Timing Statistics:

| Input Size | Algorithm | Average Time (ms) |
|------------|-----------------------|-------------------|
| 1k | Standard Sort | 0.4 |
| | HalfSelectionSort | 3 |
| | MergeSort | 1 |
| | InPlaceMerge | 1 |
| | HalfHeapSort | 0.2 |
| | QuickSelect | 0.8 |
| | QS w/ MedianofMedians | 0.6 |
| 31k | Standard Sort | 3 |
| | HalfSelectionSort | 1188 |
| | MergeSort | 44 |
| | InPlaceMerge | 27 |
| | HalfHeapSort | 7 |
| | QuickSelect | 2 |
| | QS w/ MedianofMedians | 10 |
| 1M | Standard Sort | 52 |
| | HalfSelectionSort | Input too big |
| | MergeSort | 1094 |
| | InPlaceMergeSort | 641 |
| | HalfHeapSort | 157 |
| | QuickSelect | 37 |
| | QS w/ MedianofMedians | 131 |

Worst Case Input quickSelect:

Median of 3: ~167 ms average

Median of Medians: ~5 ms average

Algorithmic Analysis:

HalfSelectionSort:

O(n^2) but more precisely we are only swapping and comparing about half as much as a full selection sort, so the constant factor will be a lot smaller, and in practice it will be faster than full selection sort.

std::sort

is O(n^2) although this is a rare case. Its average and best case is O(nlogn) and in practice it's generally faster than any other algorithm used in this project apart from quickSelect.

mergeSort and inPlaceMergeSort

Both mergeSort and inPlaceMergeSort are O(nlogn) but inPlaceMerge has a smaller space complexity.std::merge has a space complexity O(n+m) and std::inplace_merge has space complexity O(1) since it uses a constant amount of additional memory

halfHeapSort

Is still O(nlogn) like full Heap Sort. This is because you still build a full heap which is O(n) and the deleteMax phase will be O(logn)/2 which is still O(logn). In practice it'll be faster but the worst case time complexity is the same.

quickSelect

Has average time complexity O(n) and worst case time complexity $O(n^2)$ although this is quite rare and depends on the pivot selection method used and the input. Median of 3 generally avoids worst case scenarios and helps quickSelect average O(n) time. My worstCasePivot generator function causes quickSelect with the median of 3 pivot selection method to approach $O(n^2)$.

Using the **median of medians** method for pivot selection guarantees O(n) complexity although this comes with a large constant factor. In practice, median of medians was a lot slower for large inputs. However it really improved performance for quickselect on the worst case inputs (~5ms instead of ~167ms average).

What surprised me: That inPlaceMerge generally performs better than merge. Also how much better medianOfMedians performs for worst case input. 167 vs 5 is a huge difference!!