



UNIVERSITÀ
DELLA CALABRIA

The background image shows a long, empty corridor with red brick walls and metal grating floors. The corridor is brightly lit by overhead lights, creating a strong perspective effect. A dark semi-transparent rectangle is overlaid in the center of the image, containing the title and author's name.

Dynamic Programming with tree decomposition

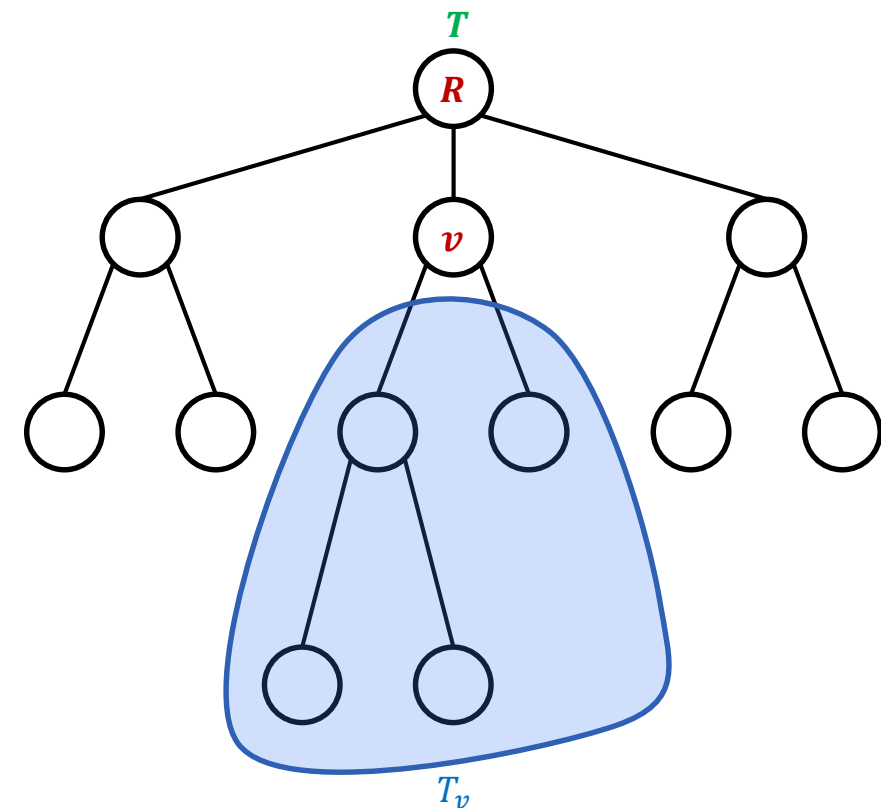
Sebastiano A. Piccolo



Recap



Maximum Independent Set on Trees



For each vertex v we compute:

► $M^+[v] = |MIS(T_v) \cup \{v\}|$

► $M^-[v] = |MIS(T_v)|$

For a vertex v with children w_1, \dots, w_d

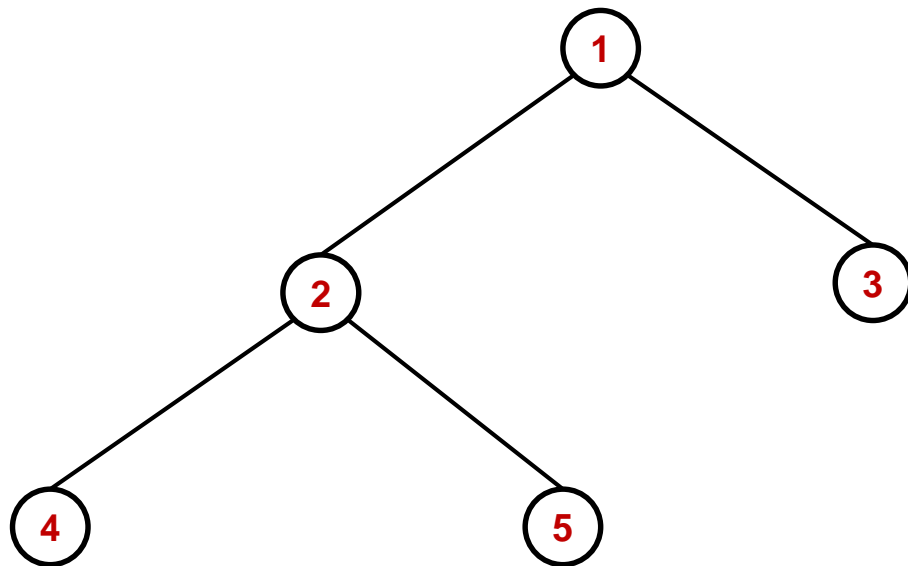
► $M^+[v] = 1 + \sum_{i=1}^d M^-[w_i]$

► $M^-[v] = \sum_{i=1}^d \max\{M^+[w_i], M^-[w_i]\}$

$$MIS(T) = \max\{M^+[R], M^-[R]\}$$

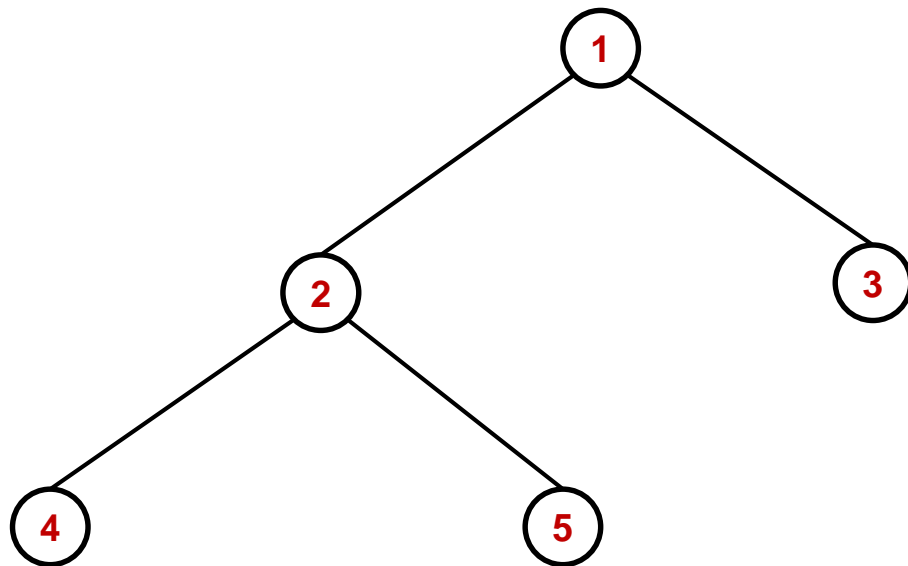


Maximum Independent Set on Trees





Maximum Independent Set on Trees

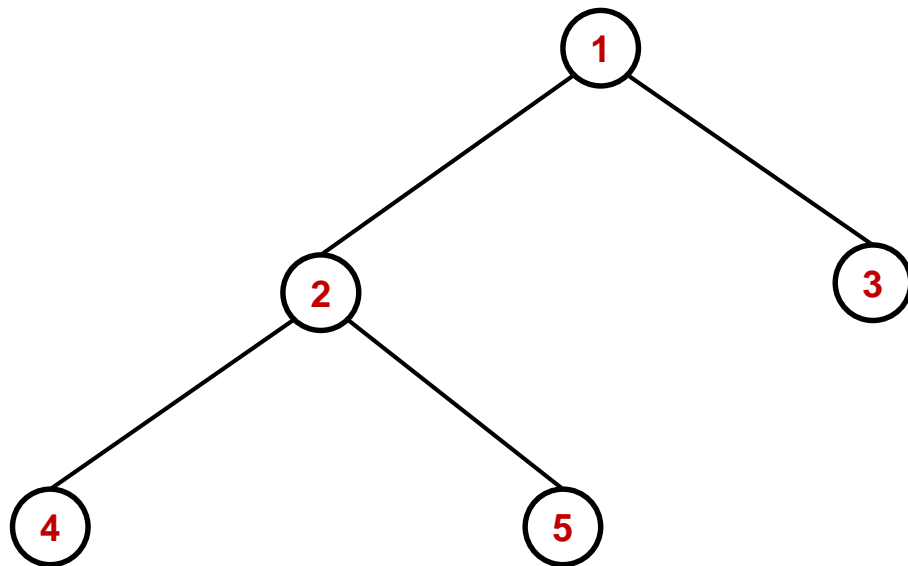


$$M^+[4] = 1$$

$$M^-[4] = 0$$



Maximum Independent Set on Trees

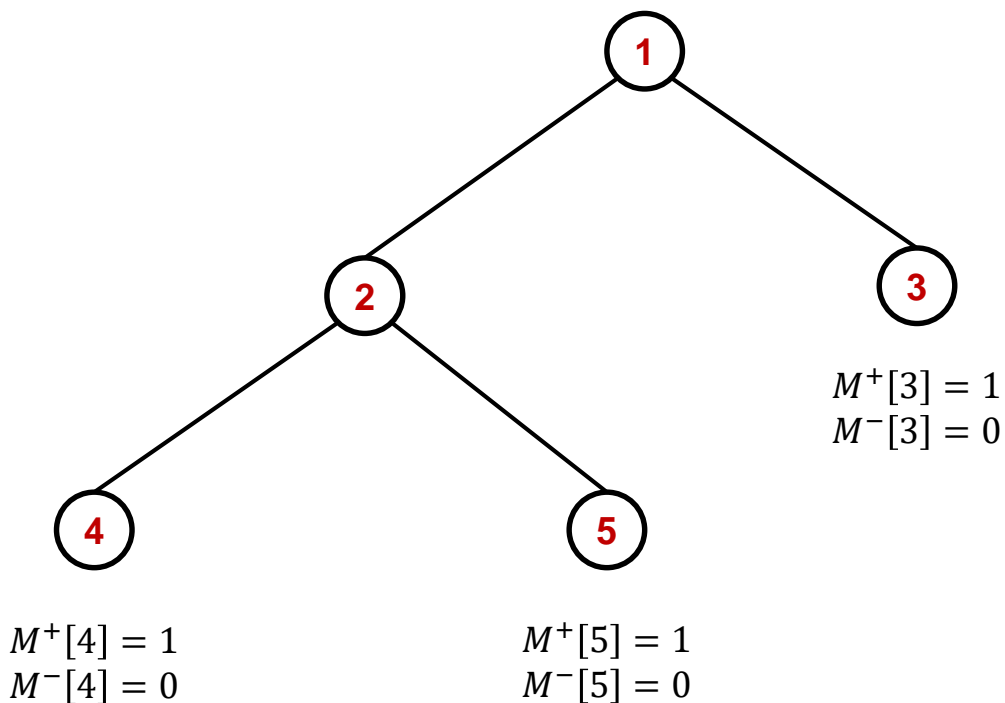


$$\begin{aligned} M^+[4] &= 1 \\ M^-[4] &= 0 \end{aligned}$$

$$\begin{aligned} M^+[5] &= 1 \\ M^-[5] &= 0 \end{aligned}$$

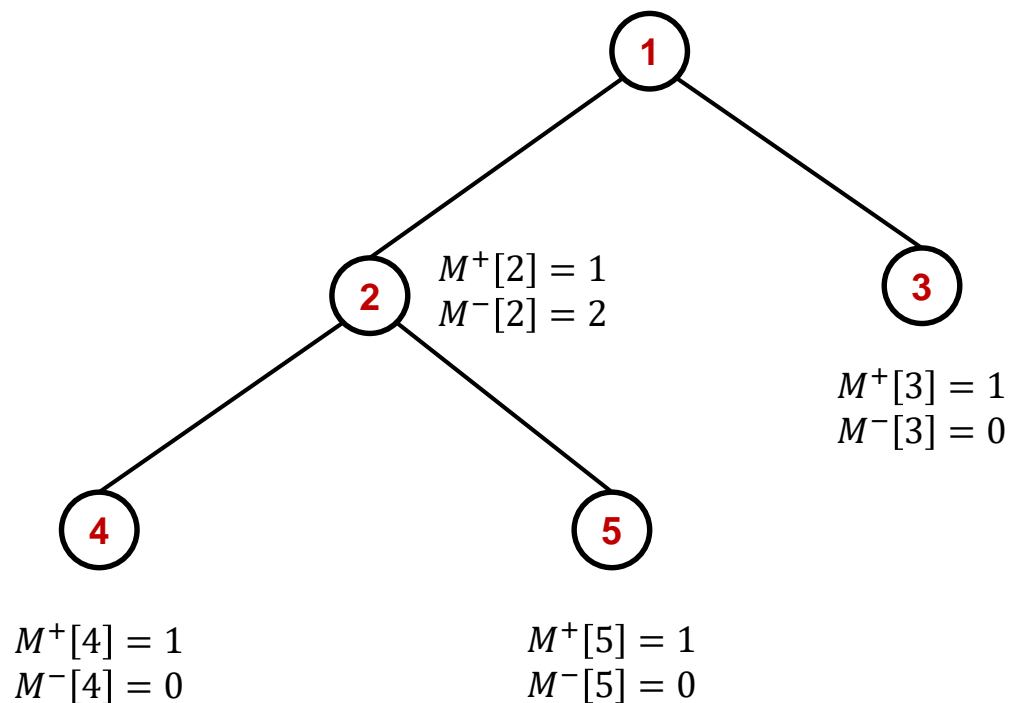


Maximum Independent Set on Trees



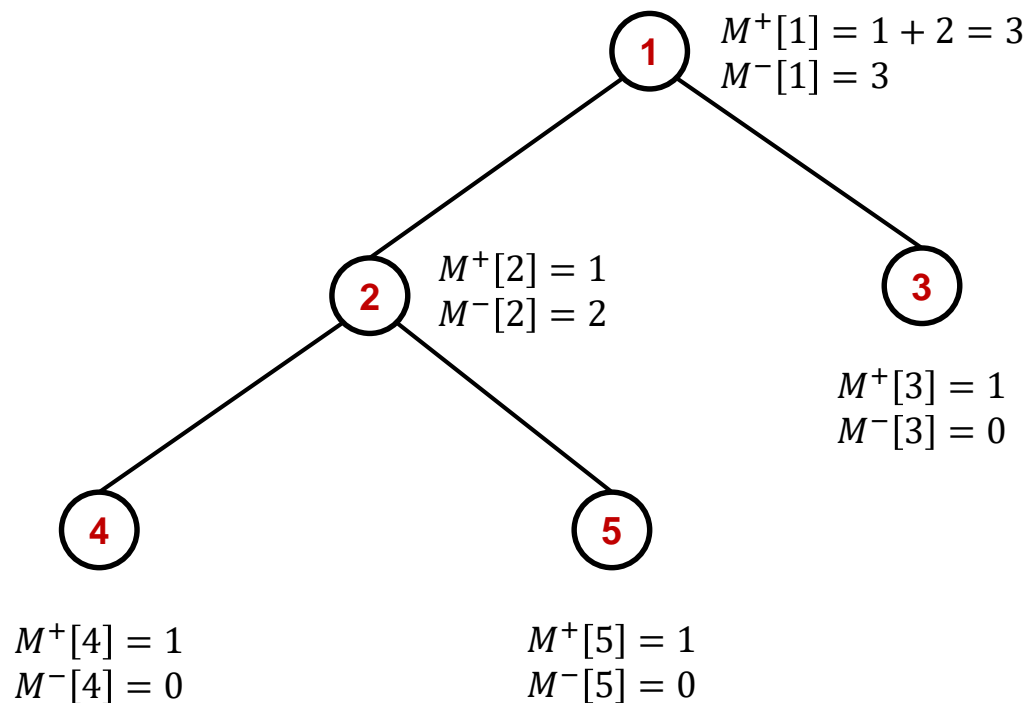


Maximum Independent Set on Trees

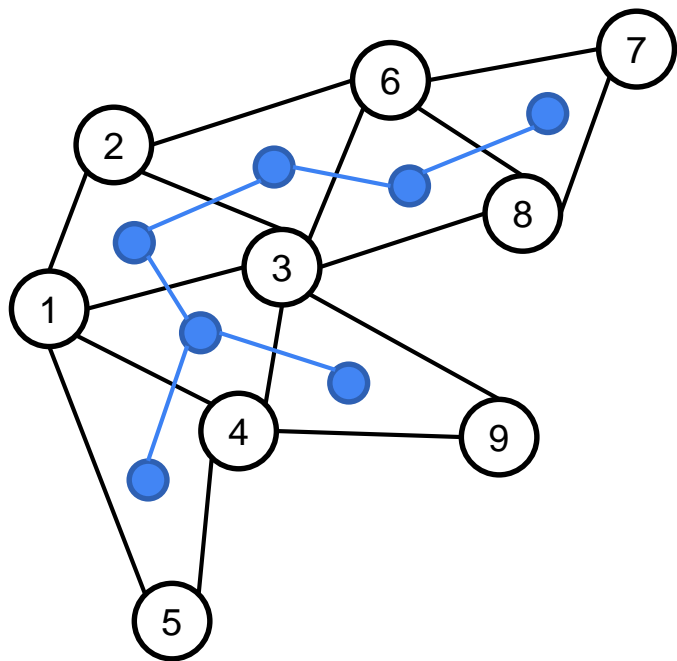




Maximum Independent Set on Trees

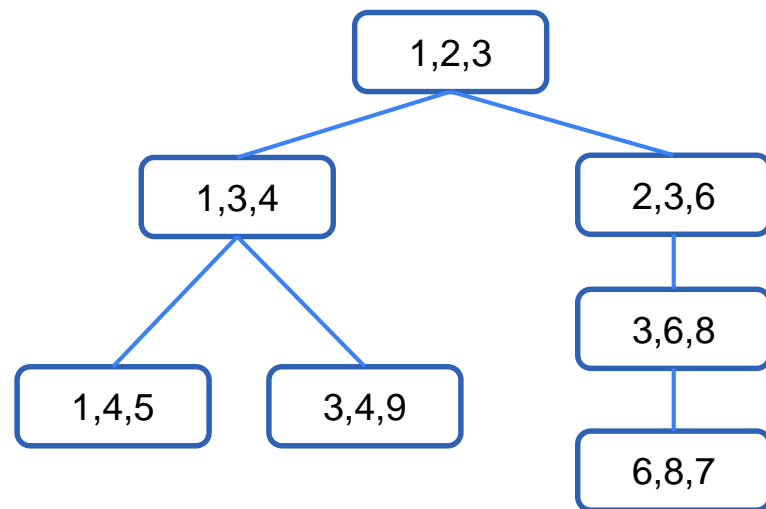
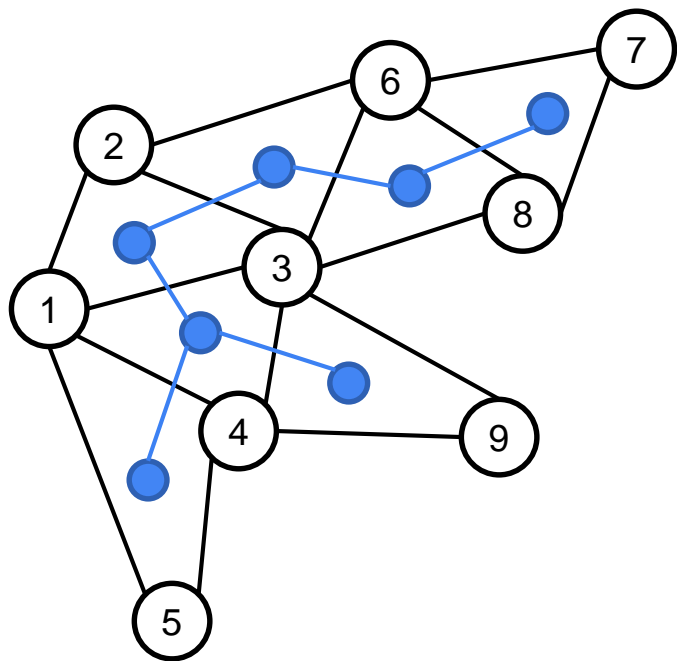



Tree decomposition: intuition





Tree decomposition

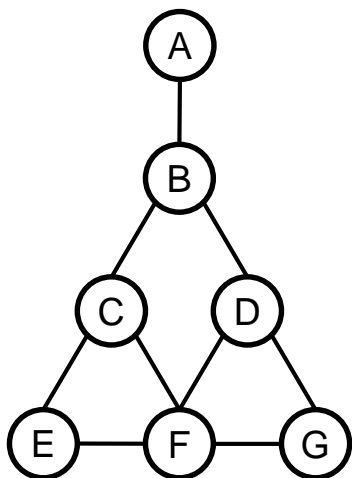




Dynamic Programming with Tree Decompositions

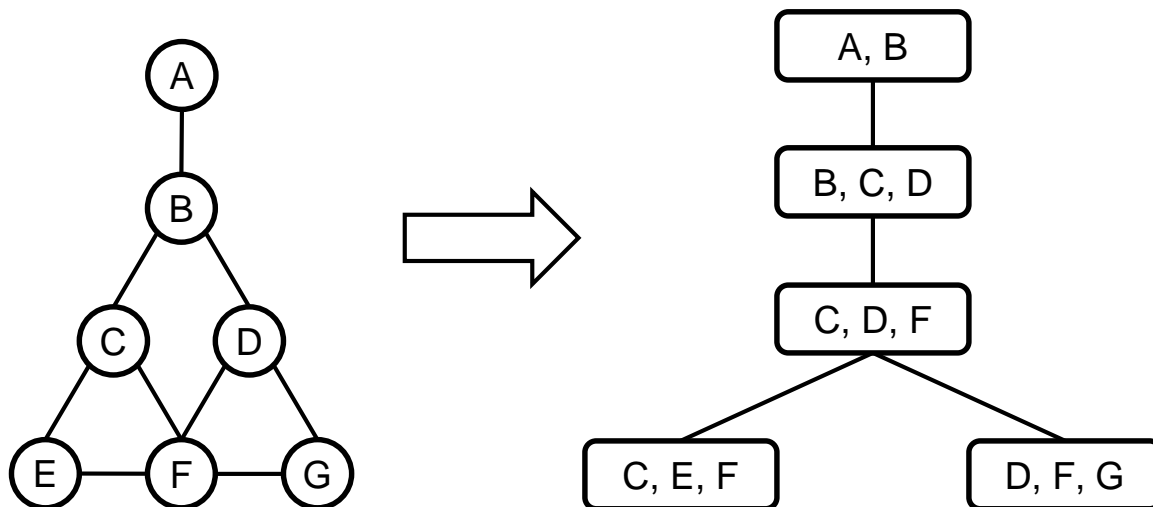


MIS on Graphs of bounded treewidth



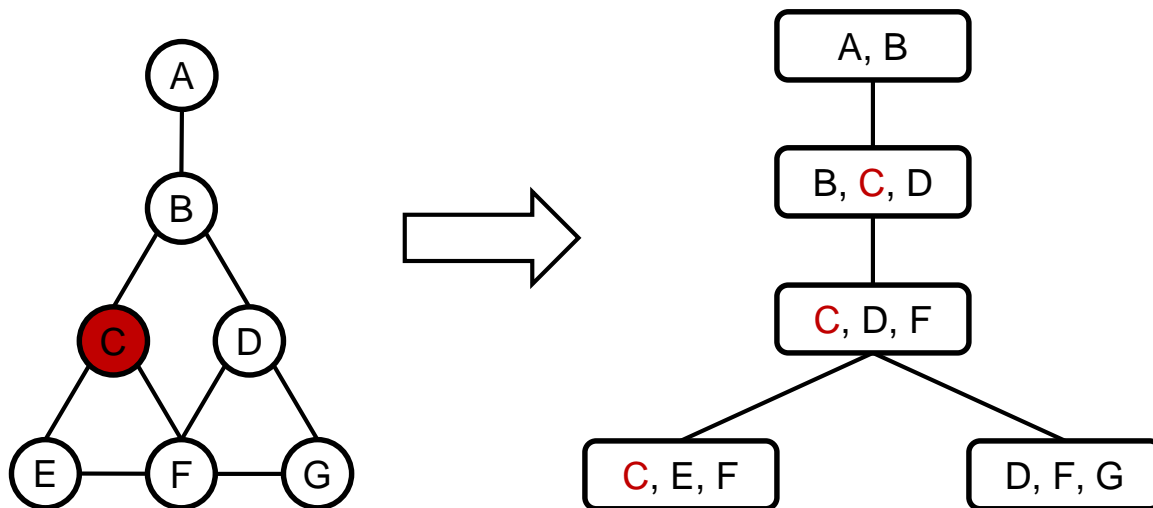


MIS on Graphs of bounded treewidth



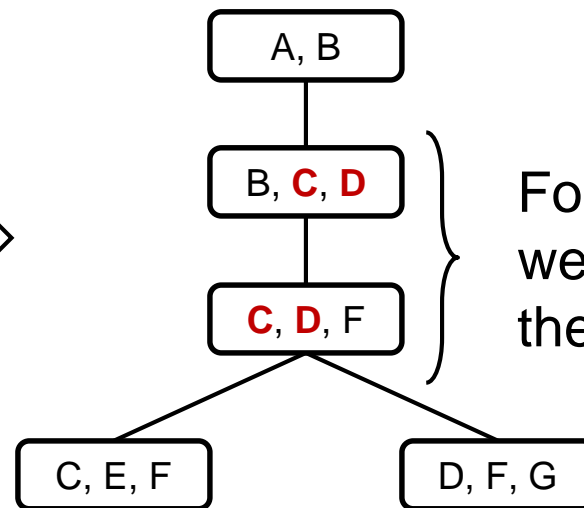
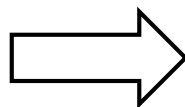
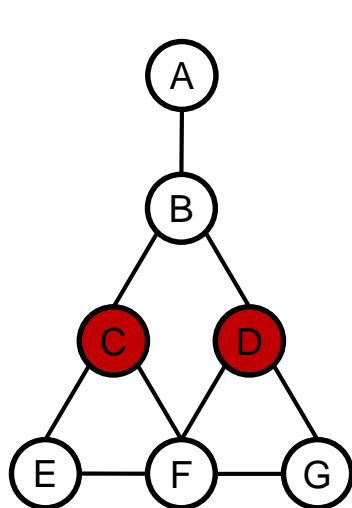


MIS on Graphs of bounded treewidth



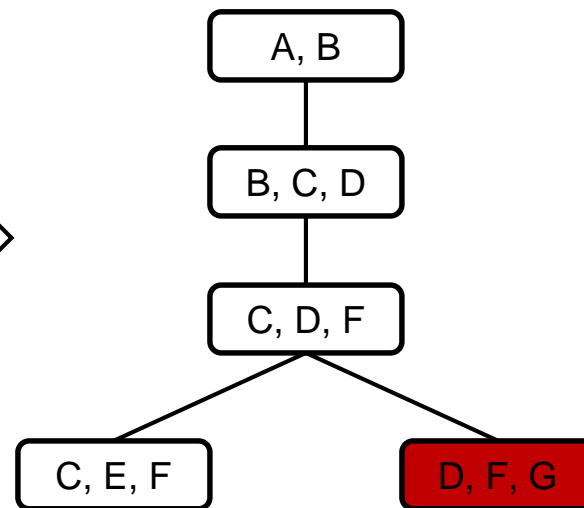
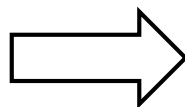
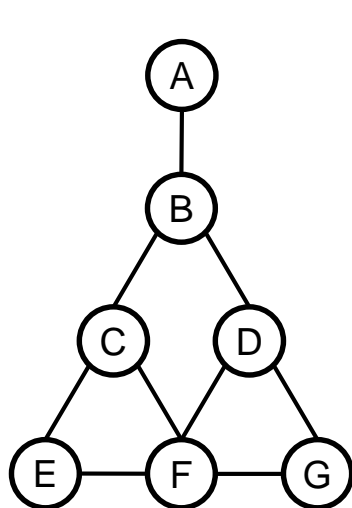


MIS on Graphs of bounded treewidth



For all parent-child bags,
we need to take care of
their intersection

MIS on Graphs of bounded treewidth

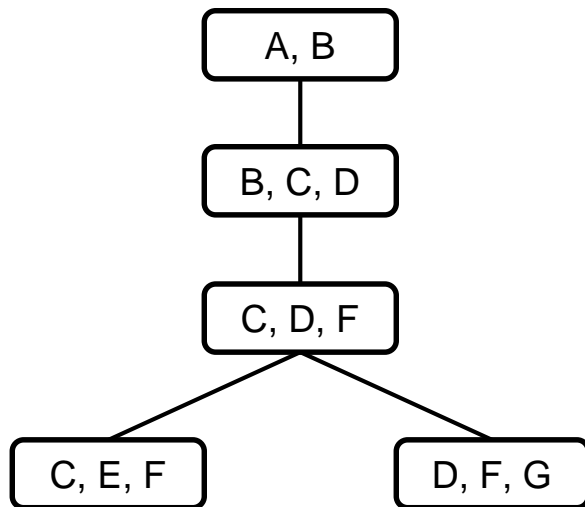


Solve the subproblem for all subsets:

$\{D\}, \{F\}, \{G\}, \{D, F\},$
 $\{D, G\}, \{F, G\}, \{D, F, G\}$



MIS on Graphs of bounded treewidth



R is the root

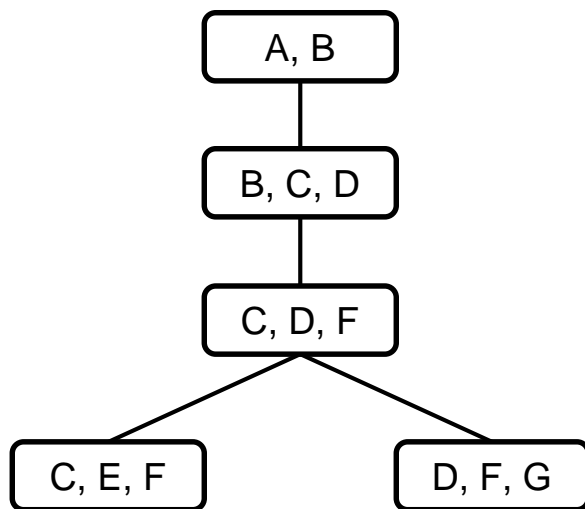
X_i is a bag

X_j is a child of X_i

$X_{k,i}$ is the parent of X_i



MIS on Graphs of bounded treewidth



R is the root

X_i is a bag

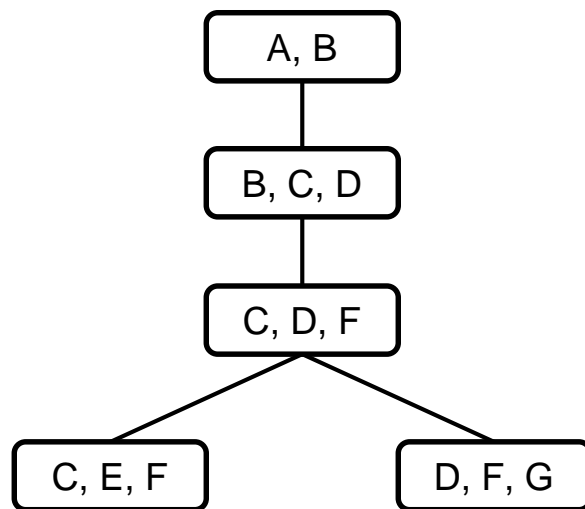
X_j is a child of X_i

$X_{k,i}$ is the parent of X_i

$$A[S, i] = MIS(S) \forall S \in X_i + \sum_j MIS(X_j)$$



MIS on Graphs of bounded treewidth



R is the root

X_i is a bag

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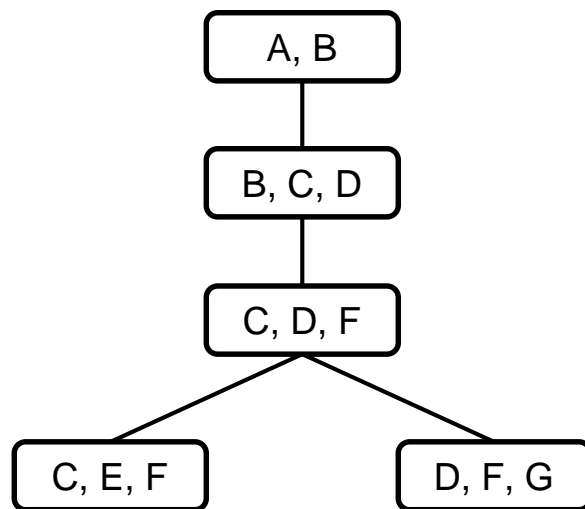
$X_{k,i}$ is the parent of X_i

$$A[S, i] = MIS(S) \forall S \in X_i + \sum_j MIS(X_j)$$

$$B[S, i] = \max_{S' \subset X_i} A[S', i], \quad S = S' \cap X_{k,i}$$



MIS on Graphs of bounded treewidth



R is the root

X_i is a bag

X_j is a child of X_i

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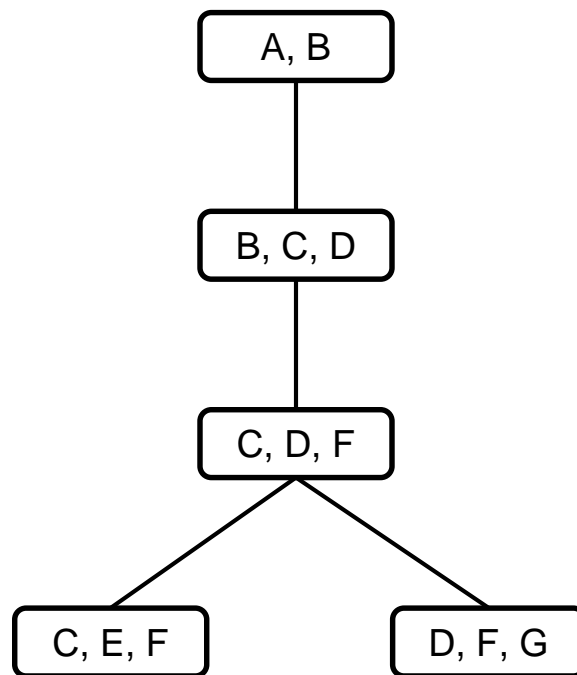
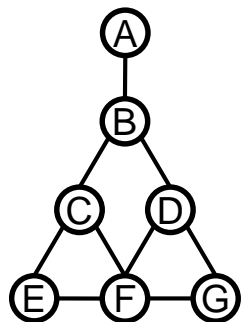
$$MIS(G) = \max_{S \subset R} B[S, R]$$



MIS on Graphs of bounded treewidth

$$A[S, i] = |S| + \sum_j B[S \cap X_j, j] - |S \cap X_j|$$

$$B[S, i] = \max_{S' \subset X_i} A[S', i], \quad S = S' \cap X_{k,i}$$

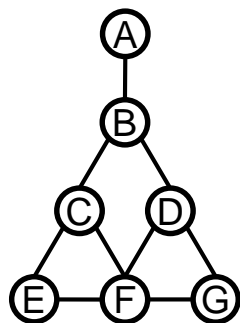




MIS on Graphs of bounded treewidth

$$A[S, i] = |S| + \sum_j B[S \cap X_j, j] - |S \cap X_j|$$

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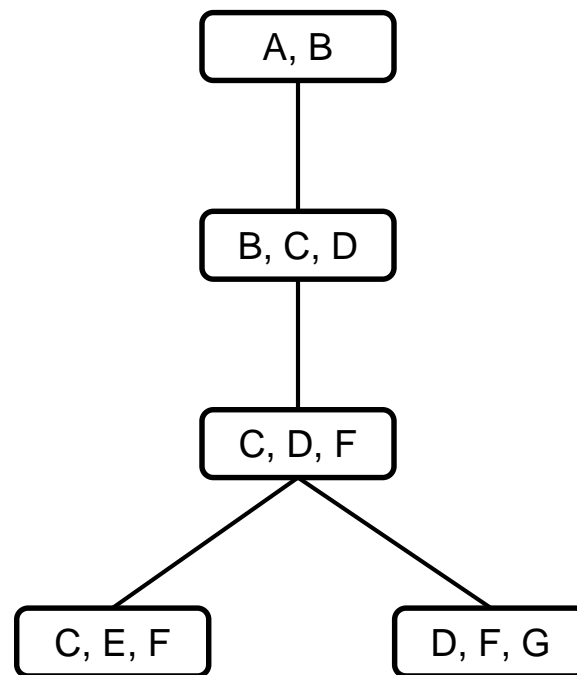


$$A[\emptyset] = 0 \quad A[E] = 1$$

$$A[C] = 1 \quad A[F] = 1$$

$$B[\emptyset] = 0 \quad B[F] = 1$$

$$B[C] = 1$$

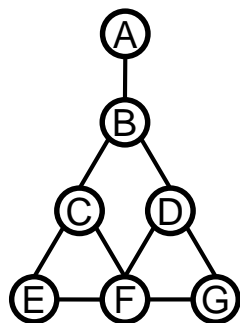




MIS on Graphs of bounded treewidth

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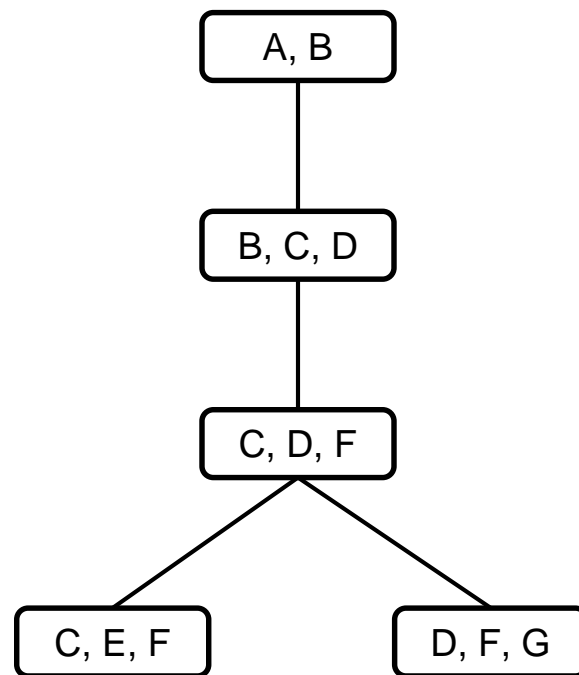


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$$B[\emptyset] = 0 \quad B[F] = 1$$

$$B[C] = 1$$



$$A[\emptyset] = 0 \quad A[F] = 1$$

$$A[D] = 1 \quad A[G] = 1$$

$$B[\emptyset] = 0 \quad B[F] = 1$$

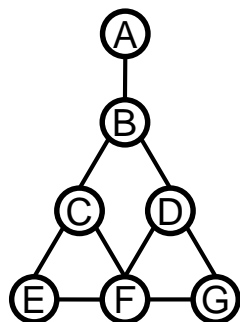
$$B[D] = 1$$



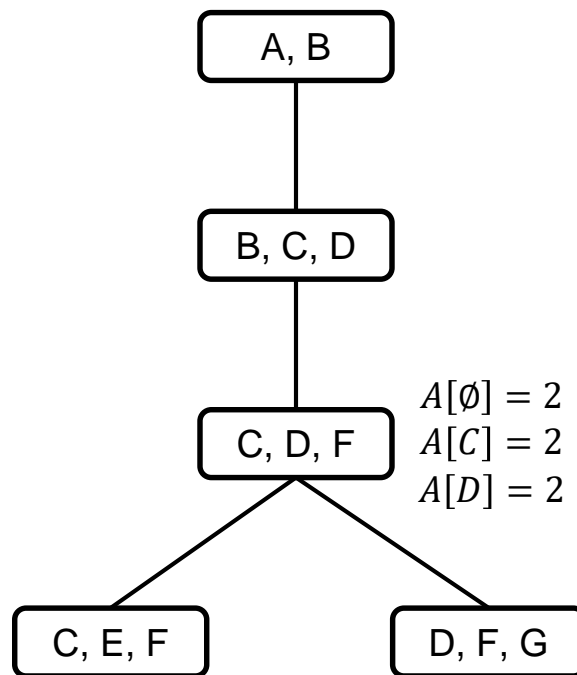
MIS on Graphs of bounded treewidth

$$A[S, i] = |S| + \sum_j B[S \cap X_j, j] - |S \cap X_j|$$

$$B[S, i] = \max_{S' \subset X_i} A[S', i], \quad S = S' \cap X_{k,i}$$



$$\begin{aligned} A[\emptyset] &= 0 & A[E] &= 1 \\ A[C] &= 1 & A[F] &= 1 \\ B[\emptyset] &= 0 & B[F] &= 1 \\ B[C] &= 1 \end{aligned}$$



$$\begin{aligned} A[\emptyset] &= 2 & A[F] &= 1 & B[\emptyset] &= 2 & B[D] &= 2 \\ A[C] &= 2 & A[C, D] &= 2 & B[C] &= 2 & B[C, D] &= 2 \\ A[D] &= 2 \end{aligned}$$

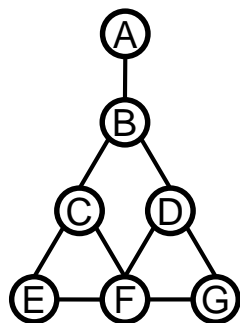
$$\begin{aligned} A[\emptyset] &= 0 & A[F] &= 1 \\ A[D] &= 1 & A[G] &= 1 \\ B[\emptyset] &= 0 & B[F] &= 1 \\ B[D] &= 1 \end{aligned}$$



MIS on Graphs of bounded treewidth

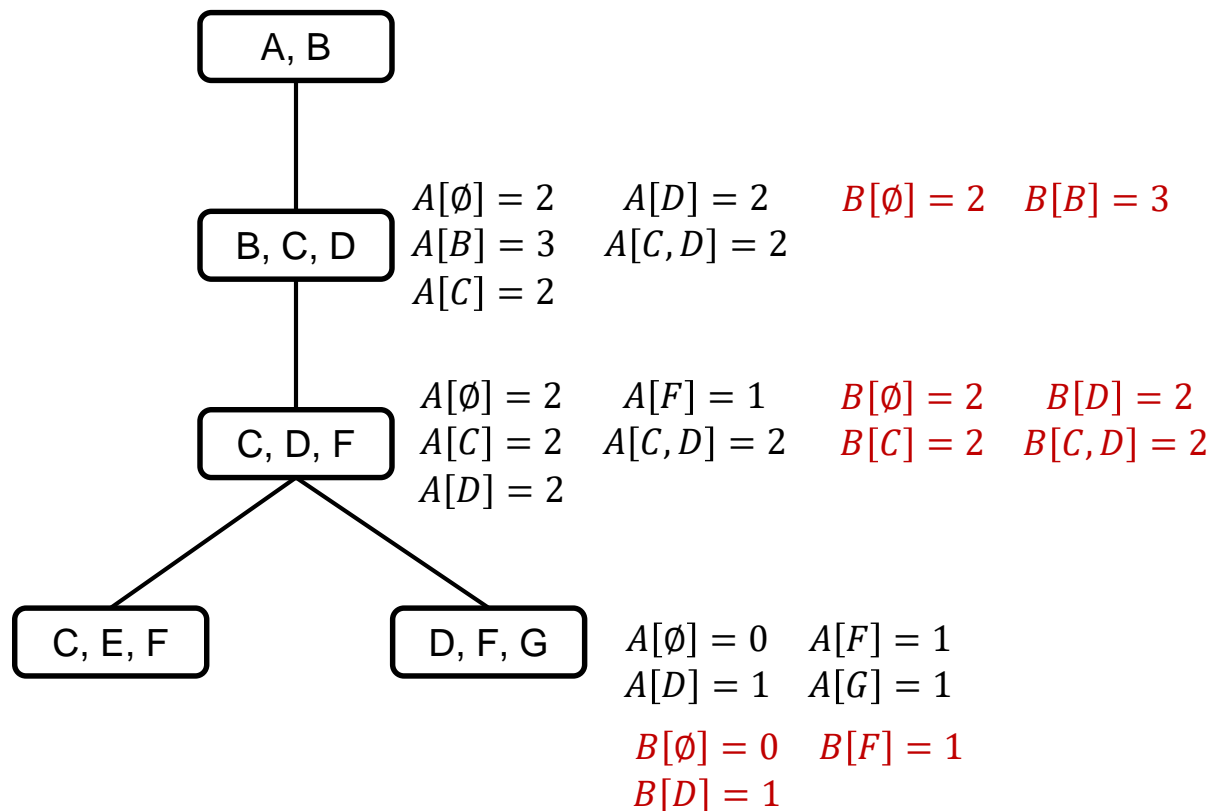
$$A[S, i] = |S| + \sum_j B[S \cap X_j, j] - |S \cap X_j|$$

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$$\begin{array}{ll} B[\emptyset] = 0 & B[F] = 1 \\ B[C] = 1 & \end{array}$$

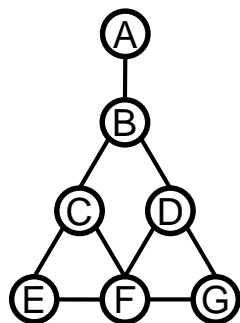




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