

Homework 8

CS 325

Daniel Kim

Problem 1: (20 pts) In the bin packing problem, items of different weights (or sizes) must be packed into a finite number of bins each with the capacity C in a way that minimizes the number of bins used. The decision version of the bin packing problem (deciding if objects will fit into $\leq k$ bins) is NP-complete. There is no known polynomial time algorithm to solve the optimization version of the bin packing problem. In this homework you will be examining three greedy approximation algorithms to solve the bin packing problem.

- **First-Fit:** Put each item as you come to it into the first (earliest opened) bin into which it fits. If there is no available bin then open a new bin.
- **First-Fit-Decreasing:** First sort the items in decreasing order by size, then use First-Fit on the resulting list.
- **Best Fit:** Place the items in the order in which they arrive. Place the next item into the bin which will leave the least room left over after the item is placed in the bin. If it does not fit in any bin, start a new bin.

a) Give pseudo code and the running time for each of the approximation algorithms.

b) Implement the algorithms in Python, C++ or C. Your program named `binpack` should read in a text file named `bin.txt` with multiple test cases as explained below and output to the terminal the number of bins each algorithm calculated for each test case. Submit a README file and your program to TEACH.

Example bin.txt: The first line is the number of test cases, followed by the capacity of bins for that test case, the number of items and then the weight of each item. You can assume that the weight of an item does not exceed the capacity of a bin for that problem.

```
3
10
6
5 10 2 5 4 4
10
20
4 4 4 4 4 4 4 4 4 6 6 6 6 6 6 6 6 6
10
4
3 8 2 7
```

Sample output:

```
Test Case 1 First Fit: 4, First Fit Decreasing: 3, Best Fit: 4
Test Case 2 First Fit: 15, First Fit Decreasing: 10, Best Fit: 15
Test Case 3 First Fit: 3, First Fit Decreasing: 2, Best Fit: 2
```

c) Randomly generate at least 20 bin packing instances. Summarize the results for each algorithm. Which algorithm performs better? How often? *Note: Submit a description of how the inputs were generated not the code used to produce the random inputs.*

a) First-Fit

For all items $i = 1, 2, 3, \dots, n$ do

 For all bins $j = 1, 2, \dots$ do

 if item i fits bin j then

 Put item i in bin j .

 Break and move to next item

 If item does not fit in any current bins,

 create a new bin

 Put item in new bin

Time complexity for First-Fit = $O(n^2)$

First-Fit-Decreasing

Sort items $i=1, \dots, n$ in decreasing order

Apply **First-Fit** algorithm (above)

Time complexity for First-Fit-Decreasing = $O(n^2)$

Best-Fit

for all items $i = 1, 2, \dots, n$ do

 for All bins $j = 1, 2, \dots$ do

 if item i fits in bin j then

 Add item and calculate remaining capacity

 Pack item i in the bin j with minimum calculated remaining capacity if item i were packed

 If no bin exists

 Create a new bin

 Add item to new bin

Time complexity for **Best-Fit** = $O(n^2)$

b) Code Submitted to Teach! Thank you!

c) Random Numbers Generated

1. Cases (iterations) = 300
2. Bin Capacity Weight (integer) from 10 to 50
3. Item Weight (integer) from 2 to Bin Capacity
4. Number of Items (integer) from 10 to 30

Results

- All methods (First-fit, First-fit-decreasing, and Best-fit) had equal performance: 67.33% of 300 iterations
- First-fit-decreasing had the best outcome 32.33% of 300 iterations
- First-fit resulted in the best outcome 0.33% of 300 iterations

Problem 2: (10 pts) An exact solution to the bin packing optimization problem can be found using 0-1 integer programming (IP) see the format on the [Wikipedia page](#).

Write an integer program for each of the following instances of bin packing and solve with the software of your choice. Submit a copy of the code and interpret the results.

a) Six items $S = \{4, 4, 4, 6, 6, 6\}$ and bin capacity of 10

b) Five items $S = \{20, 10, 15, 10, 5\}$ and bin capacity of 20

Note: The version of LINDO that you have access to on the OSU server has a limit of 50 integer variables. Therefore, LINDO will only be able to solve problems with at most 6 items.

a) MIN $Y_1 + Y_2 + Y_3 + Y_4 + Y_5 + Y_6$

ST

$$X_{11} + X_{12} + X_{13} + X_{14} + X_{15} + X_{16} = 1$$

$$X_{21} + X_{22} + X_{23} + X_{24} + X_{25} + X_{26} = 1$$

$$X_{31} + X_{32} + X_{33} + X_{34} + X_{35} + X_{36} = 1$$

$$X_{41} + X_{42} + X_{43} + X_{44} + X_{45} + X_{46} = 1$$

$$X_{51} + X_{52} + X_{53} + X_{54} + X_{55} + X_{56} = 1$$

$$X_{61} + X_{62} + X_{63} + X_{64} + X_{65} + X_{66} = 1$$

$$4X_{11} + 4X_{21} + 4X_{31} + 6X_{41} + 6X_{51} + 6X_{61} - 10Y_1 \leq 0$$

$$4X_{12} + 4X_{22} + 4X_{32} + 6X_{42} + 6X_{52} + 6X_{62} - 10Y_2 \leq 0$$

$$4X_{13} + 4X_{23} + 4X_{33} + 6X_{43} + 6X_{53} + 6X_{63} - 10Y_3 \leq 0$$

$$4X_{14} + 4X_{24} + 4X_{34} + 6X_{44} + 6X_{54} + 6X_{64} - 10Y_4 \leq 0$$

$$4X_{15} + 4X_{25} + 4X_{35} + 6X_{45} + 6X_{55} + 6X_{65} - 10Y_5 \leq 0$$

$$4X_{16} + 4X_{26} + 4X_{36} + 6X_{46} + 6X_{56} + 6X_{66} - 10Y_6 \leq 0$$

END

INT X_{11}

INT X_{12}

INT X_{13}

INT X14

INT X15

INT X16

INT Y2

INT X21

INT X22

INT X23

INT X24

INT X25

INT X26

INT Y3

INT X31

INT X32

INT X33

INT X34

INT X35

INT X36

INT Y4

INT X41

INT X42

INT X43

INT X44

INT X45

INT X46

INT X51

INT X52

INT X53

INT X54

INT X55

INT X56

INT Y6

INT X61

INT X62

INT X63

INT X64

INT X65

INT X66

```
MAX  
X  
<untitled>  
MIN Y1+Y2+Y3+Y4+Y5+Y6  
ST  
X11 + X12 + X13 + X14 + X15 + X16 = 1  
X21 + X22 + X23 + X24 + X25 + X26 = 1  
X31 + X32 + X33 + X34 + X35 + X36 = 1  
X41 + X42 + X43 + X44 + X45 + X46 = 1  
X51 + X52 + X53 + X54 + X55 + X56 = 1  
X61 + X62 + X63 + X64 + X65 + X66 = 1  
  
4X11 + 4X21 + 4X31 + 6X41 + 6X51 + 6X61 - 10Y1 <= 0  
4X12 + 4X22 + 4X32 + 6X42 + 6X52 + 6X62 - 10Y2 <= 0  
4X13 + 4X23 + 4X33 + 6X43 + 6X53 + 6X63 - 10Y3 <= 0  
4X14 + 4X24 + 4X34 + 6X44 + 6X54 + 6X64 - 10Y4 <= 0  
4X15 + 4X25 + 4X35 + 6X45 + 6X55 + 6X65 - 10Y5 <= 0  
4X16 + 4X26 + 4X36 + 6X46 + 6X56 + 6X66 - 10Y6 <= 0  
END  
  
INT X11  
INT X12  
INT X13  
INT X14  
INT X15  
INT X16  
INT Y2  
INT X21  
INT X22  
INT X23  
INT X24  
INT X25  
INT X26  
INT Y3  
INT X31  
INT X32  
INT X33  
INT X34  
INT X35  
INT X36  
INT Y4  
INT X41  
INT X42  
INT X43  
INT X44  
INT X45  
INT X46  
INT X51  
INT X52  
INT X53  
INT X54  
INT X55  
INT X56  
INT Y6  
INT X61  
INT X62  
INT X63  
INT X64  
INT X65  
INT X66
```

Reports Window		
8)	0.000000	0.100000
9)	0.000000	0.000000
10)	0.000000	0.000000
11)	0.000000	0.000000
12)	0.000000	0.100000
13)	0.000000	0.000000
NO. ITERATIONS= 21		
BRANCHES= 0 DETERM.= 1.000E 0		
LP OPTIMUM FOUND AT STEP 40		
OBJECTIVE VALUE = 3.000000000		
NEW INTEGER SOLUTION OF 3.000000000 AT BRANCH 0 PIVOT 42		
BOUND ON OPTIMUM: 3.000000		
ENUMERATION COMPLETE. BRANCHES= 0 PIVOTS= 42		
LAST INTEGER SOLUTION IS THE BEST FOUND		
RE-INSTALLING BEST SOLUTION...		
OBJECTIVE FUNCTION VALUE		
1)	3.000000	
VARIABLE	VALUE	REDUCED COST
X11	0.000000	0.400000
X12	1.000000	0.000000
X13	0.000000	0.000000
X14	0.000000	0.000000
X15	0.000000	0.400000
X16	0.000000	0.000000
Y2	1.000000	1.000000
X21	0.000000	0.400000
X22	0.000000	0.000000
X23	0.000000	0.000000
X24	0.000000	0.000000
X25	1.000000	0.400000
X26	0.000000	0.000000
Y3	0.000000	1.000000
X31	1.000000	0.400000
X32	0.000000	0.000000
X33	0.000000	0.000000
X34	0.000000	0.000000
X35	0.000000	0.400000
X36	0.000000	0.000000
Y4	0.000000	1.000000
X41	0.000000	0.600000
X42	0.000000	0.000000
X43	0.000000	0.000000
X44	0.000000	0.000000
X45	1.000000	0.600000
X46	0.000000	0.000000
X51	0.000000	0.600000
X52	1.000000	0.000000
X53	0.000000	0.000000
X54	0.000000	0.000000
X55	0.000000	0.600000
X56	0.000000	0.000000
Y6	0.000000	1.000000
X61	0.000000	0.600000
X62	0.000000	0.000000
X63	0.000000	0.000000
X64	0.000000	0.000000
X65	1.000000	0.600000
X66	0.000000	0.000000
Y1	0.400000	0.000000
Y5	1.600000	0.000000
ROW	SLACK OR SURPLUS	DUAL PRICES
2)	0.000000	0.000000
3)	0.000000	0.000000
4)	0.000000	0.000000
5)	0.000000	0.000000
6)	0.000000	0.000000
7)	0.000000	0.000000
8)	0.000000	0.100000
9)	0.000000	0.000000
10)	0.000000	0.000000
11)	0.000000	0.000000
12)	0.000000	0.100000
13)	0.000000	0.000000
NO. ITERATIONS= 42		
BRANCHES= 0 DETERM.= 1.000E 0		

b)

Reports Window							
X41	1.000000	0.000000		MIN Y1+Y2+Y3+Y4+Y5			
X42	0.000000	0.000000		ST			
X43	0.000000	0.000000		X11+	X12+	X13+	X14+ X15= 1
X44	0.000000	0.000000		X21+	X22+	X23+	X24+ X25= 1
X45	0.000000	0.000000		X31+	X32+	X33+	X34+ X35= 1
Y5	1.000000	1.000000		X41+	X42+	X43+	X44+ X45= 1
X51	0.000000	0.000000		X51+	X52+	X53+	X54+ X55= 1
X52	0.000000	0.000000		20X11 + 10X21 + 15X31 + 10X41 + 5X51 -20Y1 <=0			
X53	0.000000	0.000000		20X12 + 10X22 + 15X32 + 10X42 + 5X52 -20Y2 <=0			
X54	0.000000	0.000000		20X13 + 10X23 + 15X33 + 10X43 + 5X53 -20Y3 <=0			
X55	1.000000	0.000000		20X14 + 10X24 + 15X34 + 10X44 + 5X54 -20Y4 <=0			
				20X15 + 10X25 + 15X35 + 10X45 + 5X55 -20Y5 <=0			
ROW	SLACK OR SURPLUS	DUAL PRICES		END			
2)	0.000000	0.000000		INT Y1			
3)	0.000000	0.000000		INT Y2			
4)	0.000000	0.000000		INT Y3			
5)	0.000000	0.000000		INT Y4			
6)	0.000000	0.000000		INT Y5			
7)	0.000000	0.000000		INT X11			
8)	0.000000	0.000000		INT X12			
9)	0.000000	0.000000		INT X13			
10)	0.000000	0.000000		INT X14			
11)	0.000000	0.000000		INT X15			
NO. ITERATIONS= 48				INT X21			
BRANCHES= 0 DETERM.= 1.000E 0				INT X22			
LP OPTIMUM FOUND AT STEP 19				INT X23			
OBJECTIVE VALUE = 3.00000000				INT X24			
				INT X25			
NEW INTEGER SOLUTION OF 3.00000000 AT BRANCH 0 PIVOT 19				INT X31			
RE-INSTALLING BEST SOLUTION...				INT X32			
OBJECTIVE FUNCTION VALUE				INT X33			
1)	3.000000			INT X34			
VARIABLE	VALUE	REDUCED COST		INT X35			
Y1	1.000000	1.000000		INT X41			
Y2	1.000000	1.000000		INT X42			
Y3	1.000000	1.000000		INT X43			
Y4	0.000000	1.000000		INT X44			
Y5	0.000000	1.000000		INT X45			
X11	0.000000	0.000000		INT X51			
X12	1.000000	0.000000		INT X52			
X13	0.000000	0.000000		INT X53			
X14	0.000000	0.000000		INT X54			
X15	0.000000	0.000000		INT X55			
X21	1.000000	0.000000					
X22	0.000000	0.000000					
X23	0.000000	0.000000					
X24	0.000000	0.000000					
X25	0.000000	0.000000					
X31	0.000000	0.000000					
X32	0.000000	0.000000					
X33	1.000000	0.000000					
X34	0.000000	0.000000					
X35	0.000000	0.000000					
X41	1.000000	0.000000					
X42	0.000000	0.000000					
X43	0.000000	0.000000					
X44	0.000000	0.000000					
X45	0.000000	0.000000					
X51	0.000000	0.000000					
X52	0.000000	0.000000					
X53	1.000000	0.000000					
X54	0.000000	0.000000					
X55	0.000000	0.000000					
ROW	SLACK OR SURPLUS	DUAL PRICES					
2)	0.000000	0.000000					
3)	0.000000	0.000000					
4)	0.000000	0.000000					
5)	0.000000	0.000000					
6)	0.000000	0.000000					
7)	0.000000	0.000000					
8)	0.000000	0.000000					
9)	0.000000	0.000000					
10)	0.000000	0.000000					
11)	0.000000	0.000000					
NO. ITERATIONS= 19							
BRANCHES= 0 DETERM.= 1.000E 0							

MIN Y1+Y2+Y3+Y4+Y5

ST

X11+X12+X13+X14+X15= 1

X21+X22+X23+X24+X25= 1

X31+X32+X33+X34+X35= 1

X41+ X42+ X43+ X44+ X45= 1

X51+ X52+ X53+ X54+ X55= 1

$$20X_{11} + 10X_{21} + 15X_{31} + 10X_{41} + 5X_{51} - 20Y_1 \leq 0$$

$$20X_{12} + 10X_{22} + 15X_{32} + 10X_{42} + 5X_{52} - 20Y_2 \leq 0$$

$$20X_{13} + 10X_{23} + 15X_{33} + 10X_{43} + 5X_{53} - 20Y_3 \leq 0$$

$$20X_{14} + 10X_{24} + 15X_{34} + 10X_{44} + 5X_{54} - 20Y_4 \leq 0$$

$$20X_{15} + 10X_{25} + 15X_{35} + 10X_{45} + 5X_{55} - 20Y_5 \leq 0$$

END

INT Y1

INT Y2

INT Y3

INT Y4

INT Y5

INT X11

INT X12

INT X13

INT X14

INT X15

INT X21

INT X22

INT X23

INT X24

INT X25

INT X31

INT X32

INT X33

INT X34

INT X35

INT X41

INT X42

INT X43

INT X44

INT X45

INT X51

INT X52

INT X53

INT X54

INT X55