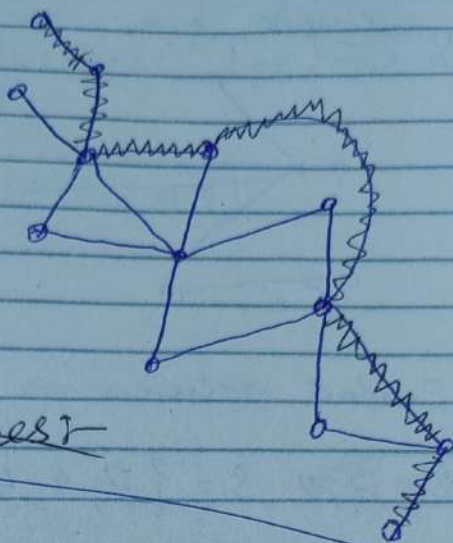


After sessional 1

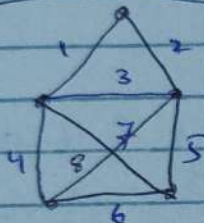
LECTURE # 8

farthest

connect
and
circuits

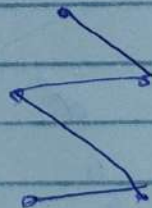
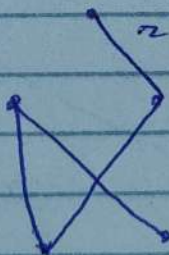
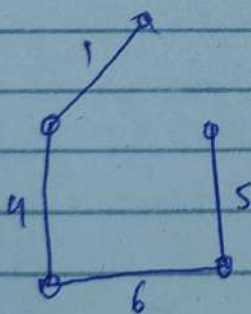


Distance of trees



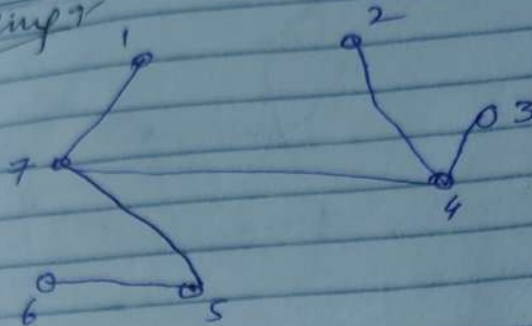
Number of different edges in two spanning trees of a graph ~~are~~ is called distance of the

edges



In case of fully connected graph the maximum number of distance of trees could be $n-1$ (maybe I think).

PRUFER CODES:- Encoding



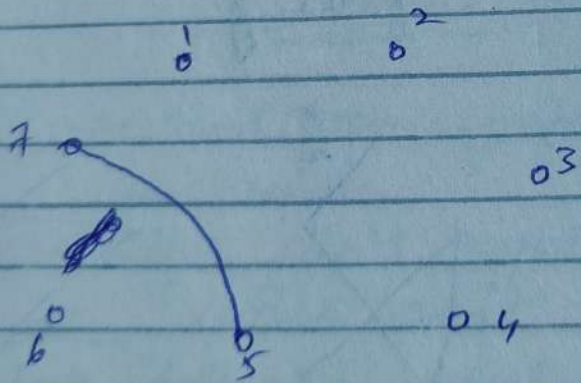
Pendant vertices $\Rightarrow 1, 2, 6, 3$

~~$S = \{7, 4, 4, 7, 5\}$~~

\rightarrow Remove the edge between 1 & 7 and add the 7 in S.

\rightarrow Same steps for all pendant vertices.

This is left:



~~$S = \{7, 4, 4, 7, 5\}$~~
 $S = \{7, 4, 4, 7, 5\}$
 $V = \{1, 2, 3, 4, 5, 6, 7\}$

\rightarrow Now remove the corresponding element from the start of S & V and connect those vertices in the graph.

\rightarrow It is ve and receive saving, because

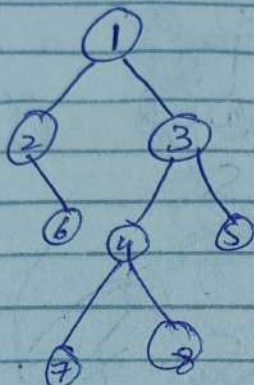
preorders

1, 2,

$\rightarrow 4a$
 $4x$

→ It is very useful for sending and receiving tree and also for saving, because it takes less space.

LECTURE #9

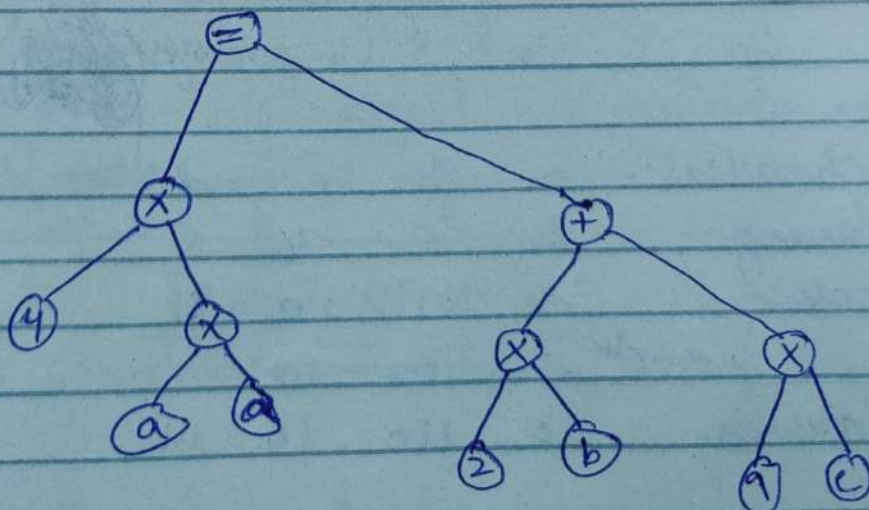


preorder:-

1, 2, 6, 3, 4, 7, 8, 5

$$\rightarrow 4a^2 = 2b + 9c$$

$$4 \times (a \times a) = (2 \times b) + (9 \times c)$$



Preorder:- $y = -x3 + y1/a \times yy$



Inorder:-

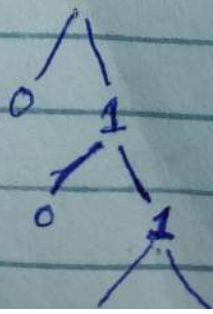
$$y = (3 \times (y + 1)) - (a / (y \times y))$$

Storage of messages upto characters:-

a, b, c, d
with frequencies
60, 5, 30, 5 (in percentages)

Characters:-	a	b	c	d
frequency:-	60	5	30	5
code:-	00	01	10	11

Optimization
Representation
0 110 10 111



Example:-
Data file

characters:- a
frequency:- 45
fixed

variable

Assignment

Columns:-

Airline
departure

→

and

Minim

- ① P
- ② K
- ③ D

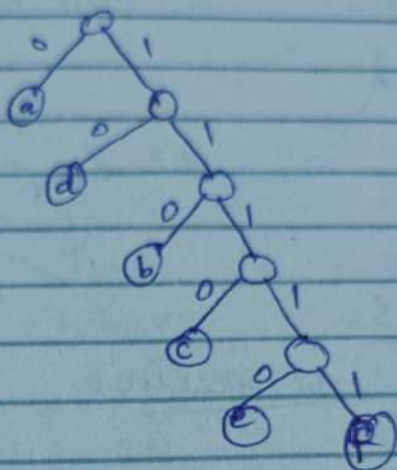
Example:-

Data File: 100,000 characters

characters:- a b c d e f
frequency:- 45 13 12 16 9 5

fixed length code: $3 \times 100,000$
 $= 300,000 \text{ bits}$

variable 0 110 1110 10 11110 11111



Assignment 1:-

Columns:-

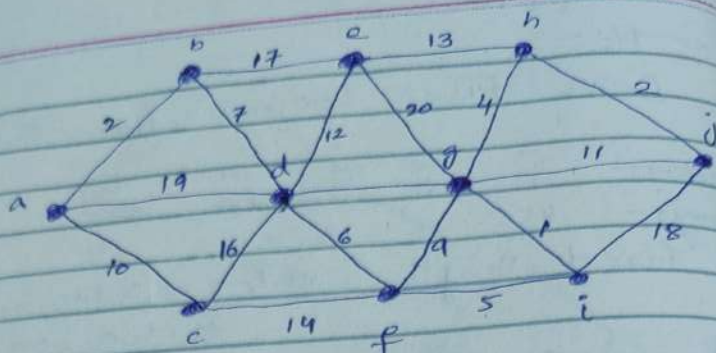
→ Airline, Origin airport, destination airport,
departure delay, arrival delay

→ ~~pick~~ pick origin airport and destination airport
and add ~~edge~~ weighted edges.

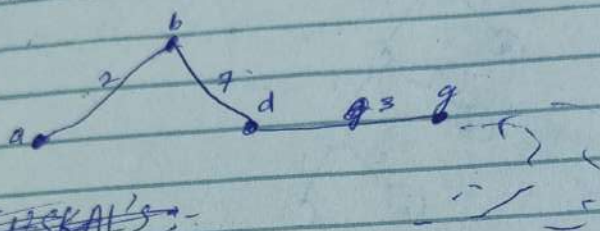
LECTURE # 10

Minimum Spanning Tree:-

- ① Prim's
- ② Kruskal's
- ③ Dijkstra's



PRIMS:
 untraversed = $\{a, b, c, d, e, f, g, h, i, j\}$



KRUSKAL'S:
PRIMS:
 programming perspectives

- ① Traversed nodes: $0 \rightarrow 10$
- ② Untraversed nodes: $10 \rightarrow 0$
- ③ Adjacent edge list: $e_1, e_2, e_3, \dots, e_2, e_3, e_6, e_7$

LECTURE # 11

Euler circuit/path/Tour

Hamiltonian path/cycle/circuit

Vertex connectivity:-

Minimum number of vertices
 must be removed to divide the
 graph into multiple components.

Q: This is a planar

Prim's Algorithm

that are connected

we are considering

Kruskal's

We

as a whole

minimum of

We also

chosen minimum

Planar Graph

Geometric

in such a way

intersect

Regions
 faces

Euler's

C

g e

To

n >

must

be

Q: This is a planar graph, show it

Lecture # 12

Prim's Algorithm:

We see only the edges that are connected with the edge we are currently at.

Kruskal's Algorithm:

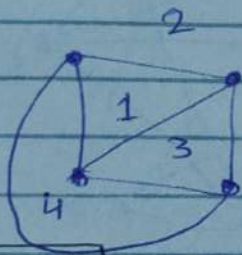
We see the whole graph as a whole and choose the minimum one.

We also check the edge we chosen must not make a cycle.

Planar Graphs:-

Geometric representation of graph in such a way that no two edges intersect each other.

Regions
faces.



Euler's Formula:-

Connected graph with n vertices & e edges has $e - n + 2$ regions.

$$n = 4, e = 6$$

$$\text{Regions} = e - n + 2 = 6 - 4 + 2$$

$$\text{Regions} = 4$$

To check graph is planar or not:-

Graph G is a simple graph with $n > 3$, then the following rule must hold in order for this graph to be planar