Database assignment 2

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1 Relational Algebra

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Exercize 1 -
\pi_{\text{mid}}\sigma_{\text{"title"}} = \text{"Not clickbait"} \land \text{"likecount"} > 9000 messages
\rho(invited, \pi_{\mathrm{id}}((\pi_{\mathrm{eid}}\sigma_{\mathrm{name}="Cheap \; \mathrm{sunglasses \; check \; description"}}event) \bowtie_{eid} invitedToEvent)) \bowtie_{\mathrm{id}} user)
\rho(users, \pi_{id}user)
\pi_{id}(invited \bowtie_{invited.id} \neq user.id \ users)
\rho(mes1, message)
\rho(mes2, message)
\rho(res, \pi_{\text{mes1.mid}}((message \times message) - \sigma_{\text{mes1.likecount} < \text{mes2.likecount}}(mes1 \times mes2)))
\pi_{\mathrm{id}}(\pi_{\mathrm{mid}}(res) \bowtie_{\mathrm{mid}} messageLikes) \bowtie_{\mathrm{id}} user
\rho(invited, \pi_{id}(\sigma_{accepted \ = \ "true"}(\pi_{eid}\sigma_{name = "Nude \ painting"}event) \bowtie_{eid} invitedToEvent)) \bowtie_{id} user)
\rho(users, \pi_{id}user)
\rho(friends, \pi_{id}(\pi_{fid}friend) \bowtie_{id} user)
\pi_{\text{name}}(invited \bowtie_{\text{invited.id}} = \text{user.id} \ users \cup friends)
\rho(friends, \pi_{id}(\pi_{fid}friend) \bowtie_{id} user)
\rho(othergender, \pi_{id}friends - \pi_{id}(\pi_{id}friends) \bowtie_{gender} user)
\rho(otherage, \pi_{id}(\pi_{id}friends) \bowtie_{friends.age < user.age} user)
\pi_{\text{name}}(friends/(othergender \cap otherage))user
\rho(one, \pi_{id}\sigma_{name} = "Crazy Cosplay Caroline" user)
\rho(onemessages, \pi_{mid}\rho_{creator=one}message)
\rho(messageliked, \pi_{id}\sigma_{mid=onemessages}messageLikes)
\pi_{\text{name}}((messageliked)/(\pi_{\text{id. name}}user))
\rho(mes1, \sigma_{likecount>999}message)
\rho(mes2, \sigma_{likecount>999}message)
\rho(mes3, \sigma_{likecount>999}message)
\rho(res1, \pi_{\text{mes1.creator}}(\sigma_{\text{mes1.mid} \neq \text{mes2.mid} \wedge \text{mes2.mid} \neq \text{mes3.mid} \wedge \text{mes1.mid} \neq \text{mes3.mid}(mes1 \times mes2 \times mes3)))
\rho(res2, \pi_{\text{mes2.creator}}(\sigma_{\text{mes1.mid}\neq \text{mes2.mid} \land \text{mes2.mid} \neq \text{mes3.mid} \land \text{mes1.mid} \neq \text{mes3.mid} (mes1 \times mes2 \times mes3)))
\rho(res3, \pi_{\text{mes3.creator}}(\sigma_{\text{mes1.mid}\neq \text{mes2.mid} \land \text{mes2.mid} \neq \text{mes3.mid} \land \text{mes1.mid} \neq \text{mes3.mid} (mes1 \times mes2 \times mes3)))
\pi_{\mathrm{id}}(\sigma_{\mathrm{res1.creator}=\mathrm{res2.creator}\wedge\mathrm{res2.creator}=\mathrm{res3.creator}\wedge\mathrm{res1.creator}=\mathrm{res3.creator}(res1 \times res2 \times res3)) \bowtie_{res.creator}=\mathrm{user.id}\ user
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2 Shema Normalization

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Exercize 2
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1.

Determine all the functional dependencies (F) that are derivable from the points mentioned above.

 $V \to TUCSD(A \text{ video ID implies all video attributes})$

 $UT \to V$ (Since an uploader cannot upload more than one video with the same title,

these attributes implie the video ID)

 $S \to C$ (Since subcategories are unique per category, there are no subcategories placed in more than one

category)

 $N \to BA$ (Since a comment number is unique, it implies all comment attributes)

 $N \to U$ (Since a comment number is unique and used by the same account used for uploading, it implies the uploader)

 $N \to V(Since every comment$ is linked via its number to a single video ID, the comment number implies the video ID)

2. Determine the key(s) in this table

 $N \to V, V \to TUCSD$

implies $N \to TUCSD(Transitivity)$

 $N \to BA, N \to U, N \to V, N \to TUCSD$

implies $N \to BAUVTUCSD(Union)$

Since N is never on the right side,

N is certainly part of all existing keys and since N is a key of itself,

N must be the only existing key.

3. Derive a minimal cover (G) for R

First, we listed all FD's:

 $\mathbf{V} \to T$

 $\mathbf{V} \to U$

 $\mathbf{V} \to C$

 $\mathbf{V} \to S$

 $V \to D$

 $\mathrm{UT} \to V$

 $\mathcal{S} \to C$

 ${\bf N} \to A$

 $\mathcal{N} \to B$

 $N \to U$

 ${\bf N} \to V$

Then we try to derive existing FD's from others:

N V, V U implies N U (Transitivity)

So, we delete N U from the list

Finally, we write down the minimal cover: $V \to TUCSDUT \to VS \to CN \to ABV$

4.

Using your minimal cover, derive a lossless join decomposition in BCNF and indicate a primary key for each table in the decomposition.

VTUCSDNBA

 $\Downarrow \Rightarrow NVBA$

VTUCSD

The resulting lossless join decomposition in BCNF is: VTUCSD, NVBA

5.

Is your decomposition also dependency preserving? Why is or isn't it? If it is not, derive a dependency preserving decomposition in 3NF.

The decomposition is not dependency preserving, because not all FD's in the minimal cover have their own table To make it dependency preserving, we add a table for all other FD's, the resulting decomposition is:

VTUCSD, NVBA, UTV, SC

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Prove, if $XY \to BandX \to YthenX \to B$

 $XY \rightarrow B$

 $X \to Y$

 $XX \rightarrow XY(AugmentationX)$

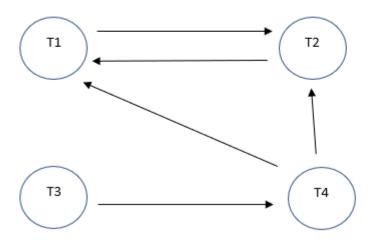
 $XX \rightarrow B(TransitivityXYB)$

 $X \rightarrow B(Union)$

Exercize 3 -

1.

T2 contains W(B) and T1 contains R(B), so there is an edge from T2 to T1 T3 contains R(A) and T4 contains W(A), so there is an edge from T3 to T4 T4 contains R(C) and T1 contains W(C), so there is an edge from T4 to T1 T1 contains R(B) and T2 contains W(B), so there is an edge from T1 to T2 T4 contains W(A) and T2 contains R(A), so there is an edge from T4 to T2



2. The graph is not conflict serializable because there exists a cycle inbetween T1 and T2.

3.															
T1:	R(C)					R(B)	W(C)	W(A)							Cmt.
T1:	S(C)					S(B)	X(C)	X(A)							Cmt.
T2:			R(B)		W(B)							W(B)	R(A)	W(C)	Cmt.
T2:			S(B)		X(B)							X(B)	S(A)	X(C)	Cmt.
T3:		R(A)		W(A)							R(A)				Cmt.
T3:		S(A)		X(A)							S(A)				Cmt.
T4:									R(C)	W(A)					Cmt.
T4:									S(C)	S(A)					Cmt.