# Lecture 5 Stack and Queue

### Our Roadmap

Stack

Queue

Stack vs. Queue

### Stack

- A stack is a sequence in which:
  - Items can be added and removed only at one end (the top)
  - You can only access the item that is currently at the top
- Stack Analogy

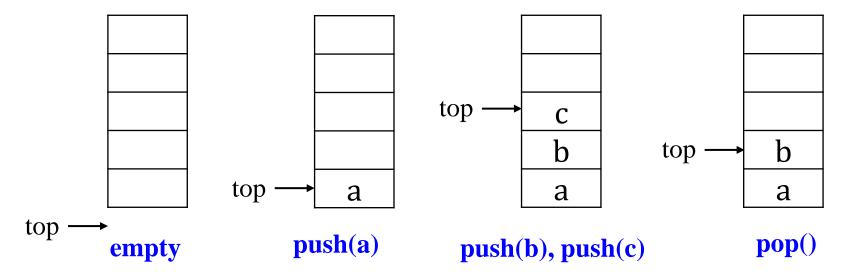






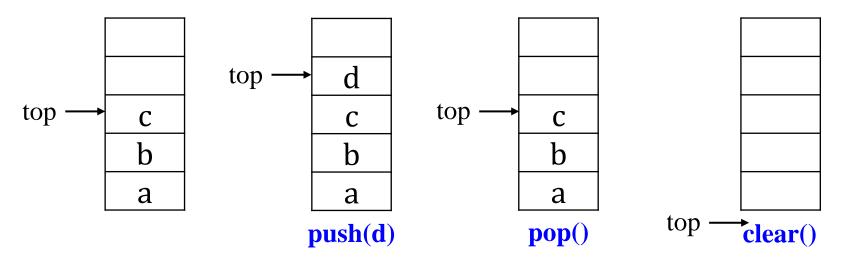
### Stack

- First In Last Out (FILO)
  - Constrained item access
- Major Operations
  - push: add an item to the top of the stack
  - pop: remove the item at the top of the stack
- Illustration



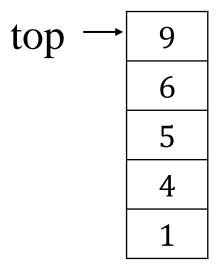
# Stack Operation

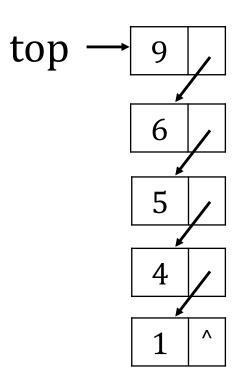
- push: add an item to the top of the stack
- pop: remove the item at the top of the stack
- top/peek: get the item at the top of the stack, but do not remove it
- isEmpty: test if the stack is empty
- isFull: test if the stack is full
- clear: clear the stack, set it as empty stack
- size: return the current size of the stack



### Implementation of Stack

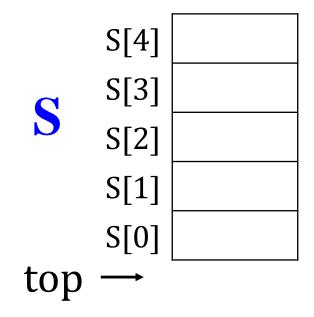
- Array-based Stack
- Linked Stack



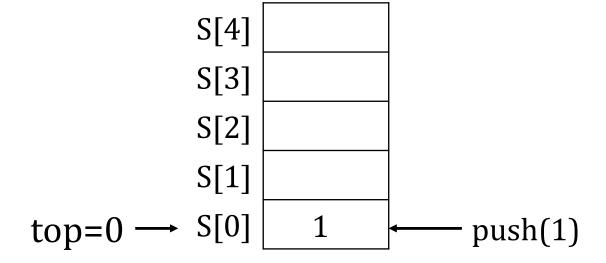


### Implementation of Stack

- Array based Stack
  - MAX\_SIZE = n // the max size of stack
  - $\bullet$  top = -1 // the current top position
  - Array S with n elements
- Example
  - ♦ MAX\_SIZE = 5
  - top = -1
  - Array S



- push(item):
  - 1. top++;
  - 2. S[top] = item
- push(1)



- push(item):
   1. top++;
   2. S[top] = item

  push(1)
  push(4)

  S[4]

push(item): 1. top++; 2. S[top] = itempush(1) push(4) S[4] push(5) S[3]  $top=2 \longrightarrow S[2]$ push(5) $top=1 \longrightarrow S[1]$ push(4)4  $top=0 \longrightarrow S[0]$ push(1)

push(item): 1. top++; 2. S[top] = itempush(1) push(4) S[4] push(5)  $top=3 \longrightarrow S[3]$ 6 push(6)push(6)  $top=2 \longrightarrow S[2]$ push(5)5  $top=1 \longrightarrow S[1]$ push(4)4  $top=0 \longrightarrow S[0]$ 

push(1)

top=5

push(9)

top=4 
$$\rightarrow$$
 S[4] 9  $\leftarrow$  push(9)  
S[3] 6  
S[2] 5  
S[1] 4  
S[0] 1

- push(10)
  - OVERFLOW
  - How to avoid that?

S[4]	9
S[3]	6
S[2]	5
S[1]	4
S[0]	1

### Push / Pop Operator

push(item):
 1. if(top == MAXSIZE-1)
 2. Stack is FULL! No push!
 3. else
 4. top++;
 5. S[top] = item

```
pop(): // should avoid underflow
1. if(top == -1)
2.    Stack is EMPTY! No pop!
3. else
4.    top--;
```

### Application of Stacks

- Making sure the delimiters (parens, brackets, etc.) are balanced:
  - Push open (i.e., left) delimiters onto a stack
  - When you encounter a close (i.e., right) delimiter, pop an item off the stack and see if it matches
- Evaluating arithmetic expressions
  - Parsing arithmetic expressions written using infix notation
- The runtime stack in memory
  - Converting a recursive algorithm to an iterative one by using a stack to emulate the runtime stack

### Brackets Balance Problem

- $\bullet$  a+{2-[b+c]\*(8\* [8+g]/[m-e]-7)-p}
- \* {[]([][])}
- Skip operators and notations
- Is the bracket expression balanced or not?
  - () Yes
  - ♦ No
  - ♦ { [()} No

### Brackets Balance Problem

- Given a bracket expression, determine whether it is balanced or not?
- \* {[]([][])}
  - Mow to solve it by using stack?
  - Bracket pairs: ( ), [ ], { }
  - Any ideas?
- Methodology
  - Employ stack store checked left bracket
  - Pop out left bracket if it is matched

# Arithmetic Expression Evaluation

- Arithmetic expression
  - operands (a, b, c), operator (+, \*)
  - a + b \* c
- Prefix expression
  - ♦ + a \* b c
- Infix expression
  - $\diamond$  a + b \* c
- Postfix expression
  - a b c \* +

### Postfix Expression

- Infix expression
  - 5 \* ((9 + 3) \* (4\*2) + 7)
- Postfix expression
  - § 5 9 3 + 4 2 \* \* 7 + \*
- Parse postfix expression is somewhat easier problem than directly parsing infix (why)
- Postfix has a nice property that parentheses are unnecessary
- Postfix Expression Evaluation
  - Convert from infix to postfix
  - Evaluate a postfix expression

### Postfix Expression

- Postfix expression
  - ⋄ 5 9 3 + 4 2 \* \* 7 + \*
- Methodology
  - Read the tokens in one at a time
  - If it is an operand, push it on the stack
  - If it is a binary operator:
    - pop top two elements from the stack,
    - apply the operator,
    - and push the result back on the stack

# Postfix Expression Evaluation

- § 593+42\*\*7+\*
- Postfix Expression Evaluation

#### Stack operations

- push(5)
- push(9)
- push(3)
- push(pop() + pop())
- push(4)
- push(2)
- push(pop() \* pop())
- push(pop() \* pop())
- push(7)
- push(pop() + pop())
- push(pop() \* pop())

#### Stack elements

5

59

593

5 12

5 12 4

5 12 4 2

5 12 8

5 96

5 96 7

5 103

515

### Our Roadmap

Stack



Queue

Stack vs. Queue

### Queue

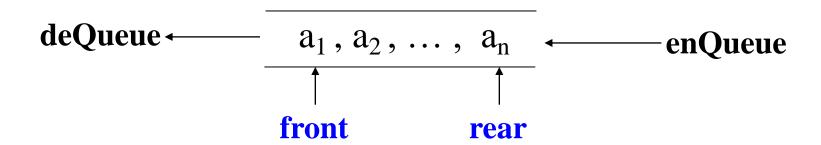
- A queue is a sequence in which:
  - items are added at the rear and removed from the front
  - You can only access the item that is currently at the front
- Queue Analogy





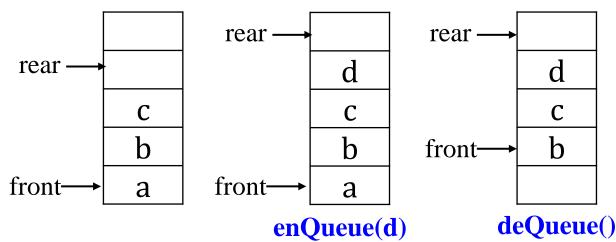
### Queue

- First In First Out (FIFO)
  - Items access constrained
- Major elements
  - front: the first element in the queue (remove)
  - rear: the last element in the queue (add)
- Illustration



# Queue Operations

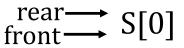
- enQueue: add an item at the rear of the queue
- deQueue: remove the item at the front of the queue
- front: get the item at the front of the queue, but do not remove it
- isEmpty: test if the queue is empty
- isFull: test the queue is full
- clear: clear the queue, set it as empty queue
- size: return the current size of the queue



### Implementation of Queue

- Array based Queue
  - MAX\_SIZE = n // the max size of stack
  - front = 0 // the current front
  - rear = 0 // the current rear
  - Array S with n elements
- Example
  - $\phi$  MAX\_SIZE = 5
  - front = 0
  - $\Rightarrow$  rear = 0
  - Array S

S[4]	

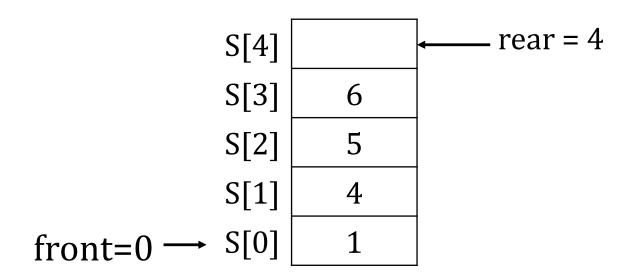


### enQueue Operator

enQueue(item):

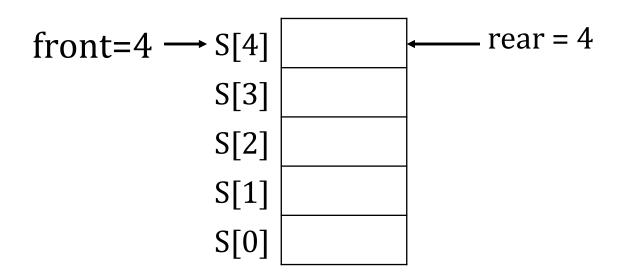
```
    if(rear < MAXSIZE)</li>
    S[rear] = item
    rear++
```

- 3. else
- 4. Queue is FULL, no enQueue
- enQueue (1), enQueue(4), enQueue(5), enQueue(6)



### deQueue Operator

- deQueue():
  - 1. if(front < rear)</pre>
  - 2. front++
  - 3. else
  - 4. Queue is empty, no deQueue
- deQueue(), deQueue(), deQueue()



### enQueue and deQueue

enQueue(9)

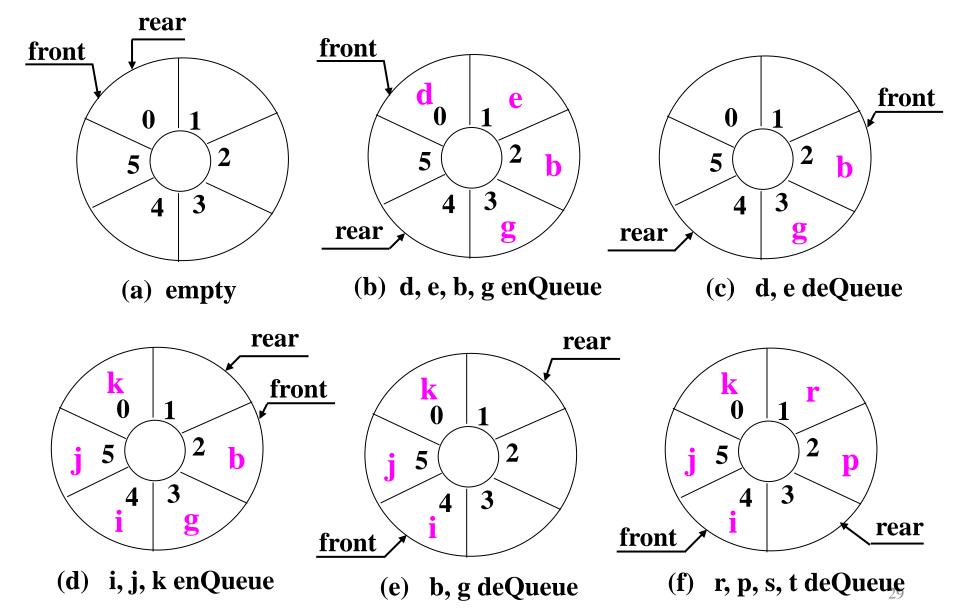
- enQueue(10)
  - \* rear >= MAXSIZE
  - Queue is FULL!!!
  - Wrong OVERFLOW
  - S[0] to S[3] is empty?
  - How to address it?



rear = 5

- front= $4 \rightarrow S[4]$  9
  - S[3]
    - S[2]
    - S[1]
    - S[0]

# Ring Queue



### Application of Queues

- First-in first-out (FIFO) inventory control
  - OS scheduling: processes, print jobs, packets, etc.
  - Breadth-first traversal of a graph or level-order traversal of a binary tree (more on these later)
- Real applications
  - iTunes playlist.
  - Data buffers (iPod, TiVo).
  - Asynchronous data transfer (file IO, pipes, sockets).
  - Dispensing requests on a shared resource (printer, processor)

### Our Roadmap

Stack

Queue

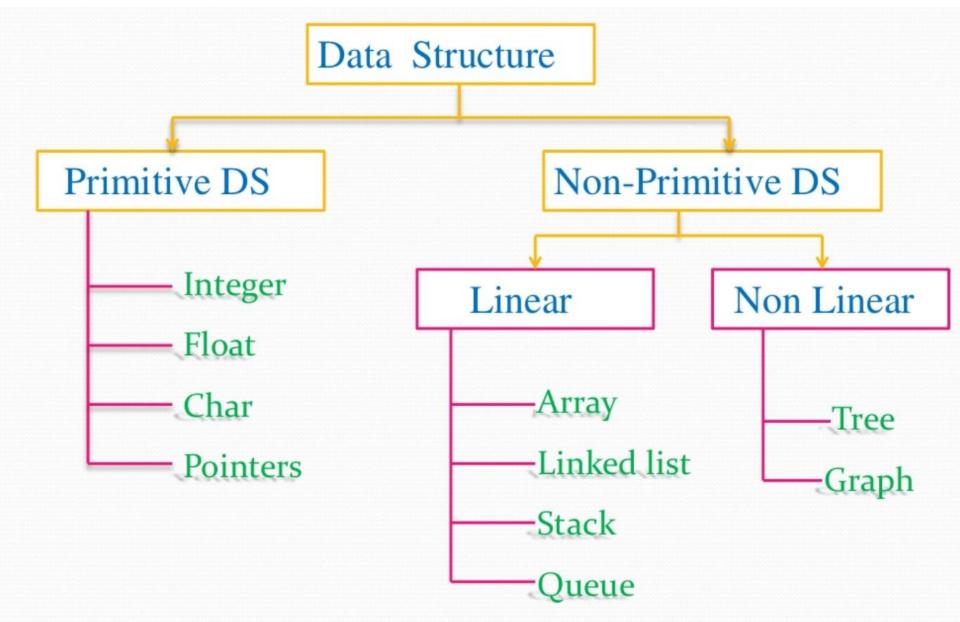


Stack vs. Queue

# Stack VS. Queue

	Stack	Queue
In-Out	FILO	FIFO
Application	function runtime	OS scheduling
Operations	push pop	enQueue, deQueue
Ops Time Complexity	0(1)	0(1)
Implementa tion	Array-based, Linked-based	Array-based, Linked-based

### Data Structure

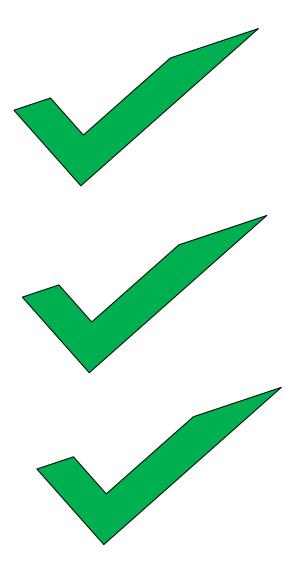


### Our Roadmap

Stack

Queue

Stack vs. Queue



### Thank You!