

CS323 Assignment1 Answer

Exercise 1

- Yes.
- Yes.
- $5 + 4 + 3 + 2 = 14$
- $2^5 = 32$
- No.

Exercise 2

$\epsilon, abc, def, abcdef, defabc, abcdefabc, defabcdef, abcdefabcdef$

Exercise 3

$\epsilon, e, de, cde, bcde, abcde$

Exercise 4

regular definition: $valid \rightarrow [0 - 9][1 - 9][0 - 9]1[0 - 9][0 - 9]2[0 - 4][0 - 9]25[0 - 5]$

regular definition for a valid IPv4 address: $valid. valid. valid. valid$

Exercise 5

Yes.

$$L_1 = L((ab)^*ac) = \{(ab)^n ac | n \geq 0\}$$

$$L_2 = L(a(ba)^*c) = \{a(ba)^n c | n \geq 0\}$$

We aim to prove that for all $n \geq 0$, the following equation holds:

$$(ab)^n ac = a(ba)^n c$$

Mathematical induction:

Base case:

For $n = 0$:

$$(ab)^0 ac = ac = a(ba)^0 c$$

Inductive Hypothesis:

Assume that for some $i \geq 0$, the statement holds. That is, we assume:

$$(ab)^i ac = a(ba)^i c$$

Inductive Step:

$$(ab)^{i+1} ac = ab(ab)^i ac = aba(ba)^i c = a(ba)^{i+1} c$$

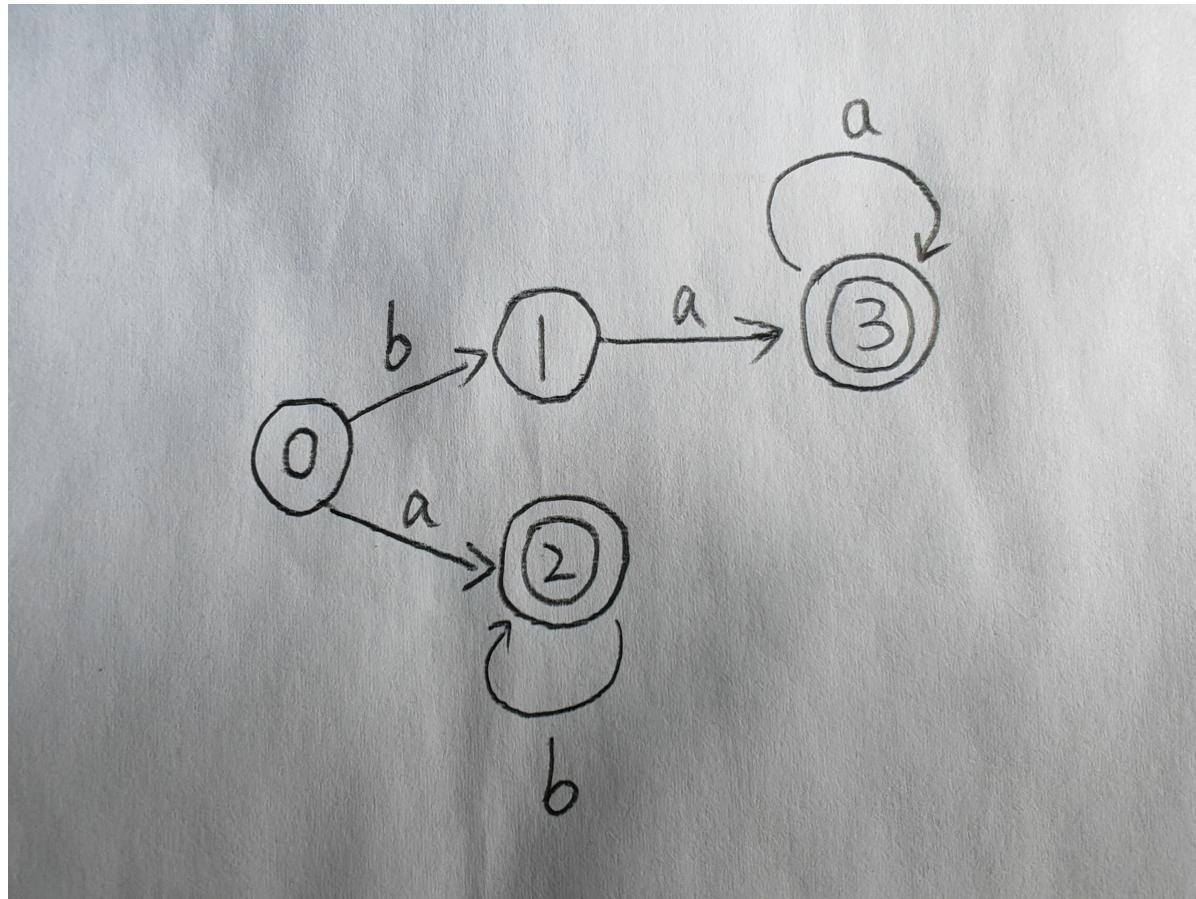
By the principle of mathematical induction, since the base case holds and the inductive step has been proven, we conclude that for all $n \geq 0$:

$$(ab)^n ac = a(ba)^n c$$

Therefore, the two regular languages L_1 and L_2 are equivalent.

Exercise 6

state transition diagram:



No.