National University of Singapore
School of Computing
CS1010S: Programming Methodology
Semester I, 2018/2019

Solutions for Recitation 2
Recursion, Iteration & Orders of Growth

Definitions

Theta (Θ) notation:

$$f(n) = \Theta(g(n)) \;\Leftrightarrow\; \exists k_1, k_2, n_0 \;.\; \; k_1 \cdot g(n) \leq f(n) \leq k_2 \cdot g(n), \text{for } n > n_0$$

Big-O notation:

$$f(n) = O(g(n)) \;\Leftrightarrow\; \exists k, n_0 \;.\; \; f(n) \leq k \cdot g(n), \text{for } n > n_0$$

Adversarial approach: For you to show that $f(n) = \Theta(g(n))$, you pick k_1, k_2, and n_0, then I (the adversary) try to pick an n which doesn't satisfy $k_1 \cdot g(n) \leq f(n) \leq k_2 \cdot g(n)$.

Implications

Ignore constants. Ignore lower order terms. For a sum, take the larger term. For a product, multiply the two terms. Orders of growth are concerned with how the effort scales up as the size of the problem increases, rather than an exact measure of the cost.

Typical Orders of Growth

- $\Theta(1)$ - Constant growth. A fixed number of simple, non-decomposable operations have constant growth.
- $\Theta(\log n)$ - Logarithmic growth. At each iteration, the problem size is scaled down by a constant amount.
- $\Theta(n)$ - Linear growth. At each iteration, the problem size is decremented by a constant amount.
- $\Theta(n \log n)$ - Nifty growth. Nice recursive solution to normally $\Theta(n^2)$ problem.
- $\Theta(n^2)$ - Quadratic growth. Computing correspondence between a set of n things, or doing something of cost n to all n things both result in quadratic growth.
- $\Theta(2^n)$ - Exponential growth. Really bad. Searching all possibilities usually results in exponential growth.

What's n?

Order of growth is *always* in terms of the size of the problem. Without stating what the problem is, and what is considered primitive (what is being counted as a "unit of work" or "unit of space"), the order of growth doesn't have any meaning.

Problems

1. Remember our point-of-sale and order-tracking system from last week? Recall that the joint only sells 4 options for combos: Classic Single Combo (hamburger with one patty), Classic Double With Cheese Combo (2 patties), and Classic Triple with Cheese Combo (3 patties), Avant-Garde Quadruple with Guacamole Combo (4 patties). We shall encode these combos as 1, 2, 3, and 4 respectively. Each meal can be *biggie-sized* to acquire a larger box of fries and drink. A *biggie-sized* combo is represented by 5, 6, 7, and 8 respectively, for combos 1, 2, 3, and 4 respectively.

   In addition, an order is a collection of combos. We'll encode an order as each digit representing a combo. For example, the order 237 represents a Double, Triple, and *biggie-sized* Triple.

   Assume that you have the following functions available:

   - `biggie_size` which when given a regular combo returns a *biggie-sized* version.
   - `unbiggie_size` which when given a *biggie-sized* combo returns a non-*biggie-sized* version.
   - `is_biggie_size` which when given a combo, returns `True` if the combo has been *biggie-sized* and `False` otherwise.
   - `combo_price` which takes a combo and returns the price of the combo.
   - `empty_order` which takes no arguments and returns an empty order which is represented by 0.
   - `add_to_order` which takes an order and a combo and returns a new order which contains the contents of the old order and the new combo. For example, `add_to_order(1,2) -> 12`.

   (a) Write a recursive function called `order_size` which takes an order and returns the number of combos in the order. For example, `order_size(237) -> 3`.

   Note: We assume that the input order is a valid one. The solutions for Recitation 1 already incorporate code that does this testing.

   ```
   def order_size(order):
       if order == 0:
           return 0
       else:
           return 1 + order_size(order // 10)
   ```

   (b) Write an iterative version of `order_size`.

   ```
   def order_size(order):
       count = 0;
       while order > 0:
           order = order // 10
           count = count + 1

       return count
   ```

   (c) Write a recursive function called `order_cost` which takes an order and returns the total cost of all the combos.

```
def order_cost(order):
    if order == 0:
        return 0.0
    else:
        return combo_price(order % 10) + order_cost(order // 10)
```

(d) Write an iterative version of `order_cost`.

```
def order_cost(order):
    cost = 0.0
    while order > 0:
        cost = cost + combo_price(order % 10)
        order = order // 10

    return cost
```

(e) **Homework:** Write a function called `add_orders` which takes two orders and returns a new order that is the combination of the two. For example, `add_orders(123,234) -> 123234`. Note that the order of the combos in the new order is not important as long as the new order contains the correct combos. `add_orders(123,234) -> 122334` would also be acceptable.

```
def add_orders(order_1, order_2):
    if order_2 == 0:
        return order_1
    else:
        return add_orders(order_1*10 + order_2%10, order_2//10)
```

2. Give order notation for the following:

(a) $5n^2 + n$

Answer: $O(n^2)$

(b) $\sqrt{n} + n$

Answer: $O(n)$

(c) $3^n n^2$

Answer: $O(3^n n^2)$

3.
```
def fact(n):
    if n == 0:
        return 1
    else:
        return n * fact(n - 1)
```

Running time: $O(n)$. Space: $O(n)$.

4. Write an iterative version of `fact`.

   Answer:

   ```
   def fact(n):
       product = 1

       for i in range(2, n+1):
           product *= i

       return product
   ```

   Running time: $O(n)$. Space: $O(1)$.

5. ```
   def find_e(n):
       if n == 0:
           return 1
       else:
           return 1/fact(n) + find_e(n - 1)
   ```

   Running time: $O(n^2)$. Space: $O(n)$. (Assume iterative `fact`)

6. Assume you have a function `is_divisible(n, x)` which returns `True` if `n` is divisible by `x`. It runs in $O(n)$ time and $O(1)$ space. Write a function `is_prime` which takes a number and returns `True` if it is prime and `False` otherwise.

   Answer:

   ```
   import math

   def is_prime(x):
       if x == 1:
           return False
       else:
           for i in range(2, int(math.sqrt(x) + 1)):
               if is_divisible(x, i):
                   return False

           return True
   ```

   Running time: $O(n^{\frac{3}{2}})$. Space: $O(1)$.

   Note: `import math` is needed in order to use `sqrt()`. Alternatively, `x**0.5` or `pow(x, 0.5)` can be used without importing `math` but `math.sqrt` is more efficient.

   In Python 3, `range()` returns an object that generates the numbers on demand. Hence the space complexity is O(1). However in Python 2, `range()` returns a list and the space complexity will be O(n) instead. `xrange()` in Python 2 is equivalent to `range()` in Python 3.
   ```

7. **Homework:** Write an iterative version of `find_e`.

```
def find_e(n):
    sum = 0

    for i in range(n+1):
        sum = sum + 1/fact(i)

    return sum
```

Running time: $O(n^2)$. Space: $O(1)$. (Assume iterative `fact`)