

# NOTES ON FFT FOR FINITE FIELDS

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What follows is an overview of Pollard's definition, as found in [1], of the Fourier transform for finite fields.

Let  $GF(p^n)$ , or  $F$  for short, be the Galois field with  $p^n$  elements, where  $p$  is a prime number and  $n$  is a positive integer. Let  $d$  be a divisor of  $p^n - 1$ , and let  $r$  be an element of order  $d$  in  $F^*$ . Then, if  $(a_i)$  is a sequence in  $F$  of length  $d$ , we define its *Fourier transform* via the rule

$$(1) \quad A_i = \sum_{j=0}^{d-1} a_j r^{ij}.$$

It has the following “convolution property”: if the sequences  $(a_i)$ ,  $(b_i)$  and  $(c_i)$  are such that their transforms  $(A_i)$ ,  $(B_i)$  and  $(C_i)$  satisfy

$$(2) \quad C_i = A_i B_i$$

then

$$(3) \quad c_i = \sum_{j=0}^{d-1} a_j b_{i-j},$$

where the indices for the terms  $b_t$  are taken modulo  $d$ .

In equation (3) we see  $(c_i)$  as the “convolution” of  $(a_i)$  and  $(b_i)$ , which allows us to rephrase the relationship between equations (2) and (3) as follows:

**Lemma 1** (Informal convolution statement). *To compute the convolution of two sequences, one may first transform them, then compute their point-wise product, and then apply the reverse transform.*

## 2. COMMENTS

**2.1.** Prove that if  $F = GF(p^n)$  then  $F^*$  is cyclic of order  $p^n - 1$ .

**2.2.** FFT, both in number fields and in  $\mathbf{C}$ , has complexity  $O(d \log d)$  versus the complexity  $O(d^2)$  of the naive way to compute.

## REFERENCES

- [1] POLLARD, J. M. The fast Fourier transform in a finite field. *Math. Comput.* 25 (1971), 365–374.

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