Meta-Circular Evaluators

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May 24, 2021

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Meta-Circular Evaluators

Problem

- Programming languages are designed to facilitate the description of certain kinds of problems.
- Not all problems can be easily described in the current programming languages.

Solution: Linguistic Abstraction

 A new programming language is invented to facilitate the description of the problem.

Evaluator

An evaluator (or interpreter) of a language is a procedure that, when applied to an expression of that language, computes the operations described in that expression.

Fundamental Ideia

An evaluator is a program that computes the meaning of a program.

Meta-Circular Evaluator

An evaluator written in the language that it evaluates.

Abstract Syntax

Our evaluator will operate in terms of an *abstract syntax*: the syntax is recognized by predicates that abstract from the concrete syntax used.

Example: Concrete Syntax for Assignments

```
a := b
a = b
a \leftarrow b
(set! a b)
```

Example: Abstract Syntax for Assignments

```
(define (set? expression)
  (eq? (car expression) 'set!))
```

Scheme Evaluator

(define (eval exp)

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (self-evaluating? exp)
  (or (number? exp) (string? exp)))
(define (repl)
  (prompt-for-input)
  (let ((input (read)))
    (let ((output (eval input)))
      (print output)))
  (repl))
(define (prompt-for-input)
  (display ">> "))
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (self-evaluating? exp)
  (or (number? exp) (string? exp)))
(define (repl)
  (prompt-for-input)
  (let ((input (read)))
    (let ((output (eval input)))
      (print output)))
  (repl))
(define (prompt-for-input)
  (display ">> "))
(define (print object)
  (write object)
  (newline))
```

Example >

Example

> (repl)

Example

> (repl)

>>

Example

- > (repl)
- >> 1

Example

```
> (repl)
```

>> 1 1

Example

```
> (repl)
```

>> 1 1 >> 2

Example

> (repl)

Example

```
> (repl)
>> 1
1
>> 2
2
>> "Hello"
```

Example

```
> (repl)
>> 1
1
>> 2
2
>> "Hello"
"Hello"
```

Example > (repl)

```
>> 1
1
>> 2
2
>> "Hello"
"Hello"
>> (+ 1 2)
```

```
Example
> (repl)
>> 1
1
>> 2
>> "Hello"
"Hello"
>> (+ 1 2)
error in "Unknown expression type -- EVAL": (+ 1 2)
backtrace:
   (eval input)
  ... "eval.scm" line 14
  1 (repl)
  ..."/dev/stdin" line 1
>
```

```
Example
> (repl)
>> 1
1
>> 2
>> "Hello"
"Hello"
>> (+ 1 2)
error in "Unknown expression type -- EVAL": (+ 1 2)
backtrace:
  0 (eval input)
  ... "eval.scm" line 14
  1 (repl)
  ..."/dev/stdin" line 1
> (repl)
```

```
Example
> (repl)
>> 1
1
>> 2
>> "Hello"
"Hello"
>> (+ 1 2)
error in "Unknown expression type -- EVAL": (+ 1 2)
backtrace:
  0 (eval input)
  ... "eval.scm" line 14
  1 (repl)
  ..."/dev/stdin" line 1
> (repl)
>>
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        ((addition? exp)
         (+ (first-operand exp)
            (second-operand exp)))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (addition? exp)
  (and (pair? exp)
       (eq? (car exp) 'plus)))
```

Scheme Evaluator (define (eval exp) (cond ((self-evaluating? exp) exp) ((addition? exp) (+ (first-operand exp) (second-operand exp))) (else (error "Unknown expression type -- EVAL" exp))))

```
(eq? (car exp) 'plus)))
(define (first-operand exp)
  (cadr exp))
```

(define (addition? exp) (and (pair? exp)

```
(define (second-operand exp)
  (caddr exp))
```

Example > (repl)

Example

- > (repl)
- >> (plus 1 2)

```
Example
```

```
> (repl)
>> (plus 1 2)
3
```

```
Example
```

```
> (repl)
>> (plus 1 2)
3
>> (plus (plus 1 2) (plus 3 4))
```

```
> (repl)
>> (plus 1 2)
>> (plus (plus 1 2) (plus 3 4))
error in +: expected number, but got (plus 1 2), as argument 1
irritants:
  ((plus 1 2) (plus 3 4))
backtrace:
  0 (+ (first-operand exp) (second-operand exp))
  ... "eval.scm" line 31
  1 (eval input)
  ... "eval.scm" line 12
  2 (repl)
  ..."/dev/stdin" line 1
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        ((addition? exp)
         (+ (eval (first-operand exp))
            (eval (second-operand exp))))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```



```
>> (plus 1 2)
3
```



```
>> (plus 1 2)
3
>> (plus (plus 1 2) (plus 3 4))
```


(error "Unknown expression type -- EVAL" exp))))

```
>> (plus 1 2)
3
>> (plus (plus 1 2) (plus 3 4))
10
```

Scheme Evaluator (define (eval exp) (cond ((self-evaluating? exp) exp) ((addition? exp) (+ (eval (first-operand exp)) (eval (second-operand exp)))) (else (error "Unknown expression type -- EVAL" exp))))

```
>> (plus 1 2)
3
   (plus (plus 1 2) (plus 3 4))
10
>> (times 2 3)
```

```
>> (plus 1 2)
3
>> (plus (plus 1 2) (plus 3 4))
10
>> (times 2 3)
error in "Unknown expression type -- EVAL": (times 2 3)
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        ((addition? exp)
         (+ (eval (first-operand exp))
            (eval (second-operand exp))))
        ((product? exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp)
  (cond ((self-evaluating? exp) exp)
        ((addition? exp)
         (+ (eval (first-operand exp))
            (eval (second-operand exp))))
        ((product? exp)
         (* (eval (first-operand exp))
            (eval (second-operand exp))))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
>> (times 2 3) 6
```

Example

```
>> (times 2 3)
6
```

Improved Syntax

Example

```
>> (times 2 3)
6
```

Improved Syntax

```
>> (+ 2 3)
5
>> (* (+ 2 3) (+ 4 5))
45
```



Naming

- Naming is the fundamental abstraction operator.
- A name abstracts a value and creates a scope where the name has a meaning.
- A name is not a first class value but denotes a first class value.
- Names are maintained in an environment: a collection of names.
- The environment is implemented as an association list:
 ((name₀ . value₀) (name₁ . value₁) ... (name_n . value_n))

```
>> (let ((pi 3.14159) (radius 5))
    (* 2 (* pi radius)))
31.4159
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Environment ;;Example:

```
;;(let ((pi 3.14159) (radius 5));; (* 2 (* pi radius)))
```

```
;;Example:
;;(let ((pi 3.14159) (radius 5))
;; (* 2 (* pi radius)))
(define (name? exp)
  (symbol? exp))
```

```
Environment
;;Example:
;;(let ((pi 3.14159) (radius 5))
;; (* 2 (* pi radius)))
(define (name? exp)
  (symbol? exp))
(define (let? exp)
  (and (pair? exp)
       (eq? (car exp) 'let)))
(define (let-names exp)
  (map car (cadr exp)))
(define (let-inits exp)
  (map cadr (cadr exp)))
```

```
Environment
;; Example:
;;(let ((pi 3.14159) (radius 5))
;; (* 2 (* pi radius)))
(define (name? exp)
  (symbol? exp))
(define (let? exp)
  (and (pair? exp)
       (eq? (car exp) 'let)))
(define (let-names exp)
  (map car (cadr exp)))
(define (let-inits exp)
  (map cadr (cadr exp)))
(define (let-body exp)
  (caddr exp))
```

```
;;Environment implementation:
;;((name-0 . value-0) (name-1 . value-1) ... (name-n . value-n))
```

```
;; Environment implementation:
::((name-0 . value-0) (name-1 . value-1) ... (name-n . value-n))
(define (eval-name name env)
  (cond ((null? env)
         (error "Unbound name -- EVAL-NAME" name))
        ((eq? name (caar env))
         (cdar env))
        (else
         (eval-name name (cdr env)))))
(define (eval-let exp env)
  (let ((values (eval-exprs (let-inits exp) env)))
    (let ((extended-environment
           (augment-environment (let-names exp)
                                 values
                                 env)))
      (eval (let-body exp) extended-environment))))
```

Environment

```
(define (eval-exprs exprs env)
  (if (null? exprs)
      (list)
      (cons (eval (car exprs) env)
            (eval-exprs (cdr exprs) env))))
;;(let ((x 1) (y (+ 1 1))) ...)
;;=> (augment-environment '(x y) '(1 2) '((z . 3)))
;; <= ((x . 1) (y . 2) (z . 3))
(define (augment-environment names values env)
  (if (null? names)
      env
      (cons (cons (car names) (car values))
            (augment-environment (cdr names) (cdr values) env))))
```

Environment

```
(define (eval-exprs exprs env)
  (if (null? exprs)
      (list)
      (cons (eval (car exprs) env)
            (eval-exprs (cdr exprs) env))))
;;(let ((x 1) (y (+ 1 1))) ...)
;;=> (augment-environment '(x y) '(1 2) '((z . 3)))
;; <= ((x . 1) (y . 2) (z . 3))
(define (augment-environment names values env)
  (if (null? names)
      env
      (cons (cons (car names) (car values))
            (augment-environment (cdr names) (cdr values) env))))
(define empty-environment (list))
```

REPL (define (repl) (prompt-for-input) (let ((input (read))) (let ((output (eval input empty-environment))) (print output))) (repl))


```
>> (let ((pi 3.14159) (radius 5))
    (* 2 (* pi radius)))
31.4159
```

```
>> (let ((pi 3.14159) (radius 5))
          (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
=> (eval 3.14159 ())
```

```
>> (let ((pi 3.14159) (radius 5))
    (* 2 (* pi radius)))

=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
=> (eval 3.14159 ())
<= 3.14159
```

```
>> (let ((pi 3.14159) (radius 5))
          (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
=> (eval 3.14159 ())
<= 3.14159
=> (eval 5 ())
<= 5
=> (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
```

```
>> (let ((pi 3.14159) (radius 5))
     (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
    (eval 3.14159 ())
=>
<= 3.14159
    (eval 5 ())
=>
<=
=>
     (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
       (eval 2 ((pi . 3.14159) (radius . 5)))
=>
<=
=>
       (eval (* pi radius) ((pi . 3.14159) (radius . 5)))
```

```
>> (let ((pi 3.14159) (radius 5))
     (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
   (eval 3.14159 ())
=>
<= 3.14159
=>
   (eval 5 ())
<=
=>
     (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
       (eval 2 ((pi . 3.14159) (radius . 5)))
=>
<=
=>
       (eval (* pi radius) ((pi . 3.14159) (radius . 5)))
         (eval pi ((pi . 3.14159) (radius . 5)))
=>
         3.14159
<=
=>
         (eval radius ((pi . 3.14159) (radius . 5)))
<=
```

```
>> (let ((pi 3.14159) (radius 5))
     (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
   (eval 3.14159 ())
=>
<= 3.14159
=>
   (eval 5 ())
<=
=>
     (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
       (eval 2 ((pi . 3.14159) (radius . 5)))
=>
<=
=>
       (eval (* pi radius) ((pi . 3.14159) (radius . 5)))
         (eval pi ((pi . 3.14159) (radius . 5)))
=>
         3.14159
<=
=>
         (eval radius ((pi . 3.14159) (radius . 5)))
<=
<=
      15.70795
```

```
>> (let ((pi 3.14159) (radius 5))
     (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
   (eval 3.14159 ())
=>
<= 3.14159
=>
   (eval 5 ())
<=
=>
     (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
       (eval 2 ((pi . 3.14159) (radius . 5)))
=>
<=
=>
       (eval (* pi radius) ((pi . 3.14159) (radius . 5)))
         (eval pi ((pi . 3.14159) (radius . 5)))
=>
         3.14159
<=
=>
         (eval radius ((pi . 3.14159) (radius . 5)))
<=
<=
      15.70795
<=
    31.4159
```

```
>> (let ((pi 3.14159) (radius 5))
     (* 2 (* pi radius)))
=> (eval (let ((pi 3.14159) (radius 5)) (* 2 (* pi radius))) ())
   (eval 3.14159 ())
=>
<= 3.14159
=>
   (eval 5 ())
<=
=>
    (eval (* 2 (* pi radius)) ((pi . 3.14159) (radius . 5)))
=>
       (eval 2 ((pi . 3.14159) (radius . 5)))
<=
=>
       (eval (* pi radius) ((pi . 3.14159) (radius . 5)))
         (eval pi ((pi . 3.14159) (radius . 5)))
=>
        3.14159
<=
=>
         (eval radius ((pi . 3.14159) (radius . 5)))
<=
<=
    15.70795
<=
    31.4159
<= 31.4159
```

Example

>> (let ((a 1))

```
(let ((b (+ a 2)))
       (let ((a (+ a b)))
(* 2 a))))
=> (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
=> (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
=> (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
```

```
>> (let ((a 1))
      (let ((b (+ a 2)))
            (let ((a (+ a b)))
            (* 2 a))))

=> (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
=> (eval 1 ()) <= 1
=> (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=> (eval (+ a 2) ((a . 1)))
=> (eval a ((a . 1))) <= 1
=> (eval 2 ((a . 1))) <= 2</pre>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
=> (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
=>
     (eval 1 ()) <= 1
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
=>
         (eval (+ a b) ((b . 3) (a . 1)))
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) <= 1
=>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
=>
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) <= 1
=>
           (eval b ((b . 3) (a . 1))) \le 3
=>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
=>
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) \le 1
=>
           (eval b ((b . 3) (a . 1))) <= 3
=>
<=
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) \le 1
=>
           (eval b ((b . 3) (a . 1))) <= 3
=>
         4
<=
         (eval (* 2 a) ((a . 4) (b . 3) (a . 1)))
=>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) <= 1
=>
           (eval b ((b . 3) (a . 1))) \le 3
=>
         4
<=
         (eval (* 2 a) ((a . 4) (b . 3) (a . 1)))
=>
           (eval 2 ((a . 4) (b . 3) (a . 1))) <= 2
=>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (* 2 a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) \le 1
=>
           (eval b ((b . 3) (a . 1))) \le 3
=>
         4
<=
         (eval (* 2 a) ((a . 4) (b . 3) (a . 1)))
=>
           (eval 2 ((a . 4) (b . 3) (a . 1))) <= 2
=>
           (eval a ((a . 4) (b . 3) (a . 1))) \le 4
=>
```

```
>> (let ((a 1))
     (let ((b (+ a 2)))
       (let ((a (+ a b)))
         (*2a))))
   (eval (let ((a 1)) (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a)))) ())
     (eval 1 ()) <= 1
=>
     (eval (let ((b (+ a 2))) (let ((a (+ a b))) (* 2 a))) ((a . 1)))
=>
       (eval (+ a 2) ((a . 1)))
=>
         (eval a ((a . 1))) <= 1
=>
         (eval 2 ((a . 1))) <= 2
=>
<=
=>
       (eval (let ((a (+ a b))) (* 2 a)) ((b . 3) (a . 1)))
         (eval (+ a b) ((b . 3) (a . 1)))
=>
           (eval a ((b . 3) (a . 1))) \le 1
=>
           (eval b ((b . 3) (a . 1))) \le 3
=>
<=
         4
         (eval (* 2 a) ((a . 4) (b . 3) (a . 1)))
=>
           (eval 2 ((a . 4) (b . 3) (a . 1))) <= 2
=>
           (eval a ((a . 4) (b . 3) (a . 1))) <= 4
=>
<=
<=
       8
<=
     8
<= 8
```

Pre-defined Names

- The REPL starts in a clean state.
- Forcing us to introduce all needed names.
- Make it better: we can provide some pre-defined names.

Pre-defined Names

- The REPL starts in a clean state.
- Forcing us to introduce all needed names.
- Make it better: we can provide some pre-defined names.

Initial Bindings

```
(define initial-bindings
(list (cons 'pi 3.14159)
(cons 'e 2.71828)
...))
```

Pre-defined Names

- The REPL starts in a clean state.
- Forcing us to introduce all needed names.
- Make it better: we can provide some pre-defined names.

Initial Bindings

REPL (define (repl) (prompt-for-input) (let ((input (read))) (let ((output (eval input empty-environment))) (print output))) (repl))

REPL (define (repl) (prompt-for-input) (let ((input (read))) (let ((output (eval input initial-environment))) (print output))) (repl))

REPL

```
(define (repl)
  (prompt-for-input)
  (let ((input (read)))
     (let ((output (eval input initial-environment)))
        (print output)))
     (repl))
```

```
>> pi
```

```
>> pi
3.14159
```

```
>> pi
3.14159
>> (+ pi e)
5.85987
```

```
REPL
(define (repl)
  (prompt-for-input)
  (let ((input (read)))
      (let ((output (eval input initial-environment)))
            (print output)))
      (repl))
```

REPL

```
(define (repl)
  (prompt-for-input)
  (let ((input (read)))
      (let ((output (eval input initial-environment)))
         (print output)))
      (repl))
```

```
Example

>> pi
3.14159
>> (+ pi e)
5.85987
>> (let ((e 10))
        (* pi e))
31.4159
>> (* (let ((e 10)) (* e e)) e)
```

```
REPL

(define (repl)
   (prompt-for-input)
   (let ((input (read)))
      (let ((output (eval input initial-environment)))
            (print output)))
      (repl))
```

Example: Add the areas of two circles with radius 10 and 100

Example: Add the areas of two circles with radius 10 and 100

- So far, we only have abstracted *values*.
- But it is also good to have abstracted expressions.
- Abstract expressions are known as functions.

Example: Add the areas of two circles with radius 10 and 100

- Abstracting an expression allow us to give it a name and to reuse it in several different places just by naming it.
- We must distinguish between names that are bound in the surrounding context and names that are bound in the function call (e.g., pi vs radius inside area-circle).
- The distinction is made by providing a parameter list.

- New naming operator: flet
 - Augment the environment to contain associations between the names and the functions.
- New call operator: call
 - Obtain the function associated with the name.
 - Evaluate all argument expressions.
 - Augment the environment to contain the associations between function parameters and the evaluated arguments.
 - Evaluate the function body in the augmented environment.

```
;(call foo 1 (+ 2 3) 4)
```

```
Functions
;(call foo 1 (+ 2 3) 4)
(define (call? exp)
  (and (pair? exp)
       (eq? (car exp) 'call)))
(define (call-operator exp)
  (cadr exp))
(define (call-operands exp)
  (cddr exp))
;(flet ((square (x) (* x x))
     (foo (a b c) (+ a (* b c))))
   ...)
```

```
Functions
;(call foo 1 (+ 2 3) 4)
(define (call? exp)
  (and (pair? exp)
       (eq? (car exp) 'call)))
(define (call-operator exp)
  (cadr exp))
(define (call-operands exp)
  (cddr exp))
;(flet ((square (x) (* x x))
      (foo (a b c) (+ a (* b c))))
(define (flet? exp)
  (and (pair? exp)
       (eq? (car exp) 'flet)))
```

```
Functions
;(flet ((square (x) (* x x))
        (foo (a b c) (+ a (* b c))))
(define (flet-names exp)
  (map car (cadr exp)))
(define (flet-functions exp)
  (map (lambda (f)
         (make-function (cadr f) (cddr f)))
       (cadr exp)))
```

```
Functions
;(flet ((square (x) (* x x))
        (foo (a b c) (+ a (* b c))))
(define (flet-names exp)
  (map car (cadr exp)))
(define (flet-functions exp)
  (map (lambda (f)
         (make-function (cadr f) (cddr f)))
       (cadr exp)))
(define (make-function parameters body)
  (cons 'function
        (cons parameters
              body)))
```

Functions ;(flet ((square (x) (* x x)) (foo (a b c) (+ a (* b c)))) ...) (define (flet-names exp) (map car (cadr exp))) (define (flet-functions exp) (map (lambda (f) (make-function (cadr f) (cddr f))) (cadr exp))) (define (make-function parameters body) (cons 'function (cons parameters body))) (define (flet-body exp) (caddr exp))

Flet Evaluation

Flet Evaluation

```
;; exp: (+ (call bar a) b)
;; env: ((bar . (function (x) (+ x 5))) ((a . 1) (b . 2)))
(define (function? obj)
  (and (pair? obj)
       (eq? (car obj) 'function)))
(define (function-parameters func)
  (cadr func))
(define (function-body func)
  (caddr func))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (let ((extended-environment
           (augment-environment (function-parameters func)
                                args
                                env)))
      (eval (function-body func) extended-environment))))
```

(else

(error "Unknown expression type -- EVAL" exp))))

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((flet? exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((flet? exp)
         (eval-flet exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval-let exp env)
  (let ((values (eval-exprs (let-inits exp) env)))
    (let ((extended-environment
           (augment-environment (let-names exp)
                                values
                                env)))
      (eval (let-body exp) extended-environment))))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (let ((extended-environment
           (augment-environment (function-parameters func)
                                args
                                env)))
      (eval (function-body func) extended-environment))))
```

```
(define (eval-let exp env)
  (let ((values (eval-exprs (let-inits exp) env)))
    (let ((extended-environment
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                                values
                                env)))
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(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (let ((extended-environment
           (augment-environment (function-parameters func)
                                args
                                env)))
      (eval (function-body func) extended-environment))))
```

```
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  (let ((values (eval-exprs (let-inits exp) env)))
    (let ((extended-environment
           (augment-environment (let-names exp)
                                values
                                env)))
      (eval (let-body exp) extended-environment))))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (let ((extended-environment
           (augment-environment (function-parameters func)
                                args
                                env)))
      (eval (function-body func) extended-environment))))
```

```
(define (eval-let exp env)
  (let ((values (eval-exprs (let-inits exp) env)))
    (let ((extended-environment
           (augment-environment (let-names exp)
                                values
                                env)))
      (eval (let-body exp) extended-environment))))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (let ((extended-environment
           (augment-environment (function-parameters func)
                                args
                                env)))
      (eval (function-body func) extended-environment))))
```

A Striking Resemblance

- The evaluation of a let form is very similar to the evaluation of a function call.
- In fact, any let form can be "translated" into a function that is immediately called.
- We will return to this subject later.

Translation Example

```
>> (let ((a 1) (b 2)) (+ a b))
```

A Striking Resemblance

- The evaluation of a let form is very similar to the evaluation of a function call.
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Translation Example

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A Striking Resemblance

- The evaluation of a let form is very similar to the evaluation of a function call.
- In fact, any let form can be "translated" into a function that is immediately called.
- We will return to this subject later.

Translation Example

Example: Add the areas of two circles with radius 10 and 100

```
>> (let ((radius-0 10) (radius-1 100))
     (flet ((square (x) (* x x)))
       (flet ((area-circle (radius)
                (* pi (call square radius))))
         (+ (call area-circle radius-0)
            (call area-circle radius-1))))
```

Example: Add the areas of two circles with radius 10 and 100

Example: Add the areas of two circles with radius 10 and 100

Problem

- There is a syntactic difference between the operators (e.g. (* 10 10)) and the functions (e.g. (call square 10))
- It will look much better if we use the same syntax (e.g. (square 10)).
- We just have to redefine the concrete syntax of function calls.

```
Evaluator
;;(square radius)
```

```
Evaluator
;;(square radius)
(define (call? exp)
  (pair? exp))
```

```
Evaluator
;;(square radius)
(define (call? exp)
  (pair? exp))
(define (call-operator exp)
  (car exp))
```

```
(;;(square radius)
(define (call? exp)
    (pair? exp))
(define (call-operator exp)
    (car exp))
(define (call-operands exp)
    (cdr exp))
```

```
Evaluator
;;(square radius)
(define (call? exp)
  (pair? exp))
(define (call-operator exp)
  (car exp))
(define (call-operands exp)
  (cdr exp))
;;Let's also add square as a pre-defined function:
```

```
;;(square radius)
(define (call? exp)
  (pair? exp))
(define (call-operator exp)
  (car exp))
(define (call-operands exp)
  (cdr exp))
;;Let's also add square as a pre-defined function:
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons 'e 2.71828)
        (cons 'square (make-function '(x) '((* x x))))))
```

Example

```
>> (square 10)
```

Example

```
>> (square 10)
100
>> (* (square (+ 1 2)) 3)
27
```

Example

>> (flet ((area-circle (radius)

31730.058999999997

(* pi (square radius)))) (+ (area-circle 10) (area-circle 100)))

Evaluator

Example >> (square 10) 100 >> (* (square (+ 1 2)) 3) 27

```
>> (flet ((square (a) "I'm a square")) (square 10))
```

```
>> (flet ((square (a) "I'm a square"))
        (square 10))
"I'm a square"
```

Example: Rebinding Names

Problem

- There is a semantic difference between operators (e.g. (+ 10 20)) and functions (e.g. (square 10))
- Names bound to functions can be shadowed.
- Names bound to operators cannot be shadowed.
- The problem can be solved by treating operators as being bound to special *primitive* operations.

```
Primitives
```

```
(define (make-primitive f)
  (list 'primitive f))
```

```
(define (make-primitive f)
  (list 'primitive f))
(define (primitive? obj)
  (and (pair? obj)
       (eq? (car obj) 'primitive)))
(define (primitive-operation prim)
  (cadr prim))
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons 'e 2.71828)
        (cons 'square (make-function '(x) '((* x x))))
        (cons '+ (make-primitive +))
        (cons '* (make-primitive *))))
```

```
(define (apply-primitive-function prim args)
  (apply (primitive-operation prim) args))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((addition? exp)
         (+ (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((product? exp)
         (* (eval (first-operand exp) env)
            (eval (second-operand exp) env)))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Example

>> (+ 3 4)

Example

>> (+ 3 4) 7

```
>> (+ 3 4)
>> (flet ((+ (x y) "I'm an addition")) (+ 10 20))
```

```
>> (+ 3 4)
>> (flet ((+ (x y) "I'm an addition"))
     (+1020))
"I'm an addition"
>> (+ 3 (flet ((+ (x y) (* x y))) (+ 4 5)))
23
>> (+ 1 2 3 4 5 6)
21
>> (let ((* 10))
    (+ * *)
20
>> (let ((+ *)
       (* +))
     (+ (* 1 2) 3))
```

23

21

20

9

>> (+ 1 2 3 4 5 6)

(* +)) (+ (* 1 2) 3))

>> (let ((* 10)) (+ * *))

>> (let ((+ *)


```
40 40 40 40 40 10 000
```

```
Example
>> (+ 3 4)
>> (flet ((+ (x y) "I'm an addition"))
     (+1020)
"I'm an addition"
>> (+ 3 (flet ((+ (x y) (* x y))) (+ 4 5)))
23
>> (+ 1 2 3 4 5 6)
21
>> (let ((* 10))
     (+ * *)
20
>> (let ((+ *)
        (* +))
     (+ (* 1 2) 3))
9
>> (flet ((+ (x y) (* x y))
      (* (x y) (+ x y)))
(+ (* 1 2) 3))
```

```
Example
>> (+ 3 4)
>> (flet ((+ (x y) "I'm an addition"))
     (+1020)
"I'm an addition"
>> (+ 3 (flet ((+ (x y) (* x y))) (+ 4 5)))
23
>> (+ 1 2 3 4 5 6)
21
>> (let ((* 10))
     (+ * *)
20
>> (let ((+ *)
        (* +))
     (+ (* 1 2) 3))
9
>> (flet ((+ (x y) (* x y))
      (* (x y) (+ x y)))
(+ (* 1 2) 3))
...infinite loop...
```

Problem

- Naming is an important abstraction tool.
- But there is more than simply attaching a name to something.
- Naming a value is different from naming a function (or is it not?).
- Using a name that abstracts a value is different from using a name that abstracts a function (or is it not?)
- Scope must also be considered.
- Shadowing must also be considered.
- Can a name abstract more than one thing at the same time?
 - He said that that that man said was false.
 - Buffalo buffalo buffalo buffalo.
- We will return to this subject later.

```
Primitives
(define initial-bindings
  (list (cons 'pi 3.14159)
          (cons 'e 2.71828)
          (cons 'square (make-function '(x) '((* x x))))
         (cons '+ (make-primitive +))
(cons '* (make-primitive *))
          ...))
```

More Primitives

```
Primitives
```

More Primitives

• We need more arithmetic operators ...

```
Primitives
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons 'e 2.71828)
        (cons 'square (make-function '(x) '((* x x))))
        (cons '+ (make-primitive +))
        (cons '* (make-primitive *))
        (cons '- (make-primitive -))
        (cons '/ (make-primitive /))
        ...))
```

More Primitives

• We need more arithmetic operators ...

Primitives

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons 'e 2.71828)
        (cons 'square (make-function '(x) '((* x x))))
        (cons '+ (make-primitive +))
        (cons '* (make-primitive *))
        (cons '- (make-primitive -))
              '/ (make-primitive /))
        ...))
```

More Primitives

- We need more arithmetic operators ...
- ...and also relational operators.



Primitives

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons 'e 2.71828)
        (cons 'square (make-function '(x) '((* x x))))
        (cons '+ (make-primitive +))
        (cons '* (make-primitive *))
        (cons '- (make-primitive -))
        (cons '/ (make-primitive /))
        (cons '= (make-primitive =))
        (cons '< (make-primitive <))</pre>
               '> (make-primitive >))
        (cons
               '<= (make-primitive <=))</pre>
        (cons
              '>= (make-primitive >=))
        ...))
```

More Primitives

- We need more arithmetic operators ...
- ...and also relational operators.



Relational Operators

Relational Operators

Relational Operators

#t

Relational Operators

```
>> (= (- 1 2) (- 3 4))
#t
>> (>= 2 5)
#f
```

Relational Operators

```
>> (= (- 1 2) (- 3 4))
#t
>> (>= 2 5)
#f
>> #t
```

Relational Operators

```
>> (= (- 1 2) (- 3 4))
#t
>> (>= 2 5)
#f
>> #t
error in "Unknown expression type -- EVAL": #t
```

Booleans

- Relational operators compute Boolean values that we pass to the evaluated language...
- ...but these are not (yet) self-evaluating forms.

Relational Operators

```
>> (= (- 1 2) (- 3 4))
#t
>> (>= 2 5)
#f
>> #t
error in "Unknown expression type -- EVAL": #t
```

Booleans

- Relational operators compute Boolean values that we pass to the evaluated language...
- ...but these are not (yet) self-evaluating forms.

Self-Evaluating Forms

```
(define (self-evaluating? exp)
  (or (number? exp) (string? exp)))
```



Relational Operators

```
>> (= (- 1 2) (- 3 4))
#t
>> (>= 2 5)
#f
>> #t
error in "Unknown expression type -- EVAL": #t
```

Booleans

- Relational operators compute Boolean values that we pass to the evaluated language...
- ...but these are not (yet) self-evaluating forms.

Self-Evaluating Forms

```
(define (self-evaluating? exp)
  (or (number? exp) (string? exp) (boolean? exp)))
```

Booleans

- What are the uses of boolean values? What are the uses of a data type that only has two members?
- Booleans are good for expressing a choice between two alternatives.
- Conditional Expressions allow such choices. They were invented in 1959 (by John McCarthy) and are available in all modern programming languages.

Conditional Expression Example

```
(if (< x y)
(+ x 5)
(* y 3))
```

```
;(if (< x y)
   (+ x 5)
(* y 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
```

```
;(if (< x y)
; (+ x 5)
; (* y 3))

(define (if? exp)
    (and (pair? exp) (eq? (car exp) 'if)))

(define (if-condition exp)
    (cadr exp))</pre>
```

```
; (if (< x y))
  (+ x 5)
  (* v 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
(define (if-condition exp)
  (cadr exp))
(define (if-consequent exp)
  (caddr exp))
(define (if-alternative exp)
  (cadddr exp))
```

```
; (if (< x y))
  (+ x 5)
  (* v 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
(define (if-condition exp)
  (cadr exp))
(define (if-consequent exp)
  (caddr exp))
(define (if-alternative exp)
  (cadddr exp))
(define (eval-if exp env)
```

```
; (if (< x y))
  (+ x 5)
  (* v 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
(define (if-condition exp)
  (cadr exp))
(define (if-consequent exp)
  (caddr exp))
(define (if-alternative exp)
  (cadddr exp))
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
```

```
; (if (< x y))
  (+ x 5)
  (* v 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
(define (if-condition exp)
  (cadr exp))
(define (if-consequent exp)
  (caddr exp))
(define (if-alternative exp)
  (cadddr exp))
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
```

```
; (if (< x y))
 (+ x 5)
  (* v 3))
(define (if? exp)
  (and (pair? exp) (eq? (car exp) 'if)))
(define (if-condition exp)
  (cadr exp))
(define (if-consequent exp)
  (caddr exp))
(define (if-alternative exp)
  (cadddr exp))
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
```

- The semantics of conditional expressions are directly inherited from the semantics of the underlying Scheme.
- But we can be more specific regarding the possible values of the condition.
- There are several possibilities:
 - The booleans are the only acceptable values (as in Pascal and Java).
 - Any value is accepted: some of them are considered false and all the other values are considered true (as in Python, Javascript and Perl).
 - Any value is accepted: one specific value is false and all the other values are considered true (as in C and Scheme).

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
```

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
;;Java
(define (true? value)
  (cond ((eq? value #t) #t)
        ((eq? value #f) #f)
        (else
         (error "Incorrect boolean value" value))))
;; Javascript
(define (true? value)
  (not (or (eq? value #f)
           (eqv? value 0)
           (eq? value ""))))
```

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
;;Java
(define (true? value)
  (cond ((eq? value #t) #t)
        ((eq? value #f) #f)
        (else
         (error "Incorrect boolean value" value))))
;; Javascript
(define (true? value)
  (not (or (eq? value #f)
           (eqv? value 0)
           (eq? value ""))))
;;Scheme
(define (true? value)
  (not (eq? value #f)))
```

Evaluator

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Evaluator

```
>> (flet ((abs (x)
             (if (< x 0)
               (- x)
               x)))
     (abs -5))
5
>> (flet ((fact (n)
             (if (= n 0)
               (* n (fact (- n 1))))))
     (fact 3))
6
>> (flet ((fact (n)
             (if (= n 0)
               (* n (fact (- n 1))))))
     (fact 50))
```

```
>> (flet ((abs (x)
            (if (< x 0)
               (-x)
              x)))
     (abs -5))
5
>> (flet ((fact (n)
            (if (= n 0)
               (* n (fact (- n 1))))))
     (fact 3))
6
>> (flet ((fact (n)
            (if (= n 0)
               (* n (fact (- n 1))))))
     (fact 50))
30414093201713378043612608166064768844377641568960512000000000000
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...)
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...)) => (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...))

=> (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1))))) ...))

=> (eval 3 ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))

<= 3
```

```
=> (eval (flat ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...))

=> (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1))))) ...))

=> (eval 3 ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1))))) ...))

<= 3

=> (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 3) (fact . \lambda) ...))
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...))
=> (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1))))) ...))
=> (eval 3 ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1))))) ...))

<= 3
=> (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 3) (fact . λ) ...))
=> (eval (= n 0) ((n . 3) (fact . λ) ...))
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...)) => (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))
       (eval 3 ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))
<=
       (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 3) (fact . \lambda) ...))
=>
        (eval (= n 0) ((n . 3) (fact . \lambda) ...))
=>
         (eval n ((n . 3) (fact . \lambda) ...))
=>
<=
=>
         (eval 0 ((n . 3) (fact . \lambda) ...))
<=
<=
=>
        (eval (* n (fact (- n 1))) ((n . 3) (fact . \lambda) ...))
=>
        (eval n ((n . 3) (fact . \lambda) ...))
<=
=>
         (eval (fact (- n 1)) ((n . 3) (fact . \lambda) ...))
          (eval (- n 1) ((n . 3) (fact . \lambda) ...))
=>
=>
           (eval n ((n . 3) (fact . \lambda) ...))
<=
=>
           (eval 1 ((n . 3) (fact . \lambda) ...))
<=
<=
          2
```

```
=> (eval (flet ((fact (n) (if (= n 0) 1 (* n (fact (- n 1)))))) (fact 3)) (...))
   (eval (fact 3) ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))
     (eval 3 ((fact . (function (n) (if (= n 0) 1 (* n (fact (- n 1)))))) ...))
<=
     (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 3) (fact . \lambda) ...))
=>
=>
       (eval (= n 0) ((n . 3) (fact . \lambda) ...))
        (eval n ((n . 3) (fact . \lambda) ...))
=>
<=
        (eval 0 ((n . 3) (fact . \lambda) ...))
=>
<=
<=
       (eval (* n (fact (- n 1))) ((n . 3) (fact . \lambda) ...))
=>
=>
       (eval n ((n . 3) (fact . \lambda) ...))
<=
=>
        (eval (fact (- n 1)) ((n . 3) (fact . \lambda) ...))
        (eval (- n 1) ((n . 3) (fact . \lambda) ...))
=>
=>
          (eval n ((n . 3) (fact . \lambda) ...))
<=
=>
          (eval 1 ((n . 3) (fact . \lambda) ...))
<=
<=
         (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 2) (n . 3) (fact . λ) ...))
=>
=>
          (eval (= n 0) ((n . 2) (n . 3) (fact . \lambda) ...))
           (eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
           (eval 0 ((n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
          #f
          (eval (* n (fact (- n 1))) ((n . 2) (n . 3) (fact . \lambda) ...))
=>
           (eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
           (eval (fact (- n 1)) ((n . 2) (n . 3) (fact . \lambda) ...))
=>
            (eval (- n 1) ((n . 2) (n . 3) (fact . \lambda) ...))
=>
```

```
=>
<=
           (eval n ((n . 2) (n . 3) (fact . \frac{\lambda}{\lambda}) ...))
           (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
         1
```

```
(eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
         (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
           0
```

```
(eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
=>
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
        (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
           (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
```

```
(eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
        (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
         #f
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...)
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
           (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n, 0) (n, 1) (n, 2) (n, 3) (fact, \lambda) ...))
=>
            (eval (= n 0) ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
```

```
(eval n ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
=>
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
        (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
         #f
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...)
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
           (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n, 0) (n, 1) (n, 2) (n, 3) (fact, \lambda) ...))
=>
            (eval (= n 0) ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
             (eval n ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
=>
             (eval 0 ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
            #t.
```

```
=>
         (eval n ((n . 2) (n . 3) (fact . \lambda) ...))
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
=>
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
        (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
         #f
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...)
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
           (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n, 0) (n, 1) (n, 2) (n, 3) (fact, \lambda) ...))
=>
            (eval (= n 0) ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
             (eval n ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
=>
             (eval 0 ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
            (eval 1 ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
```

```
=>
         (eval n ((n . 2) (n . 3) (fact . \lambda) ...))
<=
         (eval 1 ((n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
=>
        (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
        (eval (= n 0) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval 0 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
         #f
=>
         (eval (* n (fact (- n 1))) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
          (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
          (eval (fact (- n 1)) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...)
=>
=>
           (eval (- n 1) ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
            (eval n ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
=>
            (eval 1 ((n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
=>
           (eval (if (= n 0) 1 (* n (fact (- n 1)))) ((n, 0) (n, 1) (n, 2) (n, 3) (fact, \lambda) ...))
=>
            (eval (= n 0) ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
             (eval n ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
=>
             (eval 0 ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
<=
<=
            (eval 1 ((n . 0) (n . 1) (n . 2) (n . 3) (fact . \lambda) ...))
=>
<=
<=
<= 6
```

Definitions

- So far, our environments are functional: they are created, used and discarded.
- This means that the REPL has no memory of previous bindings.
- It would be good to have an incremental scope for names.

Definitions

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```
>> (def x (+ 1 2))
```

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```
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```
>> (def x (+ 1 2))
3
>> (+ x 2)
```

- So far, our environments are functional: they are created, used and discarded.
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```
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- So far, our environments are *functional*: they are created, used and discarded.
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```
>> (def x (+ 1 2))
3
>> (+ x 2)
5
>> (fdef triple (a) (+ a a a))
```

- So far, our environments are functional: they are created, used and discarded.
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```
>> (def x (+ 1 2))
3
>> (+ x 2)
5
>> (fdef triple (a) (+ a a a))
(function (a) (+ a a a))
```

- So far, our environments are functional: they are created, used and discarded.
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```
>> (def x (+ 1 2))
3
>> (+ x 2)
5
>> (fdef triple (a) (+ a a a))
(function (a) (+ a a a))
>> (triple (+ x 3))
```

Definitions

- So far, our environments are functional: they are created, used and discarded.
- This means that the REPL has no memory of previous bindings.
- It would be good to have an incremental scope for names.

```
>> (def x (+ 1 2))
3
>> (+ x 2)
5
>> (fdef triple (a) (+ a a a))
(function (a) (+ a a a))
>> (triple (+ x 3))
18
```

```
;(def x (+ 1 2))
```

```
;(def x (+ 1 2))
(define (def? exp)
   (and (pair? exp) (eq? (car exp) 'def)))
```

```
;(def (+ 1 2))
(define (def? exp)
  (and (pair? exp) (eq? (car exp) 'def)))
(define (def-name exp)
  (cadr exp))
```

```
(def x (+ 1 2))
(define (def? exp)
  (and (pair? exp) (eq? (car exp) 'def)))
(define (def-name exp)
  (cadr exp))
(define (def-init exp)
  (caddr exp))
```

```
; (def x (+ 1 2))
(define (def? exp)
  (and (pair? exp) (eq? (car exp) 'def)))
(define (def-name exp)
  (cadr exp))
(define (def-init exp)
  (caddr exp))
(define (eval-def exp env)
  (let ((value (eval (def-init exp) env)))
    (define-name! (def-name exp) value env)
    value))
```

```
; (def x (+ 1 2))
(define (def? exp)
  (and (pair? exp) (eq? (car exp) 'def)))
(define (def-name exp)
  (cadr exp))
(define (def-init exp)
  (caddr exp))
(define (eval-def exp env)
  (let ((value (eval (def-init exp) env)))
    (define-name! (def-name exp) value env)
   value))
(define (define-name! name value env)
  (let ((binding (cons name value))
        (new-pair (cons (car env) (cdr env))))
    (set-car! env binding)
    (set-cdr! env new-pair)))
```

```
;(fdef triple (a) (+ a a a))
```

```
;(fdef triple (a) (+ a a a))
(define (fdef? exp)
  (and (pair? exp) (eq? (car exp) 'fdef)))
```

```
;(fdef triple (a) (+ a a a))
(define (fdef? exp)
   (and (pair? exp) (eq? (car exp) 'fdef)))
(define (fdef-name exp)
   (cadr exp))
```

(cadr exp))

(define (fdef-parameters exp)
 (caddr exp))

```
Definitions
;(fdef triple (a) (+ a a a))
(define (fdef? exp)
   (and (pair? exp) (eq? (car exp) 'fdef)))
(define (fdef-name exp)
```

```
Definitions
;(fdef triple (a) (+ a a a))
(define (fdef? exp)
  (and (pair? exp) (eq? (car exp) 'fdef)))
(define (fdef-name exp)
  (cadr exp))
(define (fdef-parameters exp)
  (caddr exp))
(define (fdef-body exp)
  (cdddr exp))
```

```
;(fdef triple (a) (+ a a a))
(define (fdef? exp)
  (and (pair? exp) (eq? (car exp) 'fdef)))
(define (fdef-name exp)
  (cadr exp))
(define (fdef-parameters exp)
  (caddr exp))
(define (fdef-body exp)
  (cdddr exp))
(define (eval-fdef exp env)
  (let ((value
         (make-function (fdef-parameters exp)
                         (fdef-body exp))))
    (define-name! (fdef-name exp) value env)
   value))
```

```
Evaluator
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((fdef? exp)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Example

>> (def foo 1)

```
Example
>> (def foo 1)
1
```

```
Example
>> (def foo 1)
>> (+ foo (def foo 2))
```

```
Example
```

```
>> (def foo 1)
1
>> (+ foo (def foo 2))
3 ;;But it could have been 4
```

```
>> (def foo 1)
1
>> (+ foo (def foo 2))
3 ;;But it could have been 4
>> (def bar 1)
```

```
>> (def foo 1)
1
>> (+ foo (def foo 2))
3 ;;But it could have been 4
>> (def bar 1)
1
```

```
>> (def foo 1)
1
>> (+ foo (def foo 2))
3 ;;But it could have been 4
>> (def bar 1)
1
>> (+ (def bar 2) bar)
```

Example

```
>> (def foo 1)
1
>> (+ foo (def foo 2))
3  ;;But it could have been 4
>> (def bar 1)
1
>> (+ (def bar 2) bar)
4  ;;But it could have been 3
```

- The evaluation of a definition *modifies* the current environment.
- The modification occurs as a side-effect.
- Understanding side-effects requires an evaluation order:
 - Letf-to-right in Java and Common Lisp.
 - Right-to-left in APL and J.
 - Unspecified in Scheme, C, Pascal, and many others.

```
>> (def baz 3)
```

>> (def baz 3)

Example (with left-to-right evaluation order)

Problem

• An "empty" scope should not be different from a "normal" scope.

Environments

- In our current implementation, the environment is implemented as a continuous sequence of bindings without separation between different scopes.
- To add such separation, we need to partition that sequence into frames.
- Each time we augment the environment (i.e., for every let, flet and function call), we create a new frame.
- Each frame contains an association list for the bindings.
- This change require us to redefine:
 - the process of finding bindinds: eval-name
 - the process of augmenting the environment: augment-environment
 - the process of updating the environment: define-name!



```
(define (augment-environment names values env)
  (if (null? names)
        env
        (cons (car names) (car values))
             (augment-environment (cdr names) (cdr values) env))))
```

```
(define (augment-environment names values env)
  (cons (map cons names values) env))
```

Environments as Sequences of Frames

```
(define (augment-environment names values env)
  (cons (map cons names values) env))
(define (eval-name name env)
  (define (lookup-in-frame frame)
    (cond ((null? frame)
           (eval-name name (cdr env)))
          ((eq? name (caar frame))
           (cdar frame))
          (else
           (lookup-in-frame (cdr frame)))))
  (if (null? env)
      (error "Unbound name -- EVAL-NAME" name)
      (lookup-in-frame (car env))))
(define (define-name! name value env)
  (let ((binding (cons name value))
        (new-pair (cons (car env) (cdr env))))
    (set-car! env binding)
    (set-cdr! env new-pair)))
```

Environments as Sequences of Frames

```
(define (augment-environment names values env)
  (cons (map cons names values) env))
(define (eval-name name env)
  (define (lookup-in-frame frame)
    (cond ((null? frame)
          (eval-name name (cdr env)))
          ((eq? name (caar frame))
           (cdar frame))
          (else
           (lookup-in-frame (cdr frame)))))
  (if (null? env)
      (error "Unbound name -- EVAL-NAME" name)
      (lookup-in-frame (car env))))
(define (define-name! name value env)
  ((let ((binding (cons name value)))
    (set-car! env (cons binding (car env)))))
```

```
>> (def baz 3)
```

Example >> (def baz 3) 3 >> (+ (let ((x 0))

```
Example
>> (def baz 3)
3 >> (+ (let ((x 0))
       (def baz 5))
      baz)
```

(def baz 6))

8 >> (+ (let ()

baz)

```
>> (def baz 3)
>> (+ (let ((x 0))
       (def baz 5))
      baz)
8
>> (+ (let ()
       (def baz 6))
      baz)
9
    ;;Good, it's fine now!
```

Example >> (def baz 3) 3

```
>> (def baz 3)
>> (+ (let ((x 0))
       (def baz 5))
      baz)
8
>> (+ (let ()
       (def baz 6))
     baz)
  ;;Good, it's fine now!
>> (def counter 0)
0
>> (fdef incr ()
     (def counter (+ counter 1)))
(function () (def counter (+ counter 1)))
```

Example >> (def baz 3) >> (+ (let ((x 0))

(def baz 5))

(def baz 6))

;;Good, it's fine now!

(def counter (+ counter 1))) (function () (def counter (+ counter 1)))

baz)

baz)

>> (fdef incr ()

>> (incr)

>> (def counter 0)

>> (+ (let ()

8

0

```
>> (def baz 3)
>> (+ (let ((x 0))
       (def baz 5))
      baz)
8
>> (+ (let ()
       (def baz 6))
     baz)
  ;;Good, it's fine now!
>> (def counter 0)
0
>> (fdef incr ()
     (def counter (+ counter 1)))
(function () (def counter (+ counter 1)))
>> (incr)
1
```

```
Example
```

```
>> (def baz 3)
>> (+ (let ((x 0))
       (def baz 5))
      baz)
8
>> (+ (let ()
       (def baz 6))
     baz)
 ;;Good, it's fine now!
>> (def counter 0)
0
>> (fdef incr ()
     (def counter (+ counter 1)))
(function () (def counter (+ counter 1)))
>> (incr)
1
>> (incr)
```

```
>> (def baz 3)
>> (+ (let ((x 0))
       (def baz 5))
      baz)
8
>> (+ (let ()
       (def baz 6))
     baz)
 ;;Good, it's fine now!
>> (def counter 0)
0
>> (fdef incr ()
     (def counter (+ counter 1)))
(function () (def counter (+ counter 1)))
>> (incr)
1
>> (incr)
1 ;; What?
```

Assignments

- A scope protects names: all internally defined names are invisible to the outside.
- But sometimes we want to change a name defined in an outer scope.
- This requires the concept of assignment.
- Instead of defining a new binding, an assignment updates an existing binding.
- Concrete syntax: (set! name expression)

Assignments

```
;; (set! counter (+ counter 1))
```

```
Assignments
;; (set! counter (+ counter 1))
(define (set? exp)
  (and (pair? exp) (eq? (car exp) 'set!)))
```

```
Assignments

;; (set! counter (+ counter 1))

(define (set? exp)
    (and (pair? exp) (eq? (car exp) 'set!)))

(define (assignment-name exp)
    (cadr exp))
```

```
Assignments
;; (set! counter (+ counter 1))
(define (set? exp)
  (and (pair? exp) (eq? (car exp) 'set!)))
(define (assignment-name exp)
  (cadr exp))
(define (assignment-expression exp)
  (caddr exp))
```

Assignments

```
(;; (set! counter (+ counter 1))
(define (set? exp)
    (and (pair? exp) (eq? (car exp) 'set!)))
(define (assignment-name exp)
    (cadr exp))
(define (assignment-expression exp)
    (caddr exp))
(define (eval-set exp env)
    (let ((value (eval (assignment-expression exp) env)))
        (update-name! (assignment-name exp) value env)
        value))
```

Analogy

 To update a name we must locate its binding but that's what the evaluation of the name does.

Evaluating a Name

```
(define (eval-name name env)
  (define (lookup-in-frame frame)
    (cond ((null? frame)
           (eval-name name (cdr env)))
          ((eq? name (caar frame))
           (cdar frame))
          (else
           (lookup-in-frame (cdr frame)))))
  (if (null? env)
      (error "Unbound name -- EVAL-NAME" name)
      (lookup-in-frame (car env))))
```

Updating a Name

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
        ((fdef? exp)
         (eval-fdef exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
Example
```

```
>> (def counter 0)
0
>> (fdef incr ()
        (set! counter (+ counter 1)))
(function () (set! counter (+ counter 1)))
```

```
Example
```

```
>> (def counter 0)
>> (fdef incr ()
     (set! counter (+ counter 1)))
(function () (set! counter (+ counter 1)))
>> (incr)
```

```
Example
```

```
>> (def counter 0)
0
>> (fdef incr ()
          (set! counter (+ counter 1)))
(function () (set! counter (+ counter 1)))
>> (incr)
1
```

```
Example
```

```
Example
```

```
>> (def counter 0)
0
>> (fdef incr ()
          (set! counter (+ counter 1)))
(function () (set! counter (+ counter 1)))
>> (incr)
1
>> (incr)
2
```

Input/Output

- Besides assignments, there is another important source of side-effects: input/output.
- We need to include input-output operations in our evaluator.
- They must be implemented using operating system primitives.

Input/Output

```
(define initial-bindings
  (list (cons 'pi 3.14159)
    ...
        (cons '<= (make-primitive <=))
        (cons '>= (make-primitive >=))
        (cons 'display (make-primitive display))
        (cons 'newline (make-primitive newline))
        (cons 'read (make-primitive read))))
```

```
Example
(fdef print-value (text value)
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))

(fdef evaluate-all-and-return-the-last-one (v1 v2 v3 v4 v5)
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))

(fdef evaluate-all-and-return-the-last-one (v1 v2 v3 v4 v5)
    v5)
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
      (display text)
      (display " = ")
      (display value)
      (newline)
      value))

(fdef evaluate-all-and-return-the-last-one (v1 v2 v3 v4 v5)
      v5)

>> (print-value "fact(5)" (fact 5))
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
      (display text)
      (display " = ")
      (display value)
      (newline)
      value))

(fdef evaluate-all-and-return-the-last-one (v1 v2 v3 v4 v5)
      v5)

>> (print-value "fact(5)" (fact 5))
      = 120 fact(5) ;; Huh?
```

Example

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
    (display text)
  (display " = ")
    (display value)
    (newline)
    value))
```

Does it work?

- Nothing guarantees that the arguments to the function evaluate-all-and-return-the-last-one are evaluated from left to right.
- We need a specific control structure to ensure left-to-right evaluation.
- Following Scheme tradition, we will use begin.

```
;;(begin expr0 expr1 ... exprn)
```

```
;;(begin expr0 expr1 ... exprn)
(define (begin? exp)
  (and (pair? exp) (eq? (car exp) 'begin)))
```

```
(;;(begin expr0 expr1 ... exprn)
(define (begin? exp)
   (and (pair? exp) (eq? (car exp) 'begin)))
(define (begin-expressions exp)
   (cdr exp))
```

```
;;(begin expr0 expr1 ... exprn)
(define (begin? exp)
  (and (pair? exp) (eq? (car exp) 'begin)))
(define (begin-expressions exp)
  (cdr exp))
(define (eval-begin exp env)
  (define (eval-sequence expressions env)
    (if (null? (cdr expressions))
        (eval (car expressions) env)
        (begin
          (eval (car expressions) env)
          (eval-sequence (cdr expressions) env))))
  (eval-sequence (begin-expressions exp) env))
```

Evaluator

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Evaluator

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((begin? exp)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Evaluator

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        ((begin? exp)
         (eval-begin exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(fdef print-value (text value)
  (evaluate-all-and-return-the-last-one
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))
```

```
(fdef print-value (text value)
  (begin
    (display text)
    (display " = ")
    (display value)
    (newline)
    value))
```

```
(fdef print-value (text value)
  (begin
      (display text)
      (display " = ")
      (display value)
      (newline)
      value))
>> (print-value "fact(5)" (fact 5))
```

```
(fdef print-value (text value)
  (begin
      (display text)
      (display " = ")
      (display value)
      (newline)
      value))
>> (print-value "fact(5)" (fact 5))
fact(5) = 120
```

```
(fdef print-value (text value)
  (begin
        (display text)
        (display " = ")
        (display value)
        (newline)
        value))
>> (print-value "fact(5)" (fact 5))
fact(5) = 120
120
```

```
(fdef print-value (text value)
  (begin
    (display text)
(display " = ")
    (display value)
    (newline)
    value))
>> (print-value "fact(5)" (fact 5))
fact(5) = 120
120
>> (print-value "35^3" (* 35 35 35))
```

Example

```
(fdef print-value (text value)
  (begin
       (display text)
       (display " = ")
       (display value)
       (newline)
       value))
>> (print-value "fact(5)" (fact 5))
fact(5) = 120
120
>> (print-value "35^3" (* 35 35))
35^3 = 42875
42875
```

Implicit begin

• It is usual to consider that every function body, every let body, and every flet body has an *implicit* begin form.

Example (fdef print-value (text value) (begin (display text) (display " = ") (display value) (newline) value))

```
(fdef print-value (text value)
  (display text)
  (display " = ")
  (display value)
  (newline)
  value)
```

Example

```
(fdef print-value (text value)
  (display text)
  (display " = ")
  (display value)
  (newline)
  value)
```

Implicit begin

- By wrapping the list of forms in a begin form we automatically transform the list into a single expression.
- This expression will be evaluated sequentially.
- To implement the implicit begin, we have to modify the selectors that return the body of a let, the body of a flet and the body of a function.

```
;; (let (...)
;; ...)
```

```
;; (let (...)
;; ...)
(define (let-body exp)
  (caddr exp))
```

```
;; (let (...)
;; ...)
(define (let-body exp)
    (caddr exp))
;; (flet (...)
;; ...)
```

Body Selectors ;; (let (...)

```
(define (let-body exp)
  (caddr exp))
;; (flet (...)
;; ...)
(define (flet-body exp)
  (caddr exp))
```

```
;; (let (...)
(define (let-body exp)
  (caddr exp))
;; (flet (...)
(define (flet-body exp)
  (caddr exp))
;; (function (...)
```

```
;; (let (...)
(define (let-body exp)
  (caddr exp))
;; (flet (...)
;; ...)
(define (flet-body exp)
  (caddr exp))
;; (function (...)
  ...)
(define (function-body func)
  (caddr func))
```

```
;; (let (...)
(define (let-body exp)
  (caddr exp))
;; (flet (...)
;; ...)
(define (flet-body exp)
  (caddr exp))
;; (function (...)
(define (function-body func)
  (caddr func))
(define (make-begin expressions)
  (cons 'begin expressions))
```

```
;; (let (...)
;; ....)
(define (let-body exp)
  (make-begin (cddr exp)))
;; (flet (...)
;; ....)
(define (flet-body exp)
  (make-begin (cddr exp)))
;; (function (...)
;; ... ... )
(define (function-body func)
  (make-begin (cddr func)))
(define (make-begin expressions)
  (cons 'begin expressions))
```

```
>> (fdef reflexive-addition (x) (+ x x))
```

```
>> (fdef reflexive-addition (x)
          (+ x x))
(function (x) (+ x x))
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
```

```
>> (fdef reflexive-addition (x)
          (+ x x))
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
          (* x x))
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x))
(function (x) (* x x))
```

```
>> (fdef reflexive-addition (x)
          (+ x x))
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
          (* x x))
(function (x) (* x x))
>> (reflexive-product 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f x x)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f x x))
(function (f x) (f x x))
>> (reflexive-operation + 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
9
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
9
>> (reflexive-operation / 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x)
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
9
   (reflexive-operation / 3)
1
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x))
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
9
   (reflexive-operation / 3)
1
   (reflexive-operation - 3)
```

```
>> (fdef reflexive-addition (x)
     (+ x x)
(function (x) (+ x x))
>> (reflexive-addition 3)
6
>> (fdef reflexive-product (x)
     (* x x))
(function (x) (* x x))
>> (reflexive-product 3)
>> (fdef reflexive-operation (f x)
     (f \times x)
(function (f x) (f x x))
>> (reflexive-operation + 3)
6
>> (reflexive-operation * 3)
9
   (reflexive-operation / 3)
1
   (reflexive-operation - 3)
0
```

- A higher-order function is a function that accepts other functions as arguments or that computes functions as results.
- Most modern programming languages support higher-order functions and provide libraries containing higher-order functions (e.g., qsort in C, map in Scheme).
- In what concerns higher-order functions, our evaluator has the same expressive power as the C programming language:
 - We can pass functions as arguments.
 - We can apply functions contained in parameters.
 - But we cannot create anonymous functions (yet).
 - And we have other problems that C avoids by forbidding internal function definitions (we will discuss them later).

```
>> (fdef lie (truth)
   (display "it is not true that ")
   (display truth)
   (newline))
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
(display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (lie "it is not true that 1+1=2")
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (lie "it is not true that 1+1=2")
it is not true that it is not true that 1+1=2
```

```
Lies, Damned Lies
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (lie "it is not true that 1+1=2")
it is not true that it is not true that 1+1=2
>> (fdef make-it-true (lie)
     (lie lie))
(function (lie) (lie lie))
>> (make-it-true "1+1=3")
error in cadr: expected appropriate list structure, but got "1+1=3"
backtrace:
   (cadr func)
  1 (function-parameters func)
   (eval input initial-environment)
     (repl)
```

Naming Problem

- Natural languages use many Homonyms:
 - The bandage was wound around the wound.
 - 2 Don't object to the object.
 - We polish the Polish furniture.
 - The present is the best time to present the present.
- Programming languages tend to immitate natural languages and allow homonyms as long as they can be distinguished.
- In general, different namespaces are created to distinguish the different meanings (types, variables, functions, etc).
- Some languages do not allow more than one namespace and prefer to use name conventions.

Homonyms in Java

```
public class foo {
    int foo;
    foo() {
        foo = 0;
    foo(foo foo) {
        this.foo = foo.foo + 1;
    void foo(foo foo) {
        System.err.println(foo.foo);
    }
    public static void main(String args[]) {
        foo foo = new foo():
        foo:foo.foo(new foo(foo));
}
```

Homonyms in Java

```
public class foo {
    int foo;
    foo() {
        foo = 0;
    foo(foo foo) {
        this.foo = foo.foo + 1;
    void foo(foo foo) {
        System.err.println(foo.foo);
    }
    public static void main(String args[]) {
        foo foo = new foo();
        foo:foo.foo(new foo(foo));
}
```

Namespaces

- Java has namespaces for types, for constructors, for methods, for fields and variables, for labels, etc.
- C has a different namespace for structs, but a single namespace for functions and variables.
- Dylan has a single namespace and uses name conventions for distinguishing different meanings: list is a function, t> is a type and \$list is a global variable.
- Scheme has a single namespace: each name has just one meaning in each context (but recent name conventions are being proposed, e.g., exceptions must start with the letter &).
- Common Lisp has multiple namespaces: the name list designates a type, a function, a variable, a label, etc.

Lisp-1 vs Lisp-2

- A Lisp-2 is a Lisp that has one namespace for functions and another namespace for variables.
- A Lisp-1 is a Lisp that has a single namespace where a name either represents a function or a variable but not both.
- Note that Common Lisp is more than a Lisp-2: it is a Lisp-n
- Let's modify our evaluator to implement Lisp-2 semantics:
 - 1 Create a different namespace for functions.
 - Each function definition is placed in the function namespace.
 - Secondary Each function call accesses the function namespace.
- Our namespace for functions will be implemented using hidden naming conventions, but we could use more sophisticated (and efficient) approaches.

Function Namespace

```
(define (in-function-namespace name)
  (string->symbol
    (string-append
        "%function-"
        (symbol->string name))))
```

Function Namespace

```
(define (in-function-namespace name)
  (string->symbol
    (string-append
     "%function-"
     (symbol->string name))))
```

flet

Function Namespace

```
(define (in-function-namespace name)
  (string->symbol
    (string-append
        "%function-"
        (symbol->string name))))
```

flet

Function Namespace

```
(define (in-function-namespace name)
  (string->symbol
    (string-append
        "%function-"
        (symbol->string name))))
```

flet

```
fdef
```

```
;(fdef triple (a) (+ a a a))
```

```
fdef
;(fdef triple (a) (+ a a a))
(define (fdef-name exp)
    (cadr exp))
(define (fdef-name exp)
    (in-function-namespace (cadr exp)))
```

```
fdef
;(fdef triple (a) (+ a a a))
(define (fdef-name exp)
    (cadr exp))
(define (fdef-name exp)
    (in-function-namespace (cadr exp)))
```

Function Call

```
;;(square radius)
```

```
fdef
;(fdef triple (a) (+ a a a))
(define (fdef-name exp)
    (cadr exp))
(define (fdef-name exp)
    (in-function-namespace (cadr exp)))
```

```
fdef
;(fdef triple (a) (+ a a a))
(define (fdef-name exp)
   (cadr exp))
(define (fdef-name exp)
   (in-function-namespace (cadr exp)))
```

```
Function Call
```

```
;;(square radius)
(define (call-operator exp)
   (car exp))
(define (call-operator exp)
   (in-function-namespace (car exp)))
```

Initial Bindings (define initial-bindings (list (cons 'pi 3.14159) (cons '+ (make-primitive +)) (cons '* (make-primitive *)) (cons '- (make-primitive -)) (cons '/ (make-primitive /)) (cons '= (make-primitive =)) (cons '< (make-primitive <))</pre> (cons '> (make-primitive >)) (cons '<= (make-primitive <=))</pre> (cons '>= (make-primitive >=))

(cons 'read (make-primitive read))))

Initial Bindings

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons (in-function-namespace
                                       '+) (make-primitive +))
        (cons (in-function-namespace
                                           (make-primitive *))
                                          (make-primitive -))
        (cons (in-function-namespace '-)
        (cons (in-function-namespace
                                          (make-primitive /))
                                           (make-primitive =))
        (cons (in-function-namespace
        (cons (in-function-namespace
                                      '<) (make-primitive <))</pre>
        (cons (in-function-namespace
                                       '>) (make-primitive >))
        (cons (in-function-namespace
                                      '<=) (make-primitive <=))</pre>
        (cons (in-function-namespace
                                       '>=) (make-primitive >=))
        (cons (in-function-namespace 'read) (make-primitive read))))
```

```
>> (fdef lie (truth)
   (display "it is not true that ")
   (display truth)
   (newline))
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
```

```
>> (fdef lie (truth)
     (display "it is not true that ")
     (display truth)
     (newline))
(function (truth) ...)
>> (lie "1+1=2")
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (fdef make-it-true (lie)
     (lie lie))
(function (lie) (lie lie))
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (fdef make-it-true (lie)
     (lie lie))
(function (lie) (lie lie))
>> (make-it-true "1+1=3")
```

```
>> (fdef lie (truth)
    (display "it is not true that ")
    (display truth)
    (newline))
(function (truth) ...)
>> (lie "1+1=2")
it is not true that 1+1=2
>> (fdef make-it-true (lie)
     (lie lie))
(function (lie) (lie lie))
>> (make-it-true "1+1=3")
it is not true that 1+1=3
```

```
Puzzle
>> (def foo 1)
```

```
>> (def foo 1)
1
```

```
Puzzle
```

```
>> (def foo 1)
1
>> (fdef foo () foo)
```

```
Puzzle
```

```
>> (def foo 1)
>> (fdef foo () foo)
(function () foo)
```

```
>> (def foo 1)
1
>> (fdef foo () foo)
(function () foo)
>> foo
```

```
>> (def foo 1)
1
>> (fdef foo () foo)
(function () foo)
>> foo
1
```

```
Puzzle
```

```
>> (def foo 1)
1
>> (fdef foo () foo)
(function () foo)
>> foo
1
>> (foo)
```

```
>> (def foo 1)
1
>> (fdef foo () foo)
(function () foo)
>> foo
1
>> (foo)
1
```

```
Puzzle
>> (def foo 1)
>> (fdef foo () foo)
(function () foo)
>> foo
>> (foo)
1
>> (let ((foo (+ (foo) 1)))
     (flet ((foo (foo) (+ foo 1)))
       (foo (+ foo 1))))
4
>> (+ foo
      (foo)
      (let ((foo 2))
        (foo))
      (flet ((foo () foo))
        (foo)))
```

```
Puzzle
>> (def foo 1)
>> (fdef foo () foo)
(function () foo)
>> foo
>> (foo)
1
>> (let ((foo (+ (foo) 1)))
     (flet ((foo (foo) (+ foo 1)))
       (foo (+ foo 1))))
4
>> (+ foo
      (foo)
      (let ((foo 2))
        (foo))
      (flet ((foo () foo))
        (foo)))
5
```

Lisp-2 and Higher-Order Functions

```
>> (fdef reflexive-operation (f x) (f x x))
```

$Lisp-2 \ and \ Higher-Order \ Functions$

Lisp-2 and Higher-Order Functions

Lisp-2 and Higher-Order Functions

Problem

- In a Lisp-2, a name that is *not* in function call position is always treated as *not* belonging to the function namespace.
- We need extra syntax to allow a name to be treated as belonging to the function namespace.
- We will use the form (function +).
- If we allow changing the reader, we might even treat #'foo as an abbreviation for (function foo).

Evaluator

;;(function foo)

```
;;(function foo)
(define (function-reference? exp)
  (and (pair? exp) (eq? (car exp) 'function)))
```

```
;;(function foo)
(define (function-reference? exp)
  (and (pair? exp) (eq? (car exp) 'function)))
(define (function-reference-name exp)
  (cadr exp))
```

```
(;;(function foo)
(define (function-reference? exp)
   (and (pair? exp) (eq? (car exp) 'function)))
(define (function-reference-name exp)
   (cadr exp))
(define (eval-function-reference exp env)
   (eval (in-function-namespace (function-reference-name exp))
        env))
```

```
;;(function foo)
(define (function-reference? exp)
  (and (pair? exp) (eq? (car exp) 'function)))
(define (function-reference-name exp)
  (cadr exp))
(define (eval-function-reference exp env)
  (eval (in-function-namespace (function-reference-name exp))
        env))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;;(function foo)
(define (function-reference? exp)
  (and (pair? exp) (eq? (car exp) 'function)))
(define (function-reference-name exp)
  (cadr exp))
(define (eval-function-reference exp env)
  (eval (in-function-namespace (function-reference-name exp))
        env))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((function-reference? exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;;(function foo)
(define (function-reference? exp)
  (and (pair? exp) (eq? (car exp) 'function)))
(define (function-reference-name exp)
  (cadr exp))
(define (eval-function-reference exp env)
  (eval (in-function-namespace (function-reference-name exp))
        env))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((function-reference? exp)
         (eval-function-reference exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Lisp-2 and Higher-Order Functions

```
>> (fdef reflexive-operation (f x) (f x x))
```

Lisp-2 and Higher-Order Functions

```
>> (fdef reflexive-operation (f x)
          (f x x))
(function (f x) (f x x))
```

Lisp-2 and Higher-Order Functions

```
>> (fdef reflexive-operation (f x)
        (f x x))
(function (f x) (f x x))
>> (reflexive-operation (function +) 3)
error in "Unbound name -- EVAL-NAME": %function-f
```

Problem

- In a Lisp-2, a name that is in function call position is always treated as belonging to the function namespace.
- We need extra syntax to allow a name to be treated as not belonging to the function namespace.
- We will use the form (funcall f ...).

```
;;(funcall f ...)
```

```
Evaluator
```

```
;;(funcall f ...)
(define (funcall? exp)
  (and (pair? exp) (eq? (car exp) 'funcall)))
```

```
(;;(funcall f ...)
(define (funcall? exp)
  (and (pair? exp) (eq? (car exp) 'funcall)))
(define (funcall-operator exp)
  (cadr exp))
```

```
(;;(funcall f ...)
(define (funcall? exp)
  (and (pair? exp) (eq? (car exp) 'funcall)))
(define (funcall-operator exp)
  (cadr exp))
(define (funcall-operands exp)
  (cddr exp))
```

```
;;(funcall f ...)
(define (funcall? exp)
  (and (pair? exp) (eq? (car exp) 'funcall)))
(define (funcall-operator exp)
  (cadr exp))
(define (funcall-operands exp)
  (cddr exp))
(define (eval-funcall exp env)
  (let ((func (eval-name (funcall-operator exp) env))
        (args (eval-exprs (funcall-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     env)))
          (eval (function-body func) extended-environment)))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((function-reference? exp)
         (eval-function-reference exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((function-reference? exp)
         (eval-function-reference exp env))
        ((funcall? exp)
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((function-reference? exp)
         (eval-function-reference exp env))
        ((funcall? exp)
         (eval-funcall exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
>> (fdef reflexive-operation (f x) (funcall f x x))
```

```
>> (fdef reflexive-operation (f x)
        (funcall f x x))
(function (f x) (funcall f x x))
```

```
>> (fdef reflexive-operation (f x)
          (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
3025
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
3025
>> (fdef identity (x) x)
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
3025
>> (fdef identity (x) x)
(function (x) x)
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
3025
>> (fdef identity (x) x)
(function (x) x)
>> (sum (function identity) 1 100)
```

```
>> (fdef reflexive-operation (f x)
     (funcall f x x))
(function (f x) (funcall f x x))
>> (reflexive-operation (function +) 3)
6
>> (fdef sum (f a b)
     (if (> a b)
       (+ (funcall f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef cube (x) (* x x x))
(function (x) ...)
>> (sum (function cube) 1 10)
3025
>> (fdef identity (x) x)
(function (x) x)
>> (sum (function identity) 1 100)
5050
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
(if (= n 0) 1 (* x (expt x (- n 1)))))
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
      (if (= n 0) 1 (* x (expt x (- n 1)))))
(function (x n) ...)
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
        (if (= n 0) 1 (* x (expt x (- n 1))))
(function (x n) ...)
>> (fdef approx-pi (n)
        (flet ((term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
        (* 4 (sum (function term) 0 n))))
(function (n) ...)
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
      (if (= n 0) 1 (* x (expt x (- n 1)))))
(function (x n) ...)
>> (fdef approx-pi (n)
      (flet (term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
(* 4 (sum (function term) 0 n))))
(function (n) ...)
>> (approx-pi 10)
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
          (if (= n 0) 1 (* x (expt x (- n 1))))
(function (x n) ...)
>> (fdef approx-pi (n)
           (flet ((term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
           (* 4 (sum (function term) 0 n))))
(function (n) ...)
>> (approx-pi 10)
3.232315809405593
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
        (if (= n 0) 1 (* x (expt x (- n 1)))))
(function (x n) ...)
>> (fdef approx-pi (n)
        (flet ((term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
        (* 4 (sum (function term) 0 n))))
(function (n) ...)
>> (approx-pi 10)
3.232315809405593
>> (approx-pi 100)
```

Approximating π

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```
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          (if (= n 0) 1 (* x (expt x (- n 1)))))
(function (x n) ...)
>> (fdef approx-pi (n)
               (flet ((term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
               (* 4 (sum (function term) 0 n))))
(function (n) ...)
>> (approx-pi 10)
3.232315809405593
>> (approx-pi 100)
3.151493401070991
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

```
>> (fdef expt (x n)
        (if (= n 0) 1 (* x (expt x (- n 1)))))
(function (x n) ...)
>> (fdef approx-pi (n)
        (flet ((term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
        (* 4 (sum (function term) 0 n))))
(function (n) ...)
>> (approx-pi 10)
3.232315809405593
>> (approx-pi 1000)
3.151493401070991
>> (approx-pi 1000)
```

Approximating π

$$\pi = 4\sum_{i=0}^{\infty} \frac{(-1)^i}{2i+1}$$

Factorial

$$n! = 1 \times 2 \times 3 \times \ldots \times n = \prod_{i=1}^{n} i$$

Factorial

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Factorial

$$n! = 1 \times 2 \times 3 \times \ldots \times n = \prod_{i=1}^{n} i$$

Sum

Product


```
Product
```



```
Product
```

Sum

Product

Abstraction

 Sums and products are so identical that it is possible to abstract them using additional (functional) parameters.

Abstraction: accumulation (in a Lisp-2)

Lisp-2 and Functional Programming

- Functional programming in a Lisp-2 tends to spread funcal1 and function all over the code.
- This makes the code hard to read.

Abstraction: accumulation (in a Lisp-1)

Lisp-1 and Functional Programming

- Functional programming in a Lisp-1 looks much better (at the price of loosing homonyms).
- Let's change our evaluator to become a Lisp-1.



Lisp-2

```
(define (in-function-namespace name)
  (string->symbol
   (string-append
   "%function-"
    (symbol->string name))))
```

Lisp-1

```
(define (in-function-namespace name)
  name)
```

Lisp-1

```
(define (in-function-namespace name)
  name)
```

Other simplifications

- Drop function (and function-reference?, function-reference-name, eval-function-reference).
- Drop funcall (and funcall?, funcall-operator, funcall-operands, eval-funcall).
- Drop in-function-namespace.

```
>> (fdef next (i)
	(if (< i 0)
	(+ i -1)
	(+ i +1)))
(function (i) ...)
>> (next 3)
```

```
>> (fdef next (i)
	(if (< i 0)
	(+ i -1)
	(+ i +1)))
(function (i) ...)
>> (next 3)
4
>> (next -5)
-6
```

Next Integer (away from origin)

```
>> (fdef next (i)
	(if (< i 0)
	(+ i -1)
	(+ i +1)))
(function (i) ...)
>> (next 3)
4
>> (next -5)
-6
```

Conditional Expression

- There are repeated parts that could be factorized...
- ... by moving the conditional expression inside the addition.



Next Integer (away from origin)

Conditional Expression

- There are repeated parts that could be factorized...
- ... by moving the conditional expression inside the addition.



Next Integer (away from origin) - Alternative Implementation

```
>> (fdef next (i)
      (if (< i 0)
(- i 1)
         (+ i 1)))
```

Next Integer (away from origin) – Alternative Implementation

Next Integer (away from origin) – Alternative Implementation

Next Integer (away from origin) – Alternative Implementation

Next Integer (away from origin) – Alternative Implementation

```
>> (fdef next (i)
          (if (< i 0)
                (- i 1)
                (+ i 1)))
(function (i) ...)
>> (fdef next (i)
                ((if (< i 0) - +) i 1))
(function (i) ...)
>> (next 3)
error in "Unbound name -- EVAL-NAME": (if (< i 0) - +)</pre>
```

Conditional Expression in Operator Position

- We were expecting to see a name in the operator position, but we now have an expression there.
- Let's generalize the treatment of the operator.

Function Calls

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
            (eval-name exp env))
            ...
        ((call? exp)
            (eval-call exp env))
        (else
            (error "Unknown expression type -- EVAL" exp))))
```

Function Calls

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     env)))
          (eval (function-body func) extended-environment)))))
```

Function Calls

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (eval-call exp env)
  (let ((func (eval-name (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     env)))
          (eval (function-body func) extended-environment)))))
```

Function Calls

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (eval-call exp env)
  (let ((func (eval (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     env)))
          (eval (function-body func) extended-environment)))))
```

Expressions as Operators

```
(fdef next (i)
((if (< i 0) - +) i 1))
```

Lisp-1 and Functional Programming

- In a Lisp-1 it is possible to generalize the operator in a function call:
 - The operator can be an expression that evaluates to a function.
 - Using a name as operator is just a special case of the use of an expression.
- In a Lisp-2 it is also possible to generalize the operator in a function call but semantics become more complex.

Approximating π in a Lisp-1

Anonymous Functions

- The function term is not useful for other purposes.
- It seems a waste to be forced to give it a name.
- Let's introduce the concept of anonymous function.



Named Functions

Anonymous Functions

```
(fdef approx-pi (n)
(* 4 (sum (lambda (i) (/ (expt -1.0 i) (+ (* 2 i) 1)))
0 n)))
```

Named Functions

```
(fdef approx-pi (n)
  (flet (term (i) (/ (expt -1.0 i) (+ (* 2 i) 1))))
(* 4 (sum term
                0 n))))
```

Anonymous Functions

```
(fdef approx-pi (n)
  (* 4 (sum (lambda (i) (/ (expt -1.0 i) (+ (* 2 i) 1)))
            0 n)))
```

Named Functions

Anonymous Functions

```
(fdef approx-pi (n)
(* 4 (sum (lambda (i) (/ (expt -1.0 i) (+ (* 2 i) 1)))
0 n)))
```

Lambda Forms

- A lambda form evaluates to a function without requiring an explicit name for the function.
- The function is anonymous.

```
Evaluator
;;(lambda (a b)
;; (+ a b))
```

Evaluator ;;(lambda (a b)

```
Evaluator
;;(lambda (a b)
;; (+ a b))
(define (lambda? exp)
  (and (pair? exp)
       (eq? (car exp) 'lambda)))
(define (lambda-parameters exp)
  (cadr exp))
```

```
;;(lambda (a b)
;; (+ a b))
(define (lambda? exp)
  (and (pair? exp)
       (eq? (car exp) 'lambda)))
(define (lambda-parameters exp)
  (cadr exp))
(define (lambda-body exp)
  (cddr exp))
```

```
;;(lambda (a b)
;; (+ a b))
(define (lambda? exp)
  (and (pair? exp)
       (eq? (car exp) 'lambda)))
(define (lambda-parameters exp)
  (cadr exp))
(define (lambda-body exp)
  (cddr exp))
(define (eval-lambda exp env)
  (make-function (lambda-parameters exp)
                 (lambda-body exp)))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((lambda? exp)
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        ((lambda? exp)
         (eval-lambda exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp)
         (eval-let exp env))
        ((flet? exp)
         (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp)
         (eval-fdef exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
     (accumulate * 1 (lambda (i) i) 1 n))
```

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
        (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
```

Approximating *e*

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
          (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
          (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
```

Approximating *e*

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
          (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
          (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
```

Approximating *e*

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
        (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
        (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
2.7182818011463845
```

Approximating *e*

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
          (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
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          (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
2.7182818011463845
>> (approx-e 100)
```

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
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          (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
          (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
2.7182818011463845
>> (approx-e 100)
2.718281828459045
```

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
     (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
     (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
2.7182818011463845
>> (approx-e 100)
2.718281828459045
>> (approx-e 1000)
```

Approximating e

$$e = \sum_{i=0}^{\infty} \frac{1}{i!}$$

```
>> (fdef fact (n)
     (accumulate * 1 (lambda (i) i) 1 n))
(function (n) ...)
>> (fdef approx-e (n)
     (sum (lambda (i) (/ 1.0 (fact i))) 0 n))
(function (n) ...)
>> (approx-e 10)
2.7182818011463845
>> (approx-e 100)
2.718281828459045
>> (approx-e 1000)
2.718281828459045
```

A Striking Resemblance

- A def form associates a name to the evaluation of a form, e.g., (def foo (* 2 2)).
- A fdef form associates a name to a function, e.g., (fdef square (x) (* x x)).
- A lambda form evaluates to a function, e.g., (lambda (x) (* x x))
- Thus, we have the following equivalence: (fdef square (x) (* x x)) = (def square (lambda (x) (* x x)))
- \bullet Or, in more abstract terms: fdef = def + lambda
- The same equivalence occurs with flet = let + lambda.

fdef

fdef = def + lambda

From

```
(fdef square (x)
  (* x x))
```

То

```
(def square
  (lambda (x)
      (* x x)))
```

fdef

```
fdef = def + lambda
```

From

```
(fdef square (x)
(* x x))
```

То

```
(def square
(lambda (x)
(* x x)))
```

Transformation (without backquote)

fdef

```
fdef = def + lambda
```

From

```
(fdef square (x) (* x x))
```

То

```
(def square
(lambda (x)
(* x x)))
```

Transformation (with backquote)

flet

```
flet = let + lambda
```

From

То

flet

```
flet = let + lambda
```

(foo (bar 2 3)))

Transformation

let

let = lambda + function call

From

То

```
((lambda (x y)
 (- x y))
 (+ 1 2)
 (* 2 3))
```

let

let = lambda + function call

From

То

```
((lambda (x y)
(- x y))
(+ 1 2)
(* 2 3))
```

Transformation

Source to Source Transformations

- Some forms of the language can be implemented as combinations of more primitive forms.
- This implementation can be done as source-to-source transformations.
- Lambda expressions (function abstraction) and function application are two of the most fundamental forms:
 - In theory, everything can be implemented on top of these two forms (Church numerals, Church booleans, Church pairs, etc).
 - In practice, all primitive data types are implemented using low-level capabilities.
- Source-to-source transformations might also be useful for the programmer. We will return to this later.

Second Order Sum

$$\sum^{i_1} f(i)$$

where
$$f(x) = ax^2 + bx + c$$

Second Order Sum

$$\sum_{i_1}^{i_1} f(i_1)$$

where
$$f(x) = ax^2 + bx + c$$

```
>> (fdef sum (f a b)
     (if (> a b)
       0
(+ (f a)
          (sum f (+ a 1) b))))
```

Second Order Sum

$$\sum_{i=1}^{i_1} f(i)$$

where
$$f(x) = ax^2 + bx + c$$

```
>> (fdef sum (f a b)
     (if (> a b)
       (+ (f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
```

Second Order Sum

$$\sum_{i=i_k}^{i_1} f(i) \qquad \text{where} \qquad f(x) = ax^2 + bx + c$$

Second Order Sum

$$\sum_{i=1}^{i_1} f(i) \qquad \text{where} \qquad f(x) = ax^2 + bx + c$$

Second Order Sum

$$\sum_{x=-1}^{1} 10x^2 + 5x = (10 - 5) + 0 + (10 + 5) = 20$$

Second Order Sum

$$\sum_{x=-1}^{1} 10x^2 + 5x = (10 - 5) + 0 + (10 + 5) = 20$$

Second Order Sum

$$\sum_{x=-1}^{1} 10x^2 + 5x = (10 - 5) + 0 + (10 + 5) = 20$$

Dynamic Scope

- So far, we didn't think carefully about the scope of our binding constructs:
 - A def form creates a binding in the current environment.
 - A function call creates a set of bindings (for the parameters) that extends the *current* environment.
 - At the end of the call, the extended environment is discarded.
 - A name is resolved by searching the *current* environment.
- This scoping discipline is called *dynamic scope*.
- In dynamic scope, the environment grows and shrinks in parallel with the call stack.
- Dynamic scope is used in many languages such as EmacsLisp, AutoLisp, Logo, TEX, Perl (local variables) and Common Lisp (special variables).

Downwards Funarg Problem

- Passing a function with free variables as argument causes the downwards funarg problem...
- ... where the function is called in an environment where the free variables might be shadowed by other variables of the calling context.

Solving the Downwards Funarg Problem

- Functions must know the correct environment for resolving the free variables.
- One potential solution might be to use name-mangling techniques to avoid name clashes.

```
>> (fdef sum (f a b)
     (if (> a b)
       (+ (f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef second-order-sum (a b c i0 i1)
     (sum (lambda (x)
            (+ (* a x x) (* b x) c))
          i0 i1))
(function (a b c i0 i1) ...)
>> (second-order-sum 10 5 0 -1 +1)
0 ;;What???
```

Name Mangling Problems

- It makes the code hard to read.
- It doesn't really solve the problem.

Downwards Funarg Problem

- The correct solution is to associate the downward function to the scope where the function was created:
 - That scope is still active when the passed function is called...
 - ...unless it is possible to store the function somewhere and call it later.
- Some languages solve this last problem by restricting the language:
 - Pascal forbids storing functions in variables.
 - C forbids inner functions.
- Other languages (e.g., Scheme, Haskel, Common Lisp) implement environments with indefinite extent.
- But there is a more difficult problem.

$\label{thm:linear_problem} \mbox{Higher-order Functions in a Dynamically Scoped Language}$

$\label{thm:linear_problem} \mbox{Higher-order Functions in a Dynamically Scoped Language}$

$\label{thm:lighter-order} \mbox{Higher-order Functions in a Dynamically Scoped Language}$

Upwards Funarg Problem

- Returning a function with free variables as result causes the upwards funarg problem...
- ... where the function is called in an environment where the free variables are not longer bound (or are bound to different values).

Solving the Upwards Funarg Problem

- Some languages avoid the upwards funarg problem by restricting the language:
 - Pascal forbids functional returns.
 - C forbids inner functions.
 - Java forbids non-final free variables in anonymous inner classes.
- Other languages (e.g., Scheme, Haskel, Common Lisp) implement environments with indefinite extent.

```
(define (make-function parameters body)
  (cons 'function
        (cons parameters
              body)))
(define (function? obj)
  (and (pair? obj)
       (eq? (car obj) 'function)))
(define (function-parameters func)
  (cadr func))
(define (function-body func)
  (make-begin (cddr func)))
```

Evaluator (define (make-function parameters body env) (cons 'function (cons parameters (cons env body)))) (define (function? obj) (and (pair? obj) (eq? (car obj) 'function))) (define (function-parameters func) (cadr func)) (define (function-body func) (make-begin (cdddr func)))

```
(define (make-function parameters body env)
  (cons 'function
        (cons parameters
              (cons env
                    body))))
(define (function? obj)
  (and (pair? obj)
       (eq? (car obj) 'function)))
(define (function-parameters func)
  (cadr func))
(define (function-body func)
  (make-begin (cdddr func)))
(define (function-environment func)
  (caddr func))
```

```
(define (eval-lambda exp env)
  (make-function (lambda-parameters exp)
                 (lambda-body exp)))
(define (eval-call exp env)
  (let ((func (eval (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     env)))
          (eval (function-body func) extended-environment)))))
```

```
(define (eval-lambda exp env)
  (make-function (lambda-parameters exp)
                 (lambda-body exp)
                 env))
(define (eval-call exp env)
  (let ((func (eval (call-operator exp) env))
        (args (eval-exprs (call-operands exp) env)))
    (if (primitive? func)
        (apply-primitive-function func args)
        (let ((extended-environment
               (augment-environment (function-parameters func)
                                     args
                                     (function-environment func))))
          (eval (function-body func) extended-environment)))))
```

```
>> (fdef sum (f a b)
     (if (> a b)
       (+ (f a)
          (sum f (+ a 1) b))))
(function (f a b) ...)
>> (fdef second-order-sum (a b c i0 i1)
     (sum (lambda (x)
            (+ (* a x x) (* b x) c))
          i0 i1))
(function (a b c i0 i1) ...)
>> (second-order-sum 10 5 0 -1 +1)
20 ;; As expected
```

Local Recursive Functions in a Lexically Scoped Language

Local Recursive Functions in a Lexically Scoped Language

Problem

- flet creates a function and establishes a scope where a name is bound to that function.
- But the function is created outside that scope.
- As a result, the function cannot refer itself.

Local Recursive Functions in a Lexically Scoped Language

Local Recursive Functions in a Lexically Scoped Language

Solution (the Common Lisp way — labels)

- Creates the scope for the function name (temporarily binding it to some value).
- ② Creates the function inside that scope.
- Modifies the binding of the function name to point to the function.

Local Recursive Functions in a Lexically Scoped Language

Solution (the Scheme way - define)

- Creates an empty scope.
- ② Creates the function inside that scope.
- Modifies the scope, adding a new binding from the function name to the function.

```
>> (fdef kons (a b)
     (lambda (f)
       (f a b)))
(function (a b) ...)
>> (fdef kar (k)
     (k (lambda (x y) x)))
```

```
>> (fdef kons (a b)
     (lambda (f)
       (f a b)))
(function (a b) ...)
>> (fdef kar (k)
     (k (lambda (x y) x)))
(function (k) ...)
```

```
>> (fdef kons (a b)
     (lambda (f)
       (f a b)))
(function (a b) ...)
>> (fdef kar (k)
     (k (lambda (x y) x)))
(function (k) ...)
>> (fdef kdr (k)
     (k (lambda (x y) y)))
(function (k) ...)
>> (def number-pair (kons 1 2))
(function (f) ...)
>> (kar number-pair)
1
>> (kdr number-pair)
```

```
>> (fdef kons (a b)
     (lambda (f)
       (f a b)))
(function (a b) ...)
>> (fdef kar (k)
     (k (lambda (x y) x)))
(function (k) ...)
>> (fdef kdr (k)
     (k (lambda (x y) y)))
(function (k) ...)
>> (def number-pair (kons 1 2))
(function (f) ...)
>> (kar number-pair)
1
>> (kdr number-pair)
```

```
Example: Lists
```

```
>> (fdef empty-list () #f)
```

```
Example: Lists
```

```
>> (fdef empty-list ()
    #f)
(function () ...)
```

```
>> (fdef empty-list ()
    #f)
(function () ...)
>> (fdef empty-list? (1)
    (not 1))
```

```
>> (fdef empty-list ()
          #f)
(function () ...)
>> (fdef empty-list? (1)
          (not 1))
(function (1) ...)
>> (def join kons)
(function (a b) ...)
```

```
>> (fdef empty-list ()
     #f)
(function () ...)
>> (fdef empty-list? (1)
     (not 1))
(function (1) ...)
>> (def join kons)
(function (a b) ...)
>> (def first kar)
(function (k) ...)
>> (def rest kdr)
(function (k) ...)
>> (def numbers (join 1 (join 2 (join 3 (empty-list)))))
(function (f) ...)
>> (first numbers)
```

```
>> (fdef empty-list ()
     #f)
(function () ...)
>> (fdef empty-list? (1)
     (not 1))
(function (1) ...)
>> (def join kons)
(function (a b) ...)
>> (def first kar)
(function (k) ...)
>> (def rest kdr)
(function (k) ...)
>> (def numbers (join 1 (join 2 (join 3 (empty-list)))))
(function (f) ...)
>> (first numbers)
```

```
>> (fdef empty-list ()
     #f)
(function () ...)
>> (fdef empty-list? (1)
     (not 1))
(function (1) ...)
>> (def join kons)
(function (a b) ...)
>> (def first kar)
(function (k) ...)
>> (def rest kdr)
(function (k) ...)
>> (def numbers (join 1 (join 2 (join 3 (empty-list)))))
(function (f) ...)
>> (first numbers)
>> (first (rest numbers))
```

```
>> (fdef empty-list ()
     #f)
(function () ...)
>> (fdef empty-list? (1)
     (not 1))
(function (1) ...)
>> (def join kons)
(function (a b) ...)
>> (def first kar)
(function (k) ...)
>> (def rest kdr)
(function (k) ...)
>> (def numbers (join 1 (join 2 (join 3 (empty-list)))))
(function (f) ...)
>> (first numbers)
1
>> (first (rest numbers))
```

```
>> (fdef print-list (1)
      (display "(")
      (print-list-elements 1)
      (display ")")
      (newline))
```

```
>> (fdef print-list (1)
        (display "(")
        (print-list-elements 1)
        (display ")")
        (newline))
(function (1) ...)
```

```
>> (fdef print-list (1)
     (display "(")
     (print-list-elements 1)
     (display ")")
     (newline))
(function (1) ...)
>> (fdef print-list-elements (1)
     (if (empty-list? 1)
       #f
       (begin
         (display (first 1))
         (if (empty-list? (rest 1))
           #f
           (begin
             (display " ")
             (print-list-elements (rest 1)))))))
```

```
>> (fdef print-list (1)
     (display "(")
     (print-list-elements 1)
     (display ")")
     (newline))
(function (1) ...)
>> (fdef print-list-elements (1)
     (if (empty-list? 1)
       #f
       (begin
         (display (first 1))
         (if (empty-list? (rest 1))
           #f
           (begin
             (display " ")
             (print-list-elements (rest 1)))))))
(function (1) ...)
```

```
>> (fdef print-list (1)
     (display "(")
     (print-list-elements 1)
     (display ")")
     (newline))
(function (1) ...)
>> (fdef print-list-elements (1)
     (if (empty-list? 1)
       #f
       (begin
         (display (first 1))
         (if (empty-list? (rest 1))
           #f
           (begin
             (display " ")
             (print-list-elements (rest 1)))))))
(function (1) ...)
>> (print-list numbers)
```

```
>> (fdef print-list (1)
     (display "(")
     (print-list-elements 1)
     (display ")")
     (newline))
(function (1) ...)
>> (fdef print-list-elements (1)
     (if (empty-list? 1)
       #f
       (begin
         (display (first 1))
         (if (empty-list? (rest 1))
           #f
           (begin
              (display " ")
              (print-list-elements (rest 1)))))))
(function (1) ...)
>> (print-list numbers)
(1\ 2\ 3)
```

Data Structures

- So far, our language does not have any data structuring mechanisms.
- Most programming languages provide one or more such mechanisms:
 - C: structs, arrays
 - Pascal: records, arrays
 - Java: classes and arrays
 - Common Lisp: structures, classes, arrays, pairs
 - Scheme (pre R6RS): vectors, pairs
 - Scheme (pos R6RS): records, vectors, pairs
- We will implement just the fundamental data structuring mechanism: pairs.

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons '> (make-primitive >))
        (cons '<= (make-primitive <=))</pre>
        (cons '>= (make-primitive >=))
        (cons 'display (make-primitive display))
        (cons 'newline (make-primitive newline))
        (cons 'read (make-primitive read))
        (cons 'list (make-primitive list))
        (cons 'cons (make-primitive cons))
```

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons '> (make-primitive >))
        (cons '<= (make-primitive <=))</pre>
        (cons '>= (make-primitive >=))
        (cons 'display (make-primitive display))
        (cons 'newline (make-primitive newline))
        (cons 'read (make-primitive read))
        (cons 'list (make-primitive list))
        (cons 'cons (make-primitive cons))
        (cons 'car (make-primitive car))
        (cons 'cdr (make-primitive cdr))
```

```
(define initial-bindings
  (list (cons 'pi 3.14159)
        (cons '> (make-primitive >))
        (cons '<= (make-primitive <=))</pre>
        (cons '>= (make-primitive >=))
        (cons 'display (make-primitive display))
        (cons 'newline (make-primitive newline))
        (cons 'read (make-primitive read))
        (cons 'list (make-primitive list))
        (cons 'cons (make-primitive cons))
        (cons 'car (make-primitive car))
        (cons 'cdr (make-primitive cdr))
        (cons 'pair? (make-primitive pair?))
        (cons 'null? (make-primitive null?))))
```

```
>> (fdef caar (p) (car (car p)))
```

```
>> (fdef caar (p) (car (car p))) (function (p) ...)
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(function (p) ...)
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(list)
      (cons (f (car 1))
            (map f (cdr 1)))))
(function (f 1) ...)
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(function (p) ...)
>> (fdef map (f 1)
     (if (null? 1)
       (list)
       (cons (f (car 1))
             (map f (cdr 1)))))
(function (f 1) ...)
>> (map (lambda (x) (* x x x)) (list 1 2 3 4 5 6))
(1 8 27 64 125 216)
>> (fdef append (10 11)
     (if (null? 10)
       (cons (car 10)
             (append (cdr 10) 11))))
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(function (p) ...)
>> (fdef map (f 1)
     (if (null? 1)
       (list)
       (cons (f (car 1))
             (map f (cdr 1)))))
(function (f 1) ...)
>> (map (lambda (x) (* x x x)) (list 1 2 3 4 5 6))
(1 8 27 64 125 216)
>> (fdef append (10 11)
     (if (null? 10)
       (cons (car 10)
             (append (cdr 10) 11))))
(function (10 11) ...)
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(function (p) ...)
>> (fdef map (f 1)
     (if (null? 1)
       (list)
       (cons (f (car 1))
             (map f (cdr 1)))))
(function (f 1) ...)
>> (map (lambda (x) (* x x x)) (list 1 2 3 4 5 6))
(1 8 27 64 125 216)
>> (fdef append (10 11)
     (if (null? 10)
       (cons (car 10)
             (append (cdr 10) 11))))
(function (10 11) ...)
>> (append (list 1 2 3) (list 4 5 6))
```

```
>> (fdef caar (p) (car (car p)))
(function (p) ...)
>> (fdef cadr (p) (car (cdr p)))
(function (p) ...)
>> (fdef map (f 1)
     (if (null? 1)
       (list)
       (cons (f (car 1))
             (map f (cdr 1)))))
(function (f 1) ...)
>> (map (lambda (x) (* x x x)) (list 1 2 3 4 5 6))
(1 8 27 64 125 216)
>> (fdef append (10 11)
     (if (null? 10)
       (cons (car 10)
             (append (cdr 10) 11))))
(function (10 11) ...)
>> (append (list 1 2 3) (list 4 5 6))
(1 2 3 4 5 6)
```

Symbols

- Lisp was intended for symbolic processing.
- A symbol is an abstract entity that represents something, e.g., moon, john, square.
- Symbols can be used with lists to represent more complex concepts, e.g., (believes John (loves Mary Peter))
 - Two symbols are identical if they have the same name.
 - Symbols can be compared using eq?.

Symbols

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- A symbol is an abstract entity that represents something, e.g., moon, john, square.
- Symbols can be used with lists to represent more complex concepts, e.g., (believes John (loves Mary Peter))
 - Two symbols are identical if they have the same name.
 - Symbols can be compared using eq?.

Symbols

```
(define initial-bindings
(list (cons 'pi 3.14159)
...
(cons 'null? (make-primitive null?))
(cons 'eq? (make-primitive eq?))))
```

Symbols, Lists, and Quotation

- Symbols and Lists can be used as data structures...
- ...but they are also used to describe programs.
- How is it possible to distinguish between the two different purposes?
 - The evaluation rule for symbols retrieves the value associated with a symbol in the current environment.
 - We need a different evaluation rule for symbols that are to be taken literaly, i.e., that should not be evaluated.
 - Let's introduce the form quote for that purpose.
- (quote foo) evaluates to foo.
- (quote foo) can be abbreviated as 'foo.

```
;;(quote foo)
```

```
;;(quote foo)
(define (quote? exp)
   (and (pair? exp)
        (eq? (car exp) 'quote)))
```

```
;; (quote foo)
(define (quote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quote)))
(define (quoted-form exp)
  (cadr exp))
(define (eval-quote exp env)
  (quoted-form exp))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((name? exp)
         (eval-name exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;; (quote foo)
(define (quote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quote)))
(define (quoted-form exp)
  (cadr exp))
(define (eval-quote exp env)
  (quoted-form exp))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((quote? exp)
        ((name? exp)
         (eval-name exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;; (quote foo)
(define (quote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quote)))
(define (quoted-form exp)
  (cadr exp))
(define (eval-quote exp env)
  (quoted-form exp))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((quote? exp)
         (eval-quote exp env))
        ((name? exp)
         (eval-name exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
Examples
>> (quote 1)
```

```
Examples
>> (quote 1)
```

```
Examples
```

```
>> (quote 1)
1
>> (quote John)
```

```
>> (quote 1)
1
>> (quote John)
John
```

```
Examples
```

```
>> (quote 1)
>> (quote John)
.John
>> (list (quote believes)
         (quote John)
         (list (quote loves) (quote Mary) (quote Peter)))
(believes John (loves Mary Peter))
>> (list 'believes
         'John
         (list 'loves 'Mary 'Peter))
(believes John (loves Mary Peter))
>> '(believes John (loves Mary Peter))
(believes John (loves Mary Peter))
```

```
>> (quote 1)
>> (quote John)
.John
>> (list (quote believes)
         (quote John)
         (list (quote loves) (quote Mary) (quote Peter)))
(believes John (loves Mary Peter))
>> (list 'believes
         'John
         (list 'loves 'Mary 'Peter))
(believes John (loves Mary Peter))
>> '(believes John (loves Mary Peter))
(believes John (loves Mary Peter))
>> (list pi 'pi)
```

```
>> (quote 1)
>> (quote John)
.John
>> (list (quote believes)
         (quote John)
         (list (quote loves) (quote Mary) (quote Peter)))
(believes John (loves Mary Peter))
>> (list 'believes
         'John
         (list 'loves 'Mary 'Peter))
(believes John (loves Mary Peter))
>> '(believes John (loves Mary Peter))
(believes John (loves Mary Peter))
>> (list pi 'pi)
(3.14159 pi)
```

```
Examples
```

```
>> (quote 1)
>> (quote John)
.John
>> (list (quote believes)
         (quote John)
         (list (quote loves) (quote Mary) (quote Peter)))
(believes John (loves Mary Peter))
>> (list 'believes
         'John
         (list 'loves 'Mary 'Peter))
(believes John (loves Mary Peter))
>> '(believes John (loves Mary Peter))
(believes John (loves Mary Peter))
>> (list pi 'pi)
(3.14159 pi)
>> (list 'fdef 'square (list 'x) (list '* 'x 'x))
```

```
>> (quote 1)
>> (quote John)
.John
>> (list (quote believes)
         (quote John)
         (list (quote loves) (quote Mary) (quote Peter)))
(believes John (loves Mary Peter))
>> (list 'believes
         'John
         (list 'loves 'Mary 'Peter))
(believes John (loves Mary Peter))
>> '(believes John (loves Mary Peter))
(believes John (loves Mary Peter))
>> (list pi 'pi)
(3.14159 pi)
>> (list 'fdef 'square (list 'x) (list '* 'x 'x))
(fdef square (x) (* x x))
```

```
>> (fdef and (b0 b1)
        (if b0
        b1
        #f))
(function (b0 b1) ...)
>> (fdef quotient-or-false (a b)
```

```
>> (fdef and (b0 b1)
        (if b0
        b1
        #f))
(function (b0 b1) ...)
>> (fdef quotient-or-false (a b)
        (and (not (= b 0))
```

Logical Operators: and

Problem

- Function calls evaluate all their arguments...
- ...but and requires short-circuit evaluation.



Short-Circuit Evaluation

- We can implement special evaluation rules in the evaluator...
- ...but that's a job for the compiler writer.
- We want to allow the programmer to define code transformations:

From

```
(and (not (= b 0))
(/ a b))
```

То

```
(if (not (= b 0))
(/ a b)
#f)
```

Short-Circuit Evaluation

- We can implement special evaluation rules in the evaluator...
- ...but that's a job for the compiler writer.
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From

```
(and (not (= b 0))
(/ a b))
```

То

```
(if (not (= b 0))
(/ a b)
#f)
```

Transformation

```
(mdef and (b0 b1)
(list 'if b0 b1 '#f))
```

Macros

- A macro call is a form that:
 - Receives syntactical forms as arguments.
 - Computes a syntactical form as result (the macro expansion).
 - The computed form is evaluated in place of the macro call.

Function Calls vs Macro Calls

- A function call evaluates the arguments and computes a result that is not evaluated.
- A macro call does not evaluate the arguments and computes a result that is evaluated.

```
>> (fdef f-foo (x)
       (display x)
(newline)
       x)
(function (x) ...)
>> (mdef m-foo (x)
       (display x)
(newline)
       x)
(function (x) ...) >> (f-foo (+ 1 2))
3
3
>> (m-foo (+ 1 2))
(+12)
```

Functions vs Macros >> (fdef f-bar ()

```
>> (fdef f-bar ()
     (list '+
           (2)
(function () ...)
>> (mdef m-bar ()
     (list '+
           (2)
(function () ...)
```

Macro Implementation

- A macro is a function but is called differently.
- We will implement macros simply as tagged functions.
- In a call, we first check whether the function contains the tag.

```
;; (mdef quux (x) (list '+ x '1))
```

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(else

Evaluator (define (eval exp env) (cond ((self-evaluating? exp) exp) ((fdef? exp) (eval-fdef exp env)) ((begin? exp) (eval-begin exp env)) ((call? exp) (eval-call exp env))

(error "Unknown expression type -- EVAL" exp))))

((begin? exp)

((call? exp)

(else

(eval-begin exp env))

(eval-call exp env))

Evaluator

Evaluator (define (eval exp env) (cond ((self-evaluating? exp) exp) ((fdef? exp) (eval-fdef exp env)) ((mdef? exp)

(error "Unknown expression type -- EVAL" exp))))

((call? exp)

(else

(eval-call exp env))

Evaluator

(error "Unknown expression type -- EVAL" exp))))

Evaluator

- Don't evaluate the operands.
- Evaluate the returned value (in the current environment).

```
(define (eval-call exp env)
  (let ((func (eval (call-operator exp) env)))
    (if (macro? func)
          (let ((expansion (apply-function func (call-operands exp))))
                (eval expansion env))
                 (apply-function func (eval-exprs (call-operands exp) env)))))
(define (macro? obj)
    (and (function? obj)
                 (equal? (cadr (function-body obj)) "macro")))
```

- Don't evaluate the operands.
- Evaluate the returned value (in the current environment).

```
Evaluator

(define (eval-call exp env)
  (let ((func (eval (call-operator exp) env)))
      (if (macro? func)
            (let ((expansion (apply-function func (call-operands exp))))
                 (eval expansion env))
                  (apply-function func (eval-exprs (call-operands exp) env)))))

(define (macro? obj)
      (and (function? obj)
                  (equal? (cadr (function-body obj)) "macro")))
```

- Don't evaluate the operands.
- Evaluate the returned value (in the current environment).

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- Evaluate the returned value (in the current environment).

From

То

From

```
(and (not (= b 0))
(/ a b))
```

То

```
(if (not (= b 0))
(/ a b)
#f)
```

From

```
(and (not (= b 0))
(/ a b))
```

То

```
(if (not (= b 0))
(/ a b)
#f)
```

```
(mdef and (b0 b1)
  (list 'if
    b0
    b1
    '#f))
```

```
From
(and (not (= b 0))
     (/ a b))
```

```
То
(if (not (= b 0))
  (/ a b)
  #f)
```

```
(mdef and (b0 b1)
  (list 'if
    b0
    b1
    '#f))
```

From (and (not (= b 0)) (/ a b))

```
To

(if (not (= b 0))

(/ a b)

#f)
```

```
(mdef and (b0 b1)
  (list 'if
    b0
    b1
    '#f))
```

From

```
(for i 5 10
(display i))
```



#f)) (loop 5))

Evaluator

From

```
(for i 5 10
(display i))
```


(loop (+ i 1)))

From

```
(for <mark>i</mark> 5 10
(display i))
```


From

```
(for i <mark>5</mark> 10
(display i))
```

То

From (for i 5 10 (display i))

```
Τо
(let ()
  (fdef loop (i)
    (if (< i 10))
      (begin
        (display i)
        (loop (+ i 1)))
      #f))
  (loop 5))
```

```
Transformation
```

```
(mdef for (var inf sup form)
  (list 'let (list)
        (list 'fdef 'loop (list var)
          (list 'if (list '< var sup)
                (list 'begin
                      form
                       (list 'loop (list '+ var 1)))
                '#f))
        (list 'loop inf)))
```

```
From
(for i 5 10
  (display i))
```

```
Τо
(let ()
  (fdef loop (i)
    (if (< i 10))
      (begin
        (display i)
        (loop (+ i 1)))
      #f))
  (loop 5))
```

```
Transformation
```

```
(mdef for (var inf sup form)
  (list 'let (list)
        (list 'fdef 'loop (list var)
          (list 'if (list '< var sup)
                (list 'begin
                      form
                       (list 'loop (list '+ var 1)))
                '#f))
        (list 'loop inf)))
```

```
Example: or
```

```
>> (mdef or (b0 b1)
     (list 'if
       b0
       b0
       b1))
(function (b0 b1) ...)
>> (fdef in-interval (val inf sup)
    (or (and (< val inf)
             inf)
        (or (and (> val sup)
                 sup)
            val)))
(function (val inf sup) ...)
>> (in-interval 1 3 7)
3
```

```
>> (mdef or (b0 b1)
     (list 'if
       b0
       b0
       b1))
(function (b0 b1) ...)
>> (fdef in-interval (val inf sup)
    (or (and (< val inf)
             inf)
        (or (and (> val sup)
                 sup)
            val)))
(function (val inf sup) ...)
>> (in-interval 1 3 7)
3
>> (in-interval 9 3 7)
```

```
>> (mdef or (b0 b1)
     (list 'if
       b0
       b0
       b1))
(function (b0 b1) ...)
>> (fdef in-interval (val inf sup)
    (or (and (< val inf)
             inf)
        (or (and (> val sup)
                 sup)
            val)))
(function (val inf sup) ...)
>> (in-interval 1 3 7)
>> (in-interval 9 3 7)
```

```
>> (mdef or (b0 b1)
     (list 'if
       b0
       b0
       b1))
(function (b0 b1) ...)
>> (fdef in-interval (val inf sup)
    (or (and (< val inf)
             inf)
        (or (and (> val sup)
                 sup)
            val)))
(function (val inf sup) ...)
>> (in-interval 1 3 7)
3
>> (in-interval 9 3 7)
>> (in-interval 4 3 7)
```

```
>> (mdef or (b0 b1)
     (list 'if
       b0
       b0
       b1))
(function (b0 b1) ...)
>> (fdef in-interval (val inf sup)
    (or (and (< val inf)
             inf)
        (or (and (> val sup)
                 sup)
            val)))
(function (val inf sup) ...)
>> (in-interval 1 3 7)
3
>> (in-interval 9 3 7)
>> (in-interval 4 3 7)
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ a b))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ a b))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
3
>> (quotient-or-zero 6 0)
0
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0))
                (/ a b)))
         0))
(function (a b) ...)
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ab)
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
3
>> (quotient-or-zero 6 0)
0
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0))
                (/ a b)))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 0)
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ab)
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
3
>> (quotient-or-zero 6 0)
0
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0))
                (/ a b)))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ab)
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
3
>> (quotient-or-zero 6 0)
0
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0))
                (/ a b)))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
>> (quotient-or-zero 6 2)
```

```
>> (fdef quotient-or-zero (a b)
     (or (and (not (= b 0))
              (/ab)
         0))
(function (a b) ...)
>> (quotient-or-zero 6 2)
3
>> (quotient-or-zero 6 0)
0
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0))
                (/ a b)))
         0))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
>> (quotient-or-zero 6 2)
Division:Division:3 ;; What?
```

Transformation

```
(mdef or (b0 b1)
  (list 'if
    b0
    b0
    b1))
```

From

То

Transformation

```
(mdef or (b0 b1)
  (list 'if
    b0
    b0
    b1))
```

Macro Problems

- Macro calls do not evaluate their arguments.
- Macro expansions might have multiple occurrences of the same parameter.
- Thus, macro expansions might have multiple occurrences of a form that was passed as argument to the macro call.

Solution

- Never repeat a parameter in a macro expansion.
- If necessary, expand into a binding form that binds a local variable to that parameter and reuse the local variable in the expanded code.

Transformation

From (or (begin (display "Division:") (and (not (= b 0)) (/ a b))) 0)

```
То
(let ((val
     (begin
      (if val
  val
  0))
```

Transformation |

```
(mdef or (b0 b1)
  (list 'let (list (list 'val b0))
    (list 'if
      'val
      'val
      b1)))
```

```
Example: or
```

```
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0)) (/ a b)))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
>> (quotient-or-zero 6 2)
Division:3 ;; Now it's OK
>> (in-interval 1 3 7)
```

```
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0)) (/ a b)))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
>> (quotient-or-zero 6 2)
Division:3 ;; Now it's OK
>> (in-interval 1 3 7)
>> (in-interval 9 3 7)
```

Example: or

```
>> (fdef quotient-or-zero (a b)
     (or (begin
           (display "Division:")
           (and (not (= b 0)) (/ a b)))
(function (a b) ...)
>> (quotient-or-zero 6 0)
Division:0
>> (quotient-or-zero 6 2)
Division:3 ;; Now it's OK
>> (in-interval 1 3 7)
>> (in-interval 9 3 7)
error in >: expected real, but got #f, as argument 1
```

From

То

Transformation

Macro Problems

- Macro calls accept forms as arguments.
- Those forms might (will) have free variables.
- When the macro expands into binding forms, the created scope will shadow some of those free variables.
- This is an even more serious problem in a Lisp-1.

Solution

- Use name mangling techniques.
- Use gensyms.
- Use hygienic macros.

From

То

Transformation (without name mangling)

From

То

Transformation (with name mangling)

Macro Problems

- Name mangling is the same as creating a different namespace for the bindings that occur in the macro-expansion.
- But it doesn't completely solve the problem because different macros might use the same (mangled) symbols.

Solution

gensym returns a symbol that is guaranted to be unique.

gensym

Macro Problems

- Name mangling is the same as creating a different namespace for the bindings that occur in the macro-expansion.
- But it doesn't completely solve the problem because different macros might use the same (mangled) symbols.

Solution

gensym returns a symbol that is guaranted to be unique.

gensym

```
(define initial-bindings
  (list ...
        (cons 'eq? (make-primitive eq?))
        (cons 'gensym (make-primitive gensym))))
```

From

То

Transformation (without gensym)

```
(mdef or (b0 b1)
  (list 'let (list (list 'val b0))
     (list 'if
        'val
        'val
        b1)))
```

From

То

Transformation (with gensym)

```
(mdef or (b0 b1)
  (let ((val (gensym)))
    (list 'let (list (list val b0))
        (list 'if
        val
        val
        b1))))
```

Macro Problems

- Gensyms solve the problem of macro-expansions that include binding forms (e.g., lets).
- But they do nothing to solve the problem of macro-expansions that have free variables.
- This is the reverse of the problem of higher-order functions in dynamically scoped languages.

Solution

 In a lexically scoped language, functions capture the surrounding environment.

From

То

From

То

Transformation (with lambda)

```
(mdef or (b0 b1)
    (list 'or-function
        b0
        (list 'lambda (list) b1)))
```

Example

```
(fdef or-function (b0 f1)
   (if b0
        b0
        (f1)))
```

Hygiene

- The macro call expands into a function call that protects the "expanded" code from being evaluated in the scope of the macro call.
- The parameters are wrapped in lambdas that protect the macro arguments from being evaluated in any scope created by the macro expansion.
- Hygienic macros in Scheme operate differently but with the same net result.

Templates

- The macro body is an expression that, when evaluated, computes an expression that, when evaluated, computes the result of the macro call.
- In most cases, the macro expansion can be described by a template of code.
- The template, when evaluated, computes an expression that, when evaluated, computes the result of the macro call.
- To be flexible, a template must allow for some parts of it to be replaced with values computed at expansion time.

Backquote/Quasiquote

- Common Lisp provides the *backquote* form.
- Scheme provides the *quasiquote* form:
 - (quasiquote foo) means "don't evaluate foo except for unquote and unquote-splicing subforms."
 - (unquote foo) means "evaluate foo."
 - (unquote-splicing foo) means "evaluate foo and spread the resulting list."
- (quasiquote foo) can be abbreviated as `foo.
- (unquote foo) can be abbreviated as ,foo.
- (unquote-splicing foo) can be abbreviated as ,@foo.
- \bullet `(x ,x) = (list 'x x)
- (x, x, 0x) = (append (list 'x) (list x) x)



#f)) (loop 5))

Evaluator

From

```
(for i 5 10
(display i))
```


(loop (+ i 1)))

Transformation (without quasiquote)

From

```
(for i 5 10
  (display i))
```

Τо

```
(let ()
  (fdef loop (i)
    (if (< i 10))
      (begin
        (display i)
        (loop (+ i 1)))
      #f))
  (loop 5))
```

Transformation (with quasiquote)

```
(mdef for (var inf sup form)
  `(let ()
     (fdef loop (,var)
       (if (< ,var ,sup)
         (begin
           ,form
           (loop (+ ,var 1)))
         #f))
     (loop ,inf)))
```

```
Evaluator
;; (quasiquote foo)
```

```
Evaluator
;; (quasiquote foo)
(define (quasiquote? exp)
  (and (pair? exp)
          (eq? (car exp) 'quasiquote)))
```

```
Evaluator
;; (quasiquote foo)
(define (quasiquote? exp)
  (and (pair? exp)
          (eq? (car exp) 'quasiquote)))
(define (quasiquoted-form exp)
  (cadr exp))
```

```
Evaluator
;; (quasiquote foo)
(define (quasiquote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quasiquote)))
(define (quasiquoted-form exp)
  (cadr exp))
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
         (eval-quote exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
Evaluator
;; (quasiquote foo)
(define (quasiquote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quasiquote)))
(define (quasiquoted-form exp)
  (cadr exp))
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
         (eval-quote exp env))
        ((quasiquote? exp)
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
Evaluator
;; (quasiquote foo)
(define (quasiquote? exp)
  (and (pair? exp)
       (eq? (car exp) 'quasiquote)))
(define (quasiquoted-form exp)
  (cadr exp))
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
         (eval-quote exp env))
        ((quasiquote? exp)
         (eval-quasiquote exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval-quasiquote exp env)
  (eval (expand (quasiquoted-form exp)) env))
```

```
(define (splicing? form)
  (and (pair? form)
       (or (and (pair? (car form))
                (eq? (caar form) 'unquote-splicing))
           (splicing? (cdr form)))))
(define (expand-list form)
  (if (splicing? form)
      (cons 'append
            (map (lambda (subform)
                   (if (and (pair? subform)
                             (eq? (car subform) 'unquote-splicing))
                        (cadr subform)
                        (list 'list (expand subform))))
                 form))
      (cons 'list (map expand form))))
```

Evaluator (mdef and (b0 b1) (list 'if b0 b1 '#f)) (mdef or (b0 b1) (list 'or-function b0 (list 'lambda (list) b1)))

```
(mdef and (b0 b1)
  (list 'if
        b0
        b1
        '#f))
(mdef and (b0 b1)
 `(if ,b0
     ,b1
     #f))
(mdef or (b0 b1)
  (list 'or-function
        b0
        (list 'lambda (list) b1)))
```

(mdef or (b0 b1)
 `(or-function
 ,b0

(lambda () ,b1)))

Evaluator (mdef and (b0 b1) (list 'if b0 b1 '#f)) (mdef and (b0 b1) `(if ,b0 ,b1 #f)) (mdef or (b0 b1) (list 'or-function b0 (list 'lambda (list) b1)))

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
          exp)
         ((quote? exp)
          (eval-quote exp env))
         ((quasiquote? exp)
          (eval-quasiquote exp env))
         ((name? exp)
          (eval-name exp env))
        ((lambda? exp)
          (eval-lambda exp env))
        ((if? exp)
        (eval-if exp env))
((let? exp) (eval-let exp env))
         ((flet? exp) (eval-flet exp env))
         ((def? exp)
          (eval-def exp env))
         ((set? exp)
          (eval-set exp env))
        ((fdef? exp) (eval-fdef exp env))
((mdef? exp) (eval-mdef exp env))
         ((begin? exp)
          (eval-begin exp env))
         ((call? exp)
          (eval-call exp env))
           (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
          exp)
         ((quote? exp)
          (eval-quote exp env))
         ((quasiquote? exp)
          (eval-quasiquote exp env))
         ((name? exp)
          (eval-name exp env))
        ((lambda? exp)
          (eval-lambda exp env))
        ((if? exp)
        (eval-if exp env))
((let? exp) (eval-let exp env))
         ((flet? exp) (eval-flet exp env))
         ((def? exp)
          (eval-def exp env))
         ((set? exp)
          (eval-set exp env))
        ((fdef? exp) (eval-fdef exp env))
((mdef? exp) (eval-mdef exp env))
         ((begin? exp)
          (eval-begin exp env))
         ((call? exp)
          (eval-call exp env))
           (error "Unknown expression type -- EVAL" exp))))
```

Simplifying the Evaluator

- We defined and or as macros that expand into other forms.
- This means that the evaluator no longer needs to specially identify ands or ors and to provide specific semantics for their evaluation.
- The same can be said for other "special forms."

```
Macros
```

```
(mdef fdef (name parameters body)
  `(def ,name (lambda ,parameters ,body)))
```

Macros

Macros

```
(mdef fdef (name parameters body)
  `(def ,name (lambda ,parameters ,body)))
(mdef flet (binds body)
  `(let ,(map (lambda (form)
                `(,(car form)
                   (lambda , @(cdr form))))
              binds)
     ,body))
(mdef let (binds body)
 `((lambda ,(map car binds)
     ,body)
   , @(map cadr binds)))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
          exp)
         ((quote? exp)
          (eval-quote exp env))
         ((quasiquote? exp)
          (eval-quasiquote exp env))
         ((name? exp)
          (eval-name exp env))
        ((lambda? exp)
          (eval-lambda exp env))
        ((if? exp)
        (eval-if exp env))
((let? exp) (eval-let exp env))
         ((flet? exp) (eval-flet exp env))
         ((def? exp)
          (eval-def exp env))
         ((set? exp)
          (eval-set exp env))
        ((fdef? exp) (eval-fdef exp env))
((mdef? exp) (eval-mdef exp env))
         ((begin? exp)
          (eval-begin exp env))
         ((call? exp)
          (eval-call exp env))
           (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
         (eval-quote exp env))
        ((quasiquote? exp)
         (eval-quasiquote exp env))
        ((name? exp)
         (eval-name exp env))
        ((lambda? exp)
         (eval-lambda exp env))
        ((if? exp)
         (eval-if exp env))
        ((let? exp) (eval-let exp env))
        ((flet? exp) (eval-flet exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((fdef? exp) (eval-fdef exp env))
((mdef? exp) (eval-mdef exp env))
        ((begin? exp)
         (eval-begin exp env))
        ((call? exp)
         (eval-call exp env))
          (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
          (eval-quote exp env))
        ((quasiquote? exp)
(eval-quasiquote exp env))
((name? exp)
          (eval-name exp env))
        ((lambda? exp)
          (eval-lambda exp env))
        ((if? exp)
          (eval-if exp env))
        ((def? exp)
          (eval-def exp env))
        ((set? exp)
          (eval-set exp env))
        ((mdef? exp) (eval-mdef exp env))
        ((begin? exp)
         (eval-begin exp env))
        ((call? exp)
          (eval-call exp env))
        (else
           (error "Unknown expression type -- EVAL" exp))))
```

Bootstrapping Macros

Macro-defining Macro

- A macro is just a tagged function.
- An mdef call is a macro call that expands into a macro definition.
- We no longer need to specially recognize macro definitions in the evaluator.

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((quote? exp)
          (eval-quote exp env))
        ((quasiquote? exp)
(eval-quasiquote exp env))
((name? exp)
          (eval-name exp env))
        ((lambda? exp)
          (eval-lambda exp env))
        ((if? exp)
          (eval-if exp env))
        ((def? exp)
          (eval-def exp env))
        ((set? exp)
          (eval-set exp env))
        ((mdef? exp) (eval-mdef exp env))
        ((begin? exp)
         (eval-begin exp env))
        ((call? exp)
          (eval-call exp env))
        (else
           (error "Unknown expression type -- EVAL" exp))))
```

((def? exp)
 (eval-def exp env))
((set? exp)
 (eval-set exp env))
((mdef? exp) (eval-mdef exp env))

((begin? exp) (eval-begin exp env)) ((call? exp) (eval-call exp env))

(else

(error "Unknown expression type -- EVAL" exp))))

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
         ((quote? exp)
          (eval-quote exp env))
        ((quasiquote? exp)
(eval-quasiquote exp env))
((name? exp)
          (eval-name exp env))
         ((lambda? exp) (eval-lambda exp env))
         ((if? exp)
          (eval-if exp env))
         ((def? exp)
          (eval-def exp env))
         ((set? exp)
          (eval-set exp env))
         ((begin? exp)
          (eval-begin exp env))
         ((call? exp)
          (eval-call exp env))
           (error "Unknown expression type -- EVAL" exp))))
```

```
(eval
  '(begin
     (def mdef (lambda (name parameters body) "macro"
                 '(def ,name (lambda ,parameters "macro" ,body))))
     (mdef fdef (name parameters body)
           '(def ,name (lambda ,parameters ,body)))
     (mdef and (b0 b1) (if .b0 .b1 #f))
     (mdef or (b0 b1) '(or-function ,b0 (lambda () ,b1)))
     (fdef or-function (b0 f1) (if b0 b0 (f1)))
     (fdef not (b) (if b #f #t))
     (fdef map (f 1)
           (if (null? 1) (list) (cons (f (car 1)) (map f (cdr 1)))))
     (fdef append (10 11)
           (if (null? 10) 11 (cons (car 10) (append (cdr 10) 11))))
     (fdef caar (p)
           (car (car p)))
     (fdef cadr (p)
           (car (cdr p)))
     (mdef flet (binds body)
           `(let ,(map (lambda (form)
                          `(,(car form)
                           (lambda ,@(cdr form))))
                       binds)
              ,body))
     (mdef let (binds body)
           '((lambda ,(map car binds)
               , body)
             ,@(map cadr binds))))
 initial-environment)
```

```
>> (mdef define (name form)
     (if (pair? name)
       `(def ,(car name)
          (lambda, (cdr name)
            ,form))
       `(def ,name ,form)))
(function (name form) ...)
>> (define pi 3.14159)
3.14159
>> (define (fact n)
     (if (= n 0))
       (* n (fact (- n 1)))))
(function (n) ...)
>> (fact 10)
3628800
```

MetaCircular Evaluator

- We described an operational model of the language.
- The advantage of operational models is that they are executable, meaning that we can input a program and output its evaluation according to the rules of the language.
- This is useful to understand the semantics of a language but also to experiment with different semantics.
- Another important result: programs are data for the evaluator:
 - The evaluator introspects the program.
 - The evaluator manipulates the program.
 - The evaluator creates a program.
 - The evaluator interprets the program.

MetaCircular Evaluator

- But the programs do not have the same capability:
 - Although it is possible to write a program that computes, introspects or manipulates another program, we cannot interpret the generated program.
 - Macros only solve part of the problem: a macro call statically computes a new program fragment.
- To provide our programs with evaluation capabilities, we need to build an evaluator.
- Simplest solution: the language of our programs is the same as the language of the evaluator so we can simply input the evaluator to the evaluator and create another level of interpretation.
- Best solution: we can provide the evaluator as a primitive operation (but we also need to reify environments).

```
Evaluator
```

```
;;(current-environment)
```

```
Evaluator
```

```
;;(current-environment)
(define (current-environment? exp)
  (and (pair? exp)
      (eq? (car exp) 'current-environment)))
```

```
;;(current-environment)
(define (current-environment? exp)
  (and (pair? exp)
       (eq? (car exp) 'current-environment)))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((current-environment? exp)
         env)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;;(current-environment)
(define (current-environment? exp)
  (and (pair? exp)
       (eq? (car exp) 'current-environment)))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((current-environment? exp)
         env)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define initial-bindings
  (list (cons '+ (make-primitive +))
```

```
;;(current-environment)
(define (current-environment? exp)
  (and (pair? exp)
       (eq? (car exp) 'current-environment)))
(define (eval exp env)
  (cond ((self-evaluating? exp) exp)
        ((current-environment? exp)
         env)
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define initial-bindings
  (list (cons '+ (make-primitive +))
        (cons 'eval (make-primitive eval))))
```

```
>> (eval (list '+ '1 '2) (current-environment))
```

```
>> (eval (list '+ '1 '2) (current-environment))
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
(function (kons) ...)
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
(function (kons) ...)
>> (kar (kons 1 2))
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
(function (kons) ...)
>> (kar (kons 1 2))
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
(function (kons) ...)
>> (kar (kons 1 2))
>> (kdr (kons 1 2))
```

```
>> (eval (list '+ '1 '2) (current-environment))
>> (define my-env
     (let ((x 1) (y 2))
       (current-environment)))
>> (+ (eval 'x my-env) (eval 'y my-env))
>> (define (kons kar kdr)
     (current-environment))
(function (kar kdr) ...)
>> (define (kar kons)
     (eval 'kar kons))
(function (kons) ...)
>> (define (kdr kons)
     (eval 'kdr kons))
(function (kons) ...)
>> (kar (kons 1 2))
>> (kdr (kons 1 2))
```

```
>> (mdef defclass (name slots)
     `(fdef ,name ,slots
        (current-environment)))
(function (name slots) ...)
>> (fdef slot-value (instance slot)
     (eval slot instance))
(function (instance slot) ...)
>> (defclass person
     (name age sex))
(function (name age sex) ...)
>> (def john (person "John Adams" 42 "male"))
>> (def sally (person "Sally Field" 64 "female"))
```

```
>> (mdef defclass (name slots)
     `(fdef ,name ,slots
        (current-environment)))
(function (name slots) ...)
>> (fdef slot-value (instance slot)
     (eval slot instance))
(function (instance slot) ...)
>> (defclass person
     (name age sex))
(function (name age sex) ...)
>> (def john (person "John Adams" 42 "male"))
>> (def sally (person "Sally Field" 64 "female"))
>> (slot-value john 'age)
```

```
>> (mdef defclass (name slots)
     `(fdef ,name ,slots
        (current-environment)))
(function (name slots) ...)
>> (fdef slot-value (instance slot)
     (eval slot instance))
(function (instance slot) ...)
>> (defclass person
     (name age sex))
(function (name age sex) ...)
>> (def john (person "John Adams" 42 "male"))
>> (def sally (person "Sally Field" 64 "female"))
>> (slot-value john 'age)
42
```

```
>> (mdef defclass (name slots)
     `(fdef ,name ,slots
        (current-environment)))
(function (name slots) ...)
>> (fdef slot-value (instance slot)
     (eval slot instance))
(function (instance slot) ...)
>> (defclass person
     (name age sex))
(function (name age sex) ...)
>> (def john (person "John Adams" 42 "male"))
>> (def sally (person "Sally Field" 64 "female"))
>> (slot-value john 'age)
42
>> (slot-value sally 'name)
```

```
>> (mdef defclass (name slots)
     `(fdef ,name ,slots
        (current-environment)))
(function (name slots) ...)
>> (fdef slot-value (instance slot)
     (eval slot instance))
(function (instance slot) ...)
>> (defclass person
     (name age sex))
(function (name age sex) ...)
>> (def john (person "John Adams" 42 "male"))
>> (def sally (person "Sally Field" 64 "female"))
>> (slot-value john 'age)
42
>> (slot-value sally 'name)
"Sally Field"
```

```
Example
```

```
>> (mdef send (obj message)
  `(eval ',message ,obj))
```

```
>> (mdef send (obj message)
   `(eval ',message ,obj))
(function (obj message) ...)
```

```
>> (mdef send (obj message)
    `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n))
                                  (if (= n 0)
                                      (* n (fact (- n 1))))))
 (function (n) ...)
```

```
>> (mdef send (obj message)
     `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n)
                 (if (= n 0))
                   (* n (fact (- n 1))))))
(function (n) ...)
>> (send sally (fdef fact (what)
                  (begin
                    (display "It's a fact that ")
                    (display what)
                    (newline))))
(function (what) ...)
```

```
>> (mdef send (obj message)
     `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n)
                 (if (= n 0))
                   (* n (fact (- n 1))))))
(function (n) ...)
>> (send sally (fdef fact (what)
                  (begin
                    (display "It's a fact that ")
                    (display what)
                    (newline))))
(function (what) ...)
>> (send john (fact 10))
```

```
>> (mdef send (obj message)
     `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n)
                 (if (= n 0))
                   (* n (fact (- n 1))))))
(function (n) ...)
>> (send sally (fdef fact (what)
                  (begin
                    (display "It's a fact that ")
                    (display what)
                    (newline))))
(function (what) ...)
>> (send john (fact 10))
3628800
```

```
>> (mdef send (obj message)
     `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n)
                 (if (= n 0))
                   (* n (fact (- n 1))))))
(function (n) ...)
>> (send sally (fdef fact (what)
                  (begin
                    (display "It's a fact that ")
                    (display what)
                    (newline))))
(function (what) ...)
>> (send john (fact 10))
3628800
>> (send sally (fact "Lisp is great!"))
```

```
>> (mdef send (obj message)
     `(eval ',message ,obj))
(function (obj message) ...)
>> (send john (fdef fact (n)
                (if (= n 0))
                  (* n (fact (- n 1))))))
(function (n) ...)
>> (send sally (fdef fact (what)
                 (begin
                   (display "It's a fact that ")
                   (display what)
                    (newline))))
(function (what) ...)
>> (send john (fact 10))
3628800
>> (send sally (fact "Lisp is great!"))
It's a fact that Lisp is great!
```

Continuations

- There are operations whose semantics is difficult to express in the "usual" computational model:
 - error
 - throw
 - yield
 - suspend
- These operations seem to violate the semantics of function calls because either they don't return to the caller (error, throw) or they allow a function to resume its computation after returning (yield, suspend).
- We need a different computational model that allows these operations.

```
>> (define (mystery a b c d e)
	(+ a (* b (/ (- c d) e))))
...
>> (mystery 5 4 3 2 1)
```

Problem

```
>> (define (mystery a b c d e)
	(+ a (* b (/ (- c d) e))))
...
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
ERROR ;; and the evaluator stops!!!!
```

Problem

```
>> (define (mystery a b c d e)
          (+ a (* b (/ (- c d) e))))
...
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
ERROR ;; and the evaluator stops!!!!
```

```
(define (mystery a b c d e)
(+ a (* b (safe-/ (- c d) e))))
```

Problem

```
>> (define (mystery a b c d e)
	(+ a (* b (/ (- c d) e))))
...
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
ERROR ;; and the evaluator stops!!!!
```

```
(define (mystery a b c d e)
  (+ a (* b (safe-/ (- c d) e))))
(define (safe-/ x y)
  (if (= y 0)
        (error "Can't divide " (list x y))
        (/ x y)))
```

Problem

```
>> (define (mystery a b c d e)
	(+ a (* b (/ (- c d) e))))
...
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
ERROR ;; and the evaluator stops!!!!
```

```
(define (mystery a b c d e)
  (+ a (* b (safe-/ (- c d) e))))

(define (safe-/ x y)
  (if (= y 0)
        (error "Can't divide " (list x y))
        (/ x y)))

(define (error msg args)
```

Problem

```
(define (mystery a b c d e)
    (+ a (* b (safe-/ (- c d) e))))
(define (safe-/ x y)
    (if (= y 0)
          (error "Can't divide " (list x y))
          (/ x y)))
(define (error msg args)
    ???)
```

Errors

```
char *ptr = malloc(200000000UL);
if (ptr == NULL) {
 /* allocation failed, but what can I do? */
} else {
```

Signaling Errors

- Some languages (e.g., C) use unusual values as errors (a negative number when a positive number is expected, NULL when a pointer is expected, etc).
- Each caller in the call chain that caused the error must check the return of the called function.
- This is not practical.

Errors char *ptr = malloc(200000000UL); if (ptr == NULL) { throw "Allocation Error" } else { ... }

Signaling Errors

- ullet Some languages (e.g., C++) allow non-local transfers of control.
- The throw statement transfers control to the most recent try-catch statement that handles it.
- This is practical.
- Note: Java's Checked Exceptions are *not* practical.

Handling Errors

- It is generally difficult to handle exceptional situations in the callee.
- Each callee can be called from many different places.
- The callee rarely knows how to handle exceptional situations in a way that pleases all callers.
- Non-local transfers of control allow the callee to signal that an exceptional situation ocurred...
- ... and the control is transferred to the most recent handler established in the call chain.
- This is impossible to do with function calls...
- ... isn't it?



- During a function call, the caller must wait for the callee to finish its computation and return a value.
- After the return of the callee, the caller proceeds with its computation.
- This computation is the continuation of the call.
- The continuation of the call is whatever remains to be done after the return from the call.
- In most high-level languages, the continuation is implicitly stored in a stack.
- In languages with first-class (and higher-order) functions, the *continuation* can be explicitly represented with functions.

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (+ a (* b (/ r0 e)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (let ((r1 (/ r0 e)))
        (+ a (* b r1)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (+ a (* b (/ r0 e)))))
(define (mystery a b c d e)
        (let ((r0 (- c d)))
            (+ a (* b r1))))
(define (mystery a b c d e)
        (let ((r0 (- c d)))
            (+ a (* b r1))))
(define (mystery a b c d e)
        (let ((r0 (- c d)))
            (let ((r1 (/ r0 e)))
            (let ((r1 (/ r0 e)))
            (+ a r2)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (+ a (* b (/ r0 e)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (let ((r1 (/ r0 e)))
              (+ a (* b r1)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
        (let ((r1 (/ r0 e)))
              (let ((r2 (* b r1)))
              (let ((r2 (* b r1)))
              (+ a r2)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (+ a (* b (/ r0 e))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (+ a (* b r1)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (+ a r2)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (let ((r3 (+ a r2)))
          r3)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  ((\lambda (r0) (+ a (* b (/ r0 e))))
   (-cd))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (+ a (* b r1)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (+ a r2)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (let ((r3 (+ a r2)))
          r3)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  ((\lambda (r0) (+ a (* b (/ r0 e))))
   (-cd))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1) (+ a (* b r1)))
            (/ r0 e)))
   (-cd))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (+ a r2)))))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (let ((r3 (+ a r2)))
          r3)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  ((\lambda (r0) (+ a (* b (/ r0 e))))
   (- c d)))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1) (+ a (* b r1)))
             (/ r0 e)))
   (-cd))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1))((\lambda (r2)) + a r2))
                      (* b r1)))
  (- c d)))
(define (mystery a b c d e)
  (let ((r0 (- c d)))
    (let ((r1 (/ r0 e)))
      (let ((r2 (* b r1)))
        (let ((r3 (+ a r2)))
          r3)))))
```

```
(define (mystery a b c d e)
  (+ a (* b (/ (- c d) e))))
(define (mystery a b c d e)
  ((\lambda (r0) (+ a (* b (/ r0 e))))
   (-cd))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1) (+ a (* b r1)))
             (/ r0 e)))
   (-cd))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1))((\lambda (r2)) + a r2))
                       (* b r1)))
   (- c d)))
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1))((\lambda (r2))((\lambda (r3))r3))
                                 (+ a r2)))
  (- c d))) (/ r0 e))) (* b r1)))
```

Simplification

• In order to simplify the program, we will define functions that call the continuation with the value of the numeric operation.

Functions Accepting Explicit Continuations

```
(define (+c x y cont)
  (cont (+ x y)))
(define (-c x y cont)
  (cont (- x y)))
(define (*c x y cont)
  (cont (* x y)))
(define (/c x y cont)
  (cont (/ x y)))
```

```
(define (mystery a b c d e)
  ((\lambda (r0))((\lambda (r1))((\lambda (r2))((\lambda (r3))r3))
  (- c d))) (* b r1)))
                                   (+ a r2)))
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) ((\lambda (r1) ((\lambda (r2) ((\lambda (r3) r3)
                           (* b r1)))
                 (/ r0 e)))))
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                     (\lambda (r1) ((\lambda (r2) ((\lambda (r3) r3)
                                         (+ a r2)))
                                (* b r1)))))))
```

```
(define (mystery a b c d e) 
 (-c c d 
 (\lambda (r0) (/c r0 e 
 (\lambda (r1) ((\lambda (r2) ((\lambda (r3) r3) 
 (+ a r2))) 
 (* b r1))))))
```

```
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                       (\lambda (r1) ((\lambda (r2) ((\lambda (r3) r3)
                                              (+ a r2)))
                                  (* b r1)))))))
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                       (\lambda (r1) (*c b r1)
                                        (\lambda (r2) ((\lambda (r3) r3)
                                                   (+ a r2))))))))))
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                       (\lambda (r1) (*c b r1)
                                      (\lambda (r2) (+c a r2)
                                                      (\lambda (r3) r3)))))))))
```

Function mystery

```
(define (mystery a b c d e)
(+ a (* b (/ (- c d) e))))
```

Function mystery

```
(define (mystery a b c d e)

(-c c d

(\lambda (r0) (/c r0 e

(\lambda (r1) (*c b r1

(\lambda (r2) (+c a r2

(\lambda (r3) r3))))))))
```

- Each function now accepts an explicit continuation.
- In order to "return", the function calls the continuation.

Function mystery

```
(define (mystery a b c d e)
(+ a (* b (/ ( c d) e))))
```

Function mystery

```
(define (mystery a b c d e) 

(-c c d (\lambda (r0) (/c r0 e (\lambda (r1) (*c b r1 (\lambda (r2) (+c a r2 (\lambda (r3) r3))))))))
```

- Each function now accepts an explicit continuation.
- In order to "return", the function calls the continuation.

Function mystery

```
(define (mystery a b c d e)
(+ a (* b (<mark>/</mark> (- c d) e))))
```

Function mystery

```
(define (mystery a b c d e)

(-c c d

(\lambda (r0) (/c r0 e

(\lambda (r1) (*c b r1

(\lambda (r2) (+c a r2

(\lambda (r3) r3))))))))
```

- Each function now accepts an explicit continuation.
- In order to "return", the function calls the continuation.

Function mystery

```
(define (mystery a b c d e)
(+ a (* b (/ (- c d) e))))
```

Function mystery

```
(define (mystery a b c d e) 
 (-c c d 
 (\lambda (r0) (/c r0 e 
 (\lambda (r1) (*c b r1 
 (\lambda (r2) (+c a r2 
 (\lambda (r3) r3))))))))
```

- Each function now accepts an explicit continuation.
- In order to "return", the function calls the continuation.

Function mystery

```
(define (mystery a b c d e)
(+ a (* b (/ (- c d) e))))
```

Function mystery

```
(define (mystery a b c d e)
  (-c c d
        (\(\lambda\) (r0) (/c r0 e
                        (\lambda (r1) (*c b r1)
                                        (\lambda (r2) (+c a r2)
                                                         (\(\lambda\) (r3) r3)))))))))
```

- Each function now accepts an explicit continuation.
- In order to "return", the function calls the continuation.

Signalling Errors

- A "normal" function call always returns to its caller...
- ... because the continuation is implicit.
- When the continuation is explicit, it is possible to not call it, i.e., it is possible to not return to its caller.

Function /c

```
(define (/c x y cont)
  (cont (/ x y)))
```

Function /c that Signals an Error

```
(define (/c x y cont)
  (if (= y 0)
        (error "Can't divide " (list x y))
        (cont (/ x y))))
```

```
(define (mystery a b c d e)
  (-c c d
      (\lambda (r0) (/c r0 e)
                     (\lambda (r1) (*c b r1)
                                   (\lambda (r2) (+c a r2)
                                                 (\(\lambda\) (r3) r3))))))))))
(define (/c x y cont)
  (if (= v 0)
    (error "Can't divide " (list x y))
    (cont (/ x y))))
(define (error msg args)
  (display msg)
  (display args))
```

```
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                     (\lambda (r1) (*c b r1)
                                   (\lambda (r2) (+c a r2)
                                                 (\(\lambda\) (r3) r3))))))))))
(define (/c x y cont)
  (if (= y 0)
    (error "Can't divide " (list x y))
    (cont (/ x y))))
(define (error msg args)
  (display msg)
  (display args))
>> (mystery 5 4 3 2 1)
```

```
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                     (\lambda (r1) (*c b r1)
                                   (\lambda (r2) (+c a r2)
                                                 (\(\lambda\) (r3) r3))))))))))
(define (/c x y cont)
  (if (= y 0)
    (error "Can't divide " (list x y))
    (cont (/ x y))))
(define (error msg args)
  (display msg)
  (display args))
>> (mystery 5 4 3 2 1)
9
```

```
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                    (\lambda (r1) (*c b r1)
                                  (\lambda (r2) (+c a r2)
                                                (\lambda (r3) r3))))))))))
(define (/c x y cont)
  (if (= y 0)
    (error "Can't divide " (list x y))
    (cont (/ x y))))
(define (error msg args)
  (display msg)
  (display args))
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
```

```
(define (mystery a b c d e)
  (-c c d
       (\lambda (r0) (/c r0 e)
                    (\lambda (r1) (*c b r1)
                                  (\lambda (r2) (+c a r2)
                                                (\lambda (r3) r3))))))))))
(define (/c x y cont)
  (if (= v 0)
    (error "Can't divide " (list x y))
    (cont (/ x y))))
(define (error msg args)
  (display msg)
  (display args))
>> (mystery 5 4 3 2 1)
9
>> (mystery 4 3 2 1 0)
Can't divide (1 0)
```

Are we done?

>> (+ 1 (mystery 5 4 3 2 1))

Are we done?

```
>> (+ 1 (mystery 5 4 3 2 1))
```

Are we done?

```
>> (+ 1 (mystery 5 4 3 2 1))
10
>> (+ 1 (mystery 4 3 2 1 0))
```

Are we done?

```
>> (+ 1 (mystery 5 4 3 2 1))
10
>> (+ 1 (mystery 4 3 2 1 0))
ERROR ;; and the evaluator stops!!!!
```

Errors

- We converted the primitive operations to use explicit continuations.
- But we didn't convert the mystery function to use an explicit continuation.
- When the function /c calls error, the returned value of the error function is also returned from mystery.
- The expression (+ 1 (mystery ...)) then tries to add 1 to that returned value.

```
(define (mystery a b c d e) 
 (-c c d 
 (\lambda (r0) (/c r0 e 
 (\lambda (r1) (*c b r1 
 (\lambda (r2) (+c a r2 
 (\lambda (r3) r3))))))))
```

Example

```
(define (mystery-c a b c d e cont)
  (-c c d
       (\lambda (r0) (/c r0 e)
                      (\lambda (r1) (*c b r1)
                                      (\lambda (r2) (+c a r2)
                                                     (\lambda (r3) (cont r3)))))))))))
```

Including a continuation

- The function must accept a continuation.
- The function "transfers" control to (i.e., calls) the continuation with the final result.

Example

Including a continuation

- The function must accept a continuation.
- The function "transfers" control to (i.e., calls) the continuation with the final result.

Example

```
(define (mystery-c a b c d e cont)
  (-c c d
       (\(\lambda\) (r0) (/c r0 e
                      (\lambda (r1) (*c b r1)
                                      (\lambda (r2) (+c a r2)
                                                     cont))))))))
```

Including a continuation

- (lambda (x) (f x)) = f
- This is η -reduction.

Example

```
>> (+ 1 (mystery 5 4 3 2 1))
```

```
>> (mystery-c 5 4 3 2 1 (lambda (r) (+ 1 r)))
```

Example

```
>> (+ 1 (mystery 5 4 3 2 1))
10
```

```
>> (mystery-c 5 4 3 2 1 (lambda (r) (+ 1 r)))
10
```

Example

```
>> (+ 1 (mystery 5 4 3 2 1))
10
>> (+ 1 (mystery 4 3 2 1 0))
```

```
>> (mystery-c 5 4 3 2 1 (lambda (r) (+ 1 r)))
10
>> (mystery-c 4 3 2 1 0 (lambda (r) (+ 1 r)))
```

Example

```
>> (+ 1 (mystery 5 4 3 2 1))
10
>> (+ 1 (mystery 4 3 2 1 0))
ERROR ;; and the evaluator stops!!!!
```

```
>> (mystery-c 5 4 3 2 1 (lambda (r) (+ 1 r)))
10
>> (mystery-c 4 3 2 1 0 (lambda (r) (+ 1 r)))
Can't divide (1 0)
```

Example

```
>> (+ 1 (mystery 5 4 3 2 1))
10
>> (+ 1 (mystery 4 3 2 1 0))
ERROR ;;and the evaluator stops!!!!
```

Example

```
>> (mystery-c 5 4 3 2 1 (lambda (r) (+ 1 r)))
10
>> (mystery-c 4 3 2 1 0 (lambda (r) (+ 1 r)))
Can't divide (1 0)
```

Continuations

- With implicit continuations, it is impossible for the program to influence the control flow from callers to callees.
- With explicit continuations, on each function call, the program decides what to do next.

Continuations

- Implicit continuation ⇒ program in direct style.
- Explicitly continuation ⇒ program in continuation-passing style.

foo and bar in Direct Style

(foo 2 (bar 3 4))

Continuations

- Implicit continuation ⇒ program in direct style.
- Explicitly continuation ⇒ program in continuation-passing style.

foo and bar in Direct Style

```
(foo 2 (bar 3 4))
```

bar in Continuation Passing Style

```
(bar-c 3 4
(lambda (r0) (foo 2 r0)))
```

Continuations

- Implicit continuation ⇒ program in direct style.
- Explicitly continuation ⇒ program in continuation-passing style.

foo and bar in Direct Style

```
(foo 2 (bar 3 4))
```

bar in Continuation Passing Style

```
(bar-c 3 4
(lambda (r0) (foo 2 r0)))
```

foo and bar in Continuation Passing Style

```
(bar-c 3 4
(lambda (r0) (foo-c 2 r0
(lambda (r1) ...))))
```



Continuations

- Implicit continuation \Rightarrow program in *direct style*.
- Explicitly continuation \Rightarrow program in continuation-passing style.

Direct Style

```
(define (fact n)
  (if (= n 0)
      (* n (fact (- n 1)))))
```

```
(define (fact n c)
  (if (= n 0)
      (c 1)
      (fact (- n 1)
            (lambda (r) (c (* n r))))))
```

Direct Style

>> (fact 10)

Continuation Passing Style

>> (fact 10 (lambda (r) r))

Direct Style

>> (fact 10) 3628800

```
>> (fact 10 (lambda (r) r)) 3628800
```

Direct Style >> (fact 10)

```
3628800
>> (+ (fact 10)
      (fact 5))
```

```
>> (fact 10 (lambda (r) r))
3628800
>> (fact 10 (lambda (r0)
              (fact 5 (lambda (r1) (+ r0 r1)))))
```

Direct Style

```
>> (fact 10)
3628800
>> (+ (fact 10)
      (fact 5))
3628920
```

```
>> (fact 10 (lambda (r) r))
3628800
>> (fact 10 (lambda (r0)
              (fact 5 (lambda (r1) (+ r0 r1)))))
3628920
```

Direct Style

Direct Style

Direct Style

In continuation-passing style:

- The continuation is represented as a function.
- The explicit representation of a continuation is the reification of the continuation.
- There are no hidden pending computations.
- Each computed result is explicitely named.
- All calls become tail calls.
- All recursive calls become tail recursive calls.
- A stack is not required.

Continuations

• The reification of a continuation allows its capture.

```
>> (define pending-fact-computations #f)
#f
```

Continuations

• The reification of a continuation allows its capture.

Continuations

The reification of a continuation allows its capture.

Continuations

The reification of a continuation allows its capture.

Continuations

The reification of a continuation allows its capture.

Continuations

• The reification of a continuation allows its capture.

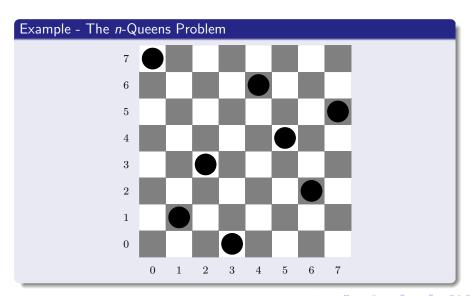
Continuations

The reification of a continuation allows its capture.

Continuations

• The reification of a continuation allows its capture.

```
>> (define pending-fact-computations #f)
#f
>> (define (fact n c)
     (if (= n 0)
       (begin
         (set! pending-fact-computations c)
         (c, 1)
       (fact (- n 1) (lambda (r) (c (* n r))))))
>> (fact 10 (lambda (r) r))
3628800
>> (pending-fact-computations 1)
3628800
>> (pending-fact-computations 10)
36288000
```



Example - The *n*-Queens Problem

- Place n Queens in a $n \times n$ chessboard so that no two queens attack each other.
- A solution requires that no two queens share the same row, column, or diagonal.

Is position (i_0, j_0) threatened by (i_1, j_1) ?

```
(define (threat? i0 j0 i1 j1)
  (or (= i0 i1)
        (= j0 j1)
        (= (+ i0 j0) (+ i1 j1))
        (= (- i0 j0) (- i1 j1))))
```

Example - The *n*-Queens Problem

- We will represent the current board by the list of occupied positions.
- We will represent each position by the pair (row. column).

Is position (i, j) attacked by the currently placed queens?

Solving the puzzle: Using backtracking

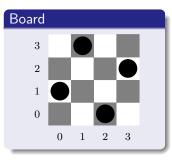
```
(define (queens n i j board)
  (if (= in)
       board
       (if (= j n)
            #f
            (if (attacked? i j board)
                 (queens n i (+ j 1) board)
                 (or (queens n (+ i 1) 0 (cons (cons i j) board))
      (queens n i (+ j 1) board)))))
```

Solving the puzzle: Using backtracking

Solution

>> (solve-queens 4)

Solution

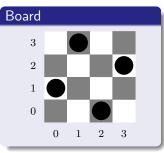


Solution

>> (solve-queens 4) ((3 . 2) (2 . 0) (1 . 3) (0 . 1))

Problem

- That's just one solution.
- Can we see the next one?

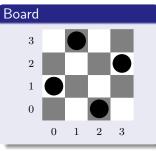


Solution

>> (solve-queens 4) ((3 . 2) (2 . 0) (1 . 3) (0 . 1))

Problem

- That's just one solution.
- Can we see the next one?



Solution

- Clarify the choices made during the process of problem solving.
- "Either α or β " \Longrightarrow "if α then α else β ".

Choices - using or

```
(or (queens n (+ i 1) 0 (cons (cons i j) board))
      (queens n i (+ j 1) board)))
```

Choices - using or

```
(or (queens n (+ i 1) 0 (cons (cons i j) board))
     (queens n i (+ j 1) board)))
```

Choices - using if

```
(let ((solution?
       (queens n (+ i 1) 0 (cons (cons i j) board))))
  (if solution?
      solution?
      (queens n i (+ j 1) board)))
```

Choices

If a non-attacked place is found:

Choices - using or

```
(or (queens n (+ i 1) 0 (cons (cons i j) board))
(queens n i (+ j 1) board)))
```

Choices - using if

Choices

If a non-attacked place is found:

• Either place a queen and try to solve the rest of the board...

Choices - using or

```
(or (queens n (+ i 1) 0 (cons (cons i j) board))
     (queens n i (+ j 1) board)))
```

Choices - using if

Choices

If a non-attacked place is found:

- Either place a queen and try to solve the rest of the board...
- ...or continue the search for a non-attacked place.



Next Solution

- The first solution is the sequence of the choices that avoided dead ends.
- In order to see the next solution, it is sufficient to consider the last successfull choice as a dead end.
- Currently, the choices are just pending computations...
- ...that disappear as soon as the first solution is returned.

Next Solution

- The first solution is the sequence of the choices that avoided dead ends.
- In order to see the next solution, it is sufficient to consider the last successfull choice as a dead end.
- Currently, the choices are just pending computations...
- ...that disappear as soon as the first solution is returned.
- But if we transform the program into continuation-passing style:
 - It becomes possible to save the sequence of choices that computed a solution.
 - It becomes possible to restore the sequence of choices using failure as an argument.

Solving the puzzle (in direct style)

Solving the puzzle (in direct style)

Transformation

• Include additional parameter for the continuation.

Solving the puzzle (in direct style)

Solving the puzzle (in direct style)

```
(define (queens n i j board c)
  (if (= i n)
     board
      (if (= j n)
          (if (attacked? i j board)
              (queens n i (+ j 1) board)
              (let ((solution?
                     (queens n (+ i 1) 0 (cons (cons i j) board))))
                (if solution?
                    solution?
                    (queens n i (+ j 1) board)))))))
```

Transformation

Each returned value becomes the argument of a continuation call.

Solving the puzzle (in direct style)

Solving the puzzle (in direct style)

Transformation

• Each tail call receives the current continuation.

(let ((solution?

(if solution?

Evaluator

(queens n i (+ j 1) board c)))))))

(queens n i (+ j 1) board c)

(c solution?)

(queens n (+ i 1) 0 (cons (cons i j) board))))

Solving the puzzle (in direct style)

```
(define (queens n i j board c)
  (if (= i n)
      (c board)
      (if (= j n)
          (c #f)
          (if (attacked? i j board)
              (queens n i (+ j 1) board c)
              (let ((solution?
                      (queens n (+ i 1) 0 (cons (cons i j) board))))
                (if solution?
                     (c solution?)
                     (queens n i (+ j 1) board c)))))))
```

Transformation

Each pending computation becomes a new continuation.

Solving the puzzle (in direct style)

Solving the puzzle (in direct style)

Transformation

• The initial continuation is the identity function.

```
(define (queens n i j board c)
  (if (= i n)
      (c board)
      (if (= j n)
           (c #f)
           (if (attacked? i j board)
               (queens n i (+ j 1) board c)
(queens n (+ i 1) 0 (cons (cons i j) board)
                         (lambda (solution?)
                           (if solution?
                                (c solution?)
                                (queens n i (+ j 1) board c))))))))
(define (solve-queens n)
  (queens n 0 0 (list) (lambda (solution) solution)))
```

Solving the puzzle (in continuation-passing style)

```
(define (queens n i j board c)
  (if (= i n)
      (c board)
      (if (= j n)
           (c #f)
           (if (attacked? i j board)
                (queens n i (+ j 1) board c)
(queens n (+ i 1) 0 (cons (cons i j) board)
                         (lambda (solution?)
                           (if solution?
                                (c solution?)
                                (queens n i (+ j 1) board c))))))))
(define (solve-queens n)
  (queens n 0 0 (list) (lambda (solution) solution)))
```

Transformation

Done!

Using Continuations

 If a solution is found, the continuation is called with the current board.

Using Continuations

• If a dead end is found, the continuation is called with false.

Using Continuations

• If the current position is attacked, we try the next one, using the same continuation.

Using Continuations

 If the current position is not attacked, we place a queen and we solve the rest of the problem with a new continuation.

Using Continuations

When the continuation receives the result ...

Solving the puzzle (in continuation-passing style)

```
(define (queens n i j board c)
  (if (= i n)
       (c board)
       (if (= j n)
           (c #f)
           (if (attacked? i j board)
                (queens n i (+ j 1) board c)
(queens n (+ i 1) 0 (cons (cons i j) board)
                          (lambda (solution?)
                            (if solution?
                                 (c solution?)
                                 (queens n i (+ j 1) board c))))))))
```

Using Continuations

we test it to see if it is a solution.

Solving the puzzle (in continuation-passing style)

```
(define (queens n i j board c)
  (if (= i n)
       (c board)
       (if (= j n)
           (c #f)
           (if (attacked? i j board)
                (queens n i (+ j 1) board c)
(queens n (+ i 1) 0 (cons (cons i j) board)
                          (lambda (solution?)
                            (if solution?
                                 (c solution?)
                                 (queens n i (+ j 1) board c))))))))
```

Using Continuations

If it is, we send it to our continuation.

Solving the puzzle (in continuation-passing style)

```
(define (queens n i j board c)
  (if (= i n)
       (c board)
       (if (= j n)
           (c #f)
           (if (attacked? i j board)
                (queens n i (+ j 1) board c)
(queens n (+ i 1) 0 (cons (cons i j) board)
                          (lambda (solution?)
                            (if solution?
                                 (c solution?)
                                 (queens n i (+ j 1) board c))))))))
```

Using Continuations

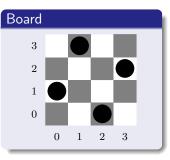
• Otherwise, we continue the search for a non-attacked postion.

Solution

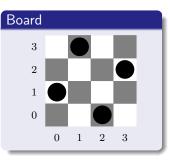
>> (solve-queens 4)

Board

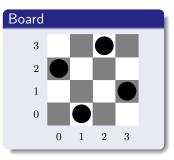
Solution



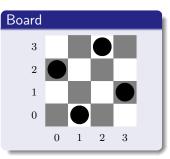
```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
```



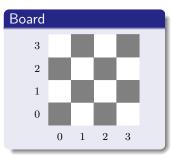
```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((3 . 1) (2 . 3) (1 . 0) (0 . 2))
```



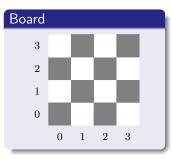
```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
```



```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
#f
```

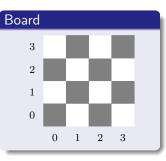


```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
#f
>> (next-solution #f)
```



Solution

```
>> (solve-queens 4)
((3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
#f
>> (next-solution #f)
#f
```



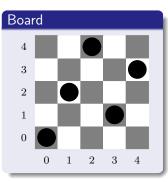
Nondeterministic Programming Languages

- Our program is now nondeterministic.
- It's just like Prolog.

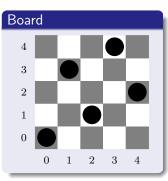
Finding Solutions for 5-Queens

>> (solve-queens 5)

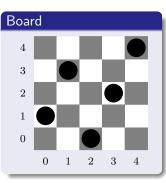
Board



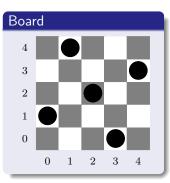
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
```



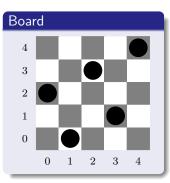
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
```



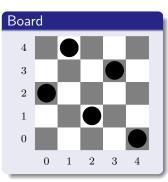
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
```



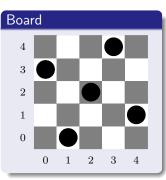
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
```



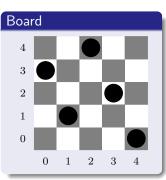
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
```



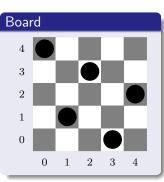
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (next-solution #f)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
```



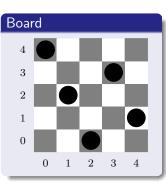
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4.4)(3.2)(2.0)(1.3)(0.1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4.0)(3.3)(2.1)(1.4)(0.2))
>> (next-solution #f)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
>> (next-solution #f)
((4 . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
```



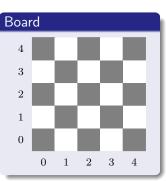
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4.0)(3.3)(2.1)(1.4)(0.2))
>> (next-solution #f)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
  (next-solution #f)
((4.0)(3.2)(2.4)(1.1)(0.3))
>> (next-solution #f)
((4 . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
```



```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
  (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4.0)(3.3)(2.
                     1) (1 . 4) (0 . 2))
>> (next-solution #f)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
  (next-solution #f)
   . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
  (next-solution #f)
((4 . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
>> (next-solution #f)
((4 . 1) (3 . 3) (2 . 0) (1 . 2) (0 . 4))
```



```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (next-solution #f)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (next-solution #f)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (next-solution #f)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
  (next-solution #f)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (next-solution #f)
((4.0)(3.3)(2.1)(1.4)(0.2))
>> (next-solution #f)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
  (next-solution #f)
   . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
  (next-solution #f)
   . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
  (next-solution #f)
((4 . 1) (3 . 3) (2 . 0) (1 . 2) (0 . 4))
>> (next-solution #f)
#f
```



Continuations

- Being forced to write programs in continuation-passing style is a huge disadvantage.
- Can we write an operator that captures the current continuation?

Continuations

- Yes, we can!
- Just convert the evaluator into continuation-passing style and implement the operator as simply "returning" the current continuation.

Continuations

- Being forced to write programs in continuation-passing style is a huge disadvantage.
- Can we write an operator that captures the current continuation?

Continuations

- Yes, we can!
- Just convert the evaluator into continuation-passing style and implement the operator as simply "returning" the current continuation.
- Or something like that.

Evaluator

```
(define (eval exp env)
  (cond ((self-evaluating? exp)
         exp)
        ((name? exp)
         (eval-name exp env))
        ((current-environment? exp)
         env)
        ((lambda? exp)
         (eval-lambda exp env))
        ((if? exp)
         (eval-if exp env))
        ((def? exp)
         (eval-def exp env))
        ((set? exp)
         (eval-set exp env))
        ((begin? exp)
         (eval-begin exp env))
        ((call? exp)
         (eval-call exp env))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
(define (eval exp env cont)
  (cond ((self-evaluating? exp)
         (cont exp))
        ((name? exp)
         (eval-name exp env cont))
        ((current-environment? exp)
         (cont env))
        ((lambda? exp)
         (eval-lambda exp env cont))
        ((if? exp)
         (eval-if exp env cont))
        ((def? exp)
         (eval-def exp env cont))
        ((set? exp)
         (eval-set exp env cont))
        ((begin? exp)
         (eval-begin exp env cont))
        ((call? exp)
         (eval-call exp env cont))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Evaluator

Evaluator

Evaluator

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
```

Evaluator

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
```

Evaluator

```
(define (eval-if exp env)
  (if (true? (eval (if-condition exp) env))
      (eval (if-consequent exp) env)
      (eval (if-alternative exp) env)))
```

Evaluator

```
(define (eval-def exp env)
  (let ((value (eval (def-init exp) env)))
    (define-name! (def-name exp) value env)
   value))
```

```
(define (eval-def exp env cont)
  (eval (def-init exp) env
        (lambda (value)
          (define-name! (def-name exp) value env)
          (cont value))))
```

Evaluator (define (eval-def exp env) (let ((value (eval (def-init exp) env))) (define-name! (def-name exp) value env) value))

```
(define (eval-def exp env cont)
 (eval (def-init exp) env
        (lambda (value)
          (define-name! (def-name exp) value env)
          (cont value))))
```

Evaluator

```
(define (eval-def exp env)
  (let ((value (eval (def-init exp) env)))
    (define-name! (def-name exp) value env)
   value))
```

```
(define (eval-def exp env cont)
  (eval (def-init exp) env
        (lambda (value)
          (define-name! (def-name exp) value env)
          (cont value))))
```

Evaluator

```
(define (eval-set exp env)
  (let ((value (eval (assignment-expression exp) env)))
     (update-name! (assignment-name exp) value env)
    value))
```

Evaluator

```
(define (eval-set exp env)
  (let ((value (eval (assignment-expression exp) env)))
    (update-name! (assignment-name exp) value env)
   value))
```

```
(define (eval-set exp env cont)
 (eval (assignment-expression exp) env
        (lambda (value)
          (update-name! (assignment-name exp) value env)
          (cont value))))
```

Evaluator

```
(define (eval-set exp env)
  (let ((value (eval (assignment-expression exp) env)))
     (update-name! (assignment-name exp) value env)
    value))
```

Evaluator

Evaluator (define (eval-begin exp env) (define (eval-sequence expressions env) (if (null? (cdr expressions)) (eval (car expressions) env) (begin (eval (car expressions) env) (eval-sequence (cdr expressions) env)))) (eval-sequence (begin-expressions exp) env))

```
(define (eval-begin exp env cont)
  (define (eval-sequence expressions env)
    (if (null? (cdr expressions))
        (eval (car expressions) env cont)
        (eval (car expressions) env
              (lambda (r) ;; ignore the result
                (eval-sequence (cdr expressions) env)))))
  (eval-sequence (begin-expressions exp) env))
```


(apply-function func (eval-exprs (call-operands exp) env)))))

Evaluator

(function-environment func))))

(augment-environment (function-parameters func) args

```
Evaluator in CPS
```

(eval (function-body func) extended-environment))))


```
Evaluator in CPS
```


Continuations

- The evaluator is now in continuation-passing style.
- The transformation preserves the semantics of the program.
- But we now have an explicit representation of the pending computations of the evaluator.
- Note that this is a procedure in the evaluating language.
- The pending computations of the evaluator are directly related to the pending computations of the evaluated program.
- It is possible to make this continuation available to the evaluated program as a procedure in the evaluated language.
- Just wrap the continuation of the evaluator into a primitive procedure of the evaluated language.

```
;;(current-continuation)
```

```
:: (current-continuation)
(define (current-continuation? exp)
  (and (pair? exp)
       (eq? (car exp) 'current-continuation)))
(define (eval exp env cont)
  (cond ((self-evaluating? exp)
         (cont exp))
        ((name? exp)
         (eval-name exp env cont))
        ((current-environment? exp)
         (cont env))
        ((lambda? exp)
         (eval-lambda exp env cont))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
:: (current-continuation)
(define (current-continuation? exp)
  (and (pair? exp)
       (eq? (car exp) 'current-continuation)))
(define (eval exp env cont)
  (cond ((self-evaluating? exp)
         (cont exp))
        ((name? exp)
         (eval-name exp env cont))
        ((current-environment? exp)
         (cont env))
        ((current-continuation? exp)
         (cont (make-primitive cont)))
        ((lambda? exp)
         (eval-lambda exp env cont))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

Testing current-continuation

>> (define foo (current-continuation))

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
>> (foo 1)
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
>> (foo 1)
1
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
>> (foo 1)
1
>> foo
```

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
>> (foo 1)
1
>> foo
1
```

Testing current-continuation

```
>> (define foo (current-continuation))
(primitive #<closure 0x975264c0>)
>> foo
(primitive #<closure 0x975264c0>)
>> (foo 1)
1
>> foo
1
```

Explanation

- When the continuation was captured, we were waiting for a value to assign to foo.
- When we called the continuation with argument 1, it means that the value we were waiting to assign to foo was the number 1.
- It means that we now are evaluating (define foo 1).



Testing current-continuation

>> (current-continuation)

```
>> (current-continuation)
(primitive #<closure 0x970f48c0>)
```

```
>> (current-continuation)
(primitive #<closure 0x970f48c0>)
>> ((current-continuation) 1)
```

```
>> (current-continuation)
(primitive #<closure 0x970f48c0>)
>> ((current-continuation) 1)
ERROR ;; Why?
```

Testing current-continuation

```
>> (current-continuation)
(primitive #<closure 0x970f48c0>)
>> ((current-continuation) 1)
ERROR ;; Why?
```

Explanation

- When the continuation was captured, we were waiting for a function to call with argument 1.
- When we called the continuation with argument 1, the current continuation was replaced with the captured continuation as if the function we were waiting to call with argument 1 was the number 1.
- It means that we now are trying to evaluate (1 1).

```
>> (let ((c (current-continuation)))
    (display "Hi!")
    (+ 1 (c (lambda (x) 1))))
```

```
>> (let ((c (current-continuation)))
     (display "Hi!")
     (+ 1 (c (lambda (x) 1))))
Hi!Hi!2
```

Testing current-continuation

```
>> (let ((c (current-continuation)))
     (display "Hi!")
     (+ 1 (c (lambda (x) 1))))
Hi!Hi!2
```

Explanation

- When the continuation was captured, we were waiting for a value to assign to c and then to display "Hi!" and then to add 1 to the invocation of c with a function that returns 1.
- We then did all the computations we were waiting to do.
- At the call to c, we went back to the assignment to c that now becomes bound to the function (lambda (x) 1) and then we did all the computations we were waiting to do until we called c that now is a function that simply returns 1.

```
>> (let ((longjmp (current-continuation)))
     (if longjmp
         (begin
           (display "Before the longimp")
           (newline)
           (longjmp #f)
           (display "Did I survive the longjmp?"))
         (begin
           (display "After the longjmp"))))
Before the longjmp
```

Testing current-continuation

Explanation

Go see the setjmp/longjmp available in C.

More Useful Continuations

- In our current implementation, the use of a continuation is usually part of the continuation itself.
- This means that (in general) the use of the continuation will re-evaluate its use when, presumably, that second use was not intended.
- It's preferable to separate the part of the program that captures
 the continuation from the part of the program that uses the
 continuation.
- That's the purpose of the function call/cc (also known as call-with-current-continuation).

```
;; (call/cc (lambda (c) ...))
```

```
;; (call/cc (lambda (c) ...))
(define (call/cc? exp)
  (and (pair? exp)
       (eq? (car exp) 'call/cc)))
(define (call/cc-expression exp)
  (cadr exp))
(define (eval exp env cont)
  (cond ...
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
:: (call/cc (lambda (c) ...))
(define (call/cc? exp)
  (and (pair? exp)
       (eq? (car exp) 'call/cc)))
(define (call/cc-expression exp)
  (cadr exp))
(define (eval exp env cont)
  (cond ...
        ((call/cc? exp)
         (eval-call/cc exp env cont))
        (else
         (error "Unknown expression type -- EVAL" exp))))
```

```
;; (call/cc (lambda (c) ...))
(define (call/cc? exp)
  (and (pair? exp)
       (eq? (car exp) 'call/cc)))
(define (call/cc-expression exp)
  (cadr exp))
(define (eval exp env cont)
  (cond ...
        ((call/cc? exp)
         (eval-call/cc exp env cont))
        (else
         (error "Unknown expression type -- EVAL" exp))))
(define (eval-call/cc exp env cont)
  (eval `(,(call/cc-expression exp)
          ',(make-primitive cont))
        env
        cont))
```

Implementing call/cc

 Note that call/cc is implemented as special syntax in our evaluator but is a "normal" function in the Scheme language.

Implementing current-continuation

```
(define (current-continuation)
  (call/cc (lambda (c) c)))
```

call/cc vs current-continuation

- call/cc is sufficient for implementing current-continuation.
- call/cc makes it easier to implement other control structures.

```
>> (define foo (call/cc (lambda (c) c)))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
1
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
1
>> foo
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
1
>> foo
1
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
1
>> foo
(primitive #<closure 0xbd721560>)
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
>> (+ 2 (* 3 (foo 1)))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
>> (+ 2 (* 3 (foo 1)))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
1
>> (+ 2 (* 3 (foo 1)))
>> (foo (+ 2 (foo (* 2 (foo 1)))))
```

```
>> (define foo (call/cc (lambda (c) c)))
(primitive #<closure 0xbd961420>)
>> foo
(primitive #<closure 0xbd961420>)
>> (foo 1)
>> foo
>> (call/cc (lambda (c) (set! foo c)))
(primitive #<closure 0xbd721560>)
>> (foo 1)
>> foo
(primitive #<closure 0xc85d7c80>)
>> (+ 2 (foo 1))
1
>> (+ 2 (* 3 (foo 1)))
>> (foo (+ 2 (foo (* 2 (foo 1)))))
```

Aborting a Computation

```
>> (define abort-top-level #f)
#f
>> (call/cc (lambda (c) (set! abort-top-level c)))
(primitive #<closure 0xbca29d20>)
```

Aborting a Computation

```
>> (define abort-top-level #f)
#f
>> (call/cc (lambda (c) (set! abort-top-level c)))
(primitive #<closure 0xbca29d20>)
```

Aborting a Computation

```
>> (define abort-top-level #f)
#f
>> (call/cc (lambda (c) (set! abort-top-level c)))
(primitive #<closure 0xbca29d20>)
```

```
(define (mystery a b c d e)
  (+ a (* b (safe-/ (- c d) e))))

(define (safe-/ x y)
  (if (= y 0)
        (error "Can't divide " (list x y))
        (/ x y)))

(define (error msg args)
  (display msg)
  (display args)
```

Aborting a Computation

```
>> (define abort-top-level #f)
#f
>> (call/cc (lambda (c) (set! abort-top-level c)))
(primitive #<closure 0xbca29d20>)
```

```
(define (mystery a b c d e)
  (+ a (* b (safe-/ (- c d) e))))
(define (safe-/ x y)
  (if (= y 0)
      (error "Can't divide " (list x y))
      (/ x y))
(define (error msg args)
  (display msg)
  (display args)
  (abort-top-level #f))
```

```
(define (palindrome n f)
  (if (= n 0)
        (f n)
        (begin
              (display n)
              (palindrome (- n 1) f)
              (display n))))
>> (palindrome 9 display)
```

```
(define (palindrome n f)
  (if (= n 0)
        (f n)
        (begin
              (display n)
              (palindrome (- n 1) f)
        (display n))))
>> (palindrome 9 display)
9876543210123456789...
```

```
(define (palindrome n f)
  (if (= n 0)
       (f n)
       (begin
         (display n)
         (palindrome (- n 1) f) (display n)))
>> (palindrome 9 display)
9876543210123456789...
>> (call/cc (lambda (return)
               (palindrome 9 return)))
9876543210
;; That was the Scheme way.
;; What about the Common Lisp way?
```

```
(define (palindrome n f)
  (if (= n 0)
      (f n)
      (begin
        (display n)
        (palindrome (- n 1) f) (display n)))
>> (palindrome 9 display)
9876543210123456789...
>> (call/cc (lambda (return)
               (palindrome 9 return)))
9876543210
;; That was the Scheme way.
;; What about the Common Lisp way?
>> (block end
     (palindrome 9 (lambda (r)
                       (return-from end r))))
9876543210
```

Common Lisp's block/return-from

Common Lisp's block/return-from

- The operators block and return-from provide a structured, lexical, non-local exit facility.
- At any point lexically contained within forms, return-from can be used with the given name to return control and values from the block form.

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

Output

>> (foo 0)

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

```
>> (foo 0)
Handling error-a 1
```

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

```
>> (foo 0)
Handling error-a 1
>> (foo 1)
```

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

```
>> (foo 0)
Handling error-a 1
>> (foo 1)
Handling error-b 3
```

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

```
>> (foo 0)
Handling error-a 1
>> (foo 1)
Handling error-b 3
>> (foo 2)
```

Java's try-catch/throw (define (foo n) (+ 1 (try (bar n) catch error-a (display "Handling error-a ") 0))) (define (bar n) (+ 1 (try (baz n) catch error-b (display "Handling error-b ") 1))) (define (baz n) (if (= n 0))(throw 'error-a) (if (= n 1)(throw 'error-b) n)))

```
>> (foo 0)
Handling error-a 1
>> (foo 1)
Handling error-b 3
>> (foo 2)
```

```
Java's try-catch/throw
(define error-handlers (list))
```

```
(define error-handlers (list))
(define (push-error-handler tag eh)
  (set! error-handlers (cons (cons tag eh) error-handlers)))
```

```
(define error-handlers (list))
(define (push-error-handler tag eh)
  (set! error-handlers (cons (cons tag eh) error-handlers)))
(define (pop-error-handler)
  (set! error-handlers (cdr error-handlers)))
```

```
(define error-handlers (list))
(define (push-error-handler tag eh)
  (set! error-handlers (cons (cons tag eh) error-handlers)))
(define (pop-error-handler)
  (set! error-handlers (cdr error-handlers)))
(define (find-error-handler tag)
  (if (null? error-handlers)
      abort-top-level
      (let ((tag-eh (car error-handlers)))
        (set! error-handlers (cdr error-handlers))
        (if (eq? tag (car tag-eh))
            (cdr tag-eh)
            (find-error-handler tag)))))
```

```
(mdef try (expr catch tag . handler)
  `(call-with-error-handler
    (lambda () ,expr)
   ',tag
    (lambda () ,@handler)))
```

```
(mdef try (expr catch tag . handler)
  `(call-with-error-handler
    (lambda () ,expr)
    ',tag
    (lambda () ,@handler)))
(define (call-with-error-handler f tag eh)
  (call/cc (lambda (c)
             (push-error-handler
               tag
               (lambda () (c (eh))))
             (let ((result (f)))
               (pop-error-handler)
               result))))
```

```
Java's try-catch/throw
(mdef try (expr catch tag . handler)
  `(call-with-error-handler
    (lambda () ,expr)
    ',tag
    (lambda () ,@handler)))
(define (call-with-error-handler f tag eh)
  (call/cc (lambda (c)
             (push-error-handler
               tag
               (lambda () (c (eh))))
             (let ((result (f)))
               (pop-error-handler)
               result))))
(define (throw tag)
  ((find-error-handler tag)))
```

Nondeterministic Programs

- Certain expressions in a nondeterministic language can have a different value in each different world.
- Each world encodes one of many possible sequences of choices, each choice corresponding to one of the possible values of a nondeterministic expression.
- In practice:
 - The evaluator computes just one of the possible values of a nondeterministic expression.
 - But remembers that there were more choices available.
 - If the choice taken does not satisfy subsequent requirements, a different choice is attempted.
 - This process is repeated until the evaluation succeeds or the evaluation runs out of choices.



A Pythagorean Triple

```
(define (a-pythagorean-triple n)
  (let ((i (an-integer-between 1 n)))
     (let ((j (an-integer-between i n)))
        (let ((k (an-integer-between j n)))
           (require (= (+ (* i i) (* j j)) (* k k)))
           (list i j k)))))
```

Nondeterministic Programs

- The function a-pythagorean-triple uses the function (an-integer-between a b) that nondeterministically returns one integer between a b.
- The function require accepts a boolean expression and requires that the computation must satisfy the expression. Otherwise, the computation fails.

An Integer Between a and b

The amb (ambiguous) operator

- Invented by John McCarthy in 1961.
- Nondeterministically returns the value of one of the expressions.

The fail operator

- Cause the current computation to fail.
- It means that the choices that were taken in the current computation were not good.



amb and fail

>> (amb 1 2)

```
amb and fail
```

```
>> (amb 1 2)
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
```

```
amb and fail
>> (amb 1 2)
```

```
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
```

```
amb and fail
```

```
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
```

```
amb and fail
>> (amb 1 2)
1
>> (fail)
```

```
// (fail)
no-more-choices
>> (+ 1 (amb 10 20))
```

>> (+ 1 (amb 10 20))
11

>> (fail)

21 >> (cons (amb 1 2) (amb 3 4))

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1 . 3)
```

>> (fail) 21

(1 . 3) >> (fail)

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
```

>> (cons (amb 1 2) (amb 3 4))

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
```

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
>> (fail)
```

>> (fail) (1 . 4) >> (fail) (2 . 3)

amb and fail >> (amb 1 2) 1 >> (fail) 2 >> (fail) no-more-choices >> (+ 1 (amb 10 20)) 11 >> (fail) 21 >> (cons (amb 1 2) (amb 3 4)) (1 . 3)

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
>> (fail)
(2.3)
>> (fail)
```

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
>> (fail)
(2.3)
>> (fail)
(2.4)
```

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
>> (fail)
(2.3)
>> (fail)
(2.4)
>> (fail)
```

```
amb and fail
>> (amb 1 2)
1
>> (fail)
2
>> (fail)
no-more-choices
>> (+ 1 (amb 10 20))
11
>> (fail)
21
>> (cons (amb 1 2) (amb 3 4))
(1.3)
>> (fail)
(1.4)
>> (fail)
(2.3)
>> (fail)
(2.4)
>> (fail)
no-more-choices
```

Implementation

- The ambiguous operator amb evaluates just one expression.
- Each of the remaining expressions is saved so that, if necessary, the entire computation can be repeated using its evaluation.

Saving Choices

```
(define choices (list))
```

Implementation

- The ambiguous operator amb evaluates just one expression.
- Each of the remaining expressions is saved so that, if necessary, the entire computation can be repeated using its evaluation.

Saving Choices

```
(define choices (list))
(define (add-choices! cs)
  (set! choices (append cs choices)))
```

Implementation

- The ambiguous operator amb evaluates just one expression.
- Each of the remaining expressions is saved so that, if necessary, the entire computation can be repeated using its evaluation.

Saving Choices

```
(define choices (list))
(define (add-choices! cs)
  (set! choices (append cs choices)))
(define (pop-choice!)
  (let ((choice (car choices)))
      (set! choices (cdr choices))
      choice))
```

amb and fail

amb and fail

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
14
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))

14
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
14
>> (fail)
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))

14
>> ((pop-choice!))
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
14
>> (fail)
16
```

amb and fail

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))

14
>> ((pop-choice!))
16
```

```
amb and fail
>> (* 2 (+ 3 (amb 4 5 6)))
14
>> (fail)
16
>> (fail)
18
```

```
amb and fail
```

```
>> (* 2 (+ 3 (call/cc (lambda (c) (add-choices! (list (lambda () (c 5)) (lambda () (c 6))))

4))))

14
>> ((pop-choice!))
16
>> ((pop-choice!))
18
```

From

```
(amb e_0 e_1 ... e_n)
```

То

From

```
(amb e_0 e_1 \ldots e_n)
```

То

The amb Macro

Failing

```
(define (fail)
  (if (null? choices)
        (abort-top-level 'no-more-choices)
        ((pop-choice!))))
```

Failing

- When a world fails it means that the computation took some wrong choice and cannot continue.
- In order to proceed, take another choice and continue the computation from there.
- What should be the meaning of (amb e)?
- What should be the meaning of (amb)?

A Pythagorean Triple

```
(define (a-pythagorean-triple n)
  (let ((i (an-integer-between 1 n)))
    (let ((j (an-integer-between i n)))
        (let ((k (an-integer-between j n)))
            (require (= (+ (* i i) (* j j)) (* k k)))
            (list i j k)))))
```

A Pythagorean Triple

The require Function

- The function require wants its argument to be true.
- If it is not true, then the computation fails.

A Pythagorean Triple

```
(define (a-pythagorean-triple n)
  (let ((i (an-integer-between 1 n)))
    (let ((j (an-integer-between i n)))
      (let ((k (an-integer-between j n)))
            (require (= (+ (* i i) (* j j)) (* k k)))
            (list i j k)))))
```

The require Function

- The function require wants its argument to be true.
- If it is not true, then the computation fails.

The require Function

```
(define (require e)
  (or e (fail)))
```

Example

>> (a-pythagorean-triple 25)

Example

```
>> (a-pythagorean-triple 25) (3 4 5)
```

Example

```
>> (a-pythagorean-triple 25)
(3 4 5)
>> (fail)
(5 12 13)
```

Example

```
>> (a-pythagorean-triple 25)
(3 4 5)
>> (fail)
(5 12 13)
>> (fail)
(6 8 10)
```

```
Example
```

```
>> (a-pythagorean-triple 25)
(3 \ 4 \ 5)
>> (fail)
(5 12 13)
>> (fail)
(6 8 10)
>> (fail)
(72425)
>> (fail)
(8 15 17)
>> (fail)
(9 12 15)
>> (fail)
(12 16 20)
>> (fail)
(15 20 25)
>> (fail)
no-more-choices
```

Solving the puzzle: Using backtracking

Solving the puzzle: Using nondeterminism

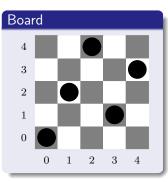
```
(define (queens n i j board)
  (if (= i n)
    board
    (let ((j (an-integer-between 0 (- n 1))))
        (require (not (attacked? i j board)))
        (queens n (+ i 1) 0 (cons (cons i j) board)))))

(define (solve-queens n)
  (queens n 0 0 (list)))
```

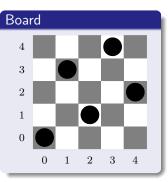
Finding Solutions for 5-Queens

>> (solve-queens 5)

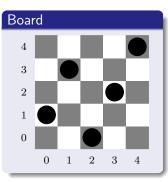
Board



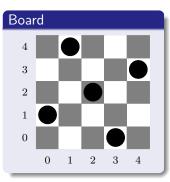
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
```



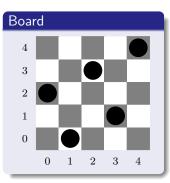
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
```



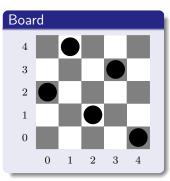
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
```



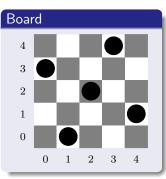
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
```



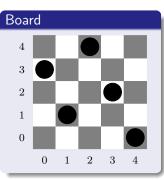
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
```



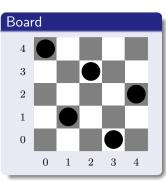
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3)
```



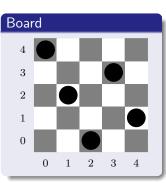
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4.3)(3.0)(2.2)(1.4)(0.1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
>> (fail)
((4.0)(3.2)(2.4)(1.1)(0.3))
```



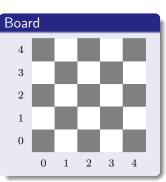
```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4.3)(3.0)(2.2)(1.4)(0.1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
>> (fail)
((4 . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
>> (fail)
((4 . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
```



```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
>> (fail)
((4 . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
>> (fail)
((4 . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
>> (fail)
((4 . 1) (3 . 3) (2 . 0) (1 . 2) (0 . 4))
```



```
>> (solve-queens 5)
((4 . 3) (3 . 1) (2 . 4) (1 . 2) (0 . 0))
>> (fail)
((4 . 2) (3 . 4) (2 . 1) (1 . 3) (0 . 0))
>> (fail)
((4 . 4) (3 . 2) (2 . 0) (1 . 3) (0 . 1))
>> (fail)
((4 . 3) (3 . 0) (2 . 2) (1 . 4) (0 . 1))
>> (fail)
((4 . 4) (3 . 1) (2 . 3) (1 . 0) (0 . 2))
>> (fail)
((4 . 0) (3 . 3) (2 . 1) (1 . 4) (0 . 2))
>> (fail)
((4 . 1) (3 . 4) (2 . 2) (1 . 0) (0 . 3))
>> (fail)
((4 . 0) (3 . 2) (2 . 4) (1 . 1) (0 . 3))
>> (fail)
((4 . 2) (3 . 0) (2 . 3) (1 . 1) (0 . 4))
>> (fail)
((4 . 1) (3 . 3) (2 . 0) (1 . 2) (0 . 4))
>> (fail)
no-more-choices
```



Puzzle

- Adrian, Bob, Charles, David, and Edward live on different floors of an apartment house that contains only five floors.
- Adrian does not live on the top floor.
- Bob does not live on the bottom floor.
- Charles does not live on either the top or the bottom floor.
- David lives on a higher floor than does Bob.
- Edward does not live on a floor adjacent to Charles's.
- Charles does not live on a floor adjacent to Bob's.
- Where does everyone live?

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4 5))
        (b (amb 1 2 3 4 5))
        (c (amb 1 2 3 4 5))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
        (e (amb 1 2 3 4 5)))
        (require (distinct? (list a b c d e)))
```

Puzzle

Adrian, Bob, Charles, David, and Edward live on different floors
of an apartment house that contains only five floors.

Puzzle

Adrian does not live on the top floor.

Puzzle

Bob does not live on the bottom floor.

The apartment-puzzle Function

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4))
        (b (amb 2 3 4 5))
        (c (amb 2 3 4))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
    (require (distinct? (list a b c d e)))
```

Puzzle

Charles does not live on either the top or the bottom floor.

Puzzle

• David lives on a higher floor than does Bob.

Puzzle

• Edward does not live on a floor adjacent to Charles's.

Puzzle

• Charles does not live on a floor adjacent to Bob's.

Puzzle

• Where does everyone live?

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4))
        (b (amb 2 3 4 5))
        (c (amb 2 3 4))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
    (require (distinct? (list a b c d e)))
    (require (> d b))
    (require (not (= (abs (- e c)) 1)))
    (require (not (= (abs (- c b)) 1)))
    (list (list 'adrian a) (list 'bob b) (list 'charles c)
          (list 'david d) (list 'emil e))))
>> (apartment-puzzle)
```

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4))
        (b (amb 2 3 4 5))
        (c (amb 2 3 4))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
    (require (distinct? (list a b c d e)))
    (require (> d b))
    (require (not (= (abs (- e c)) 1)))
    (require (not (= (abs (- c b)) 1)))
    (list (list 'adrian a) (list 'bob b) (list 'charles c)
          (list 'david d) (list 'emil e))))
>> (apartment-puzzle)
((adrian 3) (bob 2) (charles 4) (david 5) (edward 1))
```

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4))
        (b (amb 2 3 4 5))
        (c (amb 2 3 4))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
    (require (distinct? (list a b c d e)))
    (require (> d b))
    (require (not (= (abs (- e c)) 1)))
    (require (not (= (abs (- c b)) 1)))
    (list (list 'adrian a) (list 'bob b) (list 'charles c)
          (list 'david d) (list 'emil e))))
>> (apartment-puzzle)
((adrian 3) (bob 2) (charles 4) (david 5) (edward 1))
>> (fail)
```

```
(define (apartment-puzzle)
  (let ((a (amb 1 2 3 4))
        (b (amb 2 3 4 5))
        (c (amb 2 3 4))
        (d (amb 1 2 3 4 5))
        (e (amb 1 2 3 4 5)))
    (require (distinct? (list a b c d e)))
    (require (> d b))
    (require (not (= (abs (- e c)) 1)))
    (require (not (= (abs (- c b)) 1)))
    (list (list 'adrian a) (list 'bob b) (list 'charles c)
          (list 'david d) (list 'emil e))))
>> (apartment-puzzle)
((adrian 3) (bob 2) (charles 4) (david 5) (edward 1))
>> (fail)
no-more-choices
```

Conclusions

Continuations are the basis for the implementation of specialized control operators:

- Lexical transfer of control (Lisp's block&return-from, Java's break and continue, etc)
- Dynamical transfer of control (Lisp's catch&throw)
- Exceptions (Lisp's signal and error, Java's throw)
- Exception Handling (Lisp's handler-bind and handler-case and unwind-protect, Java's try-catch-finally)
- Generators (Python, Ruby and C#'s yield)
- Nondeterminism (McCarthy's amb, Prolog, Icon)
- Coroutines (cooperative multithreading)
- ..





Harold Abelson and Gerald Jay Sussman.

Structure and Interpretation of Computer Programs.

The MIT Press, Cambridge, Massachusetts, 2nd edition, July 1996.