

As always, let R be a ring. We noted

$$D(R) \cong K\text{Proj}(R).$$

To recap, a complex $P \in \mathcal{C}(\text{Proj}(R))$ is K projective we have the existing map in the following diagram:

$$\begin{array}{ccc} & & X \\ & \nearrow \exists & \downarrow \wr \\ P & \longrightarrow & Y \end{array}$$

For example,

Example. $P \in \mathcal{C}(\text{Proj}(R))$ with $P_i = 0$ for all $i \ll 0$. To sketch a proof of this claim, construct liftings one step at a time.

$$\left(\begin{array}{ccccccc} & & & & & & X \\ & & & & & \nearrow \tilde{\alpha} & \downarrow \varepsilon \\ \dots \longrightarrow P_{n+1} \longrightarrow P_n \longrightarrow \dots \longrightarrow P_{a+1} \longrightarrow P_a \longrightarrow 0 \longrightarrow \dots \end{array} \right) \xrightarrow{\alpha} Y$$

Here, $\tilde{\alpha}(S)$ should satisfy

- $\varepsilon \tilde{\alpha} = \alpha(S)$
- $\partial \tilde{\alpha}(S) = \tilde{\alpha}(\partial S)$.
- ε is surjective implies that ε is surjective on boundaries.
- $H(\varepsilon)$ is isomorphic (H is the homology functor) implies that $\varepsilon : Z(X) \twoheadrightarrow Z(Y)$
- $\ker(\varepsilon)$ is acyclic.

Once you construct these maps, the claim holds.

Note that every module has a K -projective resolution. Every module M admits a surjection $F_0 \twoheadrightarrow M$ from a free module F_0 , call this map ε . The kernel of ε then admits a surjection $F_1 \twoheadrightarrow \ker \varepsilon$ from a free module F_1 . Continuing the process one more step yields the following diagram:

$$\begin{array}{ccccc} F_2 & \xrightarrow{\partial_2} & F_1 & \xrightarrow{\partial_1} & F_0 \xrightarrow{\varepsilon} M \\ & \searrow & \nearrow & \searrow & \nearrow \\ & & \ker \partial_1 & & \ker \varepsilon \end{array}$$

If we continue this process, we get the K -projective resolution. To make the choice of resolution canonical, we choose the free modules to be freely generated by the elements of the kernels that they surject onto. This yields a functor from modules to K -projective resolutions.

Definition 1. A K -projective resolution of $M \in \mathcal{C}(R)$ is a morphism $\varepsilon : P \rightarrow M$ such that

- ε is a quasi-isomorphism, and
- P is K -projective.

We can make this functorial by taking $R \times M \twoheadrightarrow M$,

$$\varepsilon(r, m) = rm.$$

We then set

$$\bigoplus_{m \in M} R \times m = R^{\oplus M} \twoheadrightarrow M.$$

The last surjection above is the ε map.

Theorem 1. For all $M \in \mathcal{C}(R)$, there exists a surjective K -projective resolution

$$P \twoheadrightarrow M$$

which is a weak equivalence.

Definition 2. An R -complex F is semi-free if it has a filtration

$$(0) \subset F_0 \subset F_1 \subset \dots \subset \bigcup_{n \geq 0} F_n = F$$

such that

1. $F(n) \subset F$ is a subcomplex,
2. $F(n+1)/F(n)$ is a graded free module with trivial boundary maps ($\partial = 0$). In particular, $\partial(F_{n+1}) \subset F_n$.

Exercise 1

Show that semi-free implies K -projective. Also show that unions are K -projective.

Example. Take a complex of free modules

$$\dots \rightarrow F_{n+1} \rightarrow F_n \rightarrow 0.$$

Let $F(n)$ be the truncation consisting of degrees n and lower. Then $\frac{F(n+1)}{F(n)}$ is precisely $\Sigma^{n+1} F_{n+1}$.

Theorem 2. Each $M \in \mathcal{C}(R)$ has a semi-free resolution

$$F \twoheadrightarrow M.$$

(The map above is a weak equivalence.)

Corollary 1. Every K -projective is a retract of a semi-free.

Proof. (Corollary proof): Let P be K -projective. The following diagram completes the proof:

$$\begin{array}{ccc} & & F \\ & \nearrow \tilde{\alpha} & \downarrow \wr \\ P & \xlongequal{\quad} & P \end{array}$$

□

Proof. (Theorem proof): We use the small object argument. We want to construct

$$\begin{array}{ccccccc} F(1) & \hookrightarrow & F(2) & \hookrightarrow & \dots & \hookrightarrow & F(n) \hookrightarrow \dots \hookrightarrow \bigcup_{n \geq 0} F(n) \\ \varepsilon(1) \downarrow & & \swarrow \varepsilon(2) & & & \searrow \varepsilon(n) & \\ & & M & & & & \end{array}$$

such that $\frac{F(n+1)}{F(n)}$ is a graded complex of free modules with zero differential which is surjective on cycles. We need that $\varepsilon(1)$ is surjective on cycles. We have

$$\begin{array}{ccccccc} F(1) & \xlongequal{\quad} & H(F(1)) & \longrightarrow & H(F(n)) & \xrightarrow{\phi} & H(F(n+1)) \\ & & \downarrow & \xRightarrow{\quad} & \swarrow & & \\ & & H(M) & & H(\varepsilon(n)) & & \end{array}$$

for all n where the kernel of ϕ contains the kernel of $H(\varepsilon(n))$, so that $H(\varepsilon(n))$ is surjective for all n .

Next, we note $\ker(H(\varepsilon(n))) \subset H(F(n))$ goes to 0 under $H(i(n))$. We have $H(F) \xrightarrow{\sim} H(M)$.

Remark. Given $\varepsilon : X \rightarrow Y$ where $Z(\varepsilon)$ and $H(\varepsilon)$ are surjective, we have that ε is surjective. See the four lemmas applied to the following diagrams:

$$\begin{array}{ccccccccc} 0 & \longrightarrow & B(X) & \longrightarrow & Z(X) & \longrightarrow & H(X) & \longrightarrow & 0 \\ \parallel & & \downarrow & & \downarrow & & \downarrow & & \parallel \\ 0 & \longrightarrow & B(Y) & \longrightarrow & Z(Y) & \longrightarrow & H(Y) & \longrightarrow & 0 \\ \\ 0 & \longrightarrow & Z(X) & \longrightarrow & X & \longrightarrow & \varepsilon(B(X)) & \longrightarrow & 0 \\ \parallel & & \downarrow & & \downarrow & & \downarrow & & \parallel \\ 0 & \longrightarrow & Z(Y) & \longrightarrow & Y & \longrightarrow & \varepsilon(B(Y)) & \longrightarrow & 0 \end{array}$$

Say we have calculated $\varepsilon(n) : F(n) \rightarrow M$. Choose cycles $\{z_\lambda\}$ in $F(n)$ that map to the generating set of $\ker(H(\Sigma(n)))$. Pick $w_\lambda \in M$ such that $\partial w_\lambda = \varepsilon(n)(z_\lambda)$. Set

$$F(n+1) = F(n) \oplus_\lambda R \cdot e_\lambda.$$

Above, we have $|e_\lambda| = |z_\lambda| + 1$ and $\partial|_{F(n)} = \partial^{F(n)}$. (Attached are some iPad notes, where the proof continues.) \square

$$E(n+1)(e_\lambda) = w_\lambda.$$

$$E(n+1): F(n) \rightarrow M$$

- $E(n+1)|_{F(n)} = E(n)$
- $E(n+1)(e_\lambda) = w_\lambda$

Can make a functorial realization of the semi-free resolution of any complex by making free module quite big. Given any cycle, also include info on how it reads a hel.

There are other ways of constructing resolutions.

$$M$$

$$H(M) = \{H_i(M)\}_{i \in \mathbb{Z}} \text{ graded module.}$$

$$P \xrightarrow{\sim} H(M) \text{ } K\text{-projective resolution.}$$

One can perturb the differential on P to construct a K -projective resolution of M . (Adams' resolution)
(Cartan / Eilenberg resolution)

Example: P projective $\Rightarrow P_i$ projective modules.

Converse is false. (Dale)

$$R = \mathbb{Z}/4\mathbb{Z}$$

$$\Rightarrow R \xrightarrow{2} R \xrightarrow{2} R \rightarrow \dots$$

Not K -projective.

Ex: Suppose P is K -projective and acyclic. Then P is contractible, i.e. $\text{id} \simeq 0$

$$K \text{ Proj } (R) \hookrightarrow K(\text{Proj } R)$$

$$\uparrow =$$

$$D(R)$$

Fact: Say R is commutative Noetherian. Then
 $K \text{ Proj } (R) = K(\text{Proj } R) \iff R$ is regular.

Ex: $\mathbb{Z}, \mathbb{Z}[x], k[x]$ are regular rings

Derived functors: Given an R -complex.

If $P \xrightarrow{\sim} M, Q \xrightarrow{\sim} N$ K projective resolutions then $P \simeq Q$
 in $K(\text{Proj } R)$.

$$\begin{array}{ccc} & \xrightarrow{\sim} & F \\ P & \xrightarrow{\sim} & M \end{array}$$

Given any N , set

$$R\text{Hom}_R(M, N) := \text{Hom}_R(P, N) \text{ where } P \text{ is a } K\text{-proj resn.}$$

$$\text{Ext}_R^i(M, N) := H^i(R\text{Hom}_R(M, N)) = H^{-i}(\text{Hom}_R(P, N)).$$

$$R\text{Hom}_R(M, -) : \mathcal{C}(R) \rightarrow K(\mathbb{Z})$$

or $K(R)$ if R is comm.

$$\text{Ext}_R^0(M, N) = H^0(\text{Hom}_R(P, N)).$$

$$\text{Ext}_R^i(M, N) = \text{morphisms } P \rightarrow \Sigma^i N.$$

If $Q \xrightarrow{\sim} N$ is K -projective resn. then

$$\text{Hom}_R(P, Q) \xrightarrow{\sim} \text{Hom}_R(P, N).$$

is quasi-iso

— \mathcal{C}_R — $M \in \mathcal{C}(R^{op})$
 ex. of right R -mods
 $N \in \mathcal{C}(R)$

$$M \otimes_R N \quad \text{ex } (\text{of } \mathbb{Z}\text{-modules})$$

$$(M \otimes_R N)_n = \bigoplus_{i \in \mathbb{Z}} M_i \otimes N_{n-i}$$

$$\downarrow \cong \quad \downarrow \partial^m \otimes 1 = 1 \otimes \partial^n$$

$$(M \otimes_R N)_{n+1}$$

$$\partial(m \otimes n) = \partial m \otimes n + (-1)^m m \otimes \partial n.$$

Fact. $X \simeq Y$ quasi-is.

Then $P \otimes X \simeq P \otimes Y$ for any K -Proj complex P

Tensor products behave nicely w/ limits.

In absence of subscripts
 $\text{Tor}^R(M, N)$ refers to
 $\text{Tor}_*^R(M, N)$.

- Check for semi-free

$$\text{Defn: } M \otimes_R^L N = P \otimes_R N$$

$$P \simeq M$$

Is a K -proj resln.

$$\text{Tor}_i^R(M, N) = H_i(P \otimes_R N).$$

$$X \xrightarrow[\text{quasi-is}]{\simeq} Y \implies \text{Tor}^R(M, X) \simeq \text{Tor}^R(M, Y)$$