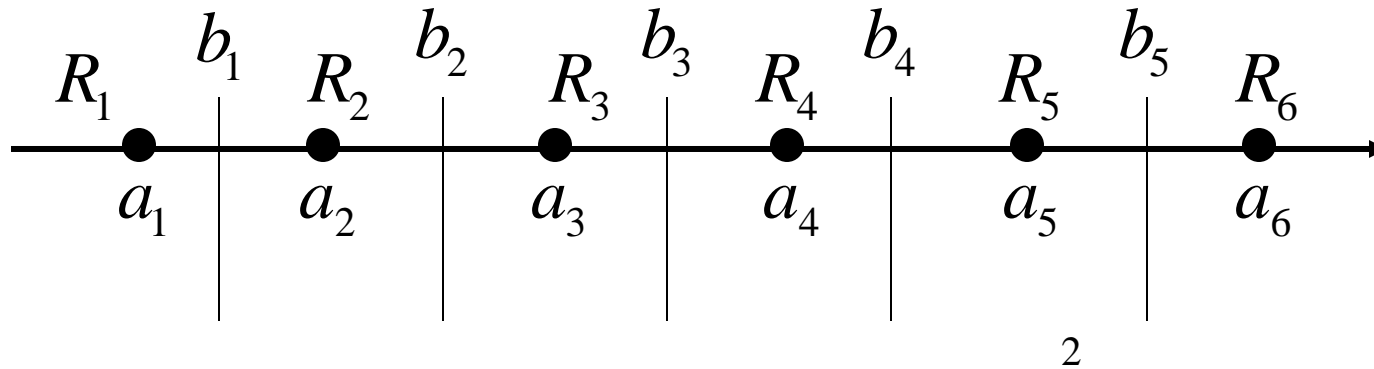
An analog sequence u_1, u_2, \dots . These real numbers are sample values of random variables U_1, U_2, \dots .

- Given a set of representation points $\{a_j\}$, how to choose $\{R_j\}$?
- Given a set of intervals $\{R_j\}$, how to choose representation points $\{a_j\}$?

Choice of intervals for given representation points

1.5

Given representation points $a_j : 1 \leq j \leq M$, determine b_j



Given u , the squared error to a_j is $u - a_j$

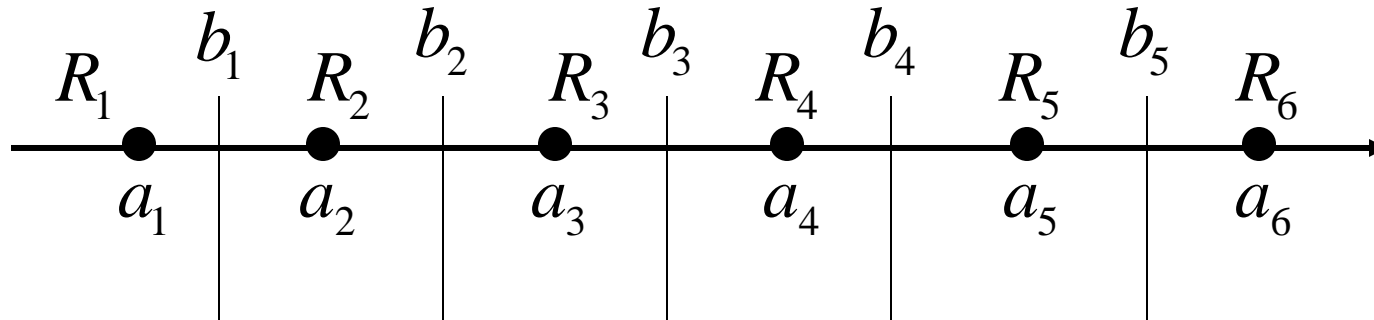
If u is between a_j and a_{j+1} then u is mapped to the closest one.

This implies that $b_j = (a_j + a_{j+1}) / 2$ and the quantization regions are determined.

Note that, in this case, the quantization result does not depend on the probabilistic model for U_1, U_2, \dots

Choice of representation points for given intervals 1 ^{1.6}

Given intervals (i.e., b_j) find representation points a_j



$$\text{MSE} = \mathbf{E} \left[(U - V)^2 \right] = \int_{-\infty}^{\infty} f_U(u) (u - v)^2 du$$

$$= \sum_{j=1}^M \int_{R_j} f_U(u) (u - a_j)^2 du \rightarrow \min \text{ over } a_j$$

Let $f_j(u)$ be the conditional pdf of U given that $u \in R_j$

Choice of representation points for given intervals 2 ^{1.7}

$$\begin{aligned}\text{MSE} &= \mathbf{E} \left[(U - V)^2 \right] = \int_{-\infty}^{\infty} f_U(u) (u - v)^2 du \\ &= \sum_{j=1}^M \int_{R_j} f_U(u) (u - a_j)^2 du \rightarrow \min \text{ over } a_j\end{aligned}$$

Let $f_j(u)$ be the conditional pdf of U given that $u \in R_j$

$$f_j(u) = \begin{cases} f_U(u) / Q_j & \text{if } u \in R_j \\ 0 & \text{otherwise} \end{cases} \quad \text{where } Q_j = \Pr(U \in R_j)$$

$$\text{Thus } \int_{R_j} f_U(u) (u - a_j)^2 du = Q_j \int_{R_j} f_j(u) (u - a_j)^2 du$$

We have to minimize $\int_{R_j} f_j(u) (u - a_j)^2 du$ w.r.t. a_j

Choice of representation points for given intervals 3 ^{1.8}

We have to minimize $\int_{R_j} f_j(u) (u - a_j)^2 du$ over a_j

$$\int_R f(u) (u - a)^2 du = \int_R f(u) u^2 du - 2a \int_R f(u) u du + a^2$$

$$\overline{(X - a)^2} = \overline{X^2} - 2a\overline{X} + a^2$$

a quadratic polynomial which is minimized over a when $a = \overline{X}$

We arrive at a set of conditions that the endpoints $\{b_j\}$ and the points $\{a_j\}$ must satisfy to achieve the MSE:

- Each b_j must be the midpoint between a_j and a_{j+1}
- Each a_j must be the mean of random variable U_j with pdf $f_j(u)$

The Lloyd-Max algorithm 1

Necessary conditions that the endpoints $\{b_j\}$ and the points $\{a_j\}$ must satisfy to achieve the MSE:

- ❑ Each b_j must be the midpoint between a_j and a_{j+1}
- ❑ Each a_j must be the mean of random variable U_j with pdf $f_j(u)$

The Lloyd-Max algorithm alternates between

1. optimising the endpoints $\{b_j\}$ for a given set $\{a_j\}$ of and then
2. optimising the points $\{a_j\}$ for the new endpoints $\{b_j\}$.

Assume that M and $f_U(u)$ are given.

1. Choose an arbitrary initial set of $a_1 < a_2 < \dots < a_M$
2. Set $b_j = (a_{j+1} + a_j)/2$
3. Set a_j equal to the conditional mean of U given $U \in (b_{j-1}, b_j]$
4. Repeat 2. and 3. until further improvement in MSE is negligible.

1. Start from an initial set of representation points a_j $j=1$ M
2. Find $M - 1$ interval endpoints $b_j = (a_j + a_{j+1}) / 2$
3. Compute the mean-squared error (MSE)

$$\text{MSE} = \sum_{k=1}^M \int_{b_{k-1}}^{b_k} (x - a_k)^2 f(x) dx$$

4. Stop if MSE changes little from last time.

5. Otherwise, update $a_j = \frac{\int_{b_{j-1}}^{b_j} x f(x) dx}{\int_{b_{j-1}}^{b_j} f(x) dx}$

6. Go to Step 2.

The Lloyd-Max algorithm: an example

1.11

$$b_j = (a_j + a_{j+1}) / 2 \quad a_j = \int_{b_{j-1}}^{b_j} x f(x) dx / \int_{b_{j-1}}^{b_j} f(x) dx$$

Consider a random variable **uniformly distributed** between 0 and 2.
Let us use the Lloyd-Max algorithm to find the 2-bit quantizer that minimizes the MSE.

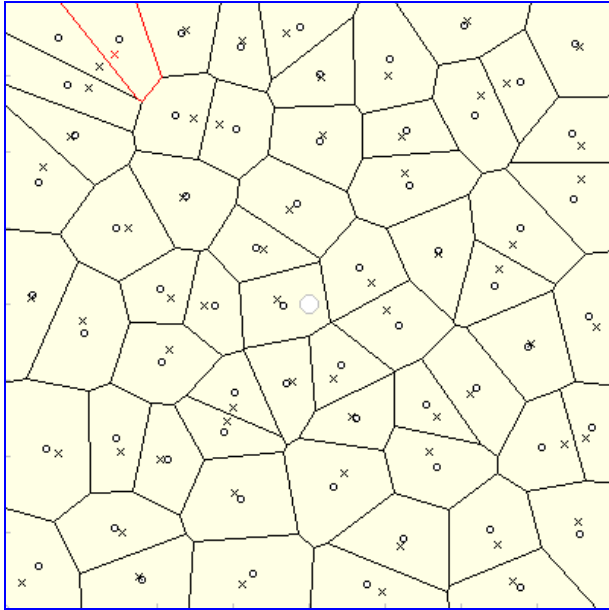
$$f(x) = \text{const}$$

$$a_j = \int_{b_{j-1}}^{b_j} x f(x) dx / \int_{b_{j-1}}^{b_j} f(x) dx = \frac{1}{2} (b_j^2 - b_{j-1}^2) / (b_j - b_{j-1}) = \frac{b_j + b_{j-1}}{2}$$

$$b_0 = 0, \quad b_1 = 0.5, \quad b_2 = 1, \quad b_3 = 1.5, \quad b_4 = 2$$

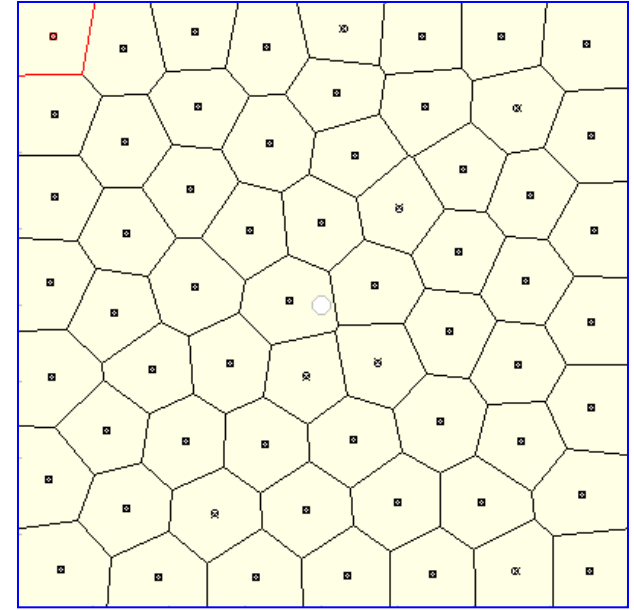
$$a_1 = 0.25, \quad a_2 = 0.75, \quad a_3 = 1.25, \quad a_4 = 1.75$$

The Lloyd-Max algorithm can be used in the multidimensional case.



Voronoi diagram / tessellation

Lloyd-Max
algorithm



Centroidal Voronoi diagram

A nice Java

applet: <http://www.raskob.de/fun/d/index-en.php>