

A Light Weight Approach to Minimize Charging Cost for Electric Bus Fleets

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Abstract—Insert abstract here

Index Terms—Insert keywords here

I. INTRODUCTION

Insert Introduction Here

II. FORMULATION

The cost objective that we desire to minimize is modeled after [1], which contains two primary elements: the cost of energy, and power demand. Energy is billed per kWh for on-peak and off-peak hours. The on-peak rate is more expensive because there is generally more demand for power during this time, whereas off-peak hours tend to be less expensive. The demand is covered in two separate chargers. The first is a facilities charge which is billed per kW for the highest 15-minute average power use over the course of the month. The second is a demand charge, which is also billed per kW, but is only billed for the highest 15-minute average power used during on-peak hours. The rates for each component are given in Table I.

Before we may compute the total monthly cost of electricity, we must define expressions for the average power and energy over time. Let each day be divided into 15-minute intervals for each bus where the average power expended for bus i during time j is denoted $p(i, j)$ as shown in Fig. 1. The resulting solution of the program we will develop will yield the average power expended by each bus during each period of time.

One constraint for which the solution must account is bus availability. When a bus is out of the station, the maximum average power for that time must be zero. For example, if bus 1 were out on route for t_5, t_6 , and t_7 , then the average power for those periods would be equal to zero as shown in Fig. 2. Let $b_{p(i, j)}$ be the average power used by bus i at time index j , and \mathbf{b} be a vector which contains $b_{p(i, j)}$ for each bus and time index. Also let $\mathcal{A} \subset i \times j$ be the set of all indices where bus i is in the station during time t_j and $\tilde{\mathcal{A}}$ be its complement. Furthermore, let p_{\max} be the maximum power that a charger can deliver.

We define a set of constraints so that buses do not use power when not in the station by letting

$$\begin{aligned} b_{p(i, j)} &= 0 \quad \forall i, j \in \tilde{\mathcal{A}} \\ b_{p(i, j)} &\leq p_{\max} \quad \forall i, j \in \mathcal{A} \\ -b_{p(i, j)} &\leq 0 \quad \forall i, j \in \mathcal{A} \end{aligned} \quad (1)$$

The constraints in Eq. 1 however to not account for buses that must charge for partial periods. For example, if the day

were divided into 15-minute time blocks, but a bus began charging at 10:07, then an average power of p_{\max} for that time slot would be inaccurate. Therefore, Eq. 1 must be modified so that the average power for each block correctly reflects partial availability. Let $\alpha(i, j)$ give the percentage of time that bus i is available during time j . Eq. 1 can be rewritten as

$$\begin{aligned} -b_{p(i, j)} &\leq 0 \quad \forall i, j \\ b_{p(i, j)} &\leq p_{\max} \cdot \alpha(i, j) \quad \forall i, j \end{aligned} \quad (2)$$

III. BATTERY

Each bus must also maintain its state of charge above acceptable levels throughout the day. When buses leave the station, each bus discharges some quantity of energy throughout the course of the route. Let $\delta(i, j)$ be the amount of charge lost by bus i at time j and let $h(i, j)$ be the state of charge of bus i at time j . The state of charge for each bus can be defined as

$$\begin{aligned} h(i, j) &= h(i, j-1) + p_b(i, j-1) \cdot \Delta T - \delta(i, j) \quad \forall i, j > 1 \\ h(i, 1) &= \eta_i \quad \forall i \end{aligned} \quad (3)$$

where η_i is the initial state of charge for bus i and ΔT is the difference in time between $t_{i, j}$ and $t_{i, j+1}$. Now that each value for the state of charge is defined, each value for h must be constrained so that it is greater than a given threshold, h_{\min} but does not exceed the maximum battery capacity h_{\max} . This yields

$$\begin{aligned} -h(i, j) &\leq -h_{\min} \quad \forall i, j \\ h(i, j) &\leq h_{\max} \quad \forall i, j. \end{aligned} \quad (4)$$

The final battery related constraint has to do with how we are planning for the bus. The expenses that come from power are computed monthly, but we desire to simulate the movements of the bus for only a day, and use this to extrapolate what the monthly cost may be. Therefore, the state of charge for a bus at the end of the day must reflect its starting value. This yields the following constraint:

$$h_{i, \text{end}} = h(i, 1) \quad \forall i$$

so that

$$h_{i, \text{end}} - h(i, 1) = 0 \quad \forall i \quad (5)$$

IV. CHARGER MANAGEMENT

Limited numbers of chargers is another limitation that many transit authorities face. Let the number of chargers be denoted n_{charger} . We desire to maintain the average cumulative power for each time step at a level that is serviceable given n_{charger} . We define a slack variable $p_c(j)$ which represents the total

TABLE I: Description of the billing structure

	On-Peak	Off-Peak	Facilities (Both)
Energy Rate	5.8282 ¢/kWh	2.9624 ¢/kWh	None
Energy Rate Symbol	μ_{e-on}	μ_{e-off}	None
Power Rate	\$ 15.73 /kW	None	\$ 4.81 /kW
Power Rate Symbol	μ_{p-on}	None	μ_{p-all}

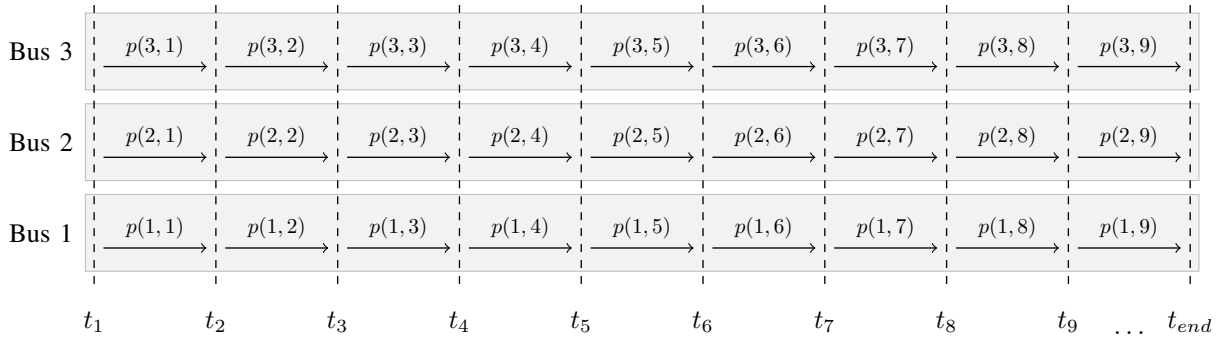


Fig. 1: Demonstrates how bus power use is conceptualized

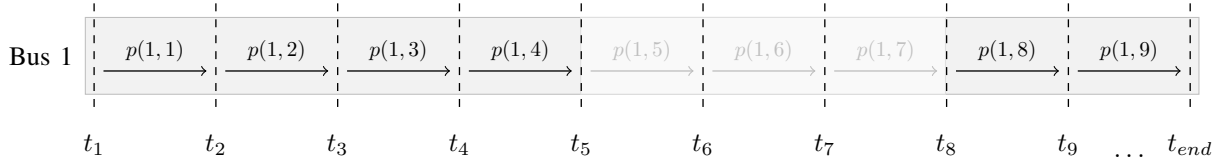


Fig. 2: Bus schedule with availability

average power consumed by all buses at time j . The variable $p_c(j)$ is computed as the sum of average bus powers so that

$$p_c(j) = \sum_i b_{p(i,j)}$$

or also as

$$p_c(j) - \sum_i b_{p(i,j)} = 0 \quad \forall j. \quad (6)$$

From a practical standpoint, we must also avoid multiple charging sessions in one visit to the station, or charger thrashing. We do this by minimising the difference in the average power values throughout the day. When a bus uses the charger for longer periods of time, the average power remains the same for multiple time periods. If a bus reconnects/disconnects multiple times, there will be larger difference in the power use from one time period to another. This can be expressed as

$$J_{\text{thrash}} = \sum_{i,j>1} |b_{p(i,j)} - b_{p(i,j-1)}|. \quad (7)$$

To minimize J_{thrash} , we define a slack variable $g_{i,j}$ so that

$$g_{i,j} = |b_{p(i,j)} - b_{p(i,j-1)}|.$$

which can be enforced by the linear constraints

$$\begin{aligned} -g_{i,j} &\leq b_{p(i,j)} - b_{p(i,j-1)} \\ g_{i,j} &\geq b_{p(i,j)} - b_{p(i,j-1)}. \end{aligned}$$

Suppose $b_{p(i,j)} - b_{p(i,j-1)} = -5$, then we have

$$-g \leq -5 \leq g$$

which minimizes g when $g = 5$. A similar expression is minimized when g is positive. The expression for $g_{i,j}$ can alternatively be expressed in standard form as

$$\begin{aligned} -g_{i,j} - b_{p(i,j)} + b_{p(i,j-1)} &\leq 0 \\ -g_{i,j} + b_{p(i,j)} - b_{p(i,j-1)} &\leq 0. \end{aligned} \quad (8)$$

V. OBJECTIVE

Now that the relevant constraints have been addressed, we must work towards computing the total objective function. We do so by first computing the total average power for the complete system. This total power is comprised of power used by the buses, and power used by external sources such as lights, ice melt, electric trains, etc which we refer to as “uncontrolled loads”, where the average power for the uncontrolled loads at time step j is denoted $u(j)$. We compute the total power as the sum of power used by the buses, $p_c(j)$ and the power consumed by uncontrolled loads $u(j)$ so that the total power, denoted $p_t(j)$ is computed as

$$p_t(j) = p_c(j) + u(j)$$

or

$$p_t(j) - p_c(j) - u(j) = 0 \quad \forall j. \quad (9)$$

The next step is to compute the fifteen minute average power use for each time step, denoted p_{15} . We do this by letting

$$p_{15}(j) = \frac{1}{n} \sum_{l \in \{j_{15}\}} p_t(l)$$

where $\{j_{15}\}$ is the set of all indices 15 minutes prior to j and n is the cardinality of $\{j_{15}\}$. Next, note that the rate schedule requires both the maximum overall average power, denoted $p_{\text{facilities}}$, and the maximum average power during on-peak hours, or p_{demand} . Let \mathcal{S}_{on} be the set of time indices belonging to on-peak hours, and recall that the max over all average power values is greater than or equal to $p_{15}(j)$ for all j . We can express this constraint is

$$p_{\text{facilities}} \geq p_{15}(j) \quad \forall j$$

or alternatively as

$$p_{15}(j) - p_{\text{facilities}} \leq 0 \quad \forall j \quad (10)$$

Because $p_{\text{facilities}}$ will be used in the objective function, the value for $p_{\text{facilities}}$ will be minimised until it is equal to the largest value in p_{15} . Following a similar logic, we also define a set of constraints for the maximum average on-peak power, p_{demand} so that

$$p_{15}(i) - p_{\text{demand}} \leq 0 \quad \forall i \in \mathcal{S}_{\text{on}} \quad (11)$$

The next step in computing the objective function is to compute the total *energy* consumed during on and off-peak hours respectively. Let e_{on} be the total energy consumed during on-peak hours and e_{off} be the energy consumed during off-peak hours. We can compute energy as the product of average power and time. In our case, we compute this as

$$\begin{aligned} e_{\text{on}} &= \Delta T \cdot \sum_{i \in g} p_t(i) \\ e_{\text{off}} &= \Delta T \cdot \sum_{i \in \tilde{g}} p_t(i) \end{aligned} \quad (12)$$

where \tilde{g} is the complement of g . We can now compute the total monthly charge as

$$J_{\text{cost}} = \begin{bmatrix} e_{\text{on}} \\ e_{\text{off}} \\ p_{\text{facilities}} \\ p_{\text{demand}} \end{bmatrix}^T \begin{bmatrix} \mu_{e\text{-on}} \\ \mu_{e\text{-off}} \\ \mu_{p\text{-all}} \\ \mu_{p\text{-on}} \end{bmatrix} \quad (13)$$

Which yields a complete objective function of

$$J_{\text{all}} = J_{\text{cost}} + J_{\text{thrash}} \quad (14)$$

VI. SCALABILITY

Before the bus charge problem can be solved, we need to address how the problem will scale. Historically, when buses are routed to chargers it requires a program which will evaluate all possible conflicts and selects a contention-free schedule, which tends to increase on the order of $0(n^2)$ and is NP-hard==insert citation of knapsack problem==. Before the bus problem can scale, we need a method to separate buses into groups to manage the problem's scope.

The group-selection problem separates buses into n_{group} groups, where group m is allocated n_{charger}^m chargers and n_{bus}^m buses. Each group must have sufficient chargers to fill it's needs and prefers buses with dissimilar routes.

We know that the number of cross-terms in future problems will be reduced when each group has the same number of buses. Therefore, let n_{bus}^m be described as

$$\begin{aligned} n_{\text{bus}}^m &\geq \left\lfloor \frac{n_{\text{bus}}}{n_{\text{group}}} \right\rfloor \\ n_{\text{bus}}^m &\leq \left\lceil \frac{n_{\text{bus}}}{n_{\text{group}}} \right\rceil \end{aligned}$$

or alternatively as

$$\begin{aligned} -n_{\text{bus}}^m &\leq -\left\lfloor \frac{n_{\text{bus}}}{n_{\text{group}}} \right\rfloor \\ n_{\text{bus}}^m &\leq \left\lceil \frac{n_{\text{bus}}}{n_{\text{group}}} \right\rceil. \end{aligned} \quad (15)$$

The next set of constraints ensures that there are sufficient charging resources in each group. First, we define a slack variable which gives the total power used in group m at time step j as p_j^m . Recall, we also know the expected power use for each bus as this is a result of the linear program described in sections II through V as $b_{p(i,j)}$, which allows us to describe the total power for any one group as

$$p_{j,m} = \sum_m \beta_{i,m} b_{p(i,j)}$$

or alternatively as

$$p_{j,m} - \sum_m \beta_{i,m} b_{p(i,j)} = 0 \quad \forall i, j, m \quad (16)$$

where $\beta_{i,m}$ is a binary indicator variable that is one when bus i is in group m and serves as part of the solution to this problem.

Next, we know that the total load of each group must be less than or equal to the collective capability of that group's chargers, which can be expressed as

$$n_{\text{charger}}^m \cdot p_{\text{max}} \leq p_{j,m} \quad \forall j, m$$

or in standard form as

$$n_{\text{charger}}^m \cdot p_{\text{max}} - p_{j,m} \leq 0 \quad \forall j, m \quad (17)$$

so that the number of chargers is sufficient to charge the collective load of the group.

We also desire to group buses together who's routes are most orthogonal. If two buses contain no overlap, they will be easiest to schedule and the inner product of their schedules from the first program will be equal to zero. Let

$$w_{i,i'} = \mathbf{b}_i^T \mathbf{b}_{i'},$$

where \mathbf{b}_i is the charge schedule for bus i as computed in the first program. We desire to minimize the total cross terms $w_{i,i'}$ for all buses in the same group. Define a slack variable $v_{i,i',g}$ which is equal to $w_{i,i'}$ if buses i and i' are in the same group, and zero otherwise so that

$$\begin{cases} v_{i,i',g} = w_{i,i'} & \beta_{i,m} = 1, \beta_{i',m} = 1 \\ v_{i,i',g} = 0 & \text{otherwise} \end{cases}$$

which can also be expressed by letting

$$\begin{aligned} v_{i,i',g} &\leq w_{i,i'} \\ v_{i,i',g} &\geq w_{i,i'} - M(2 - \beta_{i,m} - \beta_{i',m}) \\ v_{i,i',g} &\leq 0 + M\beta_{i,m} \\ v_{i,i',g} &\leq 0 + M\beta_{i',m} \\ v_{i,i',g} &\geq 0 \end{aligned}$$

or alternatively as

$$\begin{aligned} v_{i,i',g} &\leq w_{i,i'} \\ -v_{i,i',g} + M\beta_{i,m} + M\beta_{i',m} &\leq -w_{i,i'} + 2M \\ v_{i,i',g} - M\beta_{i,m} &\leq 0 \\ v_{i,i',g} - M\beta_{i',m} &\leq 0 \\ -v_{i,i',g} &\leq 0. \end{aligned} \quad (18)$$

The final objective can then be expressed as

$$J_{\text{select}} = \sum_{i,i',g} v_{i,i',g}. \quad (19)$$

Because this objective is more of a preference than a requirement for the bus charging problem, it is not always necessary to solve completely. Experimental results have shown that increasing the MIP Gap to upwards of 50% can be sufficient for this problem and helps reduce the computational complexity of this step.

VII. CHARGE SCHEDULES

The results from the Linear Program defined in the previous section give us a general estimate of how much and when buses should charge, however we must still address two primary issues. The first is defining concrete continuous time start and stop times for each charge session and the second is limiting the charge sessions to a finite number of chargers. After solving the first program, we take the results and derive *preliminary* intervals and average power consumptions for each charge session.

For example, consider a solution to a three bus, two charger scenario given in Fig. 3. Note that there appears to be three buses charging at the same time from t_5 to t_6 even though there are only two chargers. We can reformulate this solution in terms of continuous start and stop variables and a variable charge rate so that the *duration* of each charge session may be relaxed. The objective is to store the given energy in the corresponding bus within the given charge interval.

Note how few of the charge sessions utilize the chargers to full capacity. This implies that there exists a smaller charge window in which equivalent power can be delivered. This allow us to use the charge durations from the solution from Fig. 3 as bounds on *allowable* charge windows instead of absolute truth.

An example of how Fig. 3 may be reformulated is given in Fig. 5. Note how the actual charge sessions don't necessarily need to take up all the time they were initially allocated in the first solution and that these times can fluctuate if the average charge rate is less than the maximum charger capacity. In this example, we assume a maximum charge capacity of 350kW.

Note how the third charge session does have to be exactly where it was scheduled because the average is equal to the maximum charge rate. If we exam in just the schedule for Bus 1, we note that there are four essential variables for the corresponding charge session: $a(i, r)$, $b(i, r)$, $f(i, r)$ and $d(i, r)$ which represent the minimum start time, actual start time, actual end time, and maximum end time respectively.

The problem we must now solve is one of arranging these "rectangles" such that each one is larger than it's minimum width (or charge time). We must also account for the number of chargers. It can be helpful to view the problem as a bin packing problem, where each session must fit within the "swim lane" of a charger. For example, taking the charge sessions given in Fig. 5 and arranging them so that there is no overlap between sessions will yield a valid solution as shown in Fig. 4.

From Fig. 6, we know that $a(i, r)$, $b(i, r)$, $f(i, r)$ and $d(i, r)$ must be such that

$$\begin{aligned} a(i, r) &\leq b(i, r) \\ b(i, r) &\leq f(i, r) \\ f(i, r) &\leq d(i, r). \end{aligned}$$

or alternatively as

$$\begin{aligned} -b(i, r) &\leq -a(i, r) \\ b(i, r) - f(i, r) &\leq 0 \\ f(i, r) &\leq d(i, r) \end{aligned} \quad (20)$$

Where $a(i, r)$ and $d(i, r)$ are known from the previous optimization problem, and $b(i, r)$ and $f(i, r)$ are optimization variables.

We must differentiate between chargers and so, define $\sigma_{i,r,k}$ as a binary selector variable which is one if charger k services bus i for session r and zero otherwise. We know that only one charger can charge each bus at a time. We also know that each charge session *must* be serviced, which implies that

$$\sum_k \sigma_{i,r,k} = 1 \quad \forall i, r. \quad (21)$$

Next, we also know that during each session a certain amount of energy must be transferred from the charger to the battery. The amount of energy that must be transferred to bus i during session r can be computed from the results of the first linear program and is denoted $e(i, r)$. We can compute a minimum time window from this value as

$$w(i, r)_{\min} = \frac{e(i, r)}{p_{\max}}.$$

Because this is the minimum time window, we must ensure that the difference between the start and stop times is at least this large so that

$$f(i, r) - b(i, r) \geq w(i, r) \quad \forall i, r$$

or alternatively,

$$b(i, r) - f(i, r) \leq -w(i, r) \quad \forall i, r. \quad (22)$$

The final set of constraints deals with contention so that no charger can be scheduled for two sessions that overlap. let $\mathcal{L} = \{(i, r) \times (i', r')\}$ where charge sessions i, r and i', r'

Bus 3	0	0	0	350	350	350	350	0	0
Bus 2	175	175	175	175	70	0	0	0	0
Bus 1	0	35	105	105	140	0	0	0	0
	t_1	t_2	t_3	t_4	t_5	t_6	t_7	t_8	$t_9 \dots t_{end}$

Fig. 3: An example solution to a 3-bus, 2-charger scenario from the first QP

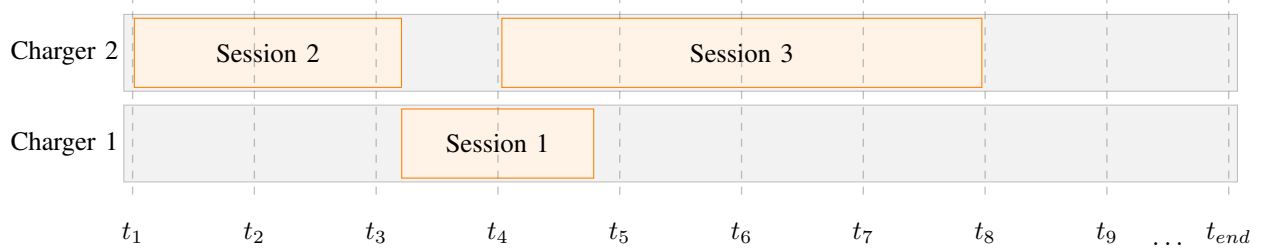


Fig. 4: Demonstrates the solution to the second program

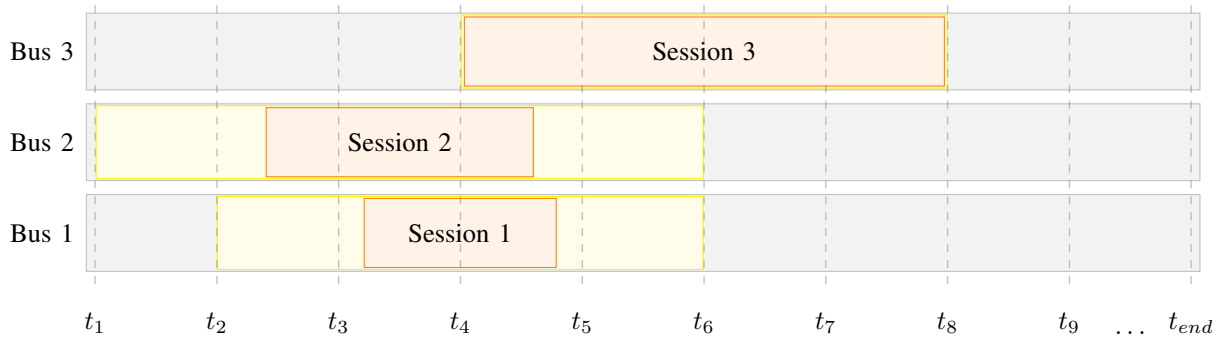


Fig. 5: Demonstrates how the problem from the first QP can be reformulated in terms of continuous variables

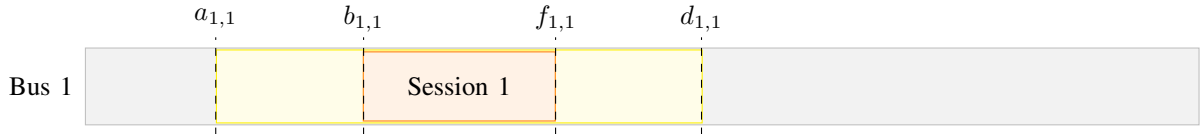


Fig. 6: Gives variables of optimization for the second program

have the potential to overlap. Before we can prevent overlap, we must define a binary variable $l_{(i,r),(i',r')}$ which is equal to one when session i, r is scheduled before session i', r' and zero otherwise so that

$$\begin{cases} f \leq b' & l = 1 \\ f' \leq b & l = 0 \end{cases}$$

Here we can expand this thought through use of the “big-M” technique. Let M be large. In this case, we can set it equal to the number of seconds in a day. We know what the top constraint must be trivially satisfied when $l = 0$ and the

bottom must also when $l = 1$. This leads to a reformulation so that

$$\begin{aligned} f - b' &\leq M(1 - l) \\ f' - b &\leq lM \end{aligned}$$

However, this constraint *only* needs to hold when sessions i, r and i', r' are scheduled to charge on the same charger or that $\sigma_{i,r,k} = \sigma_{i',r',k} = 1$. We can reformulate the above constraint to satisfy this condition by letting

$$\begin{aligned} f - b' &\leq M(3 - \sigma_{i,r,k} - \sigma_{i',r',k} - l) \\ f' - b &\leq M(2 - \sigma_{i,r,k} - \sigma_{i',r',k} + l) \end{aligned}$$

which can also be expressed as

$$\begin{aligned} f - b' + M\sigma_{i,r,k} + M\sigma_{i',r',k} + Ml_{(i,r),(i',r')} &\leq 3M \\ f' - b + M\sigma_{i,r,k} + M\sigma_{i',r',k} - Ml_{(i,r),(i',r')} &\leq 2M. \end{aligned} \quad (23)$$

Finally, we desire the power profile to closely match the profile given in the first linear program, which would occur if each charge session matched the durations given in the first solution. Because these durations form an upper bound on how long a bus *could* charge, we essentially desire to maximize the collective duration of all the charge sessions. This can be written in linear form as

$$\min_{f,b} \sum_{i,r} b(i,r) - f(i,r) \quad (24)$$

and forms the objective over which we maximize.

VIII. MULTI-RATE CHARGING

Discuss final program here

ACKNOWLEDGMENT

This material is based in part upon work supported by the National Science Foundation through the ASPIRE Engineering Research Center under Grant No. EEC-1941524, the Department of Energy through a prime award with ABB under Grant No. DE-EE0009194, and PacifiCorp under contract number 3590. Any opinions, findings, and conclusions or recommendations expressed in this material are those of the authors and do not necessarily reflect the views of the National Science Foundation, the Department of Energy, or Pacificorp.

Variable	Description	Range	Variable	Description	Range
Indices					
i	Bus index	\mathbb{N}	j	Time Index	\mathbb{N}
k	Charger index	\mathbb{N}	r	Route Index	
Formulation					
n_{bus}	The number of buses in the optimization framework.	\mathbb{Z}	n_{time}	The number of time indices in a day.	\mathbb{Z}^+
$b_{p(i,j)}$	The average power consumed by bus i during time period j .	\mathbb{R}	t_j	The time at time index j . This paper also refers to the period of time from t_j to t_{j+1} as “period t_j ”.	\mathbb{R}
\mathbf{b}	A vector containing each value for $b_{p(i,j)}$.	$\mathbb{R}^{n_{\text{bus}} \times n_{\text{time}}}$	$\alpha(i,j)$	the percentage of time bus i spends in the station from t_j to t_{j+1} .	$[0, 1]$
\mathcal{A}	The set of all $i \times j$ elements where bus i can charge at time index j	$i \times j$	p_{max}	The maximum power a bus charger can deliver to a bus in kW. This paper assumes a value of 350 for most examples and results.	\mathbb{R}^+
$\tilde{\mathcal{A}}$	The complement of \mathcal{A} .	$i \times j$			
Battery					
h_{min}	The minimum allowable state of charge	$(0, h_{\text{max}})$	h_{max}	The maximum state of charge	\mathbb{R}^+
η_i	The beginning state of charge for bus i	$(h_{\text{min}}, h_{\text{max}})$	$h(ij)$	The state of charge for bus i at time t_j .	$(h_{\text{min}}, h_{\text{max}})$
ΔT	The change in time from t_j to t_{j+1}	\mathbb{R}^+	\mathbf{h}	A vector containing all state of charge values.	$\mathbb{R}_+^{n_{\text{bus}} \times n_{\text{time}}}$
$\delta(ij)$	The battery discharge for bus i during time period j .	\mathbb{R}_+	$h(i, \text{end})$	Bus i 's final state of charge.	$(h_{\text{min}}, h_{\text{max}})$

Charger Management

n_{charger}	The time index for the start of bus i 's j^{th} stop	\mathbb{Z}_+	$p_c(j)$	The average power consumed by all buses during time period j .	\mathbb{R}
p_c	A vector containing all values of $p_c(j)$.	$\mathbb{R}_+^{n_{\text{time}}}$	J_{thrash}	A secondary objective function which penalizes multiple plug-in instances per charge session.	\mathbb{R}_+

Objective

$\mu_{\text{e-on}}$	On-Peak Energy Rate	\mathbb{R}_+	$\mu_{\text{e-off}}$	Off-Peak Energy Rate	\mathbb{R}_+
$\mu_{\text{p-on}}$	On-Peak Demand Power Rate	\mathbb{R}_+	$\mu_{\text{p-all}}$	Facilities Power Rate	\mathbb{R}_+
\mathcal{S}_{on}	The set of on-peak time indices	$\{1, \dots, n_{\text{time}}\}$	p_{demand}	Maximum average power during on-peak periods	\mathbb{R}
$p_{\text{facilities}}$	Maximum average power over all time instances.	\mathbb{R}_+	$p_t(j)$	The total average power consumed by both the bus chargers and the uncontrolled loads.	$\mathbb{R}_+^{n_{\text{time}}}$
$u(j)$	The average power over time j consumed by the uncontrolled loads	$\mathbb{R}_+^{n_{\text{time}}}$	p_t	a vector containing $p_t(i)$ for all i .	$\mathbb{R}_+^{n_{\text{time}}}$
e_{on}	The total amount of energy consumed by the bus chargers and uncontrolled loads during off-peak hours.	\mathbb{R}_+	e_{off}	The total energy consumed by the bus chargers and uncontrolled loads during on-peak hours.	\mathbb{R}_+
J_{cost}	The section of the objective function pertaining to the fiscal expense of charging buses.	\mathbb{R}	J_{all}	The expression for the complete objective function.	\mathbb{R}

Charge Schedules

$a(i, r)$	The beginning of the allowable charge interval for bus i 's r^{th} charge session.	\mathbb{R}_+	$b(i, r)$	The commanded start time for bus i 's r^{th} charge session	\mathbb{N}
$f(i, r)$	The commanded end time for bus i 's r^{th} charge session.	\mathbb{R}_+	$d(i, r)$	The end time of the allowable charge interval for bus i 's r^{th} charge session.	\mathbb{R}_+
$\sigma(i, r, k)$	A selector variable which is one when bus i charges at charger k for session r .	$\{0, 1\}$	M	The number of seconds in a day	\mathbb{Z}_+

A selector variable which is one when
 $l_{(i,r),(i',r')}$ bus i charges before bus i' during the $(i,r) \times (i',r')$
 r and r' sessions respectively.

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